

GRAPHS: DFS AND BFS BASED ALGORITHMS

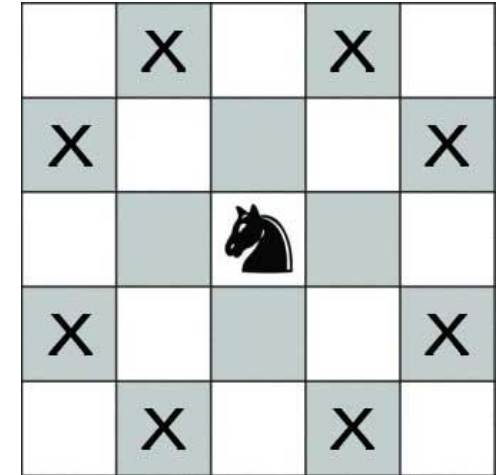
DATA STRUCTURES AND ALGORITHMS



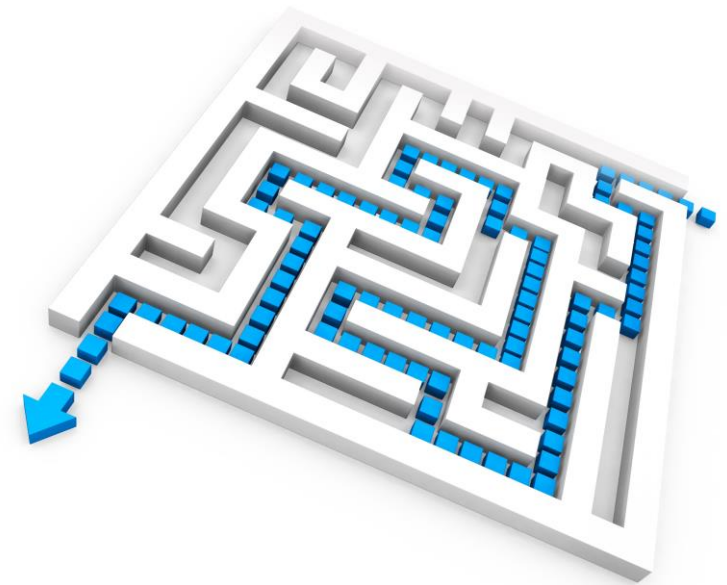
GRAPHS DATA STRUCTURE AND ALGORITHMS

DFS and BFS based Algorithms

- Connected Components
- Cycle finding
 - For Directed
 - Undirected
- Connected Component
- Topological sorting
- All paths finding
- Shortest path (not weighted graph)
- Bipartite graph finding
- Edges Classification
- Articulation points finding
- Bridges Finding



ROCK CRUSHES LIZARD
 SCISSORS DECAPITATE LIZARD
 LIZARD EATS PAPER
 LIZARD POISONS SPOCK
 PAPER DISPROVES SPOCK
 SPOCK VAPORIZES ROCK
 SPOCK BENDS SCISSORS



CONNECTED COMPONENTS

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CONNECTED COMPONENTS

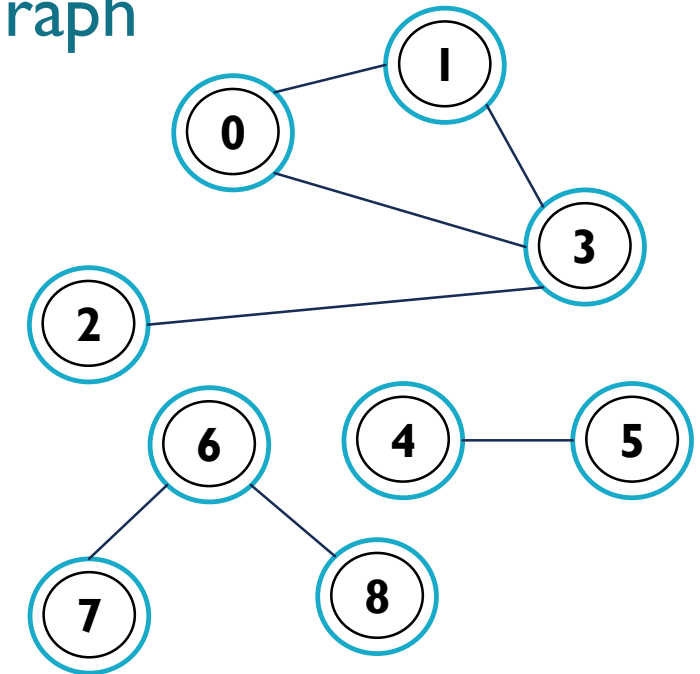
Definition

Connected component of an undirected graph is a subarray in which any two vertices are connected to each other by paths

restart DFS (or BFS) from one of the remaining unvisited vertices to find the next connected component. This process is repeated until all vertices have been visited and has an overall time complexity of $O(V + E)$.

```
int cc_num = 0; //connected components
visited.assign(N, false);
for (int i = 0; i < N; i++){
    if(!visited[i]){
        cout<<"connected component "<<
        ++cc_num<<":";
        dfs(i);
        cout<<endl;
    }
}
```

Graph



Output

```
connected component 1: 0 1 3 2
connected component 2: 4 5
connected component 3: 6 7 8
```


CYCLE FINDING ALGORITHM

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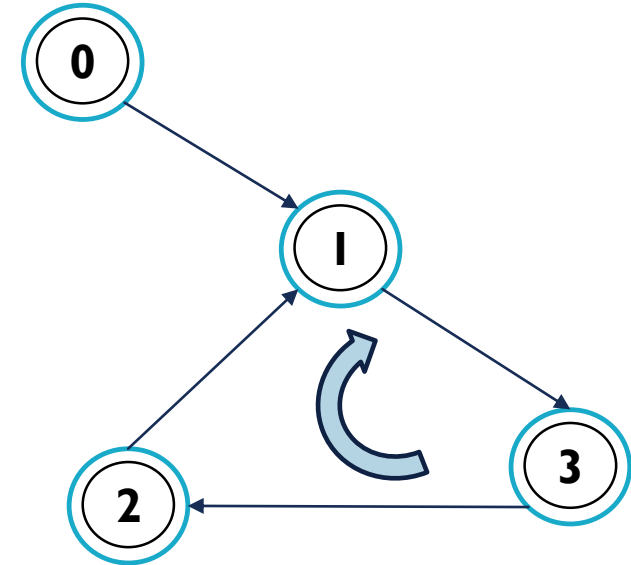


CYCLE FINDING IN DIRECTED GRAPH

```
vector<vector<int>> graph;
```

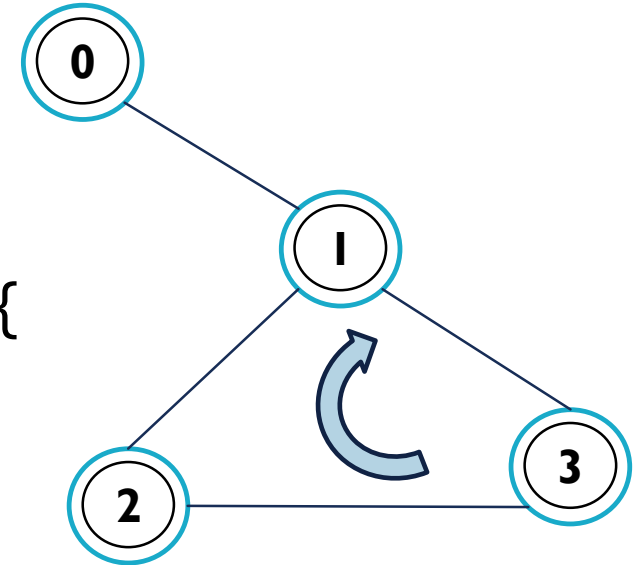
```
vector<bool> visited;
```

```
bool dfs(int u){  
    visited[u] = true;  
    for (int i = 0; i < graph[u].size(); i++){  
  
        int v = graph[u][i];  
        if(visited[v])  
            return true;  
        else  
            dfs(v);  
    }  
    return false;  
}
```



CYCLE FINDING IN UNDIRECTED GRAPH

```
vector<vector<int>> graph;  
vector<bool> visited;  
  
bool dfs(int u, int p = -1){  
    visited[u] = true;  
    for (int i = 0; i < graph[u].size(); i++){  
        int v = graph[u][i];  
        if(visited[v] && p != u)  
            return true;  
        else  
            dfs(v, u);  
    }  
    return false;  
}
```



TOPOLOGICAL SORTING

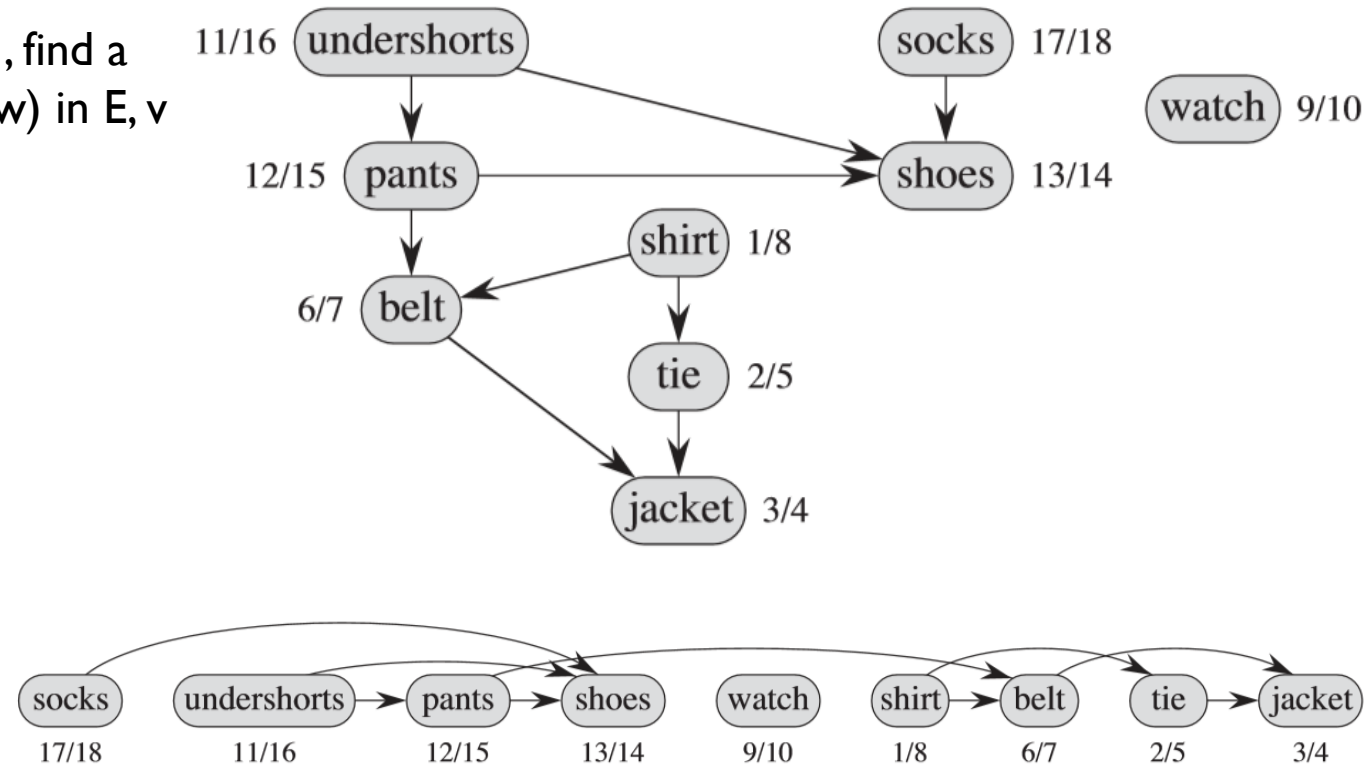
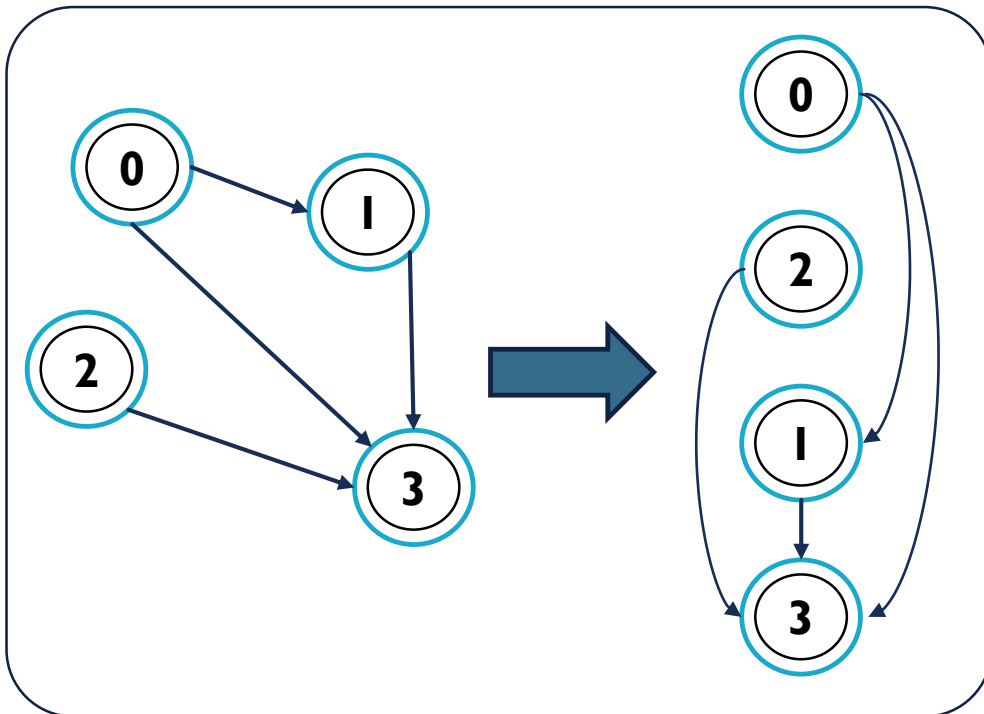
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TOPOLOGICAL SORTING

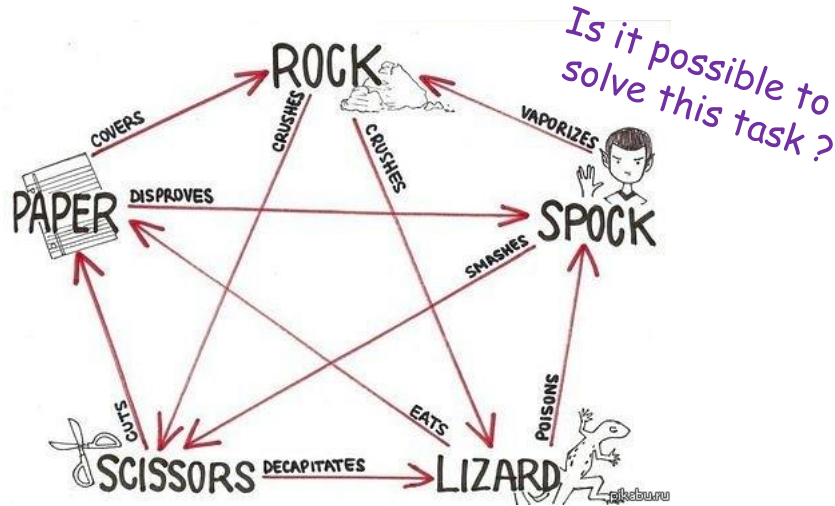
Definition

Topological sorting problem: given digraph $G = (V, E)$, find a linear ordering of vertices such that: for any edge (v, w) in E , v precedes w in the ordering



TOPOLOGICAL SORTING

Algorithm



A topological sorting is possible if and only if the graph has no directed cycles i.e., it is a Directed Acyclic Graph (DAG). It is always possible to find a topological order for a DAG and this can be done in linear time.

Topological sorting of a DAG can be found with DFS. During the depth-first traversal of the graph, just when a vertex finishes expanding (i.e., all its outlinks have been visited), add it to a stack. The order of vertices in the stack represents the topological order of the DAG.

Let L be an empty list

Let all vertices be initially unmarked

while there are unmarked vertices:

 select an unmarked vertex u

 dfs(u)

dfs(vertex u):

 mark u

 foreach edge u -> v

 if v is unmarked:

 dfs(v)

 add u to head of L

L represents the topological order of the DAG

TOPOLOGICAL SORTING IMPLEMENTATION

```
#include<iostream>
#include<vector>
using namespace std;

vector<vector<int>> g;
vector<bool> visited;
vector<int> ans;

void dfs(int v){
    visited[v] = true;
    for (int i=0; i<g[v].size(); ++i) {
        int to = g[v][i];
        if (!visited[to])
            dfs (to);
    }
    ans.push_back (v); // reverse path
}
```

```
void topological_sort(){
    for (int i=0; i<visited.size(); ++i)
        visited[i] = false;
    ans.clear();
    for (int i=0; i<visited.size(); ++i)
        if (!visited[i])
            dfs (i);
    reverse (ans.begin(), ans.end());
}
```


ALL PATHS FINDING

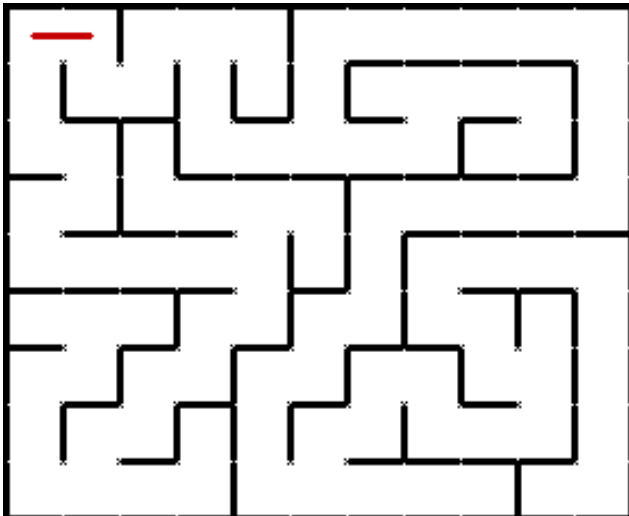
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DFS ALL PATHS FINDING

Algorithm

- Additional parameter for storing the path
- Use the DFS with two additional parts
 - Source – Destination producing (printing or storing)
 - Backtracking



DFS

dfs(curr, dest)

If curr == dest
print path from curr to dest

visited [cur] = true;
from cur vertex to all adjuts
if not visited[adjust]
dfs(curr, dest)

backtrack the path.

DFS ALL PATHS FINDING

```
#include<iostream>
#include<vector>
using namespace std;
int n,m, a,b;
vector <vector<int>> adj;
vector <int> visited, path;

void dfs(int curr, int dest){
    if(curr == dest) {
        for(int i = 0; i < path.size(); i++) {
            cout << path[i] << ' ';
        }
        cout << dest << '\n';
        return;
    }
    path.push_back(curr);
    visited[curr] = true;
    for (int i = 0; i < adj[curr].size(); ++i) {
        int now = adj[curr][i];
        if(!visited[now])
            dfs(now, dest);
    }
    path.pop_back();
    visited[curr] = false;
}
```

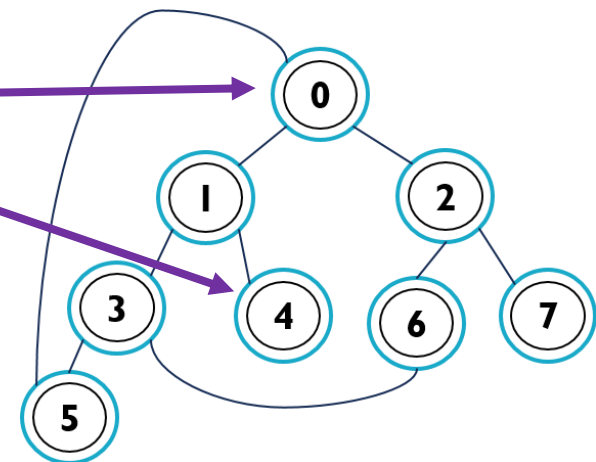
```
int main(){

    cin >> n >> m;
    adj = vector <vector <int>>(n);
    visited = vector <int> (n, 0);
    for (int i = 0; i < m; ++i) {
        cin >> a >> b;
        adj[a].push_back(b);
        adj[b].push_back(a);
    }
    cout<<endl;
    dfs(0, 4);
    system("pause");
    return 0;
}
```

```
c:\users\sul\documents\vis...
8 9
0 1
0 2
1 3
1 4
2 6
2 7
3 6
3 5
0 5

0 1 4
0 2 6 3 1 4
0 5 3 1 4
Press any key to continue . . .
```

All paths from
0 to 4



EDGES CLASSIFICATION

DATA STRUCTURES AND ALGORITHMS

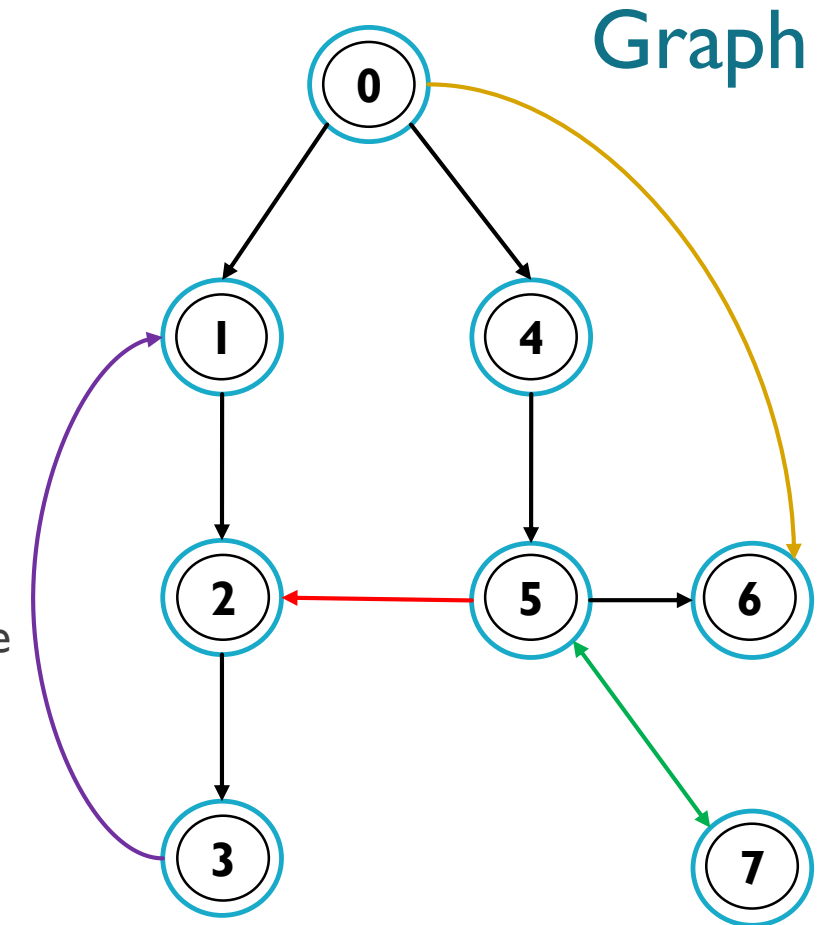
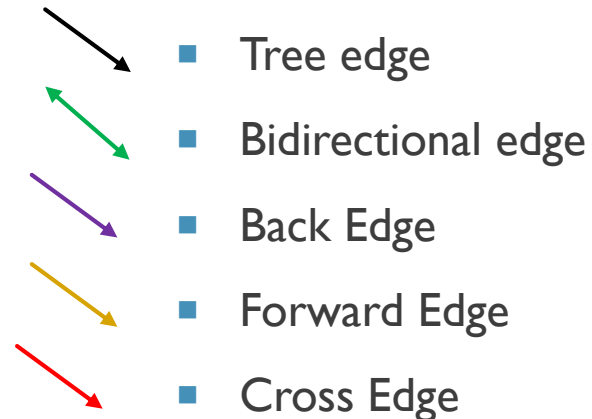


EDGES CLASSIFICATION

Definition

- If v is visited for the first time as we traverse the edge (u, v) , then the edge is a tree edge
- Else, v has been visited
 - If v is an ancestor of u , then edge (u, v) is a back edge
 - Else, if v is a descendant of u , then edge (u, v) is a forward edge
 - Else, if v is neither an ancestor or descendant of u , then edge (u, v) is a cross edge

Notation



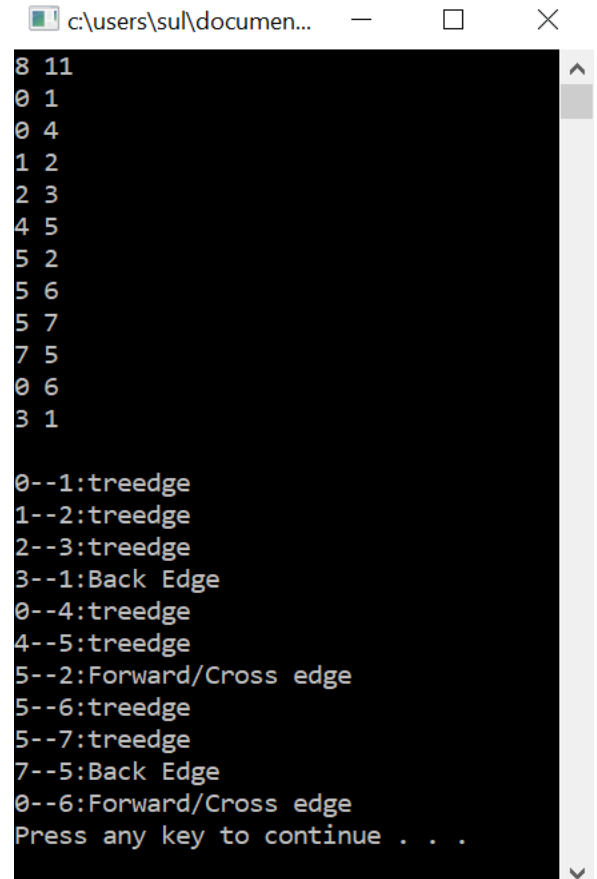
EDGES CLASSIFICATION

```
#include<iostream>
#include<vector>
#define white 0
#define gray 1
#define black 2
using namespace std;
int n,m, a,b;
vector <vector<int>> adj;
vector <int> color;
vector<int> parent;

void dfs(int curr);

int main(){
    cin >> n >> m;
    adj = vector < vector <int> > (n);
    color = vector <int> (n, 0);
    parent = vector<int> (n, -1);
    for (int i = 0; i < m; ++i) {
        cin >> a >> b;
        adj[a].push_back(b);
    }
    cout<<endl;
    dfs(0);
    system("pause");
    return 0;
}
```

```
void dfs(int curr){
    color[curr] = gray;
    for (int i = 0; i < adj[curr].size(); ++i) {
        int now = adj[curr][i];
        if(color[now] == white){
            parent[now] = curr;
            cout<<curr<<"--"<<now<<":treedge\n";
            dfs(now);
        }
        else if(color[now] == gray){
            if(now == parent[now])
                cout<<curr<<"--"<<now<<":Bidirectional\n";
            else
                cout<<curr<<"--"<<now<<":Back Edge\n";
        }
        else if(color[now] == black)
            cout<<curr<<"--"<<now<<":Forward/Cross edge\n";
        }
        color[curr] = black;
    }
}
```



```
c:\users\sul\documen...
8 11
0 1
0 4
1 2
2 3
4 5
5 2
5 6
5 7
7 5
0 6
3 1

0--1:treedge
1--2:treedge
2--3:treedge
3--1:Back Edge
0--4:treedge
4--5:treedge
5--2:Forward/Cross edge
5--6:treedge
5--7:treedge
7--5:Back Edge
0--6:Forward/Cross edge
Press any key to continue . . .
```

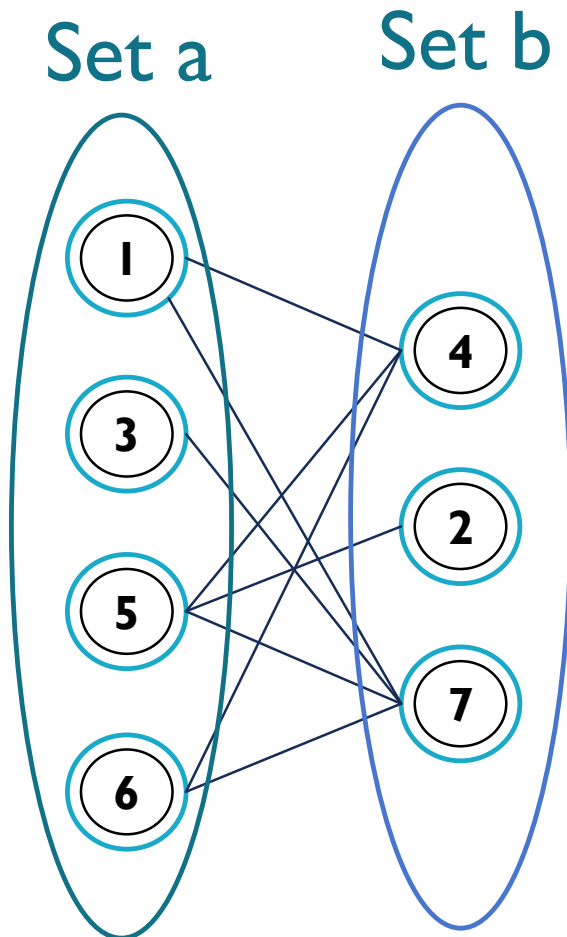

BIPARTED GRAPH FINDING

DATA STRUCTURES AND ALGORITHMS



BIPARTED GRAPH IMPLEMENTATION

```
vector<vector<int>> graph;
vector<int> color;
```



```
bool bfs(int s){
    fill(color.begin(), color.end(), -1);
    color[s] = 1;
    queue<int> q;
    q.push(s);
    while (!q.empty()){
        int u = q.front(); q.pop();
        for (int i = 0; i < graph[u].size(); i++){
            int v = graph[u][i];

            if(v == u) // self loop
                return false;
            if(color[v] == -1){
                color[v] = 1 - color[u];
                q.push(v);
            }else if(color[v] == color[u])
                return false;
        }
    }
    return true;
}
```


ARTICULATION POINT FINDING

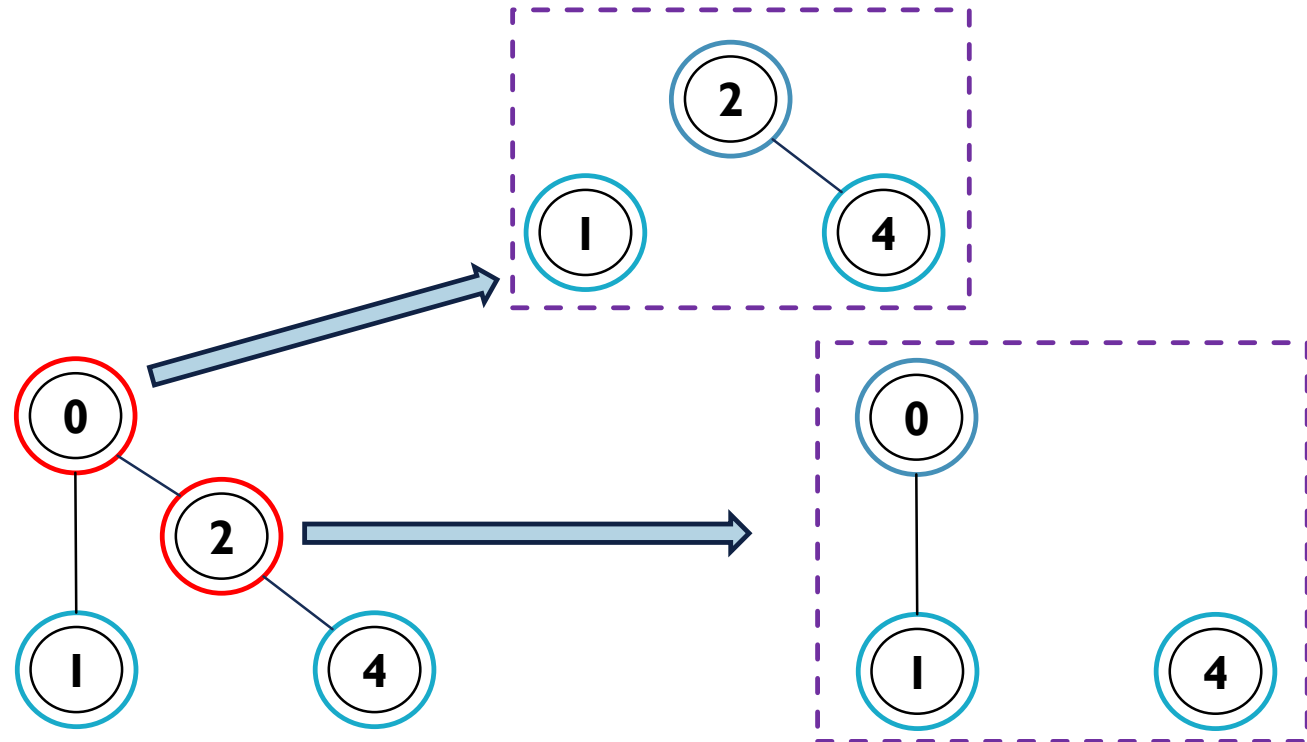
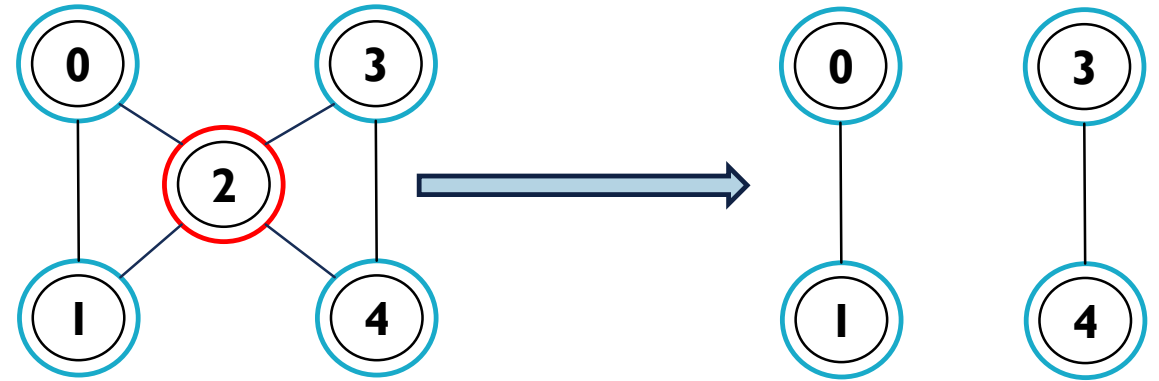
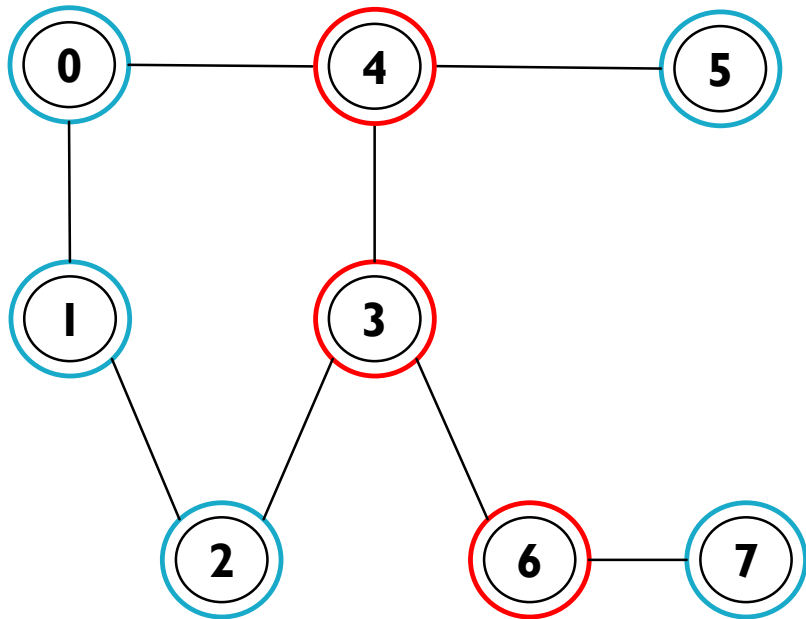
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ARTICULATION POINT

Definition

A vertex in an undirected connected graph is an articulation point iff removing it disconnects the graph

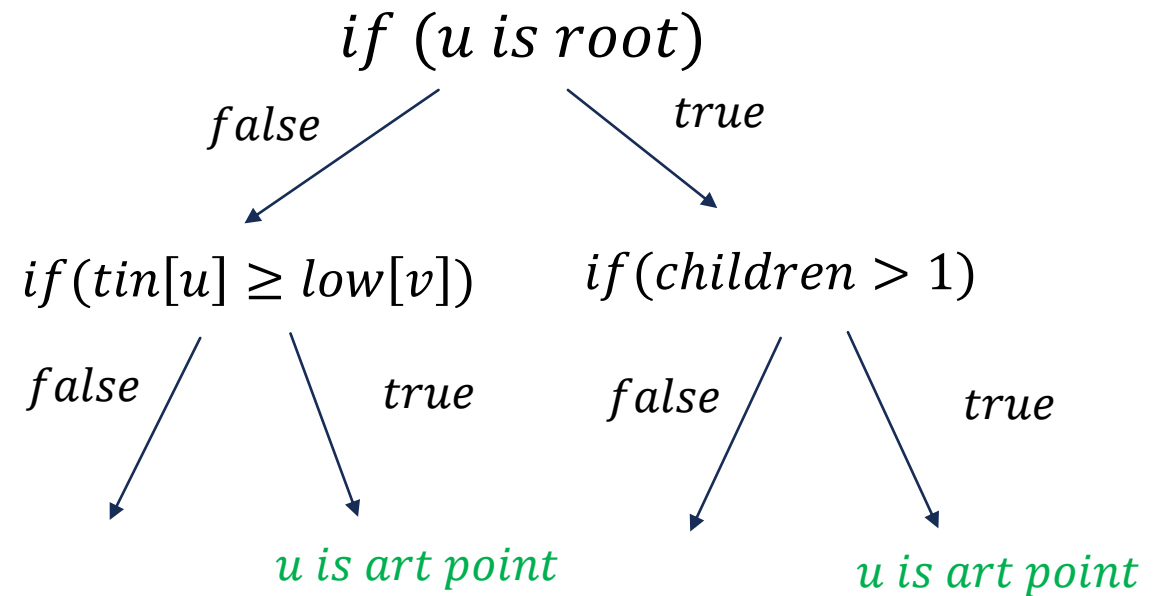


ARTICULATION POINT

Algorithm

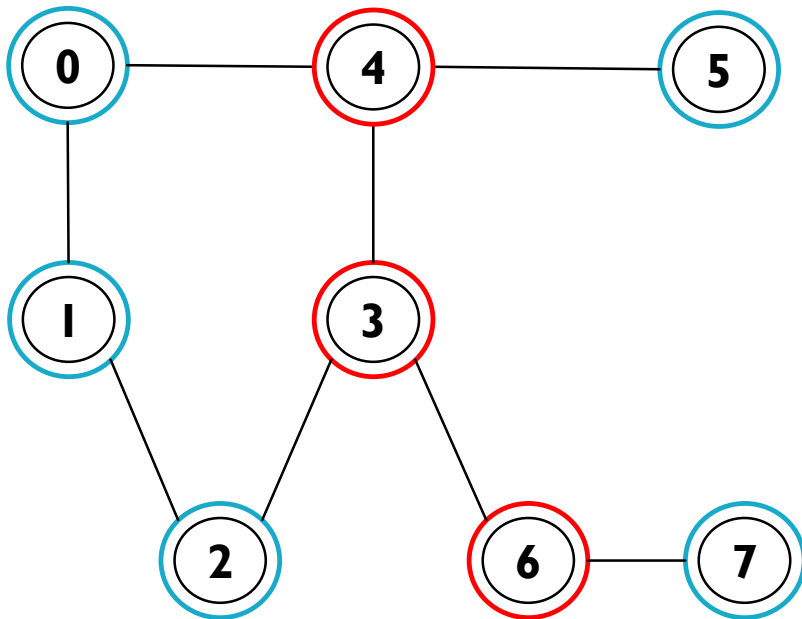
- u is root of DFS tree and it has more than one children
- u is not the root
 - If it has a child v such that no vertex in subtree rooted with v has a back edge to one of the ancestors in of u

$$low[u] = \begin{cases} u \rightarrow v \text{ tree edge: } \min(low[u], low[v]) \\ u \rightarrow v \text{ back edge: } \min(low[u], tin[v]) \end{cases}$$



ARTICULATION POINT IMPLEMENTATION

```
#include<iostream>
#include<vector>
#include<algorithm>
using namespace std;
vector<vector<int>> graph;
vector<bool> visited;
vector<int> tin, low;
int timer;
```



```
void dfs(int u, int p = -1){
    low[u] = tin[u] = timer++;
    visited[u] = true;
    int child = 0;
    for (int i = 0; i < graph[u].size(); i++)
    {
        int v = graph[u][i];
        if(v == p)
            continue;
        if(visited[v])
            low[u] = min(low[u], tin[v]); // back edge
        else{
            dfs(v, u);
            low[u] = min(low[u], low[v]); // tree edge
            if( low[v] >= tin[u] && p != -1)
                cout<<"art point: "<<u<<endl;
            child++;
        }
    }
    if(child > 1 && p == -1)
        cout<<"art point: "<<u<<endl;
}
```


BRIDGES FINDING

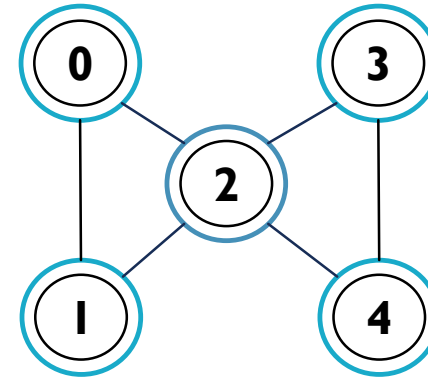
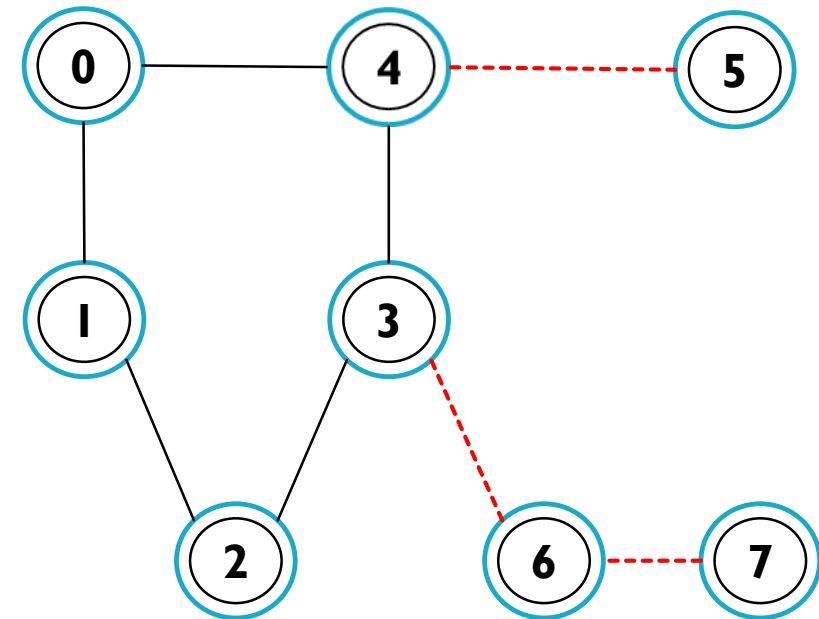
DATA STRUCTURES AND ALGORITHMS



BRIDGES

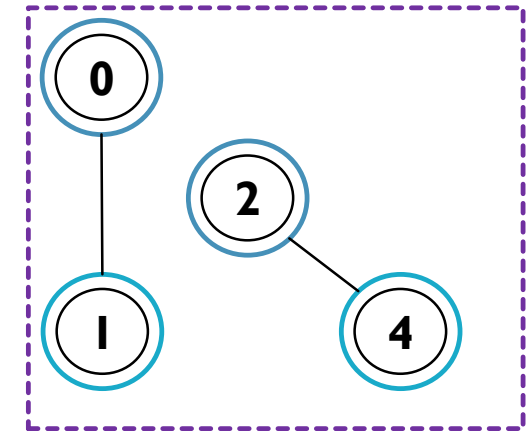
Definition

An edge in an undirected connected graph is the bridge iff removing it disconnects the graph

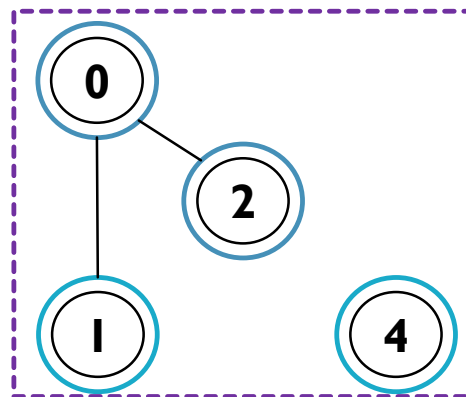


No Bridges,
but has an articulation point

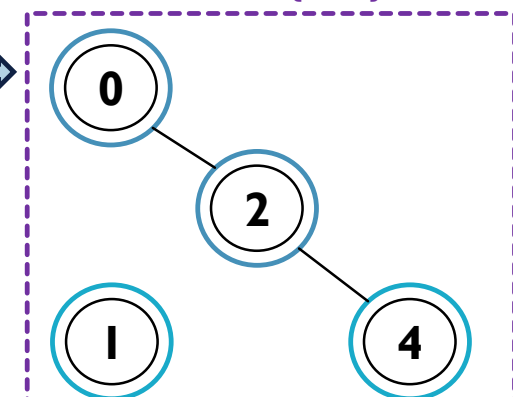
Remove $\{0, 2\}$



Remove $\{2, 4\}$

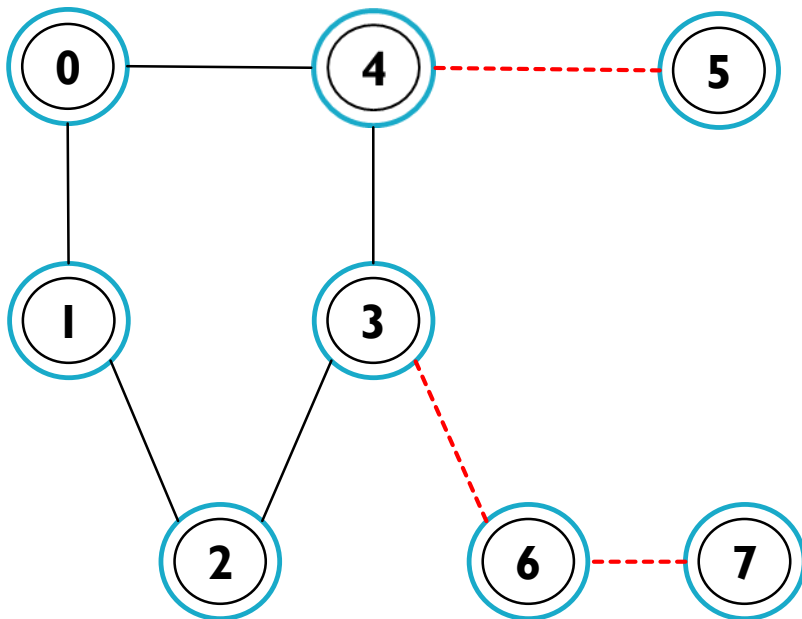


Remove $\{1, 0\}$



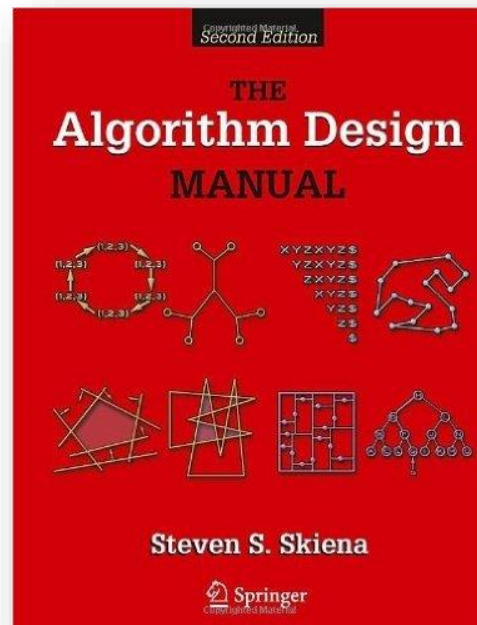
BRIDGE FINDING IMPLEMENTATION

```
#include<iostream>
#include<vector>
#include<algorithm>
using namespace std;
vector<vector<int>> graph;
vector<bool> visited;
vector<int> tin, low;
int timer;
```

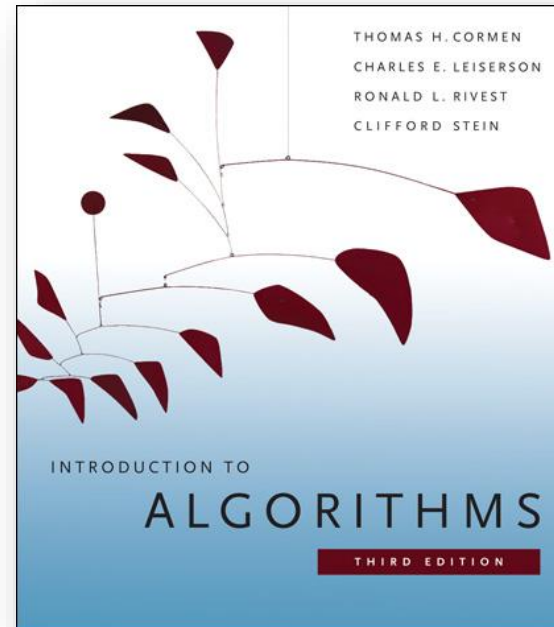


```
void dfs(int u, int p = -1){
    low[u] = tin[u] = timer++;
    visited[u] = true;
    for (int i = 0; i < graph[u].size(); i++)
    {
        int v = graph[u][i];
        if(v == p)
            continue;
        if(visited[v])
            low[u] = min(low[u], tin[v]); // back edge
        else{
            dfs(v, u);
            low[u] = min(low[u], low[v]); // tree edge
            if( low[v] > tin[u] && p != -1)
                cout<<u<<"- - -"<<v<<endl;
        }
    }
}
```

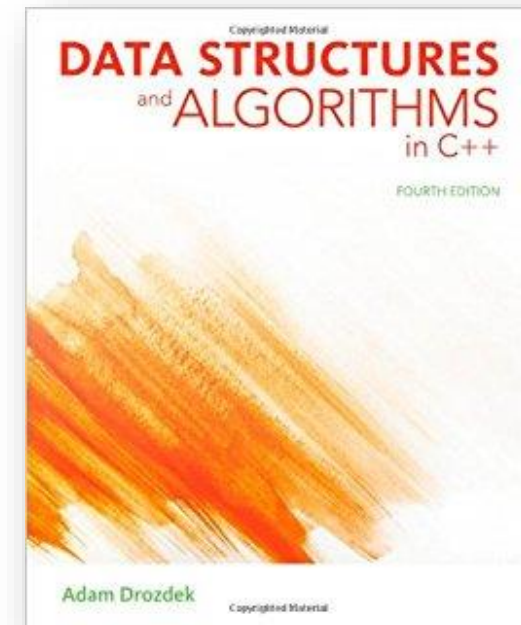

LITERATURE



Stieven Skienna
Algorithms design manual
Chapter 5: Graph Traversal
Page 145



Thomas H. Cormen
Introduction to Algorithms
Chapter VI Graph Algorithms
Page 587.



Adam Drozdek
Data structures and Algorithms in C++
Chapter 8: Graphs
Page 391