CASH Approach for Capacity Sharing

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1 Current Scheduler

Changes in task 4 extended the scheduler to EDF scheduling with CBS. CBS provides a basis on which we will add an overrun control.

2 Modified Scheduler

Aperiodic Task: A task that does not run on a fix period

Task Server: A mechanism to schedule aperiodic tasks

Constant Bandwidth Server: A dynamic task server that supports soft-real time scheduling

Overrun: The amount of time a given task has exceeded its deadline

Capacity Sharing for Overrun Control: A policy that allows multiple task servers, in particular CBS, to share available, but unused budget

Required outcome from this task: extend the scheduler to support CASH

More on CASH: Whenever an aperiodic task runs on the CBS, there is a possibility that the task finishes before the budget is used up. When dealing with multiple CBS, we can implement a queue that keeps track of available, but unused budget. This allows task servers to consume unused resources first, increasing efficiency in resource consumption, and lowering overall execution times for aperiodic tasks.

Based on the requirements, we need to add a reclaiming mechanism to support CASH. We need to first implement a queue to store unused resources, and we need to store the budget and deadline of these resources. Instead of a queue, we have extended the TCB to support place holders for the budget and deadline (Figure 1).

In $nrk_scheduler.c$, we implement 3 mechanisms to adjust and update unused budget to cover 3 different scenarios. See figure 2, 3, and 4 for code.

1. We need to check and update the validity of unused budget. First, we determine whether the deadline of CBS has passed. If the deadline has passed, we can remove the resource, since we can no longer use them. This is done by setting cash and cash_period in the TCB to 0. Secondly, we also update the time values, specifically cash_period to reflect the new relative deadline each time the scheduler is called. We check to prevent the available budget from being larger than the relative deadline to avoid using non-existing budget.

```
TASK CONTROL BLOCK

"TASK Pointer to current top of stack "/

"TASK CONTROL BLOCK

"TASK Pointer to current top of stack "/

"TASK CONTROL BLOCK

"TASK CONTROL BLOCK

"TASK POINTER TO EVENT TO EV
```

Figure 1: Updated TCB

- 2. We need to start keeping track of unused budget. Whenever an CBS task finishes running, we would like store the remaining budget of a the CBS_TASK as cash and update cash_period in the TCB to store the relative deadline of the resource.
- 3. We need to check for unused budget to consume whenever a new CBS task is run. This can be determined by evaluating the *cash* and *cash_period* of all TCBs. An *cash* value larger than 0 indicate unused budget is available to be reclaim (Figure 5)). If unused resources are available, then the task can consume the resource held by the CBS, and adjust the budget to reflect the usage by decrementing the budget. As per the algorithm of CASH, we choose the CBS with the earliest deadline. We also handle cases when a single task requires more execution time than a single CBS remaining budget can provide; the algorithm allows a single running task to use up unused budgets from multiple CBS. In the case that there are no unused resources available, the task will follow the default protocol.

Figure 2: Code snippet: Adds remaining budget to CASH

```
// .iddd (cash )
// .id
```

Figure 3: Code snippet: Updates and maintains CASH and relative deadline values

Figure 4: Code snippet: Tasks use CASH first

3 Testing

The main function as part of main.c, is responsible for setting up and initialization of the system. It also holds the task set required to verify the accuracy of the scheduler.

Table 1: Task Set for CBS			
Task	e_i	P_i	Type
1	3 (Budget)	8	CBS
2	1	5	Basic
3	2 (Budget)	9	CBS

We chose the execution time of aperiodic tasks to be around 1.3 seconds in order to simulate task completion prior to budget depletion shown in figure 6. This task will help us demonstrate the functionality, explained in the next section.

Figure 5: Code snippet: Helper function to get the *task_ID* of the smallest CASH remaining

Figure 6: Code snippet: CBS task simulation

4 Expected Results

The expected behaviour of the scheduler for a new CBS task is to check and use any remaining budget, before using CPU reserves. The output of the simulation should reflect the checking and usage, and decreasing the CASH budget accordingly. We should also expect to see unused budget get added to the CASH budget when a CBS task completes.

5 Results

The screen captions below captures two scenarios. In figure 7, we can observe a running CBS task and its completion. Upon completion, we can see that it takes the remaining budget of 2071, and updates the CASH to 2071, as well as setting the relative deadline of the reserve. In figure 8, we can observe a

CBS task using CASH reserves. The budget is originally at a value of 2047. The unused budget is 4142, and is added to the current reserve, resulting in a new budget of 5939. These results successfully shows us the increment and decrement of the CASH budget.

Note: Further results are recorded in testResult.txt, submitted as part of this assignment.

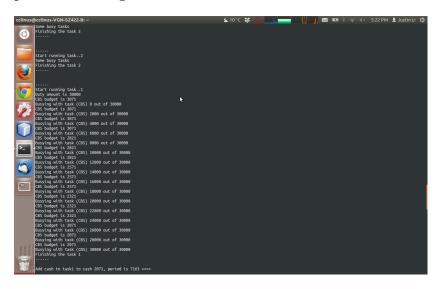


Figure 7: Handling remaining budget

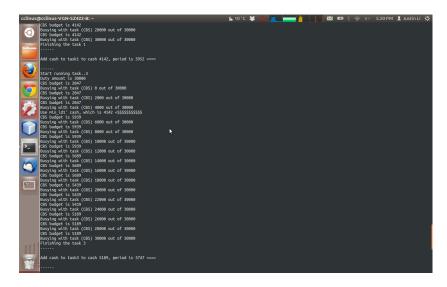


Figure 8: Using CASH reserves