### Team Notebook

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### 1 Data Structures

### 1.1 Dynamic Convex Hull Trick

```
const ld is_query = -(1LL << 62);</pre>
struct Line {
    ld m. b:
    mutable std::function<const Line *()> succ;
    bool operator<(const Line &rhs) const {</pre>
       if (rhs.b != is query) return m < rhs.m;</pre>
       const Line *s = succ();
       if (!s) return 0:
       1d x = rhs.m:
       return b - s->b < (s->m - m) * x;
};
struct HullDvnamic : public multiset<Line> { // dvnamic
     upper hull + max value querv
    bool bad(iterator y) {
       auto z = next(y);
       if (y == begin()) {
           if (z == end()) return 0;
           return v->m == z->m && v->b <= z->b:
       auto x = prev(v):
       if (z == end()) return y->m == x->m && y->b <= x->b;
       return (x-b - y-b) * (z-m - y-m) >= (y-b - z-b)
             * (v->m - x->m):
    void insert_line(ld m, ld b) {
       auto y = insert({m, b});
       v->succ = [=] { return next(v) == end() ? 0 : &*next(
            y); };
       if (bad(y)) {
           erase(v);
           return:
       while (next(y) != end() && bad(next(y))) erase(next(y))
       while (y != begin() && bad(prev(y))) erase(prev(y));
    ld best(ld x) {
       auto 1 = *lower_bound((Line) {x, is_query});
       return 1.m * x + 1.b:
};
```

### 1.2 Heavy Light

```
const int N = 2000*100 + 10:
const int L = 20;
int par[N][L], h[N], fath[N], st[N], en[N], sz[N];
vector<int> c[N]: //Adjacency List
int dsz(int s, int p) {
sz[s] = 1:
for(int xt = 0: xt < (int)c[s].size(): xt++) {
 int x = c[s][xt];
 if( x != p ) {
  sz[s] += dsz(x, s);
  if(sz[x] > sz[c[s][0]])
   swap( c[s][0], c[s][xt] );
return sz[s];
void dfs(int s, int p) {
static int ind = 0;
st[s] = ind++:
for(int k = 1: k < L: k++)
 par[s][k] = par[par[s][k-1]][k-1];
for(int xt = 0: xt < (int)c[s].size(): xt++) {</pre>
 int x = c[s][xt];
 if( x == p ) continue;
 fath[x] = x:
 if( xt == 0 ) fath[x] = fath[s];
 h[x] = h[s] + 1:
 par[x][0] = s:
 dfs(x, s);
en[s] = ind;
void upset(int u, int w, int qv) {
int stL = max( st[w] , st[fath[u]] );
set( stL, st[u] + 1 , qv , 0, n , 1 ); //l,r,val,s,e,id
if( stL == st[w] ) return;
upset( par[fath[u]][0] , w , qv );
```

### 1.3 Implicit Treap

```
const int MAX=1e6+9;

typedef struct item * pitem;
struct item {
  int prior, value, cnt;
```

```
bool rev:
pitem 1, r;
int cnt (pitem it) {
return it ? it->cnt : 0:
void upd_cnt (pitem it) {
if (it)
 it->cnt = cnt(it->1) + cnt(it->r) + 1;
void push (pitem it) {
if (it && it->rev) {
 it->rev = false:
 swap (it->1, it->r);
 if (it->1) it->1->rev ^= true;
 if (it->r) it->r->rev ^= true;
void merge (pitem & t, pitem 1, pitem r) {
push (1);
push (r);
if (!1 || !r)
 t = 1 ? 1 : r:
else if (l->prior > r->prior)
 merge (1->r, 1->r, r), t = 1;
 merge (r->1, 1, r->1), t = r;
upd cnt (t):
void split (pitem t, pitem & 1, pitem & r, int key, int add
    = 0) {
if (!t)
 return void( l = r = 0 ):
int cur_key = add + cnt(t->1);
if (key <= cur_key)</pre>
 split (t->1, 1, t->1, key, add), r = t;
 split (t->r, t->r, r, key, add + 1 + cnt(t->1)), 1 = t;
upd cnt (t):
void reverse (pitem t, int 1, int r) {
pitem t1. t2. t3:
split (t, t1, t2, 1);
split (t2, t2, t3, r-l+1);
t2->rev ^= true:
merge (t, t1, t2);
```

```
3
```

```
merge (t, t, t3);
}
void output (pitem t) {
   if (!t) return;
   push (t);
   output (t->1);
   printf ("%d ", t->value);
   output (t->r);
}
```

### 1.4 Link-Cut tree

```
Node x[N];
struct Node {
int sz. label: /* size. label */
Node *p, *pp, *1, *r; /* parent, path-parent, left, right
     pointers */
Node() { p = pp = 1 = r = 0; }
void update(Node *x) {
x->sz = 1:
if(x->1) x->sz += x->1->sz;
if(x->r) x->sz += x->r->sz:
void rotr(Node *x) {
Node *y, *z;
y = x->p, z = y->p;
if((y->1 = x->r)) y->1->p = y;
x->r = y, y->p = x;
if((x->p = z)) {
 if(y == z->1) z->1 = x;
 else z->r = x:
x->pp = y->pp;
y->pp = 0;
update(y);
void rotl(Node *x) {
Node *y, *z;
y = x->p, z = y->p;
if((y->r = x->1)) y->r->p = y;
x->1 = y, y->p = x;
if((x->p = z)) {
 if(y == z->1) z->1 = x;
 else z->r = x:
x \rightarrow pp = y \rightarrow pp;
y->pp = 0;
update(y);
```

```
void splay(Node *x) {
Node *y, *z;
while(x->p) {
 y = x->p;
 if(y->p == 0) {
  if(x == y->1) rotr(x);
  else rotl(x);
 else {
  z = y - p;
  if(v == z->1) {
   if(x == y -> 1) rotr(y), rotr(x);
   else rotl(x), rotr(x);
  else { if(x == y->r) rotl(y), rotl(x);
   else rotr(x). rotl(x):
  }
 }
}
update(x);
Node *access(Node *x) {
splay(x);
if(x->r) {
 x->r->pp = x;
 x->r->p = 0;
 x->r = 0:
 update(x);
Node *last = x;
while(x->pp) {
 Node *y = x->pp;
 last = v;
 splay(y);
 if(y->r) {
  y->r->pp = y;
 y->r->p = 0;
 y->r = x;
 x->p = y;
 x->pp = 0;
 update(v):
 splay(x);
return last:
Node *root(Node *x) {
access(x):
while(x \rightarrow 1) x = x \rightarrow 1;
```

```
splav(x):
return x;
void cut(Node *x) {
access(x):
x->1->p = 0;
x->1 = 0;
update(x);
void link(Node *x, Node *y) {
access(x):
access(v):
x->1 = y;
y->p = x;
update(x);
Node *lca(Node *x, Node *y) {
access(x):
return access(v):
int depth(Node *x) {
access(x):
return x->sz - 1;
void init(int n) {
for(int i = 0; i < n; i++) {</pre>
 x[i].label = i:
 update(&x[i]);
```

### 1.5 Ordered Set

```
4
```

```
cout<<*X.find_by_order(2)<<end1; // 4
cout<<*X.find_by_order(4)<<end1; // 16
cout<<(end(X)==X.find_by_order(6))<<end1; // true

cout<<X.order_of_key(-5)<<end1; // 0
cout<<X.order_of_key(1)<<end1; // 0
cout<<X.order_of_key(3)<<end1; // 2
cout<<X.order_of_key(4)<<end1; // 2
cout<<X.order_of_key(4)<<end1; // 5</pre>
```

### 1.6 Seg Persistant

```
typedef pair<pair<int,int>,int > ANS;
#define MX second
#define LE first.first
#define RT first.second
const int MAXN=1e5+9.LOG=22:
int root[MAXN], le[LOG * MAXN * 2], ri[LOG * MAXN * 2], sz,
    lleft[LOG * MAXN * 2]. rright[LOG * MAXN * 2]:
int maxi[LOG * MAXN * 2];
int n, q;
int h[MAXN]. vec[MAXN]:
pair<int, int> sec[MAXN];
int build(int b, int e){
int id = sz++:
if (e - b == 1) return
 id:
int mid = (b + e)/2;
le[id] = build(b, mid):
ri[id] = build(mid, e):
return id;
void merge(int id, int b, int e, int mid){
maxi[id] = max(maxi[le[id]], maxi[ri[id]]);
maxi[id] = max(maxi[id], rright[le[id]] + lleft[ri[id]]):
lleft[id] = lleft[le[id]];
if (lleft[id] == (mid - b))
 lleft[id] += lleft[ri[id]]:
rright[id] = rright[ri[id]]:
if (rright[id] == (e - mid)) rright[id] += rright[le[id]];
int modify(int id, int b, int e, int x){
int nid = sz++:
if (e - b == 1){}
 lleft[nid] = rright[nid] = maxi[nid] = 1;
 return nid:
```

```
int mid = (b+e)/2:
le[nid]=le[id]:
ri[nid]=ri[nid];
if (x<mid)</pre>
 le[nid]=modify(le[nid],b,mid,x);
 ri[nid] = modify(ri[nid], mid , e, x);
merge(nid,b,e,mid);
return nid:
ANS mg(ANS a. ANS b. int s1. int s2){
ANS ret:
ret.MX = max(a.MX, b.MX);
ret.MX = max(ret.MX, a.RI + b.LE):
ret.LE = a.LE:
if (a.LE == s1) ret.LE += b.LE;
ret.RI = b.RI:
if (b.RI == s2) ret.RI += a.RI;
return ret:
ANS get(int id, int b, int e, int 1, int r){
if (1 <= b && e <= r)
 return {{maxi[id], lleft[id]}, rright[id]};
if (r \le b \mid l \in \le 1)
 return {{0, 0}, 0};
int mid = (b + e)/2:
return mg(get(le[id], b, mid, l, r), get(ri[id], mid, e, l,
   ), min(mid - b, max(0, mid - 1)), min(e - mid, max(0, r -
     mid))):
void init(){
copv(h, h + n, vec);
sort(vec. vec + n):
reverse(vec, vec + n);
for (int i = 0; i < n; i++)</pre>
 sec[i] = {h[i], i}:
sort(sec. sec + n):
reverse(sec, sec + n);
root[0] = build(0, n):
for (int i = 0; i < n; i++)</pre>
 root[i + 1] = modify(root[i], 0, n, sec[i].second);
int main(){
ios::svnc with stdio(false):
cin.tie(0):
cin >> n:
```

```
for (int i = 0; i < n; i++)
  init();
  cin >> q;
  while (q--){
    int 1, r, w;
    cin >> 1 >> r >> w;
    l--;
    int b = 0, e = n, mid, ret = n;
    while (b <= e){
        mid = (b + e) / 2;
        if (get(root[mid], 0, n, 1, r).MX >= w){
        ret = mid;
        e = mid - 1;
    }
    else
        b = mid + 1;
    }
    cout << vec[ret - 1] << "\n";
}
    return 0;
}</pre>
```

### 1.7 Segment Beats

```
struct JiDriverSegmentTree {
   static const int T = (1 << 20);
   static const int INF = 1e9 + 7;
   struct Node {
       int max:
      long long sum;
   } tree[T]:
   void updateFromChildren(int v) {
       tree[v].sum = tree[2 * v].sum + tree[2 * v + 1].sum:
       tree[v].max = max(tree[2 * v].max. tree[2 * v + 1].
           max):
   void build(int v, int 1, int r, const vector<int>&
        inputArray) {
      if (1 + 1 == r) {
          tree[v].max = tree[v].sum = inputArray[1];
          int mid = (r + 1) / 2;
          build(2 * v, 1, mid, inputArray);
          build(2 * v + 1, mid, r, inputArray):
          updateFromChildren(v);
   void build(const vector<int>& inputArray) {
```

```
n = inputArrav.size();
       build(1, 0, n, inputArray);
   void updateModEq(int v, int 1, int r, int q1, int qr, int
       if (gr <= 1 || r <= gl || tree[v].max < val) return;</pre>
       if (1 + 1 == r) {
           tree[v].max %= val:
           tree[v].sum = tree[v].max:
       int mid = (r + 1) / 2:
       updateModEq(2 * v, 1, mid, ql, qr, val);
       updateModEq(2 * v + 1, mid, r, ql, qr, val);
       updateFromChildren(v);
   void updateModEq(int ql, int qr, int val) {
       updateModEq(1, 0, n, ql, qr, val);
   void updateEq(int v, int l, int r, int qi, int val) {
       if (1 + 1 == r) {
           tree[v].max = tree[v].sum = val;
           return:
       int mid = (1 + r) / 2;
       if (qi < mid) {</pre>
           updateEq(2 * v, 1, mid, qi, val);
           updateEq(2 * v + 1, mid, r, qi, val);
       updateFromChildren(v);
   void updateEq(int qi, int val) {
       updateEq(1, 0, n, qi, val);
   long long findSum(int v, int l, int r, int ql, int qr) {
       if (qr <= 1 || r <= q1) {</pre>
          return 0;
       if (ql <= 1 && r <= qr) {</pre>
           return tree[v].sum:
       int mid = (r + 1) / 2;
       return findSum(2 * v, 1, mid, ql, qr) + findSum(2 * v
             + 1, mid, r, ql, qr);
   long long findSum(int ql, int qr) {
       return findSum(1, 0, n, ql, qr);
} segTree;
```

### 1.8 Treap

```
struct item {
 int key, prior;
item * 1, * r;
 item() { }
item (int key, int prior) : key(key), prior(prior), 1(NULL)
      . r(NULL) { }
typedef item * pitem;
void split (pitem t. int kev. pitem & l. pitem & r) {
if (!t)
 1 = r = NULL;
else if (key < t->key)
 split (t->1, kev, 1, t->1), r = t;
 split (t\rightarrow r, key, t\rightarrow r, r), l = t;
void insert (pitem & t, pitem it) {
 if (!t)
 t = it;
else if (it->prior > t->prior)
 split (t, it->key, it->l, it->r), t = it;
 insert (it->key < t->key ? t->l : t->r, it);
void merge (pitem & t, pitem 1, pitem r) {
if (!1 || !r)
 t = 1 ? 1 : r:
 else if (1->prior > r->prior)
 merge (1->r, 1->r, r), t = 1;
 merge (r->1, 1, r->1), t = r:
void erase (pitem & t. int kev) {
 if (t->key == key)
 merge (t, t->1, t->r);
 erase (key < t->key ? t->l : t->r, key);
pitem unite (pitem 1, pitem r) {
if (!1 || !r) return 1 ? 1 : r:
if (1->prior < r->prior) swap (1, r);
 pitem lt, rt;
 split (r, 1->key, lt, rt);
1->1 = unite (1->1, 1t):
1->r = unite (1->r, rt);
return 1:
pitem root = NULL;
```

```
int main()
{
  ios_base::sync_with_stdio(false),cin.tie(0);
  item a = item(10,20);
  item b = item(10,20);
  insert(root, &a);
  insert(root, &b);
  return 0;
}
```

### 2 Dp Optimizations

### 2.1 Convex Hull Trick

```
#define F first
#define S second
#define pii pair <int, int>
#define pb psh_back
typedef long long 11;
vector <pair <11, 11> > cv;
11 barkhord(pair<11, 11> p1, pair<11, 11> p2) { //Make sure
    m1 > m2:
return (p2.S - p1.S + p1.F - p2.F - 1) / (p1.F - p2.F);
11 get(11 t)
int lo = -1. hi = cv.size() - 1:
while(hi - lo > 1)
 int mid = (lo + hi)/2;
 if(barkhord(cv[mid + 1], cv[mid]) <= t) lo = mid;</pre>
 else hi = mid:
return t * cv[hi].F + cv[hi].S:
//\{m, h\} in points.
void build(vector <pair <11, 11> > points) {
sort(points.begin(), points.end(), cmp); //Make them
     increasing in m and decreasing in h.
for (auto X : points)
```

### 2.2 Knuth

Knuth Optimization is applicable if  $C_{i,j}$  satisfied the following 2 conditions:

- 1- Quadrangle Inequality:  $C_{a,c} + C_{b,d} \le C_{a,d} + C_{b,c}$  for  $a \le b \le c \le d$ 
  - 2- Monotonicity:  $C_{b,c} \leq C_{a,d}$  for  $a \leq b \leq c \leq d$

Then if the smallest k that gives optimal answer in  $dp_{i,j} = dp_{i-1,k} + C_{k,j}$  equals to  $A_{i,j}$  we have:

$$A_{i,j-1} \le A_{i,j} \le A_{i+1,j}$$

### 3 Geometry

### 3.1 Convex Hull 3D

```
pt cross(pt u.pt v){return pt(u.Y*v.Z-u.Z*v.Y.u.Z*v.X-u.X*v
     .Z,u.X*v.Y-u.Y*v.X); }
pt cross(pt o,pt p,pt q){return _cross(p-o,q-o);}
ld dot(pt p,pt q){return p.X*q.X+p.Y*q.Y+p.Z*q.Z;}
pt shift(pt p) {return pt(p.Y,p.Z,p.X);}
pt norm(pt p)
if(p.Y<p.X || p.Z<p.X) p=shift(p);</pre>
if(p.Y<p.X) p=shift(p);</pre>
return p;
const int MAX=1000;
pt P[MAX];
vector<pt>ans:
queue<pair<int,int> >Q;
set<pair<int.int> >mark:
int main()
 cin>>n;
 int mn=0:
 for(int i=0;i<n;i++)</pre>
 cin>>P[i].X>>P[i].Y>>P[i].Z;
 if(P[i]<P[mn]) mn=i:</pre>
int nx=(mn==0):
for(int i=0;i<n;i++)</pre>
 if(i!=mn && i!=nx && cross2d(P[nx]-P[mn],P[i]-P[mn])>0)
 Q.push({mn,nx});
 while(!Q.empty())
 int v=Q.front().first,u=Q.front().second;
  Q.pop();
  if(mark.count({v,u})) continue;
  mark.insert({v.u}):
  int p=-1:
  for(int q=0;q< n;q++)
  if(a!=v && a!=u)
   if(p==-1 || dot(cross(P[v],P[u],P[p]),P[q]-P[v])<0)</pre>
  ans.push_back(norm(pt(v,u,p)));
 Q.push({p,u});
  Q.push(\{v,p\});
 sort(ans.begin(),ans.end());
```

```
ans.resize(unique(ans.begin(),ans.end())-ans.begin());
for(int i=0;i<ans.size();i++)
  cout<<ans[i].X<<" "<<ans[i].Y<<" "<<ans[i].Z<<end1;
}</pre>
```

### 3.2 Delaunay Triangulation N\*N

```
struct Delaunav{
vector <pt> p;
vector <int> to, nxt, perm:
int add_edge(int q, int bef=-1){
 int cnt = sz(to);
 to.pb(q);
 nxt.pb(-1);
 if (bef != -1){
 nxt[bef] = cnt:
  to.pb(to[bef]);
 nxt.pb(-1);
 return cnt;
bool onconvex(int e){
 if (nxt[nxt[e]]] != e) return true;
 if (dir(p[to[e^1]], p[to[e]], p[to[nxt[e]]]) < 0) return</pre>
 return false:
int before(int e){
 int cur = e, last = -1;
 do{
  last = cur:
  cur = nxt[cur^1]:
 }while (cur != e);
 return last^1:
void easy_triangulate(){
 to.clear():
 nxt.clear();
 perm = vector<int>(sz(p));
 for (int i = 0; i < sz(p); i++)</pre>
  perm[i] = i:
 sort(perm.begin(), perm.end(), [&] (int i, int j){
   return p[i] < p[i]; });</pre>
 sort(p.begin(), p.end());
 if (dir(p[0], p[1], p[2]) > 0){
  swap(p[1], p[2]);
  swap(perm[1], perm[2]);
 int to0 = add_edge(0), to0c = add_edge(2),
```

```
to1 = add \ edge(1), to1c = add \ edge(0).
 to2 = add_edge(2), to2c = add_edge(1);
nxt[to1] = to2; nxt[to2] = to0;
nxt[to0] = to1: nxt[to0c] = to2c:
nxt[to2c] = to1c; nxt[to1c] = to0c;
int e = to0:
bool D2 = true;
for (int i = 3; i < sz(p); i++){
 pt q = p[i];
 if (D2){
  int edge = e:
   if (dir(q, p[to[edge^1]], p[to[edge]])){
    D2 = false:
    break:
   edge = nxt[edge]:
  } while (edge != e);
 vector <int> vis:
 if (D2){
  while (p[to[e^1]] < p[to[e]])</pre>
   e = nxt[e];
  vis.pb(e);
  e = nxt[e];
 }
 else{
  while (dir(q, p[to[e^1]], p[to[e]]) <= 0 || dir(q, p[to[e
       ^1]], p[to[before(e)^1]]) < 0)
   e = nxt[e]:
  while (dir(q, p[to[e^1]], p[to[e]]) > 0){
   vis.pb(e);
   e = nxt[e]:
 }
 int b = before(vis[0]):
 int ex = add_edge(i, b);
 int last = ex^1;
 for (int edge : vis){
  nxt[last] = edge;
  int eq = add edge(i, edge):
  nxt[edge] = eq;
  nxt[eq] = last;
  last = eq^1;
 nxt[ex] = last:
 nxt[last] = e;
bool incircle(pt a, pt b, pt c, pt d){
```

```
if (dir(a, b, c) < 0)
 swap(b, c):
return a.z() * (b.x * (c.y - d.y) - c.x * (b.y - d.y)
  + d.x * (b.v - c.v))
 -b.z() * (a.x * (c.y - d.y) - c.x * (a.y - d.y) + d.x *
      (a.v - c.v))
 + c.z() * (a.x * (b.y - d.y) - b.x * (a.y - d.y) + d.x *
      (a.v - b.v))
 -d.z() * (a.x * (b.y - c.y) - b.x * (a.y - c.y) + c.x *
      (a.v - b.v)) > 0;
bool locallv(int e){
pt a = p[to[e^1]], b = p[to[e]], c = p[to[nxt[e]]], d = p[
     to[nxt[e^1]]]:
if (onconvex(e)) return true:
if (onconvex(e^1)) return true;
if (incircle(a, b, c, d)) return false:
if (incircle(b, a, d, c)) return false;
return true:
void flip(int e){
int a = nxt[e], b = nxt[a],
c = nxt[e^1], d = nxt[c];
nxt[d] = a:
nxt[b] = c:
to[e] = to[c]:
 nxt[a] = e:
nxt[e] = d:
 to[e^1] = to[a];
nxt[c] = e^1:
nxt[e^1] = b;
void delaunav triangulate(){
if (sz(to) == 0)
 easy_triangulate();
bool *mark = new bool[sz(to)]:
fill(mark, mark + sz(to), false);
vector <int> bad:
for (int e = 0; e < sz(to); e++){}
 if (!mark[e/2] && !locallv(e)){
  bad.pb(e):
  mark[e/2] = true;
 }
while (sz(bad)){
 int e = bad.back();
 bad.pop_back();
 mark[e/2] = false;
 if (!locallv(e)){
```

```
int to_check[4] = {nxt[e], nxt[nxt[e]], nxt[e^1], nxt[nxt
   flip(e);
   for (int i = 0: i < 4: i++)
    if (!mark[to_check[i]/2] && !locally(to_check[i])){
     bad.pb(to check[i]):
     mark[to_check[i]/2] = true;
 for (int e = 0: e < sz(to): e++)
  assert(locallv(e)):
vector <tri> get_triangles(){
 vector <tri> res:
 bool *mark = new bool[sz(to)];
 fill(mark, mark + sz(to), false):
 for (int e = 0; e < sz(to); e++){</pre>
  if (mark[e]) continue:
  if (onconvex(e)) continue:
  pt a = p[to[e^1]], b = p[to[e]], c = p[to[nxt[e]]];
  mark[e] = mark[nxt[e]] = mark[nxt[nxt[e]]] = true;
  res.pb(tri(perm[to[e^1]], perm[to[e]], perm[to[nxt[e]]]))
 return res;
vector <pair<ls, pt>> get_voronoi_edges(){
 vector <pair<ls, pt>> res;
 for (int e = 0; e < sz(to); e++){</pre>
 pt a = p[to[e^1]], b = p[to[e]], c = p[to[nxt[e]]], d = p
       [to[nxt[e^1]]];
  if (onconvex(e^1)){
   pt o1 = center(a, b, c),
      o2 = (a+b)/2:
   pt ab = (b-a):
   pt per(ab.v, -ab.x);
   o2 = o2 + per*100000; //infinity
   res.pb({{o1, o2}, a});
   continue:
  if (onconvex(e)) continue;
  if (e&1) continue:
  res.pb({{center(a, b, c), center(b, a, d)}, a});
 return res;
Delaunay(vector <pt> &p):p(p){}
};
```

### 3.3 Delaunay Triangulation NlogN

```
typedef long long 11;
bool ge(const ll& a, const ll& b) { return a >= b; }
bool le(const ll& a, const ll& b) { return a <= b; }
bool eg(const ll& a, const ll& b) { return a == b; }
bool gt(const ll& a, const ll& b) { return a > b; }
bool lt(const ll& a, const ll& b) { return a < b; }
int sgn(const ll& a) { return a >= 0 ? a ? 1 : 0 : -1; }
struct pt {
    11 x, v;
    pt() { }
    pt(11 _x, 11 _y) : x(_x), y(_y) { }
    pt operator-(const pt& p) const {
       return pt(x - p.x, y - p.y);
    11 cross(const pt& p) const {
       return x * p.y - y * p.x;
    ll cross(const pt& a, const pt& b) const {
       return (a - *this).cross(b - *this):
    11 dot(const pt% p) const {
       return x * p.x + y * p.y;
    11 dot(const pt& a, const pt& b) const {
       return (a - *this).dot(b - *this):
    11 sqrLength() const {
       return this->dot(*this):
    bool operator==(const pt& p) const {
        return eq(x, p.x) && eq(y, p.y);
};
const pt inf_pt = pt(1e18, 1e18);
struct QuadEdge {
    pt origin;
    QuadEdge* rot = nullptr;
    QuadEdge* onext = nullptr;
    bool used = false;
    QuadEdge* rev() const {
       return rot->rot;
    QuadEdge* lnext() const {
       return rot->rev()->onext->rot;
```

```
QuadEdge* oprev() const {
        return rot->onext->rot;
    pt dest() const {
        return rev()->origin:
};
QuadEdge* make_edge(pt from, pt to) {
    QuadEdge* e1 = new QuadEdge:
    QuadEdge* e2 = new QuadEdge:
    QuadEdge* e3 = new QuadEdge;
    QuadEdge* e4 = new QuadEdge;
    e1->origin = from;
    e2->origin = to;
    e3->origin = e4->origin = inf pt:
    e1->rot = e3:
    e2 \rightarrow rot = e4:
    e3 \rightarrow rot = e2:
    e4 \rightarrow rot = e1;
    e1->onext = e1:
    e2 - onext = e2;
    e3 \rightarrow onext = e4:
    e4 \rightarrow onext = e3:
    return e1;
void splice(QuadEdge* a, QuadEdge* b) {
    swap(a->onext->rot->onext. b->onext->rot->onext):
    swap(a->onext, b->onext);
void delete_edge(QuadEdge* e) {
    splice(e, e->oprev());
    splice(e->rev(), e->rev()->oprev());
    delete e->rev()->rot:
    delete e->rev():
    delete e->rot;
    delete e:
QuadEdge* connect(QuadEdge* a. QuadEdge* b) {
    QuadEdge* e = make_edge(a->dest(), b->origin);
    splice(e, a->lnext());
    splice(e->rev(), b);
    return e;
bool left_of(pt p, QuadEdge* e) {
```

```
return gt(p.cross(e->origin, e->dest()), 0);
bool right_of(pt p, QuadEdge* e) {
   return lt(p.cross(e->origin, e->dest()), 0);
template <class T>
T det3(T a1, T a2, T a3, T b1, T b2, T b3, T c1, T c2, T c3)
   return a1 * (b2 * c3 - c2 * b3) - a2 * (b1 * c3 - c1 * b3
          a3 * (b1 * c2 - c1 * b2);
bool in_circle(pt a, pt b, pt c, pt d) {
// If there is int128, calculate directly.
// Otherwise, calculate angles.
#if defined( LP64 ) || defined( WIN64)
   __int128 det = -det3<__int128>(b.x, b.y, b.sqrLength(), c
        .x, c.v,
                                c.sqrLength(), d.x, d.y, d.
                                    sqrLength());
   det += det3<__int128>(a.x, a.y, a.sqrLength(), c.x, c.y,
        c.sqrLength(), d.x,
                       d.v, d.sqrLength());
   det -= det3<__int128>(a.x, a.y, a.sqrLength(), b.x, b.y,
        b.sarLength(), d.x.
                       d.v, d.sqrLength());
   det += det3<__int128>(a.x, a.y, a.sqrLength(), b.x, b.y,
        b.sqrLength(), c.x,
                       c.y, c.sqrLength());
   return det > 0:
#else
   auto ang = [](pt 1, pt mid, pt r) {
      11 x = mid.dot(1, r):
      11 v = mid.cross(1, r);
       long double res = atan2((long double)x, (long double)
       return res:
   long double kek = ang(a, b, c) + ang(c, d, a) - ang(b, c, d)
         d) - ang(d, a, b):
   if (kek > 1e-8)
       return true;
   else
       return false;
#endif
```

```
pair<QuadEdge*. QuadEdge*> build tr(int 1, int r, vector<pt</pre>
    } (a %<
   if (r - 1 + 1 == 2) {
       QuadEdge* res = make_edge(p[1], p[r]);
       return make_pair(res, res->rev());
   if (r - 1 + 1 == 3) {
       QuadEdge *a = make_edge(p[1], p[1 + 1]), *b =
            make_edge(p[l + 1], p[r]);
       splice(a->rev(), b);
       int sg = sgn(p[1].cross(p[1 + 1], p[r]));
       if (sg == 0)
          return make_pair(a, b->rev());
       QuadEdge* c = connect(b, a);
       if (sg == 1)
          return make_pair(a, b->rev());
          return make_pair(c->rev(), c);
   int mid = (1 + r) / 2:
   QuadEdge *ldo, *ldi, *rdo, *rdi;
   tie(ldo, ldi) = build_tr(l, mid, p);
   tie(rdi, rdo) = build_tr(mid + 1, r, p);
   while (true) {
       if (left_of(rdi->origin, ldi)) {
          ldi = ldi->lnext();
          continue:
       if (right_of(ldi->origin, rdi)) {
          rdi = rdi->rev()->onext:
          continue;
       break:
   QuadEdge* basel = connect(rdi->rev(), ldi);
   auto valid = [&basel](QuadEdge* e) { return right of(e->
        dest(), basel); };
   if (ldi->origin == ldo->origin)
       ldo = basel->rev();
   if (rdi->origin == rdo->origin)
       rdo = basel:
   while (true) {
       QuadEdge* lcand = basel->rev()->onext:
       if (valid(lcand)) {
          while (in_circle(basel->dest(), basel->origin,
               lcand->dest().
                          lcand->onext->dest())) {
              QuadEdge* t = lcand->onext:
              delete_edge(lcand);
              lcand = t:
```

```
QuadEdge* rcand = basel->oprev();
      if (valid(rcand)) {
          while (in_circle(basel->dest(), basel->origin,
               rcand->dest().
                          rcand->oprev()->dest())) {
              QuadEdge* t = rcand->oprev();
              delete_edge(rcand);
              rcand = t;
      }
      if (!valid(lcand) && !valid(rcand))
          break:
      if (!valid(lcand) ||
          (valid(rcand) && in_circle(lcand->dest(), lcand->
                                   rcand->origin, rcand->
                                        dest())))
          basel = connect(rcand, basel->rev());
       else
          basel = connect(basel->rev(), lcand->rev());
   return make_pair(ldo, rdo);
vector<tuple<pt, pt, pt>> delaunay(vector<pt> p) {
   sort(p.begin(), p.end(), [](const pt& a, const pt& b) {
      return lt(a.x, b.x) || (eq(a.x, b.x) && lt(a.y, b.y))
   auto res = build_tr(0, (int)p.size() - 1, p);
   QuadEdge* e = res.first;
   vector<QuadEdge*> edges = {e};
   while (lt(e->onext->dest().cross(e->dest(), e->origin),
        0))
       e = e->onext:
   auto add = [&p, &e, &edges]() {
       QuadEdge* curr = e;
      do {
          curr->used = true:
          p.push_back(curr->origin);
          edges.push_back(curr->rev());
          curr = curr->lnext();
      } while (curr != e);
   }:
   add();
   p.clear();
   int kek = 0:
   while (kek < (int)edges.size()) {</pre>
```

```
if (!(e = edges[kek++])->used)
          add();
}
vector<tuple<pt, pt, pt>> ans;
for (int i = 0; i < (int)p.size(); i += 3) {
          ans.push_back(make_tuple(p[i], p[i + 1], p[i + 2]));
}
return ans;</pre>
```

### 3.4 Find Polynomial from it's Points

```
P(x) = \sum_{i=1}^{n} y_i \prod_{j=1, j \neq i}^{n} \frac{x - x_j}{x_i - x_j}
```

### 3.5 Geometry Duality

duality of point (a, b) is y = ax - b and duality of line y = ax + b is (a, -b)Properties:

- 1. p is on l iff l\* is in p\*
- 2. p is in intersection of l1 and l2 iff l1\* and l2\* lie on p\*
- 3. Duality preserve vertical distance
- 4. Translating a line in primal to moving vertically in dual
- 5. Rotating a line in primal to moving a point along a non-vertical line
- 6.  $li \cap lj$  is a vertex of lower envelope  $\iff$  (li\*, lj\*) is an edge of upper hull in dual

### 3.6 Half Planes

```
typedef int T;
typedef long long T2;
typedef long long T4; // maybe int128_t

const int MAXLINES = 100 * 1000 + 10;
const int INF = 20 * 1000 * 1000;
```

```
typedef pair<T, T> point;
typedef pair<point, point> line;
#define X first
#define Y second
#define A first
#define B second
// REPLACE ZERO WITH EPS FOR DOUBLE
point operator - (const point &a, const point &b) {
return point(a.X - b.X, a.Y - b.Y):
T2 cross(point a, point b) {
return ((T2)a.X * b.Y - (T2)a.Y * b.X):
}
bool cmp(line a, line b) {
 bool aa = a.A < a.B;
 bool bb = b.A < b.B:
 if (aa == bb) {
 point v1 = a.B - a.A;
 point v2 = b.B - b.A;
 if (cross(v1, v2) == 0)
  return cross(b.B - b.A, a.A - b.A) > 0;
  return cross(v1, v2) > 0;
 else
 return aa;
bool parallel(line a, line b) {
return cross(a.B - a.A, b.B - b.A) == 0;
pair<T2, T2> alpha(line a, line b) {
return pair<T2, T2>(cross(b.A - a.A, b.B - b.A),
  cross(a.B - a.A. b.B - b.A)):
bool fcmp(T4 f1t, T4 f1b, T4 f2t, T4 f2b) {
 if (f1b < 0) {
 f1t *= -1:
 f1b *= -1;
 if (f2b < 0) {
 f2t *= -1:
```

```
f2b *= -1:
return f1t * f2b < f2t * f1b; // check with eps</pre>
bool check(line a, line b, line c) {
bool crs = cross(c.B - c.A, a.B - a.A) > 0;
pair<T2, T2> a1 = alpha(a, b);
pair<T2, T2> a2 = alpha(a, c);
bool alp = fcmp(a1.A, a1.B, a2.A, a2.B);
return (crs ^ alp):
bool notin(line a, line b, line c) { // is intersection of a
     and b in ccw direction of c?
if (parallel(a, b))
 return false:
if (parallel(a, c))
 return cross(c.B - c.A. a.A - c.A) < 0:
if (parallel(b, c))
 return cross(c.B - c.A, b.A - c.A) < 0;
return !(check(a, b, c) && check(b, a, c));
void print(vector<line> lines) {
cerr << " @ @ @ " << endl:
for (int i = 0: i < lines.size(): i++)</pre>
 cerr << lines[i].A.X << " " << lines[i].A.Y << " -> " <<
      lines[i].B.X << " " << lines[i].B.Y << endl;
cerr << " @ @ @ " << endl<< endl:
line dq[MAXLINES];
vector<line> half_plane(vector<line> lines) {
lines.push back(line(point(INF, -INF), point(INF, INF)));
lines.push_back(line(point(-INF, INF), point(-INF, -INF)));
lines.push_back(line(point(-INF, -INF), point(INF, -INF)));
lines.push_back(line(point(INF, INF), point(-INF, INF)));
sort(lines.begin(), lines.end(), cmp);
int ptr = 0:
for (int i = 0; i < lines.size(); i++)</pre>
 if (i > 0 &&
   (lines[i-1].A < lines[i-1].B) == (lines[i].A < lines[
        il.B) &&
   parallel(lines[i - 1], lines[i]))
  continue;
 else
 lines[ptr++] = lines[i];
lines.resize(ptr):
```

```
if (lines.size() < 2)</pre>
 return lines:
//print(lines);
int f = 0, e = 0:
dq[e++] = lines[0];
da[e++] = lines[1]:
for (int i = 2; i < lines.size(); i++) {</pre>
 while (f < e - 1 \&\& notin(dq[e - 2], dq[e - 1], lines[i]))
 //print(vector<line>(dg + f, dg + e));
 if (e == f + 1) {
  T2 crs = cross(dq[f].B - dq[f].A, lines[i].B - lines[i].A
       );
  if (crs < 0)
   return vector<line>():
  else if (crs == 0 && cross(lines[i].B - lines[i].A, dq[f
       l.B - lines[i].A) < 0
   return vector<line>();
 while (f < e - 1 \&\& notin(dq[f], dq[f + 1], lines[i]))
 dq[e++] = lines[i];
while (f < e - 1 \&\& notin(dq[e - 2], dq[e - 1], dq[f]))
while (f < e - 1 \&\& notin(dq[f], dq[f + 1], dq[e - 1]))
vector<line> res:
res.resize(e - f);
for (int i = f: i < e: i++)</pre>
res[i - f] = dq[i];
return res;
int main() {
int n:
cin >> n:
vector<line> lines;
for (int i = 0; i < n; i++) {</pre>
 int x1, v1, x2, v2;
 cin >> x1 >> y1 >> x2 >> y2;
 lines.push_back(line(point(x1, y1), point(x2, y2)));
lines = half_plane(lines);
cout << lines.size() << endl:</pre>
for (int i = 0; i < lines.size(); i++)</pre>
 cout << lines[i].A.X << " " << lines[i].A.Y << " " <<</pre>
      lines[i].B.X << " " << lines[i].B.Y << endl:</pre>
```

### 3.7 Minimum Enclosing Circle

```
const int N = 1000*100 + 10:
struct point {
   11 x, y, z;
typedef vector<point> circle;
bool ccw(point a, point b, point c) {
   return (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y)
         (v) >= 0:
bool incircle( circle a , point p ) {
   if( sz(a) == 0 ) return false;
   if(sz(a) == 1)
       return a[0].x == p.x && a[0].y == p.y;
   if(sz(a) == 2) {
       point mid = \{a[0].x+a[1].x, a[0].y+a[1].y\};
       return sq(2*p.x-mid.x) + sq(2*p.y-mid.y) <= sq(2*a
            [0].x-mid.x) + sq(2*a[0].y-mid.y);
   if( !ccw(a[0], a[1], a[2]) )
       swap(a[0], a[2]):
   return incircle(a[0],a[1],a[2], p) >= 0;
point a[N]:
circle solve(int i, circle curr) {
   assert(curr.size() <= 3):</pre>
   if(i == 0)
       return curr:
   circle ret = solve(i-1, curr);
   if( incircle(ret, a[i-1]) )
       return ret:
   curr.pb(a[i-1]):
   return solve(i-1, curr);
}
int n;
void gg(circle c) {
   if(sz(c) == 1) {
       cout << ld(a[0].x) << " " << ld(a[0].y) << endl;</pre>
       cout << 0.1 << endl:
       return:
   if(sz(c) == 2) {
       point mid = \{c[0].x+c[1].x, c[0].y+c[1].y\};
       1d ret = sqrt(sq(2*c[0].x-mid.x) + sq(2*c[0].y-mid.y)
       cout << ld(mid.x) / 2 << " " << ld(mid.y) /2 << endl;</pre>
       cout << ret << endl:</pre>
   } else {
       1pt a[3];
```

```
for(int i = 0; i < 3; i++)
           a[i] = lpt(c[i].x, c[i].y);
       lpt A = 1d(0.5) * (a[0] + a[1]), C = 1d(0.5) * (a[1])
       lpt B = A + (a[1] - a[0]) * lpt(0, 1), D = C + (a[2])
            - a[1]) * lpt(0, 1):
       lpt center = intersection( A , B , C , D );
       ld ret = abs(a[0] - center);
       cout << center.real() << " " << center.imag() << endl</pre>
       cout << ret << endl:</pre>
   }
int main(){
   cin >> n:
   for(int i = 0; i < n; i++) {</pre>
       cin >> a[i].x >> a[i].y;
       a[i].z = sq(a[i].x) + sq(a[i].y);
   srand(time(NULL)):
   for(int i = 1; i < n; i++)</pre>
       swap(a[i], a[rand()%(i+1)]);
   circle ans = solve(n, circle());
   cout << fixed << setprecision(3) ;</pre>
   gg(ans);
   return 0;
```

### 3.8 Points Inside Polygon

```
S = I + B / 2 - 1
```

### 3.9 Primitives

```
typedef long double ld;
typedef complex<ld> pt;
typedef vector<pt> poly;
#define x real()
#define y imag()
typedef pair<pt, pt> line;
// +, -, * scalar well defined
const ld EPS = 1e-12;
const ld PI = acos(-1);
const int ON = 0, LEFT = 1, RIGHT = -1, BACK = -2, FRONT =
2, IN = 3, OUT = -3;
inline bool Lss(ld a, ld b){ return a - b < -EPS; }</pre>
```

```
inline bool Grt(ld a, ld b){ return a - b > +EPS: }
inline bool Leg(ld a, ld b){ return a - b < +EPS; }</pre>
inline bool Geq(ld a, ld b){ return a - b > -EPS; }
inline bool Equ(ld a. ld b){ return abs(a-b) < EPS: }</pre>
bool bvX(const pt &a. const pt &b)
if (Equ(a.x, b.x)) return Lss(a.v, b.v);
return Lss(a.x, b.x);
bool byY(const pt &a, const pt &b){
if (Equ(a.v, b.v)) return Lss(a.x, b.x):
return Lss(a.y, b.y);
struct cmpXY{ inline bool operator ()(const pt &a, const pt
    &b)const { return byX(a, b); } };
struct cmpYX{ inline bool operator ()(const pt &a, const pt
    &b)const { return byY(a, b); } };
bool operator < (const pt &a, const pt &b){ return byX(a, b)
    : }
istream& operator >> (istream& in, pt p){ld valx,valy; in>>
    valx>>valy; p={valx,valy}; return in;}
ostream& operator << (ostream& out, pt p){out<<p.x<<' ', '<p.y
    ; return out;}
ld dot(pt a, pt b){return (conj(a) * b).x;}
ld cross(pt a, pt b){return (coni(a) * b).v:}
ld disSQ(pt a, pt b){return norm(a - b);}
ld dis(pt a, pt b){return abs(a - b);}
ld angleX(pt a, pt b){return arg(b - a);}
ld slope(pt a, pt b){return tan(angleX(a,b));}
//polar(r,theta) -> cartesian
pt rotate(pt a, ld theta){return a * polar((ld)1, theta);}
pt rotatePiv(pt a, ld theta, pt piv){return (a - piv) *
    polar((ld)1, theta) + piv:}
ld angleABC(pt a, pt b, pt c){return abs(remainder(arg(a-b)
    - arg(c-b), 2.0 * PI));}
pt proj(pt p, pt v){return v * dot(p,v) / norm(v);}
pt projPtLine(pt a, line 1){return proj(a - 1.first, 1.second
    -l.first)+l.first:}
ld disPtLine(pt p, line l){return dis(p-1.first, proj(p-1.
    first.l.second-l.first));}
int relpos(pt a, pt b, pt c) //c to a-b
b = b-a, c = c-a;
if (Grt(cross(b,c), 0)) return LEFT;
if (Lss(cross(b,c), 0)) return RIGHT:
if (Lss(dot(b,c), 0)) return BACK:
```

```
if (Lss(dot(b,c), abs(b))) return FRONT:
return ON:
}
int relpos(line 1, pt b){return relpos(l.first, l.second, b)
pair<pt,bool> intersection(line a, line b)
ld c1 = cross(b.first - a.first, a.second - a.first);
ld c2 = cross(b.second - a.first, a.second - a.first);
if (Equ(c1,c2))
 return {{-1,-1},false}:
return {(c1 * b.second - c2 * b.first) / (c1 - c2), true};
bool intersect(line a. line b)
pair<pt, bool> ret = intersection(a,b);
if (!ret.second) return false:
if (relpos(a, ret.first) == ON and relpos(b, ret.first) ==
 return true;
return false:
bool isconvex(poly &pl)
int n = pl.size();
bool neg = false, pos = false;
for (int i=0:i<n:i++)</pre>
 int rpos = relpos(pl[i], pl[(i+1)%n], pl[(i+2)%n]);
 if (rpos == LEFT) pos = true;
 if (rpos == RIGHT) neg = true;
return !(neg&pos);
int crossingN(poly &pl, pt a)
int n = pl.size();
pt b = a:
for (pt p:pl)
 b.real(max(b.x,p.y));
int cn = 0:
for (int i=0:i<n:i++)</pre>
 pt p = pl[i], q=pl[(i+1)%n];
 if (intersect({a,b},{p,q}) && (relpos({p,q},a)!= RIGHT ||
      relpos({p,q},b) != RIGHT))
  cn ++:
return cn;
```

```
int pointInPoly(poly &pl, pt p)
int n = pl.size():
for (int i=0;i<n;i++)</pre>
 if (relpos(pl[i], pl[(i+1)%n], p) == ON)
return crossingN(pl,p)%2? IN : OUT;
poly getHull(poly &pl, bool lower)
sort(pl.begin(), pl.end(), byX);
poly res;
int n = res.size();
for (auto p : pl)
 while (n \ge 2 \&\& relpos(res[n-2], res[n-1], p) == (lower?
      RIGHT : LEFT))
  res.pop_back(), n--;
 res.push_back(p), n++;
return res;
pair<pt, pt> nearestPair(poly &pl)
int n = pl.size():
sort(pl.begin(), pl.end(), byX);
multiset<pt. cmpYX> s:
ld rad = abs(pl[1] - pl[0]);
pair<pt, pt> res = {pl[0], pl[1]};
int 1 = 0, r = 0:
for (int i=0;i<n;i++)</pre>
 while (1<r && Geq(pl[i].x - pl[l].x, rad))</pre>
  s.erase(pl[1++]);
 while (r<1 && Leq(pl[r].x, pl[i].x))</pre>
  s.insert(pl[r++]);
 for (auto it = s.lower_bound(pt(pl[i].x, pl[i].y-rad)); it
       != s.end(): it++)
  if (Grt(it->y, pl[i].y+rad))
   break:
  ld cur = abs(pl[i] - (*it));
  if (Lss(cur, rad))
   rad = cur, res = {*it, pl[i]};
return res;
```

```
typedef struct circle{
pt c:
ld r;
} cir:
//number of common tangent lines
int tangentCnt(cir c1, cir c2)
ld d= abs(c1.c-c2.c):
if (Grt(d, c1.r+c2.r)) return 4: //outside
if (Equ(d, c1.r+c2.r)) return 3; //tangent outside
if (Lss(d, c1.r+c2.r) && Grt(d, abs(c1.r-c2.r))) return 2;
     //interfere
if (Equ(d, abs(c1.r-c2.r))) return 1; //tangent inside
return 0://inside
line intersection(line 1, cir c)
ld dis = disPtLine(c.c. 1):
ld d = sqrt(c.r*c.r - dis*dis);
pt p = projPtLine(c.c, 1);
pt vec = (1.second-1.first)/abs(1.second - 1.first);
return {p + d * vec, p - d * vec};
 0 = other is inside this, zero point
 1 = other is tangent inisde of this, one point
  2 = other is intersect with this, two point
  3 = other is tangent outside of this, one point
  4 = other is outside of this, zero point
pair<int. vector<pt> > intersect(cir c. cir other) {
1d r = c.r:
pt o = c.c;
vector<pt> v;
ld sumr = other.r + r;
ld rr = r - other.r:
ld d = dis(o, other.c);
ld a = (r*r - other.r*other.r + d*d)/(2*d):
ld h = sqrt(r*r-a*a);
pt p2 = a * (other.c - o) / d;
if(Equ(sumr - d, 0)) {
 v.push_back(p2);
 return make_pair(3, v);
if(Equ(rr - d, 0)) {
```

```
13
```

```
v.push back(p2):
 return make_pair(1, v);
 if(d <= rr)
 return make_pair(0, v);
 if(d >= sumr)
 return make_pair(4, v);
 pt p3(p2.x + h*(other.c.y - o.y)/d, p2.y - h*(other.c.x - o
 pt p4(p2.x - h*(other.c.y - o.y)/d, p2.y + h*(other.c.x - o.y)/d
      (x)/d:
 v.push back(p3):
 v.push_back(p4);
 return make_pair(2, v);
ld arcarea(ld l, ld r, ld R){//circle with radius(r)
    intersect with circle with radius (R) and distance
    between centers equal to (d)
 ld cosa = (1*1 + r*r - R*R)/(2.0*r*1):
ld a = acos(cosa):
 return r*r*(a - \sin(2*a)/2);
```

### 3.10 Rotating Calipers

```
vector<pair<pt, pt>> get_antipodals(poly &p)
int n = p.size();
sort(p.begin(), p.end(), byX);
vector <pt> U, L;
for (int i = 0; i < n; i++){</pre>
 while (U.size() > 1 && relpos(U[U.size()-2], U[U.size()
      -1], p[i]) != LEFT)
  U.pop back():
 while (L.size() > 1 && relpos(L[L.size()-2], L[L.size()
      -1], p[i]) != RIGHT)
  L.pop back():
 U.push_back(p[i]);
 L.push back(p[i]):
vector <pair<pt, pt>> res;
int i = 0, j = L.size()-1;
while (i+1 < (int)U.size() || j > 0){
 res.push_back({U[i], L[j]});
 if (i+1 == (int)U.size())
 else if (j == 0)
 i++:
 else if (cross(L[i]-L[i-1], U[i+1]-U[i]) >= 0) i++;
```

```
else
  j--;
}
return res;
}
```

pt bary(pt A, pt B, pt C, ld a, ld b, ld c) {

return (A\*a + B\*b + C\*c) / (a + b + c):

### 3.11 Triangles

```
pt centroid(pt A, pt B, pt C) {
   // geometric center of mass
   return bary(A, B, C, 1, 1, 1);
pt circumcenter(pt A, pt B, pt C) {
   // intersection of perpendicular bisectors
   double a = norm(B - C), b = norm(C - A), c = norm(A - B);
   return bary(A, B, C, a*(b+c-a), b*(c+a-b), c*(a+b-c));
pt incenter(pt A, pt B, pt C) {
   // intersection of internal angle bisectors
   return bary(A, B, C, abs(B-C), abs(A-C), abs(A-B));
pt orthocenter(pt A, pt B, pt C) {
   // intersection of altitudes
   double a = norm(B - C), b = norm(C - A), c = norm(A - B);
   return bary(A, B, C, (a+b-c)*(c+a-b), (b+c-a)*(a+b-c), (c
        +a-b)*(b+c-a));
pt excenter(pt A, pt B, pt C) {
   // intersection of two external angle bisectors
   double a = abs(B - C), b = abs(A - C), c = abs(A - B):
   return bary(A, B, C, -a, b, c);
   //// NOTE: there are three excenters
   // return bary(A, B, C, a, -b, c);
   // return bary(A, B, C, a, b, -c);
```

### 3.12 Useful Geometry Facts

```
Area of triangle with sides a, b, c: sqrt(S *(S-a)*(S-b)*(S-c)) where S = (a+b+c)/2
```

```
Area of equilateral triangle: s^2 * sqrt(3) / 4 where is side lenght

Pyramid and cones volume: 1/3 area(base) * height

if p1=(x1, x2), p2=(x2, y2), p3=(x3, y3) are points on circle, the center is

x = -((x2^2 - x1^2 + y2^2 - y1^2)*(y3 - y2) - (x2^2 - x3^2 + y2^2 - y3^2)*(y1 - y2)) / (2*(x1 - x2)*(y3 - y2) - 2*(x3 - x2)*(y1 - y2))

y = -((y2^2 - y1^2 + x2^2 - x1^2)*(x3 - x2) - (y2^2 - y3^3 + x2^2 - x3^2)*(x1 - x2)) / (2*(y1 - y2)*(x3 - x2) - 2*(y3 - y2)*(x1 - x2))
```

### 4 Graph

### 4.1 2-SAT

```
vector\langle int \rangle adj[2 * N], jda[2 * N], top;
bool mark[2 * N];
int c[2 * N]:
void add_clause(int x, int y) {
adj[x ^ 1].pb(y);
adj[y ^ 1].pb(x);
jda[y].pb(x^1);
jda[x].pb(y ^ 1);
void dfs(int u) {
mark[u] = 1:
for(auto v : adj[u]) if(!mark[v]) dfs(v);
top.pb(u);
void sfd(int u. int col) {
c[u] = col:
for(auto v : jda[u]) if(!c[v]) sfd(v, col);
vector<int> two sat(int n) {
memset(mark, 0, sizeof mark):
memset(c, 0, sizeof c):
top.clear():
for(int i = 2; i < 2 * n + 2; i++) if(!mark[i]) dfs(i);</pre>
int cnt = 1:
while(top.size()) {
 int x = top.back(); top.pop_back();
 if(!c[x]) sfd(x, cnt++);
vector<int> ans, ans1;
```

```
ans1.pb(-1);
for(int i = 1; i <= n; i++) {
  if(c[2 * i] == c[2 * i + 1]) return ans1;
  if(c[2 * i] > c[2 * i + 1]) ans.pb(i);
}
return ans;
}
```

### 4.2 Biconnected-Component

```
vector<int> adi[N]:
bool vis[N]:
int dep[N], par[N], lowlink[N];
vector<vector<int> > comp;
stack<int> st;
void dfs(int u, int depth = 0, int parent = -1){
 vis[u] = true:
 dep[u] = depth;
 par[u] = parent;
 lowlink[u] = depth:
 st.push(u);
 for (int i = 0; i < adj[u].size(); i++){</pre>
 int v = adj[u][i];
 if (!vis[v])
  dfs(v, depth + 1, u);
  lowlink[u] = min(lowlink[u], lowlink[v]):
 }
  lowlink[u] = min(lowlink[u], dep[v]);
 if (lowlink[u] == dep[u] - 1){
 comp.push_back(vector<int>());
 while (st.top() != u)
  comp.back().push_back(st.top());
  st.pop();
 comp.back().push_back(u);
 st.pop();
 comp.back().push_back(par[u]);
}
void bicon(int n){
for (int i = 0; i < n; i++)</pre>
 if (!vis[i])
  dfs(i);
}
```

### 4.3 Directed Minimum Spanning Tree MN

/\*

```
GETS:
 call make_graph(n) at first
 you should use add_edge(u,v,w) and
 add pair of vertices as edges (vertices are 0..n-1)
 output of dmst(v) is the minimum arborescence with root v
      in directed graph
 (-1 if it hasn't a spanning arborescence with root v)
O(mn)
const int INF = 2e7:
struct MinimumAborescense{
int n:
struct edge {
 int src, dst;
 int weight:
vector<edge> edges;
void make_graph(int _n) {
 n=_n;
 edges.clear();
void add_edge(int u, int v, int w) {
 edges.push_back({u, v, w});
int dmst(int r) {
 int N = n:
 for (int res = 0; ;) {
  vector<edge> in(N, {-1,-1,(int)INF});
  vector<int> C(N, -1):
  for (auto e: edges)
  if (in[e.dst].weight > e.weight)
   in[e.dst] = e;
  in[r] = \{r, r, 0\};
  for (int u = 0; u < N; ++u) { // no comming edge ==> no
       aborescense
   if (in[u].src < 0) return -1:</pre>
   res += in[u].weight:
  vector<int> mark(N, -1); // contract cycles
  int index = 0;
  for (int i = 0: i < N: ++i) {</pre>
   if (mark[i] != -1) continue;
   int u = i:
   while (mark[u] == -1) {
   mark[u] = i;
```

```
u = in[u].src:
   if (mark[u] != i || u == r) continue;
   for (int v = in[u].src: u != v: v = in[v].src) C[v] =
   C[u] = index++:
  if (index == 0) return res; // found arborescence
  for (int i = 0: i < N: ++i) // contract</pre>
   if (C[i] == -1) C[i] = index++;
  vector<edge> next:
  for (auto &e: edges)
   if (C[e.src] != C[e.dst] && C[e.dst] != C[r])
    next.push_back({C[e.src], C[e.dst], e.weight - in[e.dst
         ].weight});
  edges.swap(next);
  N = index; r = C[r];
}
};
```

### 

```
call make_graph(n) at first
you should use add_edge(u,v,w) and
add pair of vertices as edges (vertices are 0..n-1)
output of dmst(v) is the minimum arborescence with root v in
     directed graph
(INF if it hasn't a spanning arborescence with root v)
O(mlogn)
const int INF = 2e7;
struct MinimumAborescense{
struct edge {
 int src, dst, weight;
struct union_find {
 vector<int> p;
 union_find(int n) : p(n, -1) { };
 bool unite(int u, int v) {
 if ((u = root(u)) == (v = root(v))) return false;
  if (p[u] > p[v]) swap(u, v);
  p[u] += p[v]; p[v] = u;
```

```
return true:
bool find(int u, int v) { return root(u) == root(v); }
int root(int u) { return p[u] < 0 ? u : p[u] = root(p[u]);</pre>
int size(int u) { return -p[root(u)]: }
};
struct skew_heap {
struct node {
 node *ch[2];
 edge kev:
 int delta:
} *root;
skew_heap() : root(0) { }
void propagate(node *a) {
 a->key.weight += a->delta;
 if (a->ch[0]) a->ch[0]->delta += a->delta:
 if (a->ch[1]) a->ch[1]->delta += a->delta;
 a \rightarrow delta = 0:
}
node *merge(node *a, node *b) {
 if (!a || !b) return a ? a : b;
 propagate(a); propagate(b);
 if (a->key.weight > b->key.weight) swap(a, b);
 a->ch[1] = merge(b, a->ch[1]);
 swap(a->ch[0], a->ch[1]);
 return a:
 void push(edge key) {
 node *n = new node():
 n->ch[0] = n->ch[1] = 0;
 n->key = key; n->delta = 0;
 root = merge(root, n);
void pop() {
 propagate(root):
 node *temp = root:
 root = merge(root->ch[0], root->ch[1]);
edge top() {
 propagate(root):
 return root->key;
bool empty() {
 return !root;
void add(int delta) {
 root->delta += delta:
void merge(skew_heap x) {
```

```
root = merge(root, x.root):
 }
 };
 vector<edge> edges:
 void add_edge(int src, int dst, int weight) {
 edges.push back({src. dst. weight}):
 int n:
 void make_graph(int _n) {
 n = _n;
 edges.clear():
 int dmst(int r) {
 union find uf(n):
 vector<skew_heap> heap(n);
  for (auto e: edges)
  heap[e.dst].push(e);
  double score = 0:
 vector<int> seen(n, -1);
  seen[r] = r:
  for (int s = 0; s < n; ++s) {
  vector<int> path:
   for (int u = s; seen[u] < 0;) {</pre>
   path.push_back(u);
   seen[u] = s:
   if (heap[u].empty()) return INF;
   edge min_e = heap[u].top();
   score += min e.weight:
   heap[u].add(-min_e.weight);
   heap[u].pop():
   int v = uf.root(min_e.src);
   if (seen[v] == s) {
    skew heap new heap:
    while (1) {
     int w = path.back();
     path.pop back():
     new_heap.merge(heap[w]);
     if (!uf.unite(v, w)) break;
    heap[uf.root(v)] = new_heap;
    seen[uf.root(v)] = -1:
   u = uf.root(v):
 return score:
};
```

### 4.5 Ear Decomposition

Solution:

- 1- Find a spanning tree of the given graph and choose a root for the tree.
- 2- Determine, for each edge uv that is not part of the tree, the distance between the root and the lowest common ancestor of u and v.
- 3- For each edge uv that is part of the tree, find the corresponding "master edge", a non-tree edge wx such that the cycle formed by adding wx to the tree passes through uv and such that, among such edges, w and x have a lowest common ancestor that is as close to the root as possible (with ties broken by edge identifiers).
- 4- Form an ear for each non-tree edge, consisting of it and the tree edges for which it is the master, and order the ears by their master edges' distance from the root (with the same tie-breaking rule).

### 4.6 Edmond-Blossom

```
// Order: M * Sqrt(N)
// Edges of 1-based. add_edge for adding edges and calc for
    calculating matching
// output is in match array (match[i] = 0 if i isn't in
    matching)
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
    count()):
template<int SZ> struct UnweightedMatch {
int match[SZ], N;
vector<int> adi[SZ]:
void add_edge(int u, int v) {
 adi[u].pb(v):
 adj[v].pb(u);
queue<int> q;
int par[SZ], vis[SZ], orig[SZ], aux[SZ];
void augment(int u, int v) { // toggle edges on u-v path
 while (1) { // one more matched pair
  int pv = par[v], nv = match[pv];
  match[v] = pv; match[pv] = v;
  v = nv: if (u == pv) return:
```

```
}
int lca(int u, int v) { // find LCA of supernodes in O(dist
static int t = 0;
for (++t::swap(u.v)) {
 if (!u) continue;
 if (aux[u] == t) return u; // found LCA
 aux[u] = t; u = orig[par[match[u]]];
}
void blossom(int u, int v, int a) { // go other way
for (; orig[u] != a; u = par[v]) { // around cycle
 par[u] = v; v = match[u]; // treat u as if vis[u] = 1
 if (vis[v] == 1) vis[v] = 0, q.push(v);
 orig[u] = orig[v] = a; // merge into supernode
}
bool bfs(int u) { // u is initially unmatched
for(int i = 0: i < N + 1: i++)</pre>
 par[i] = 0, vis[i] = -1, orig[i] = i;
q = queue<int>();
vis[u] = 0;
q.push(u);
while (q.size()) { // each node is pushed to q at most
 int v = q.front(); q.pop(); // 0 -> unmatched vertex
 for (int x : adi[v]) {
  if (vis[x] == -1) \{ // neither of x, match[x] visited
   vis[x] = 1; par[x] = v;
   if (!match[x])
    return augment(u,x),1;
   vis[match[x]] = 0;
   g.push(match[x]);
  } else if (vis[x] == 0 && orig[v] != orig[x]) {
   int a = lca(orig[v],orig[x]); // odd cycle
   blossom(x,v,a), blossom(v,x,a);
  } // contract O(n) times
}
return 0:
int calc(int N) { // rand matching -> constant improvement
for(int i = 0: i <= N: i++)</pre>
 match[i] = aux[i] = 0:
int ans = 0; vector<int> V(N); iota(V.begin(), V.end(),1);
```

```
shuffle(V.begin(), V.end(),rng); // find rand matching
for (int x : V) {
   if (!match[x]) {
      for (int y : adj[x]) {
        if (!match[y]) {
        match[x] = y, match[y] = x; ++ans;
        break;
      }
   }
   }
}
for(int i = 1; i <= N; i++)
   if (!match[i] && bfs(i))
   ++ans;
   return ans;
}
</pre>
```

### 4.7 Flow-Dinic

```
//Order : General: mn^2. Bipartite: mn^0.5. Zero-One: mn
    ^(2/3)
const int maxN = 1000, maxE = 2 * 1e5 + 10:
int from[maxE], to[maxE], cap[maxE], prv[maxE], head[maxN],
    pt[maxN], ec;
void addEdge(int u, int v, int uv, int vu = 0){
from[ec] = u, to[ec] = v, cap[ec] = uv, prv[ec] = head[u],
head[u] = ec++:
from[ec] = v. to[ec] = u. cap[ec] = vu. prv[ec] = head[v].
head[v] = ec++:
int lv[maxN]. a[maxN]:
bool bfs(int source, int sink){
memset(lv, 31, sizeof(lv));
int h = 0, t = 0:
lv[source] = 0;
a[t++] = source:
while (t-h){
 int v = a[h++]:
 for (int e = head[v]; ~e; e = prv[e])
  if (cap[e] && lv[v] + 1 < lv[to[e]]){</pre>
   lv[to[e]] = lv[v] + 1:
   q[t++] = to[e];
```

```
return lv[sink] < 1e8:
int dfs(int v, int sink, int f = 1e9){
if (v == sink || f == 0)
return f:
int ret = 0:
for (int &e = pt[v]; ~e; e = prv[e])
 if (lv[v]+1 == lv[to[e]]){
 int x = dfs(to[e], sink, min(f, cap[e]));
  cap[e] -= x;
  cap[e^1] += x:
  ret += x:
  f = x;
  if (!f)
  break:
return ret:
int dinic(int source, int sink){
memset(prv, -1, sizeof prv):
memset(head, -1, sizeof head);
int ret = 0;
while (bfs(source, sink)){
 memcpy(pt, head, sizeof(head));
 ret += dfs(source, sink);
return ret:
```

### 4.8 Gomory-Hu

```
bool mark[N];
int p[N], w[N];
void gfs(int u) {
mark[u] = 1;
for(int e = head[u]; e != -1; e = prv[e])
if(!mark[to[e]] && cap[e])
 gfs(to[e]):
//edges is one-directed. Order: O(n * flow)
vector<pair<int, pii>> gomory_hu(int n, vector<pair<int, pii</pre>
    >> edges) {
for(int i = 1; i <= n; i++) p[i] = 1;</pre>
memset(w, 0, sizeof w):
p[1] = 0;
for(int i = 2; i <= n; i++) {</pre>
 memset(head, -1, sizeof head):
 ec = 0:
```

```
for(auto u : edges) add_edge(u.S.F, u.S.S, u.F);
w[i] = dinic(i, p[i]);
memset(mark, 0, sizeof mark);
gfs(i);
for(int j = i + 1; j <= n; j++)
    if(mark[j] && p[j] == p[i])
    p[j] = i;
    if(p[p[i]] && mark[p[p[i]]]) {
        int pi = p[i];
        swap(w[i], w[pi]);
    p[i] = p[pi];
    p[pi] = i;
    }
}
vector<pair<int, pii>> tree;
for(int i = 1; i <= n; i++) if(p[i]) tree.pb({w[i], {i, p[i]}};
    return tree;
}</pre>
```

### 4.9 Hungarian

```
const int N = 2002:
const int INF = 1e9:
int hn, weight[N][N]; //hn should contain number of vertices
     in each part. weight must be positive.
int x[N], y[N]; //initial value doesn't matter.
int hungarian() // maximum weighted perfect matching O(n^3)
int n = hn:
int p, q;
vector<int> fx(n, -INF), fy(n, 0);
fill(x, x + n, -1):
fill(y, y + n, -1);
for (int i = 0: i < n: ++i)
 for (int j = 0; j < n; ++j)
 fx[i] = max(fx[i], weight[i][j]);
for (int i = 0: i < n: ) {
 vector<int> t(n, -1), s(n+1, i);
 for (p = 0, q = 1; p < q && x[i] < 0; ++p) {
  int k = s[p];
  for (int i = 0: i < n && x[i] < 0: ++i)
   if (fx[k] + fy[j] == weight[k][j] && t[j] < 0) {
    s[q++] = y[j], t[j] = k;
    if (y[j] < 0) // match found!</pre>
     for (int p = j; p \ge 0; j = p)
```

```
y[j] = k = t[j], p = x[k], x[k] = j;
 }
 if (x[i] < 0) {
  int d = INF:
  for (int k = 0: k < a: ++k)
  for (int j = 0; j < n; ++j)
   if (t[i] < 0) d = min(d, fx[s[k]] + fy[i] - weight[s[k]]
  for (int j = 0; j < n; ++j) fy[j] += (t[j] <0? 0: d);
  for (int k = 0: k < q: ++k) fx[s[k]] -= d:
 } else ++i:
int ret = 0:
for (int i = 0; i < n; ++i) ret += weight[i][x[i]];</pre>
int main() {
int n. e: cin >> n >> e:
for (int i=0; i<e; i++)</pre>
 int u, v; cin >> u >> v;
 --u: --v:
 cin >> weight[u][v];
cout << hungarian() << '\n':</pre>
return 0;
```

### 4.10 Min-Cost-Max-Flow

```
const int N = 810, E = N * N, INF = 1e9;
int n, ed = 0, from[E], to[E], cap[E], head[N], nex[E], par[
    N];
ld dis[N], cost[E];

void add_edge(int u, int v, int c, ld co)
{
  from[ed] = u, to[ed] = v, cap[ed] = c, cost[ed] = co , nex[
    ed] = head[u], head[u] = ed ++;
  from[ed] = v, to[ed] = u, cap[ed] = 0, cost[ed] = -co, nex[
    ed] = head[v], head[v] = ed ++;
}

pair<int, ld> spfa(int sink, int source)
{
```

```
for(int i=0: i<N: i++)dis[i] = INF:</pre>
memset(mark, 0, sizeof mark);
memset(par, -1, sizeof par);
queue<int> q;
dis[source] = 0. mark[source] = true:
q.push(source);
while(q.size())
 int v = q.front(); q.pop();
 mark[v] = false:
 for(int e = head[v]: e != -1: e = nex[e])
  if(cap[e] && dis[to[e]] > dis[v] + cost[e])
   dis[to[e]] = dis[v] + cost[e];
   par[to[e]] = e:
   if(!mark[to[e]])q.push(to[e]), mark[to[e]] = true;
 }
}
int curr = sink:
if(dis[curr] == INF)return make_pair(0, 0);
ld res = 0:
int flow = INF:
while(curr != source)
 flow = min(flow, cap[par[curr]]);
 curr = from[par[curr]]:
curr = sink:
while(curr != source)
 res += cost[par[curr]]:
 cap[par[curr]] -= flow;
 cap[par[curr] ^ 1] += flow:
 curr = from[par[curr]];
return make_pair(flow, res);
pair<int, ld> MinCostMaxFlow(int sink, int source)
```

```
18
```

```
{
  int flow = 0;
  pair<int, ld> f = {INF, 0};
  ld Cost = 0;

while(f.F)
  {
    f = spfa(sink, source);
    flow += f.F;
    Cost += f.F * f.S;
}

return make_pair(flow, Cost);
}
```

### 5 Number Theory

### 5.1 Chineese Reminder Theorem

```
#define lcm LLLCCM
11 GCD(11 a, 11 b) { return (b == 0) ? a : GCD(b, a % b); }
inline 11 LCM(11 a. 11 b) { return a / GCD(a. b) * b: }
inline 11 normalize(11 x, 11 mod) { x %= mod; if (x < 0) x</pre>
     += mod: return x: }
struct GCD_type { 11 x, y, d; };
GCD_type ex_GCD(11 a, 11 b){
if (b == 0) return {1, 0, a};
 GCD_type pom = ex_GCD(b, a % b);
 return {pom.v. pom.x - a / b * pom.v. pom.d}:
const int N = 2:
11 r[N], n[N], ans, 1cm;
// t: number of equations.
// r: reminder array, n: mod array
// returns {reminder. lcm}
pair <11, 11> CRT(11* r, 11 *n, int t) {
 for(int i = 0: i < t: i++)</pre>
 normalize(r[i], n[i]);
 ans = r[0];
 lcm = n[0]:
 for(int i = 1; i < t; i++){</pre>
 auto pom = ex_GCD(lcm, n[i]);
 11 x1 = pom.x;
```

### 5.2 Miller Robin

```
//with probability (1/4) iter, we might make mistake in our
//we have false positive here.
using u64 = uint64 t:
using u128 = __uint128_t;
using namespace std;
u64 binpower(u64 base, u64 e, u64 mod) {
1164 \text{ result} = 1:
base %= mod;
while (e) {
 if (e & 1)
 result = (u128)result * base % mod:
 base = (u128)base * base % mod:
 e >>= 1:
return result:
bool check_composite(u64 n, u64 a, u64 d, int s) {
u64 x = binpower(a, d, n);
if (x == 1 | | x == n - 1)
 return false;
for (int r = 1: r < s: r++) {
 x = (u128)x * x % n:
 if (x == n - 1)
 return false:
return true;
bool MillerRabin(u64 n, int iter=5) { // returns true if n
    is probably prime, else returns false.
if (n < 4)
```

```
return n == 2 || n == 3;
int s = 0;
u64 d = n - 1;
while ((d & 1) == 0) {
    d >>= 1;
    s++;
}

for (int i = 0; i < iter; i++) {
    int a = 2 + rand() % (n - 3);
    if (check_composite(n, a, d, s))
    return false;
}
return true;
}</pre>
```

### 5.3 Most Divisors

```
<= 1e2: 60 with 12 divisors
<= 1e3: 840 with 32 divisors
<= 1e4: 7560 with 64 divisors
<= 1e5: 83160 with 128 divisors
<= 1e6: 720720 with 240 divisors
<= 1e7: 8648640 with 448 divisors
<= 1e8: 73513440 with 768 divisors
<= 1e9: 735134400 with 1344 divisors
<= 1e10: 6983776800 with 2304 divisors
<= 1e11: 97772875200 with 4032 divisors
<= 1e12: 963761198400 with 6720 divisors
<= 1e13: 9316358251200 with 10752 divisors
<= 1e14: 97821761637600 with 17280 divisors
<= 1e15: 866421317361600 with 26880 divisors
<= 1e16: 8086598962041600 with 41472 divisors
<= 1e17: 74801040398884800 with 64512 divisors
<= 1e18: 897612484786617600 with 103680 divisors
```

### 5.4 Number of Primes

```
30: 10

60: 17

100: 25

1000: 168

10000: 1229

100000: 9592

1000000: 78498

10000000: 664579
```

### 6 Numerical

### 6.1 Base Vector **Z**2

```
const int maxL = 61;
struct Base{
ll a[maxL] = {}:
ll eliminate(ll x){
 for(int i=maxL-1; i>=0; --i) if(x >> i & 1) x ^= a[i];
 return x:
void add(ll x){
 x = eliminate(x):
 if(x == 0) return ;
 for(int i=maxL-1: i>=0: --i) if(x >> i & 1) {
  a[i] = x:
  return ;
 }
int size(){
 int cnt = 0:
 for(int i=0; i<maxL; ++i) if(a[i]) ++cnt;</pre>
 return cnt:
11 get_mx() {
 11 x = 0:
 for (int i=maxL-1; i>=0; i--) {
 if(x & (1LL << i)) continue;</pre>
  else x ^= a[i]:
 return x;
};
```

### 6.2 Extended Catalan

number of ways for going from 0 to A with k moves without going to -B:

$$\binom{k}{\frac{A+k}{2}} - \binom{k}{\frac{2B+A+k}{2}}$$

### 6.3 FFT

const int LG = 20; // IF YOU WANT TO CONVOLVE TWO ARRAYS OF
 LENGTH N AND M CHOOSE LG IN SUCH A WAY THAT 2LG > n + m

```
const int MAX = 1 << LG:</pre>
#define M PI acos(-1)
struct point{
double real, imag:
point(double _real = 0.0, double _imag = 0.0){
 real = real:
 imag = _imag;
point operator + (point a, point b){
return point(a.real + b.real, a.imag + b.imag);
point operator - (point a, point b){
return point(a.real - b.real, a.imag - b.imag);
point operator * (point a, point b){
return point(a.real * b.real - a.imag * b.imag, a.real * b.
 imag + a.imag * b.real);
void fft(point *a, bool inv){
for (int mask = 0; mask < MAX; mask++){</pre>
 for (int i = 0; i < LG; i++)</pre>
  if ((1 << i) & mask)</pre>
   rev |= (1 << (LG - 1 - i)):
 if (mask < rev)</pre>
  swap(a[mask], a[rev]);
for (int len = 2; len <= MAX; len *= 2){</pre>
 double ang = 2.0 * M_PI / len;
 if (inv)
  ang *= -1.0;
 point wn(cos(ang), sin(ang));
 for (int i = 0: i < MAX: i += len){</pre>
  point w(1.0, 0.0):
  for (int j = 0; j < len / 2; j++){</pre>
   point t1 = a[i + j] + w * a[i + j +
    len / 21:
   point t2 = a[i + i] - w * a[i + i +
    len / 2];
   a[i + i] = t1:
   a[i + j + len / 2] = t2;
   w = w * wn;
 }
 for (int i = 0; i < MAX; i++){</pre>
```

```
a[i].real /= MAX;
a[i].imag /= MAX;
}
```

### 6.4 Gaussian Elimination

```
const int N = 505, MOD = 1e9 + 7:
typedef vector <11> vec;
11 pw(ll a, ll b) {
if(!b)
return 1:
11 x = pw(a, b/2);
return x * x % MOD * (b % 2 ? a : 1) % MOD:
11 inv(11 x) { return pw(x, MOD - 2); }
//matrix * x = ans
vec solve(vector<vec> matrix, vec ans) {
int n = matrix.size(), m = matrix[0].size();
for (int i=0; i<n; i++)</pre>
 matrix[i].pb(ans[i]);
vector <int> ptr;
ptr.resize(n):
int i = 0, j =0;
while(i < n and j < m) {</pre>
 int ind = -1;
 for(int row = i: row < n: row++)</pre>
 if(matrix[row][i])
  ind = row;
 if(ind == -1) {
  j++;
  continue :
 matrix[i].swap(matrix[ind]);
 ll inverse = inv(matrix[i][i]):
 for(int row = i + 1: row < n: row++) {
  11 z = matrix[row][j] * inverse % MOD;
  for(int k = 0; k <= m; k++)</pre>
   matrix[row][k] = (matrix[row][k] % MOD - matrix[i][k]*z %
         MOD + MOD) % MOD:
 ptr[i] = j;
i ++;
```

```
j ++;
vector <11> sol:
if(i != n) {
 for (int row=i; row<n; row++)</pre>
  if(matrix[row][m] != 0)
   return sol; //without answer;
sol.resize(m):
for (int i=0: i<m: i++)</pre>
 sol[i] = 0;
for (int row=i-1: row>=0: row--){
 int j = ptr[row];
 sol[j] = matrix[row][m] * inv(matrix[row][j]) % MOD;
 for (int c=row-1; c>=0; c--)
  matrix[c][m] += (MOD - sol[i] * matrix[c][i] % MOD).
       matrix[c][m] %= MOD:
return sol:
int main() {
int n, m; cin >> n >> m;
vector <vec> A:
for (int i=0: i<n: i++)</pre>
 vec B:
 for (int j=0; j<m; j++)</pre>
  11 x: cin >> x:
  B.push_back(x);
 A.push back(B):
for (int i=0: i<n: i++)</pre>
 ll v; cin >> v;
 ans.pb(y);
vec sol = solve(A. ans):
for (auto X : sol)
 cout << X << ' ':
cout << endl;</pre>
```

### 6.5 General Linear Recursion

```
const int maxL = 20: // IF YOU WANT TO CONVOLVE TWO ARRAYS
    OF LENGTH N AND M CHOOSE LG IN SUCH A WAY THAT 2LG > n
const int maxN = 1 << maxL. MOD = 998244353:</pre>
#define M PI acos(-1)
int root [maxL + 2] = {0,998244352,86583718,372528824,
69212480.87557064.15053575.57475946.15032460.
4097924,1762757,752127,299814,730033,227806,
42058,44759,8996,2192,1847,646,42};
int bpow(int a, int b){
int ans = 1:
while (b){
 if (b & 1)
 ans = 1LL * ans * a % MOD:
 a = 1LL * a * a % MOD;
return ans;
void ntt(vector<int> &a, bool inv){
int LG = 0, z = 1, MAX = a.size();
while(z != MAX) z *= 2, LG ++;
int ROOT = root[LG]:
for (int mask = 0; mask < MAX; mask++){</pre>
 int rev = 0:
 for (int i = 0: i < LG: i++)</pre>
 if ((1 << i) & mask)</pre>
  rev |= (1 << (LG - 1 - i));
 if (mask < rev)</pre>
  swap(a[mask], a[rev]);
for (int len = 2; len <= MAX; len *= 2){</pre>
 int wn = bpow(ROOT, MAX / len);
 if (inv)
  wn = bpow(wn, MOD - 2);
 for (int i = 0; i < MAX; i += len){</pre>
  int w = 1;
  for (int j = 0; j < len / 2; j++){
  int l = a[i + i]:
   int r = 1LL * w * a[i + j + len / 2] %
   a[i + j] = (1 + r);
   a[i + j + len / 2] = 1 - r + MOD;
```

```
if (a[i + i] >= MOD)
    a[i + i] -= MOD:
   if (a[i + j + len / 2] >= MOD)
    a[i + j + len / 2] -= MOD;
   w = 1LL * w * wn % MOD;
 }
}
if (inv){
 int x = bpow(MAX, MOD - 2);
 for (int i = 0: i < MAX: i++)</pre>
 a[i] = 1LL * a[i] * x % MOD:
int ans[maxN]. bb[maxN]:
//ans[i] = sum i=1^i b i * ans[i - i], ans[0] = 1:
void solve(int 1. int r) {
if(r - 1 == 1) return ;
int mid = (1 + r)/2:
solve(1, mid);
vector <int> a, b;
for (int i=1; i<r; i++) {</pre>
if(i < mid) a.pb(ans[i]);</pre>
 else a.pb(0);
 b.pb(bb[i-l+1]);
for (int i=1; i<r; i++) {</pre>
 a.pb(0);
 b.pb(0);
ntt(a, false);
ntt(b, false);
vector <int> c;
c.resize(a.size()):
for (int i=0; i<2*r-2*1; i++)</pre>
 c[i] = 1LL * a[i] * b[i] % MOD:
ntt(c, true):
for (int i=0: i<r-mid: i++)</pre>
 ans[mid + i] += c[mid - l - 1 + i], ans[mid + i] %= MOD;
```

```
solve(mid, r);
}
int main() {
  int n, m; cin >> n >> m;
  for (int i=1; i<=m; i++)
    cin >> bb[i];
  int k = 1;
  while(k < n) k = 2 * k;

ans[0] = 1;
  solve(0, k);
  for (int i=0; i<n; i++)
    cout << ans[i] << ' ';
  cout << endl;
}</pre>
```

### 6.6 LP Duality

primal: Maximize  $c^T x$  subject to  $Ax \le b, x \ge 0$  dual: Minimize  $b^T y$  subject to  $A^T y \ge c, y \ge 0$ 

### 6.7 NTT

```
const int MOD = 998244353:
const int LG = 16; // IF YOU WANT TO CONVOLVE TWO ARRAYS OF
    LENGTH N AND M CHOOSE LG IN SUCH A WAY THAT 2LG > n + m
const int MAX = (1 << LG):</pre>
const int ROOT = 44759: // ENSURE THAT ROOT2(LG - 1) = MOD -
     1
int bpow(int a, int b){
int ans = 1;
while (b){
 if (b & 1)
  ans = 1LL * ans * a % MOD;
 b >>= 1:
 a = 1LL * a * a % MOD;
return ans;
void ntt(int *a, bool inv){
for (int mask = 0: mask < MAX: mask++){</pre>
 int rev = 0;
 for (int i = 0; i < LG; i++)</pre>
  if ((1 << i) & mask)</pre>
   rev |= (1 << (LG - 1 - i));
```

```
if (mask < rev)
 swap(a[mask], a[rev]);
for (int len = 2: len <= MAX: len *= 2){
int wn = bpow(ROOT, MAX / len);
if (inv)
 wn = bpow(wn, MOD - 2);
 for (int i = 0; i < MAX; i += len){</pre>
 int w = 1:
 for (int j = 0; j < len / 2; j++){
  int l = a[i + i]:
  int r = 1LL * w * a[i + i + len / 2] %
  a[i + j] = (1 + r);
  a[i + j + len / 2] = 1 - r + MOD;
  if (a[i + j] >= MOD)
   a[i + j] -= MOD;
  if (a[i + j + len / 2] >= MOD)
   a[i + i + len / 2] -= MOD:
  w = 1LL * w * wn % MOD:
 }
}
if (inv){
int x = bpow(MAX, MOD - 2);
for (int i = 0; i < MAX; i++)</pre>
 a[i] = 1LL * a[i] * x % MOD;
```

### 6.8 Popular LP

### BellmanFord:

```
maximize X_n X_1 = 0 and for eache edge (v - > u and weight w): X_n - X_v < w
```

### Flow:

```
maximize \Sigma f_{out} (where out is output edges of vertex 1) for each vertex (except 1 and n): \Sigma f_{in} - \Sigma f_{out} = 0 (where in is input edges of v and out is output edges of v)
```

```
Dijkstra(IP):
```

```
minimize \Sigma z_i * w_i for each edge (v->u and weight w): 0 \le z_i \le 1 and for each ST-cut which vertex 1 is in S and vertex n is in T: \Sigma z_e \ge 1 (for each edge e from S to T)
```

### 6.9 Simplex

```
typedef vector <ld> vd;
typedef vector <int> vi;
const ld Eps = 1e-9:
// ax <= b, max(cTx), x >= 0
// O(nm^2)
vd simplex(vector <vd> a, vd b, vd c) {
int n = a.size(), m = a[0].size() + 1, r = n, s = m - 1;
vector \langle vd \rangle d(n + 2, vd(m + 1, 0)); vd x(m - 1);
vi ix(n + m); iota(ix.begin(), ix.end(), 0);
for(int i = 0: i < n: i ++) {
 for(int j = 0; j < m - 1; j ++) d[i][j] = -a[i][j];</pre>
 d[i][m-1] = 1;
 d[i][m] = b[i];
 if(d[r][m] > d[i][m])
 r = i:
for(int j = 0; j < m - 1; j ++) d[n][j] = c[j];
d[n + 1][m - 1] = -1:
while(true) {
 if(r < n) {
  vd su:
  swap(ix[s], ix[r + m]); d[r][s] = 1 / d[r][s];
  for(int j = 0; j <=m; j ++) if(j != s) {</pre>
   d[r][j] *= -d[r][s]; if(d[r][j]) su.pb(j);
  for(int i = 0; i <= n + 1; i ++) if(i != r) {</pre>
   for(int j = 0; j < su.size(); j ++)</pre>
    d[i][su[j]] += d[r][su[j]] * d[i][s];
   d[i][s] *= d[r][s];
 for(int j = 0; j < m; j ++) if(s < 0 || ix[s] > ix[j])
  if(d[n + 1][j] > Eps || d[n + 1][j] > -Eps &&
    d[n][j] > Eps) s = j; if(s < 0) break;
```

```
for(int i = 0; i < n; i ++) if(d[i][s] < -Eps) {
   if(r < 0) {
      r = i;
      continue;
   }
   double e = d[r][m] / d[r][s] - d[i][m] / d[i][s];
   if(e < -Eps || e < Eps && ix[r + m] > ix[i + m]) r = i;
   }
   if(r < 0)
   {return vd();} // Unbounded
}
if(d[n + 1][m] < -Eps) {return vd();}// No solution
for(int i = m; i < n + m; i ++)
   if(ix[i] < m - 1) x[ix[i]] = d[i - m][m];
return x;
}</pre>
```

### 6.10 Stirling Cycle

```
const int mod = 998244353;
const int root = 15311432:
const int root_1 = 469870224;
const int root_pw = 1 << 23;</pre>
const int N = 400004:
vector<int> v[N];
11 modInv(ll a, ll mod = mod){
11 \times 0 = 0, \times 1 = 1, \times 0 = \text{mod}, \times 1 = a;
 while(r1){
 11 q = r0 / r1;
 x0 = q * x1; swap(x0, x1);
 r0 = q * r1: swap(r0, r1):
 return x0 < 0 ? x0 + mod : x0:
void fft (vector<int> &a. bool inv) {
 int n = (int) a.size();
 for (int i=1, j=0; i<n; ++i) {</pre>
 int bit = n \gg 1:
 for (; j>=bit; bit>>=1)
  j -= bit;
 j += bit;
 if (i < j)
  swap (a[i], a[j]);
 for (int len=2; len<=n; len<<=1) {</pre>
 int wlen = inv ? root 1 : root:
 for (int i=len; i<root_pw; i<<=1)</pre>
```

```
wlen = int (wlen * 111 * wlen % mod):
  for (int i=0: i<n: i+=len) {</pre>
  int w = 1:
  for (int j=0; j<len/2; ++j) {</pre>
   int u = a[i+j], v = int (a[i+j+len/2] * 111 * w % mod);
   a[i+i] = u+v < mod ? u+v : u+v-mod:
   a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
   w = int (w * 111 * wlen % mod):
 }
}
if(inv) {
 int nrev = modInv(n, mod);
 for (int i=0: i<n: ++i)</pre>
 a[i] = int (a[i] * 111 * nrev % mod);
void pro(const vector<int> &a, const vector<int> &b, vector<</pre>
    int> &res)
vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
while (n < (int) max(a.size(), b.size())) n <<= 1;</pre>
n <<= 1:
fa.resize (n), fb.resize (n);
fft(fa, false), fft (fb, false);
for (int i = 0; i < n; ++i)</pre>
 fa[i] = 1LL * fa[i] * fb[i] % mod:
fft (fa, true);
res = fa:
int S(int n, int r) {
int nn = 1:
while(nn < n) nn <<= 1;</pre>
for(int i = 0: i < n: ++i) {</pre>
 v[i].push back(i):
 v[i].push_back(1);
for(int i = n; i < nn; ++i) {</pre>
 v[i].push back(1):
for(int j = nn; j > 1; j >>= 1){
 int hn = i \gg 1:
 for(int i = 0; i < hn; ++i){</pre>
 pro(v[i], v[i + hn], v[i]);
/*for (int k=0: k<=r: k++)
 cout << v[0][k] << ' ': cout << '\n':*/
```

```
return v[0][r]:
int fac[N], ifac[N], inv[N];
void prencr(){
fac[0] = ifac[0] = inv[1] = 1;
for(int i = 2: i < N: ++i)</pre>
 inv[i] = mod - 1LL * (mod / i) * inv[mod % i] % mod;
for(int i = 1; i < N; ++i){fac[i] = 1LL * i * fac[i - 1] %</pre>
 ifac[i] = 1LL * inv[i] * ifac[i - 1] % mod;
int C(int n, int r){
return (r \ge 0 \&\& n \ge r)? (1LL * fac[n] * ifac[n - r] %
  * ifac[r] % mod) : 0;
int main(){
prencr();
int n. k:
cin >> n >> k;
cout << S(n, k) << endl; //Also you have S(n, t) for all t.
```

### 6.11 Stirling

$$\left\{\begin{array}{c} \mathbf{n} \\ \mathbf{k} \end{array}\right\} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

### 7 String

### 7.1 Aho Corrasick

```
int nxt[N][C];
int f[N], q[N], vcnt;
vector<int> adj[N];

int add(string s)
{
  int cur = 0;
  for(auto ch : s)
  {
    ch -= 'a';
    if(!nxt[cur][ch]) nxt[cur][ch] = ++vcnt;
    cur = nxt[cur][ch];
}
```

```
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```

```
return cur;
}

void aho()
{
  int hi = 0, lo = 0;
  for(int i = 0; i < C; i++) if(nxt[0][i]) q[hi++] = nxt[0][i
         ];
  while(hi != lo)
  {
    int x = q[lo++];
    adj[f[x]].pb(x);
    for(int i = 0; i < C; i++)
    {
        if(nxt[x][i])
        {
            q[hi++] = nxt[x][i];
            f[nxt[x][i]] = nxt[f[x]][i];
        }
        else nxt[x][i] = nxt[f[x]][i];
    }
}</pre>
```

### 7.2 Palindromic

```
int n, last, sz;
char s[N]:
int len[N], link[N], cnt[N];
map<short, int> to[N];
void init() {
n = 0: last = 0:
for(int i = 0; i < N; i++) to[i].clear();</pre>
 s[n++] = -1:
link[0] = 1;
len[1] = -1:
sz = 2:
int get link(int v) {
while(s[n - len[v] - 2] != s[n - 1]) v = link[v]:
return v:
void add_letter(int c) {
 s[n++] = c:
 last = get link(last):
 if(!to[last][c]) {
 len [sz] = len[last] + 2;
 link[sz] = to[get_link(link[last])][c];
 to[last][c] = sz++;
```

```
}
last = to[last][c];
cnt[last] = cnt[link[last]] + 1;
}
```

### 7.3 Suffix Array

int rank[LOG][N], n, lg;

string s:

```
pair<pair<int, int>, int> sec[N]:
int sa[N];
int lc[N];
int lcp(int a, int b)
int _a = a;
for(int w = lg - 1; ~w && max(a, b) < n; w--)</pre>
 if(max(a, b) + (1 \ll w) \ll rank[w][a] == rank[w][b])
 a += 1 << w, b += 1 << w;
return a - _a;
int cnt[N]:
pair<pii, int> gec[N];
void srt()
memset(cnt, 0, sizeof cnt);
for(int i = 0: i < n: i++) cnt[sec[i].F.S+1]++:
for(int i = 1; i < N; i++) cnt[i] += cnt[i - 1];</pre>
for(int i = 0; i < n; i++) gec[--cnt[sec[i].F.S+1]] = sec[i</pre>
     1:
memset(cnt, 0, sizeof cnt):
for(int i = 0; i < n; i++) cnt[gec[i].F.F+1]++;</pre>
for(int i = 1: i < N: i++) cnt[i] += cnt[i - 1]:</pre>
for(int i = n - 1; ~i; i--) sec[--cnt[gec[i].F.F+1]] = gec[
     i];
void build()
n = s.size():
 int cur = 1; lg = 0;
 while(cur < n)</pre>
  lg++;
  cur <<= 1;
 lg++;
```

```
for(int i = 0; i < n; i++) rank[0][i] = s[i];</pre>
for(int w = 1: w < lg: w++)</pre>
for(int i = 0: i < n: i++)
 if(i + (1 << w - 1) >= n)
  sec[i] = \{\{rank[w-1][i], -1\}, i\};
  sec[i] = \{\{rank[w-1][i], rank[w-1][i+(1<< w-1)]\}, i\};
 srt():
 rank[w][sec[0].S] = 0:
 for(int i = 1; i < n; i++)</pre>
 if(sec[i].F == sec[i - 1].F)
  rank[w][sec[i].S] = rank[w][sec[i-1].S]:
  rank[w][sec[i].S] = i:
for(int i = 0: i < n: i++)</pre>
sa[rank[lg-1][i]] = i;
for(int i = 0; i + 1 < n; i++)
lc[i] = lcp(sa[i], sa[i + 1]);
```

### 7.4 Suffix Automata

```
const int maxn = 2 e5 + 42: // Maximum amount of states
map < char , int > to [ maxn ]; // Transitions
int link [ maxn ]; // Suffix links
int len [ maxn ]; // Lengthes of largest strings in states
int last = 0: // State corresponding to the whole string
int sz = 1; // Current amount of states
void add letter ( char c ) { // Adding character to the end
int p = last ; // State of string s
last = sz ++; // Create state for string sc
len [ last ] = len [ p ] + 1;
for (; to [ p ][ c ] == 0; p = link [ p ]) // (1)
 to [p][c] = last; // Jumps which add new suffixes
if ( to [ p ][ c ] == last ) { // This is the first
     occurrence of
 c if we are here
 link [ last ] = 0;
 return :
int q = to [ p ][ c ];
if ( len [ q ] == len [ p ] + 1) {
link [ last ] = q ;
 return ;
```

```
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```

```
}
// We split off cl from q here
int cl = sz ++;
to [ cl ] = to [ q ]; // (2)
link [ cl ] = link [ q ];
len [ cl ] = len [ p ] + 1;
link [ last ] = link [ q ] = cl ;
for (; to [ p ][ c ] == q ; p = link [ p ]) // (3)
to [ p ][ c ] = cl ; // Redirect transitions where needed
}
```

### 7.5 Suffix Tree

```
#define fpos adla
const int inf = 1e9;
const int maxn = 1e4; //maxn = number of states of suffix
    tree
char s[maxn]:
map<int, int> to[maxn]; //edges of tree
int len[maxn], fpos[maxn], link[maxn];
//len[i] is the length of the inner edge of v
//fpos[i] is start position of inner edge in string s
int node, pos:
int sz = 1, n = 0;
int make_node(int _pos, int _len) {
fpos[sz] = _pos;
len [sz] = _len;
return sz++;
void go_edge() {
while(pos > len[to[node][s[n - pos]]]) {
 node = to[node][s[n - pos]];
 pos -= len[node];
}
}
void add_letter(int c) {
s[n++] = c;
pos++;
int last = 0;
while(pos > 0) {
 go_edge();
 int edge = s[n - pos];
 int &v = to[node][edge];
 int t = s[fpos[v] + pos - 1];
 if(v == 0) {
  v = make_node(n - pos, inf);
```

```
link[last] = node:
last = 0:
} else if(t == c) {
link[last] = node:
return;
} else {
int u = make_node(fpos[v], pos - 1);
to[u][c] = make_node(n - 1, inf);
to[u][t] = v;
 fpos[v] += pos - 1;
 len [v] -= pos - 1;
v = u:
link[last] = u;
last = u:
}
if(node == 0)
pos--;
else
node = link[node]:
```

### 8 Useful Fact and Constants

### 8.1 C(2n,n)

```
1: 2

2: 6

3: 20

4: 70

5: 252

6: 924

7: 3432

8: 12870

9: 48620

10: 184756

11: 705432

12: 2704156

13: 10400600

14: 40116600

15: 155117520
```

### 8.2 Factorials

```
1: 1
```

```
2: 2
3: 6
4: 24
5: 120
6: 720
7: 5040
8: 40320
9: 362880
10: 362880
11: 39916800
12: 479001600
13: 6227020800
14: 87178291200
15: 1307674368000
```

### 8.3 Long Long Integer

```
__int128 x;
unsigned __int128 y;
//Cin and Cout must be implemented
//Constants doesn't work
```

### 8.4 Power of 3

```
1: 3
2: 9
3: 27
4: 81
5: 243
6: 729
7: 2187
8: 6561
9: 19683
10: 59049
11: 177147
12: 531441
13: 1594323
14: 4782969
15: 14348907
16: 43046721
17: 129140163
18: 387420489
19: 1162261467
20: 3486784401
```

### Useful formulas

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 $\binom{n}{k} = \frac{n!}{k!(n-k)!}$  — number of ways to choose k objects out of n  $\binom{n+k-1}{k-1}$  — number of ways to choose k objects out of n with repetitions

permutations of n elements with k cycles  $\binom{n+1}{m} = n \binom{n}{m} + \binom{n}{m-1}$  $\frac{n}{m}$ ] — Stirling numbers of the first kind; number of

$${\binom{n+1}{m}} = n {\binom{n}{m}} + {\binom{n}{m-1}}$$

$$(x)_n = x(x-1)\dots x - n + 1 = \sum_{k=0}^n (-1)^{n-k} {n\brack k} x^k$$

of partitions of set  $1, \ldots, n$  into k disjoint subsets.  ${n+1 \brace m} = k \begin{Bmatrix} n \end{Bmatrix} + \begin{Bmatrix} n \cr k-1 \end{Bmatrix}$  ${n \choose m}$  — Stirling numbers of the second kind; number

$${\binom{n+1}{m}} = k {\binom{n}{k}} + {\binom{n}{k-1}}$$

$$\sum_{k=0}^{n} {n \brace k}(x)_k = x^n$$

$$C_n = \frac{1}{n+1} {2n \choose n} - \frac{\text{Catalan numbers}}{C(x)}$$

$$C(x) = \frac{1-\sqrt{1-4x}}{2x}$$

## Binomial transform

If 
$$a_n = \sum_{k=0}^{n} {n \choose k} b_k$$
, then  $b_n = \sum_{k=0}^{n} (-1)^{n-k} {n \choose k} a_k$ 

• 
$$a = (1, x, x^2, ...), b = (1, (x+1), (x+1)^2, ...)$$

• 
$$a_i = i^k, b_i = {n \brace i}!$$

### Burnside's lemma

shifts of array, rotations and symmetries of  $n \times n$ Let G be a group of action on set X (Ex.: cyclic

action f that transforms x to y: f(x) = y. Call two objects x and y equivalent if there is an

calculated as follows: CThe number of equivalence classes then can be lculated as follows:  $C = \frac{1}{|G|} \sum_{f \in G} |X^f|$ , where  $X^f$ 

is the set of fixed points of  $f: X^f = \{x | f(x) = x\}$ 

## Generating functions

 $a_0, a_1, \dots, a_n, \dots$  is  $A(x) = \sum_{i=1}^{\infty} a_i x^i$ Ordinary generating function (o.g.f.) for sequence

sequence  $a_0, a_1, \dots, a_n, \dots$  is  $A(x) = \sum_{n=0}^{\infty} a_i x^i$ Exponential generating function (e.g.f.)

$$B(x) = A'(x), b_{n-1} = n \cdot a_n$$

$$c_n = \sum_{k=0}^n a_k b_{n-k} \text{ (o.g.f. convolution)}$$

$$c_n = \sum_{k=0}^n \binom{n}{k} a_k b_{n-k} \text{ (e.g.f. convolution, compute}$$
with FFT using  $\widetilde{a_n} = \frac{a_n}{n!}$ )

# General linear recurrences

algorithm in  $O(n \log^2 n)$ . also can compute all  $a_n$  with Divide-and-Conquer If  $a_n =$  $\sum_{k=1}^{n} b_k a_{n-k}, \text{ then } A(x) =$ 

# Inverse polynomial modulo x'

Given A(x), find B(x) s  $A(x)B(x) = 1 + x^{l} \cdot Q(x) \text{ for some } Q(x)$ 

1. Start with  $B_0(x) = \frac{1}{a_0}$ 

 $B_{k+1}(x) = (-B_k(x)^2 A(x) + 2B_k(x)) \mod x^{2^{k+1}}$ 

# Fast subset convolution

Given array  $a_i$  of size  $2^k$ , calculate  $b_i =$ 

## Hadamard transform

size  $2 \times 2 \times \ldots \times 2$ , calculate FFT of that array: Treat array a of size  $2^k$  as k-dimentional array