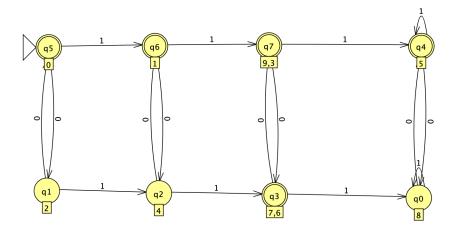
Week 2 HW, credit tasks

Task 1)

Using the NFA from P-2e, we perform product construction to obtain this. However the DFA was first obtained by building intuitively,



minimal DFA of P-2e and C-2a

The strategy here is to use the upper row of states as even parity of 0s, the bottom row being odd parity of states. So it is easy to assign, according to condition

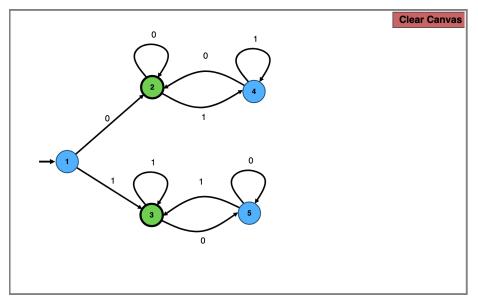
- All of top rows can be regarded as final states.
- The third state of bottom row (q3) is an accepting state for exactly two 1s.

This is achieved as the number of 1s in the sequence is the number of horizontal transitions taken from the initial state.

Additionally, the old DFA of previous submission can be *minimized* by subset partitioning to the same above diagram.

Task 2)

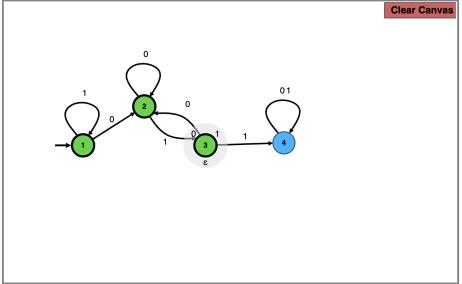
This DFA is very intuitive. We split initial state into case 0 and 1. The loops formed by (2,4) and (3,5) were constructed to ensure that all incoming transition to the final states are the same to the corresponding starting symbol.



 $DFA\ of\ same\ symbols$

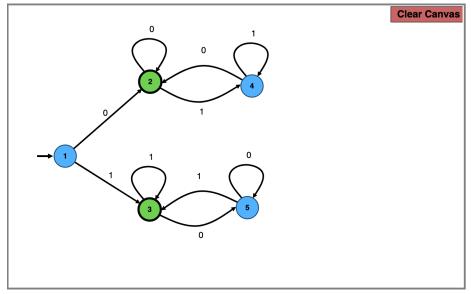
Task 3)

This DFA is sourced from the NFA constructed for P-2f, by inverting the accepting and non-accepting states from a DFA that only accepts substring 011.



inverted DFA of substring 011

Task 4) The NFA is constructed the same as that of the DFA of Task 2



 $DFA\ of\ same\ symbols$