

# DFA (Deterministic finite automata)

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## Formal definition of DFA:

An DFA can be represented by a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$  where

- ❖  $Q$  is a finite set of states.
- ❖  $\Sigma$  is a finite set of symbols called the alphabets.
- ❖  $\delta$  is the transition function
- ❖  $q_0$  is the initial state from where any input is processed ( $q_0 \in Q$ ) (DFA have must only one initial state)
- ❖  $F$  is a set of final state/states of  $Q$  ( $F \subseteq Q$ ) (DFA have one final state, but it also possible that it have more than 1 final states)

## DFA have must the following properties:

- No empty transition
- **No. of arrows** on each state **depend upon alphabets** (it means if there is 1 alphabets in language like  $\Sigma = \{a\}$ , then there must be only one arrow on each state and if there are two alphabets in language like  $\Sigma = \{a, b\}$ , then there must be two arrows on each state and so on...)

**OR**

- In DFA, state can read only one letter at a time. (It means that every state has strictly one transition for each alphabet. Suppose there are two alphabets in the language  $L = \{a, b\}$ , then each state has strictly had two transitions. One for alphabet "a" and one transition for alphabet "b")
- Reject state sometimes.

## No empty transition means?

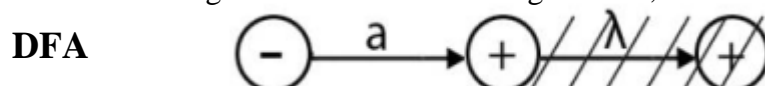
It means that null (empty) transition is not allowed in DFA.

Let's clear it with example:

We have Regular Expression "a" if  $\Sigma = \{a, b\}$



Both above diagrams have same meaning in NFA, because NFA allow empty transition.



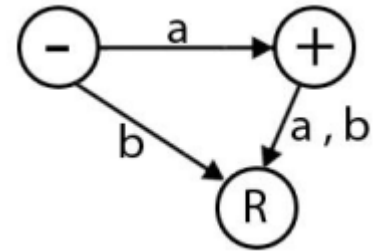
But DFA don't allow empty transition.

### Reject state means?

A state which not allow us to reach **final state**.

Simply we can say that on this particular transition **letter will be skip / ignore**.

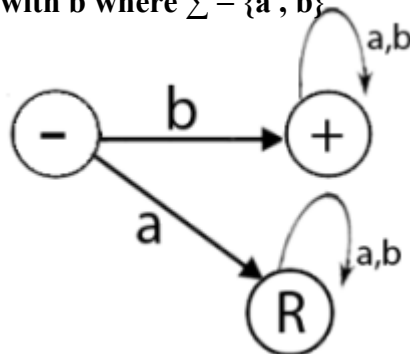
In the above DFA it will only read letter "a". *Letter "b" & "a,b" will be reject.*



### Examples of DFA:

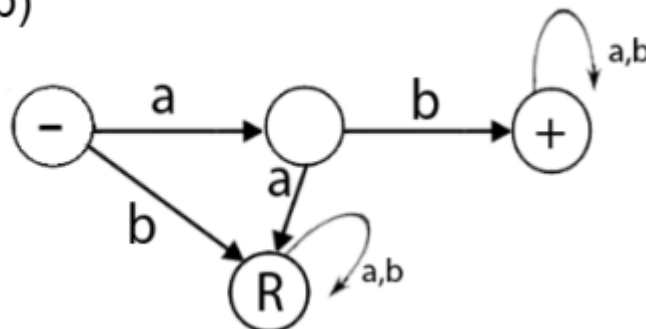
1. Draw DFA that starts with b where  $\Sigma = \{a, b\}$

R.E =  $b(a+b)^*$



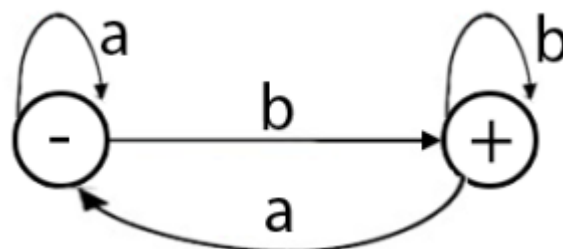
2. Draw DFA that starts with ab where  $\Sigma = \{a, b\}$

R.E =  $ab(a+b)^*$



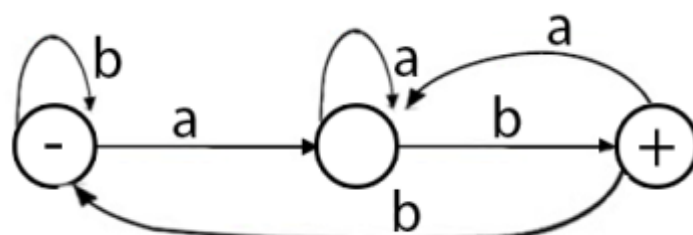
3. Draw DFA that ends with b where  $\Sigma = \{a, b\}$

RE =  $(a+b)^*b$



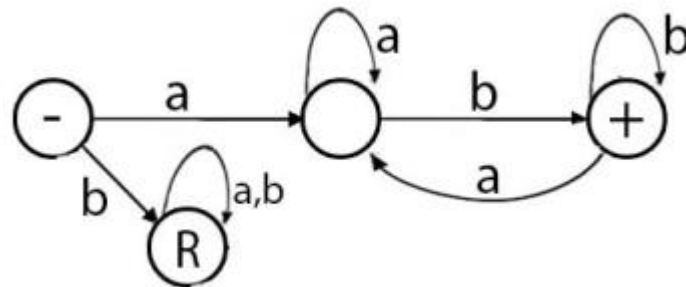
4. Draw DFA that ends with ab where  $\Sigma = \{a, b\}$

RE =  $(a+b)^*ab$



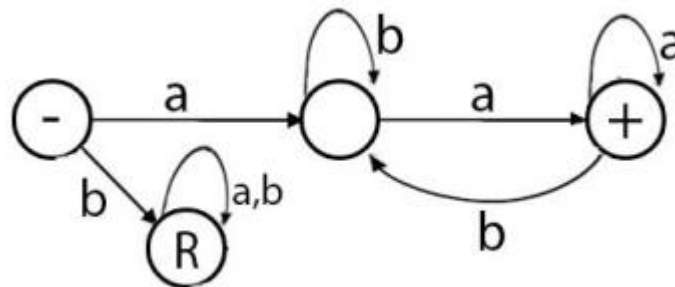
5. Draw DFA that starts with a & end with b where  $\Sigma = \{a, b\}$

RE =  $a(a+b)^*b$



6. Draw DFA that starts & end with a where  $\Sigma = \{a, b\}$

RE =  $a(a+b)^*a$



### Difference between NFA & DFA:

NFA	DFA
NFA is a generalize type of Finite Automata	DFA is a special type of Finite Automata
NFA have redundant transition. (It means that the transition from a state can be to multiple next states for each input symbol. Hence it is called non-deterministic)	DFA have no redundant transition. (It means that the transition from a state is to a single particular next state for each input symbol. Hence it is called deterministic)
NFA allows empty string transitions.	DFA don't allows empty string transitions.
NFA requires less space.(small set of strings)	DFA requires more space.(Large set of strings)
Regular expression can be easily converted to NFA.	Conversion of regular expression to DFA is complex.
NFA can't be easily implemented	DFA can be easily implemented

