

MATH280

Sammendrag av det viktige.

Markov Chains

A stochastic model describing a sequence of possible events in which the probability of each event depends only on the state attained in the previous event. Here the matrix A is the probability of going from one state to another

$$\mathbf{x}_{k+1} = A\mathbf{x}_k$$

Example

The entries of the matrix A is the probabilities of going from one state to another (migration matrix):

$$A = \begin{array}{cc} & \begin{array}{cc} \text{From} \\ \text{City} & \text{Suburb} \end{array} \\ \begin{bmatrix} .95 & .03 \\ .05 & .97 \end{bmatrix} & \begin{array}{cc} \text{City} \\ \text{Suburb} \end{array} \end{array} \quad \begin{array}{c} \text{To} \end{array}$$

If we have a population vector, we can predict the coming population:

$$A\mathbf{x}_k = \mathbf{x}_{k+1} = \begin{bmatrix} .95 & .03 \\ .05 & .97 \end{bmatrix} \begin{bmatrix} 600,000 \\ 400,000 \end{bmatrix} = \begin{bmatrix} 582,000 \\ 418,000 \end{bmatrix}$$

PageRank is a specific case of Markov Chains.

Power method

For an $n \times n$ matrix A with an eigenvalue λ_1 bigger than all the others, then

$$\frac{1}{\lambda_1^k} A^k \mathbf{x} \rightarrow c_1 \mathbf{v}_1$$

where \mathbf{v}_1 is the eigenvector of λ_1 and c_1 is a scalar.

For a large k , a scalar multiple of $A^k \mathbf{x}$ determines the same direction as the eigenvector $c_1 \mathbf{v}_1$.

This means that $A^k \mathbf{x}$ approaches the eigenspace spanned by \mathbf{v}_1 (as $k \rightarrow \infty$).

We do not know λ_1 , therefore we have must add the constant c_1 .

If we scale each $A^k \mathbf{x}$ so that its largest entry is a 1, then \mathbf{x}_k converges to a multiple of \mathbf{v}_1 whose largest entry is 1.

When \mathbf{x}_k is close to an eigenvector for λ_1 then $A\mathbf{x}_k$ is close to $\lambda_1 \mathbf{x}_k$

1. Initial vector \mathbf{x}_0 whose largest entry is 1.
2. for $k = 0, 1, \dots$:
 - a. Compute $A\mathbf{x}_k$.
 - b. Let μ_k be the largest (absolute) value in $A\mathbf{x}_k$.
 - c. Compute $\mathbf{x}_{k+1} = \frac{1}{\mu_k} A\mathbf{x}_k$.
3. For almost all choices of \mathbf{x}_0 , μ_k approaches the dominant eigenvalue and \mathbf{x}_k the eigenvector.

Jacobi- and Gauss-Seidel iterative methods

We have

$$A\mathbf{x} = \mathbf{b}$$

where A is very large, and has a lot of zeros in it. We want to solve for \mathbf{x} , but doing Gaussian eliminations is expensive.

Given that it has a unique solution \mathbf{x}^* and the diagonal is nonzero.

1. Initial matrix A whose diagonal entries are nonzero.
2. Write $A\mathbf{x}$ as a set of equations.
3. Solve each equation for its own x_i .
4. Guess a value for a "random" \mathbf{x} : for instance x_1 , so that the equations has a solution.
5. Loop until difference in \mathbf{x} is within a tolerance:
 - a. Insert the "guessed" \mathbf{x} , and solve the set of equations.
 - b. The solution will yield a new \mathbf{x} .
 - c. Input this new \mathbf{x} into the set of equations.
 - d. Check the difference in the new and old \mathbf{x} .
 - e. Repeat until the difference is below a specified tolerance.

Bases and coordinates

Let $\mathcal{B} = \{\mathbf{b}_1, \dots, \mathbf{b}_n\}$ be a basis for V . Then each \mathbf{x} in V can be expressed as a linear combination

$$\mathbf{x} = c_1\mathbf{b}_1 + \dots + c_n\mathbf{b}_n$$

The coordinates of \mathbf{x} relative to \mathcal{B} (or the \mathcal{B} -coordinates of \mathbf{x}), $[\mathbf{x}]_{\mathcal{B}}$, are the weights c_1, \dots, c_n such that $\mathbf{x} = c_1\mathbf{b}_1 + \dots + c_n\mathbf{b}_n$.

$$[\mathbf{x}]_{\mathcal{B}} = \begin{bmatrix} c_1 \\ \vdots \\ c_n \end{bmatrix}$$

Change of basis

Let $\mathcal{B} = \{\mathbf{b}_1, \dots, \mathbf{b}_n\}$ and $\mathcal{C} = \{\mathbf{c}_1, \dots, \mathbf{c}_n\}$ be bases of a vector space V . Then there is a unique $n \times n$ matrix ${}_{\mathcal{C} \leftarrow \mathcal{B}}P$ such that

$$[\mathbf{x}]_{\mathcal{C}} = {}_{\mathcal{C} \leftarrow \mathcal{B}}P [\mathbf{x}]_{\mathcal{B}}$$

The columns of ${}_{\mathcal{C} \leftarrow \mathcal{B}}P$ are the \mathcal{C} -coordinate vectors of the vectors in the basis \mathcal{B} :

$${}_{\mathcal{C} \leftarrow \mathcal{B}}P = [[\mathbf{b}_1]_{\mathcal{C}} \ \dots \ [\mathbf{b}_n]_{\mathcal{C}}].$$

The matrix is the change-of-coordinates matrix from \mathcal{B} to \mathcal{C} . That means that multiplying by ${}_{\mathcal{C} \leftarrow \mathcal{B}}P$

converts \mathcal{B} -coordinates into \mathcal{C} -coordinates. It also applies that $(\begin{smallmatrix} P \\ \mathcal{C} \leftarrow \mathcal{B} \end{smallmatrix})^{-1} = \begin{smallmatrix} P \\ \mathcal{B} \leftarrow \mathcal{C} \end{smallmatrix}$.

$$[\mathbf{c}_1 \ \mathbf{c}_2 \vdots \mathbf{b}_1 \ \mathbf{b}_2] \rightarrow [I \vdots \begin{smallmatrix} P \\ \mathcal{C} \leftarrow \mathcal{B} \end{smallmatrix}]$$

Linear transformations

When multiplying a vector \mathbf{x} with a matrix A , we get a new vector \mathbf{b} .

$$A\mathbf{x} = \mathbf{b} \quad \Leftrightarrow \quad x_1\mathbf{a}_1 + \cdots + x_n\mathbf{a}_n = \mathbf{b}$$

This is equivalent to saying that multiplication by A transforms \mathbf{x} into \mathbf{b} .

For each \mathbf{x} in \mathbb{R}^n , $T(\mathbf{x})$ is computed as $A\mathbf{x}$ where $A \in \mathbb{R}^{m \times n}$. This can be denoted by $\mathbf{x} \mapsto A\mathbf{x}$ (\mathbf{x} maps to $A\mathbf{x}$).

The range of T is the same as $\text{Col}(A)$.

A transformation (or mapping) T is *linear* if:

1. $T(\mathbf{u} + \mathbf{v}) = T(\mathbf{u}) + T(\mathbf{v})$ for all \mathbf{u}, \mathbf{v} in the domain of T ($\text{Col}(A)$).
2. $T(c\mathbf{u}) = cT(\mathbf{u})$ for all scalars c and all \mathbf{u} in the domain of T ($\text{Col}(A)$).

Projections

Normality and orthogonality

The norm of a vector is

$$\|\mathbf{v}\| = \sqrt{\mathbf{v} \cdot \mathbf{v}} \quad \& \quad \|\mathbf{v}\|^2 = \mathbf{v} \cdot \mathbf{v}$$

The distance between two vectors are

$$\text{dist}(\mathbf{v}, \mathbf{u}) = \|\mathbf{v} - \mathbf{u}\| = \sqrt{(v_1 - u_1)^2 + \cdots + (v_n - u_n)^2}$$

The orthogonal compliment is the set of all vectors orthogonal to a subspace.

$$\text{Row}(A)^\perp = \text{Nul}(A) \quad \& \quad \text{Col}(A)^\perp = \text{Nul}(A^T)$$

Proof

The definition of the nullspace $\text{Nul}(A)$ is the set of vectors \mathbf{x} such that $A\mathbf{x} = \mathbf{0}$. This is saying that the dot-product of these \mathbf{x} 's and rows are all zero, and is therefore orthogonal.

Orthogonal projections

Let \mathbf{y} be a vector you want to project down to a subspace \mathcal{W} with orthogonal basis $\mathcal{B} = \{\mathbf{u}_1, \dots, \mathbf{u}_p\}$, then

$$\hat{\mathbf{y}} = \text{Proj}_{\mathcal{W}} \mathbf{y} = \frac{\mathbf{y} \cdot \mathbf{u}_1}{\mathbf{u}_1 \cdot \mathbf{u}_1} \mathbf{u}_1 + \dots + \frac{\mathbf{y} \cdot \mathbf{u}_p}{\mathbf{u}_p \cdot \mathbf{u}_p} \mathbf{u}_p$$

Orthonormal bases

If $\mathbf{U} = \{\mathbf{u}_1, \dots, \mathbf{u}_p\}$ form an orthonormal basis, then

$$\hat{\mathbf{y}} = \text{Proj}_{\mathcal{W}} \mathbf{y} = \mathbf{U} \mathbf{U}^T \mathbf{y}$$

Gram-Schmidt process

Given a basis $\{\mathbf{x}_1, \dots, \mathbf{x}_p\}$ for a nonzero subspace \mathcal{W} of \mathbb{R}^n , then one can form an orthogonal basis for the same subspace through these steps:

$$\begin{aligned} \mathbf{v}_1 &= \mathbf{x}_1 \\ \mathbf{v}_2 &= \mathbf{x}_2 - \frac{\mathbf{x}_2 \cdot \mathbf{v}_1}{\mathbf{v}_1 \cdot \mathbf{v}_1} \mathbf{v}_1 \\ \mathbf{v}_3 &= \mathbf{x}_3 - \frac{\mathbf{x}_3 \cdot \mathbf{v}_1}{\mathbf{v}_1 \cdot \mathbf{v}_1} \mathbf{v}_1 - \frac{\mathbf{x}_3 \cdot \mathbf{v}_2}{\mathbf{v}_2 \cdot \mathbf{v}_2} \mathbf{v}_2 \\ &\vdots \\ \mathbf{v}_p &= \mathbf{x}_p - \frac{\mathbf{x}_p \cdot \mathbf{v}_1}{\mathbf{v}_1 \cdot \mathbf{v}_1} \mathbf{v}_1 - \frac{\mathbf{x}_p \cdot \mathbf{v}_2}{\mathbf{v}_2 \cdot \mathbf{v}_2} \mathbf{v}_2 - \dots - \frac{\mathbf{x}_p \cdot \mathbf{v}_{p-1}}{\mathbf{v}_{p-1} \cdot \mathbf{v}_{p-1}} \mathbf{v}_{p-1} \end{aligned}$$

By doing this one creates an orthogonal basis $\{\mathbf{v}_1, \dots, \mathbf{v}_p\}$ for \mathcal{W} .

QR-factorisation

If $A \in \mathbb{R}^{m \times n}$ with linearly independent columns, then A can be factored as $A = QR$ where $Q \in \mathbb{R}^{m \times n}$ whose columns form an orthonormal basis for $\text{Col}(A)$ and $R \in \mathbb{R}^{n \times n}$ upper triangular invertible matrix with positive diagonal entries.

The columns of Q can be found through the Gram-Schmidt process, so that

$$\text{Span}\{\mathbf{x}_1, \dots, \mathbf{x}_k\} = \text{Span}\{\mathbf{u}_1, \dots, \mathbf{u}_k\} \quad \Rightarrow \quad Q = [\mathbf{u}_1 \dots \mathbf{u}_n]$$

This means that \mathbf{x}_k can be written as a linear combinations of \mathbf{U} :

$$\mathbf{x}_k = r_{1,k} \mathbf{u}_1 + \dots + r_{k,k} \mathbf{u}_k + 0 \mathbf{u}_{k+1} + \dots + 0 \mathbf{u}_n$$

This shows that \mathbf{x}_k is a linear combination of the columns of Q using as weights the entries in the vector

$$\mathbf{r}_k = \begin{bmatrix} r_{1,k} \\ \vdots \\ r_{k,k} \\ 0 \\ \vdots \\ 0 \end{bmatrix} \quad \& \quad \mathbf{x}_k = Q \mathbf{r}_k \quad \& \quad R = [\mathbf{r}_1 \dots \mathbf{r}_n] = Q^T A$$

In order to find R one can use the fact that Q form an orthonormal set:

$$R = Q^T A \quad \Leftarrow \quad Q^T A = Q^T (QR) = IR = R \quad \Leftarrow \quad Q^T Q = I$$

Least-squares

If $A \in \mathbb{R}^{m \times n}$ and $\mathbf{b} \in \mathbb{R}^m$, a least squares solution of $A\mathbf{x} = \mathbf{b}$ is an $\hat{\mathbf{x}} \in \mathbb{R}^n$ which minimizes $\|\mathbf{b} - A\hat{\mathbf{x}}\|$.

The solution to this is to approximate \mathbf{b} onto $\text{Col}(A)$

$$\hat{\mathbf{b}} = \text{proj}_{\text{Col}(A)} \mathbf{b}$$

Which makes us able to solve for the best approximation of \mathbf{x}

$$A\hat{\mathbf{x}} = \hat{\mathbf{b}}$$

Proof

Since $\mathbf{b} - \hat{\mathbf{b}}$ is orthogonal to $\text{Col}(A)$, then $\mathbf{b} - A\hat{\mathbf{x}}$ is orthogonal to each column of A. If \mathbf{a}_j is a column of A, then $\mathbf{a}_j \cdot (\mathbf{b} - A\hat{\mathbf{x}}) = 0$.

$$\begin{aligned} A^T(\mathbf{b} - A\hat{\mathbf{x}}) &= \mathbf{0} \\ A^T\mathbf{b} - A^T A\hat{\mathbf{x}} &= \mathbf{0} \\ A^T A\hat{\mathbf{x}} &= A^T\mathbf{b} \end{aligned}$$

Which show that the set of least-squares solutions are found through the normal equations

$$A^T A\mathbf{x} = A^T\mathbf{b}$$

Spectral theorem

Properties of a $n \times n$ symmetric matrix A:

1. A has n real eigenvalues, counting duplicates.
2. The dimension of the eigenspace for eigenvalue λ is equal to the number of eigenvalues that has the same value.
3. The eigenspace for eigenvalues with different eigenvalues are orthogonal. **NOTE:** This does not apply for eigenvectors with the same eigenvalues. We can find a orthogonal space using the Gram-Schmidt method.
4. A is orthogonally diagonalizable.

Given a symmetric matrix A, we know it is orthogonally diagonalizable. And because eigenvectors will be orthogonal (if all eigenvalues are distinct) we can normalize (ending up with orthonormality). This gives us $\mathbf{A} = \mathbf{PDP}^{-1}$, where the columns of P are the normalized eigenvectors. This results in

$\mathbf{P}^{-1} = \mathbf{P}^T$, which gives:

$$\mathbf{A} = \mathbf{PDP}^{-1} = \underbrace{\begin{bmatrix} \uparrow & & \uparrow \\ \mathbf{u}_1 & \cdots & \mathbf{u}_n \\ \downarrow & & \downarrow \end{bmatrix}}_{\text{Orthonormal eigenvectors}} \underbrace{\begin{bmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_n \end{bmatrix}}_{\text{Eigenvalues}} \underbrace{\begin{bmatrix} \leftarrow & \mathbf{u}_1^T & \rightarrow \\ & \vdots & \\ \leftarrow & \mathbf{u}_n^T & \rightarrow \end{bmatrix}}_{\text{Transposed orthonormal eigenvectors}} = \begin{bmatrix} \uparrow & & \uparrow \\ \lambda_1 \mathbf{u}_1 & \cdots & \lambda_n \mathbf{u}_n \\ \downarrow & & \downarrow \end{bmatrix} \begin{bmatrix} \leftarrow & \mathbf{u}_1^T & \rightarrow \\ & \vdots & \\ \leftarrow & \mathbf{u}_n^T & \rightarrow \end{bmatrix}$$

This can be written as:

$$\mathbf{A} = \lambda_1 \mathbf{u}_1 \mathbf{u}_1^T + \cdots + \lambda_n \mathbf{u}_n \mathbf{u}_n^T$$

Which is called a *spectral decomposition* of \mathbf{A} because it breaks up \mathbf{A} into pieces determined by the spectrum (eigenvalues) of \mathbf{A} . This means that each matrix $\mathbf{u}_i \mathbf{u}_i^T$ is the projection matrix of a \mathbf{x} onto the subspace spanned by \mathbf{u}_i .

Singular Value Decomposition (SVD)

If we have a matrix $\mathbf{A} \in \mathbb{R}^{m \times n}$ with rank r . There is a matrix $\Sigma \in \mathbb{R}^{m \times n}$ where the first r diagonal entries are the singular values ($\sigma_1 \geq \cdots \geq \sigma_r > 0$) of \mathbf{A} .

$$\Sigma = \begin{bmatrix} D & 0 \\ 0 & 0 \end{bmatrix} = \begin{bmatrix} \sigma_1 & & 0 & 0 \\ & \ddots & & 0 \\ 0 & & \sigma_r & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Σ has zeros filled in for $m - r$ rows and $n - r$ columns.

There also exist two orthogonal matrices in $\mathbb{R}^{m \times m}$ (\mathbf{U}) and $\mathbb{R}^{n \times n}$ (\mathbf{V}) such that

$$\mathbf{A} = \mathbf{U}\Sigma\mathbf{V}^T = \underbrace{\begin{bmatrix} \uparrow & & \uparrow \\ \mathbf{u}_1 & \cdots & \mathbf{u}_r \\ \downarrow & & \downarrow \end{bmatrix}}_{\substack{\text{Eigenvectors of} \\ \mathbf{A}\mathbf{A}^T}} \underbrace{\begin{bmatrix} \sigma_1 & & 0 & 0 \\ & \ddots & & 0 \\ 0 & & \sigma_r & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}}_{\substack{\text{Singular values } (\sqrt{\lambda_n}, \\ \text{with } \lambda_n \text{ eigenvalues} \\ \text{of } \mathbf{A}\mathbf{A}^T \text{ and } \mathbf{A}^T\mathbf{A})}} \underbrace{\begin{bmatrix} \leftarrow & \mathbf{v}_1^T & \rightarrow \\ & \vdots & \\ \leftarrow & \mathbf{v}_r^T & \rightarrow \end{bmatrix}}_{\text{Eigenvectors of } \mathbf{A}^T\mathbf{A}}$$

$$\mathbf{U} = \begin{array}{ccccccc} \uparrow & & \uparrow & \uparrow & & \uparrow \\ \mathbf{u}_1 & \cdots & \mathbf{u}_r & \mathbf{u}_{r+1} & \cdots & \mathbf{u}_n \\ \downarrow & & \downarrow & \downarrow & & \downarrow \\ \text{Col}(\mathbf{A}) & & \text{Nul}(\mathbf{A}^T) & & & \end{array} \quad \mathbf{V} = \begin{array}{ccccccc} \uparrow & & \uparrow & \uparrow & & \uparrow \\ \mathbf{v}_1 & \cdots & \mathbf{v}_r & \mathbf{v}_{r+1} & \cdots & \mathbf{v}_n \\ \downarrow & & \downarrow & \downarrow & & \downarrow \\ \text{Row}(\mathbf{A}) & & \text{Nul}(\mathbf{A}) & & & \end{array}$$

Weighted least-squares

Aims to minimise the sum of squared errors.

$$\mathbf{W}\mathbf{y} = \begin{bmatrix} w_1 & & 0 \\ & \ddots & \\ 0 & & w_n \end{bmatrix} \begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix} = \begin{bmatrix} w_1 y_1 \\ \vdots \\ w_n y_n \end{bmatrix}$$

We then apply the normal equation, in order to solve the least squares problem.

$$(WA)^T W A \mathbf{x} = (WA)^T W \mathbf{y} \quad (W A \mathbf{x} = W \mathbf{y})$$

Minimum norm least squares

When the null space of A contains more than the zero vector, we have infinitely many solutions to a least squares problem:

$$A^T A \mathbf{x} = A^T \mathbf{b} \quad \text{is valid for any } \mathbf{x} = \mathbf{x}^* + \mathbf{q} \text{ for any } \mathbf{q} \in \text{Nul}(A).$$

In order to solve the normal equations for this kind of a problem, we use the pseudoinverse A^\dagger .

$$A^\dagger = V_r \Sigma_r^{-1} U_r^T$$

where r is the rank of A.

$$\mathbf{x}^\dagger = A^\dagger \mathbf{b}$$

Where \mathbf{x}^\dagger belongs to $\text{Nul}(A)^\perp$.
(See pdf for more information and proofs.)

Tikhonov regularisation

Is a modification of the normal equations with an added regularisation parameter α .

$$(A^T A + \alpha I) \mathbf{x}_\alpha = A^T \mathbf{b}$$

where

$$\mathbf{x}_\alpha = \sum_{i=1}^n \frac{\sigma_i}{\sigma_i^2 + \alpha} (\mathbf{u}_i \cdot \mathbf{b}) \mathbf{v}_i \quad \lim_{\alpha \rightarrow 0} \quad \mathbf{x}^\dagger = \sum_{i=1}^n \frac{\mathbf{u}_i \cdot \mathbf{b}}{\sigma_i} \mathbf{v}_i$$