

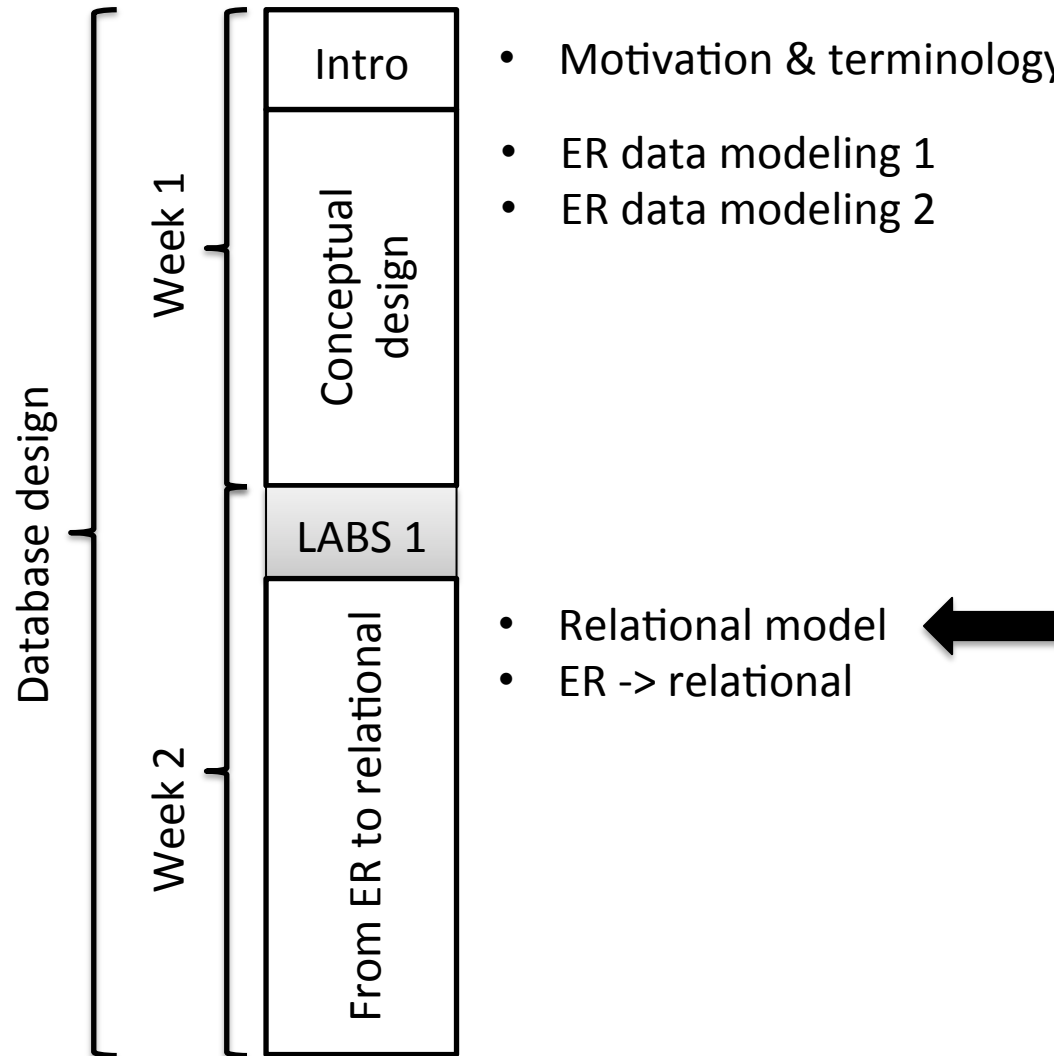


# The Relational Model

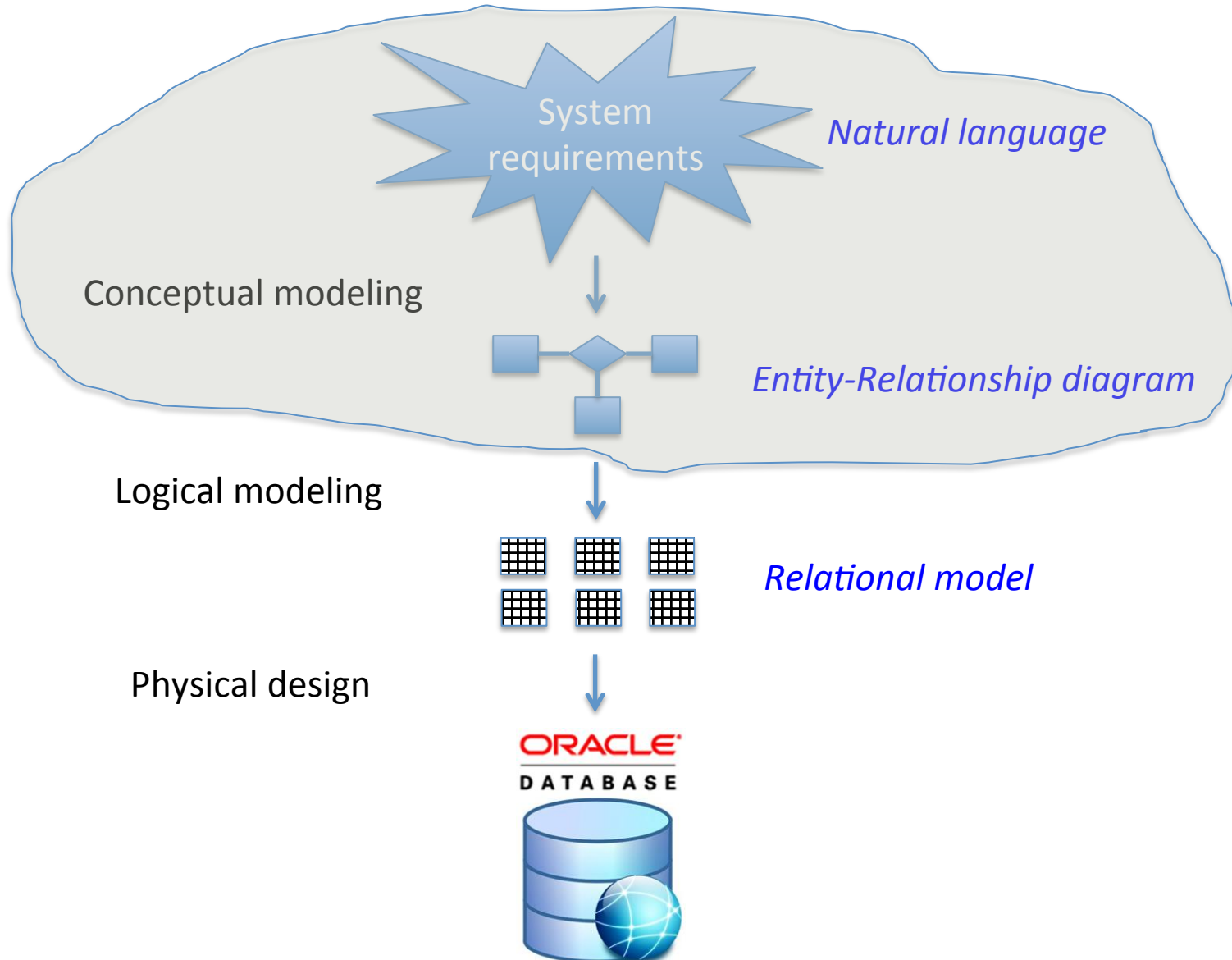
Lecturer: Neena Thota

[neena.thota@it.uu.se](mailto:neena.thota@it.uu.se)

# Where are we?



# Database design: full picture



# Intended Learning Outcomes

- Describe the main concepts of the relational model.
  - Relations, attributes, tuples.
  - First Normal Form.
  - Null values.
  - Superkeys, candidate keys, primary key.
  - Foreign keys.
- Create the corresponding structures using a DBMS.

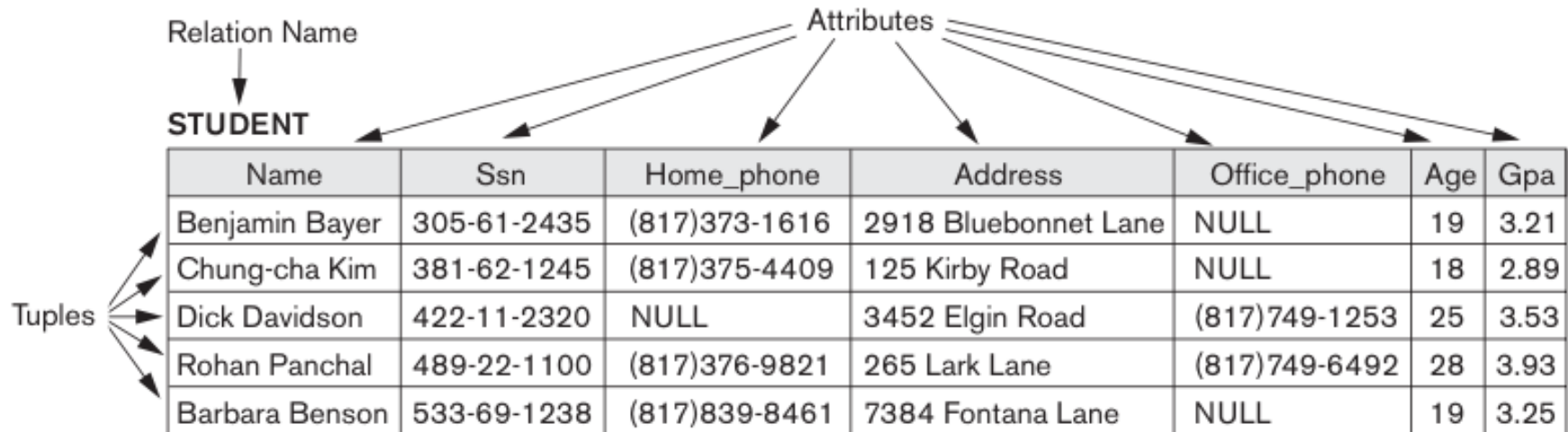
# Some history



- The relational model was introduced by Dr. Edgar (Ted) F. Codd (1924-2003) in 1970.
  - A mathematician from Oxford (UK), working as a researcher at the IBM San Jose Research Laboratory (USA).
- Many DBMSs are based on the relational data model.

# Relational Model Concepts

- Represents the database as a collection of relations.



# Basic concepts: from relations to tables (DB)

- $A_1, A_2, \dots, A_n$  are attributes/columns.
- $R(A_1, A_2, \dots, A_n)$  is a relation schema on these attributes.
  - Employee(EmpID, SSN, Name, Surname, Salary)
- $r$  is a relation on the relation schema  $R$ 
  - That is, the set of tuples/rows.
- The current values (*relation instance*) of a relation are specified by a table.

Relation → EMPLOYEE

Attribute/column ↓

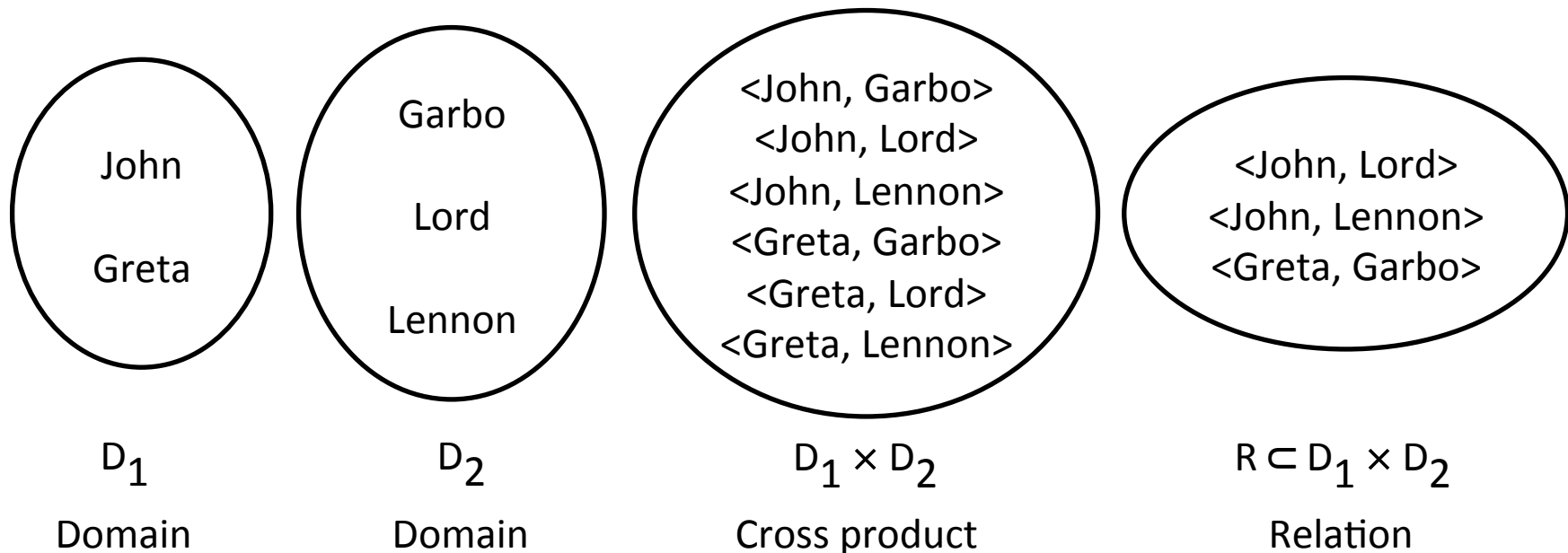
EmpID	SSN	Name	Surname	Salary
001	18160421	Charlotte	Brontë	13 000
002	18180730	Emily	Brontë	12 500
003	18200117	Anne	Brontë	11 000

Instance [

Tuple/row ←

# Basic concepts: Relations (math)

- In set theory, a relation is defined as a **subset** of the **product set** (*Cartesian* or *cross product*) of a number of **domains** (value sets).
- Members of a relation are called **tuples**.
- **Degree** of a relation is number of attributes.
- If the relation is of **degree**  $n$ , the tuples are called *n-tuples*.





# Data types of domains

- **Data types in SQL:**
  - **char(n).** Fixed length character string, with user-specified length n.
  - **varchar(n).** Variable length character strings, with user-specified maximum length n.
  - **int.** Integer (a finite subset of the integers that is machine-dependent).
  - **smallint.** Small integer (a machine-dependent subset of the integer domain type).
  - **numeric(p,d).** Fixed point number, with user-specified precision of p digits, with d digits to the right of decimal point.
  - **real, double precision.** Floating point and double-precision floating point numbers, with machine-dependent precision.
  - **float(n).** Floating point number, with user-specified precision of at least n digits.
  - **date.** Dates, containing a (4 digit) year, month and date.
  - **time.** Time of day, in hours, minutes and seconds.
  - **etc.**

# Data types and domains

- Domain is a set of **atomic** values.
- Data **type or format** is specified for each domain

Example:

STUDENT (Name: string, Ssn: string, Home-phone: string, Address: string, Office-phone: string, Age: integer, Gpa: real)

Name	Ssn	Home_phone	Address	Office_phone	Age	Gpa
Benjamin Bayer	305-61-2435	(817)373-1616	2918 Bluebonnet Lane	NULL	19	3.21
Chung-cha Kim	381-62-1245	(817)375-4409	125 Kirby Road	NULL	18	2.89
Dick Davidson	422-11-2320	NULL	3452 Elgin Road	(817)749-1253	25	3.53
Rohan Panchal	489-22-1100	(817)376-9821	265 Lark Lane	(817)749-6492	28	3.93
Barbara Benson	533-69-1238	(817)839-8461	7384 Fontana Lane	NULL	19	3.25

# Creation of a new relation (SQL)

SQL CREATE statement

Table name

Attributes

Domains

```
CREATE TABLE Employee (  
    EmpID      int,  
    SSN        int,  
    Name       varchar(100),  
    Surname    varchar(100),  
    Salary     int,  
    Department int,  
    Supervisor int  
)
```

# Insertion of new data into relation

```
INSERT INTO Employee  
VALUES (1,19851,'John','McEnroe',43000,4,13)
```

```
CREATE TABLE Employee (  
    EmpID          int,  
    SSN            int,  
    Name           varchar(100),  
    Surname        varchar(100),  
    Salary         int,  
    Department     int,  
    Supervisor     int  
)
```

# Basic concepts: First Normal Form (1NF)

## First Normal Form

- Only **simple/atomic** values are allowed in the relational model.
  - E.g., attributes are not allowed to have composite or multiple values, like an attribute “phone” with values {00461254362, 0039338746384} for a single employee.

EmpID	SSN	Name	Surname	Salary
001	18160421	Charlotte	Brontë	13 000
002	18180730	Emily	Brontë	12 500
003	18200117	Anne	Brontë	11 000

# Feature 1: Null values

- A special value, **null** or  $\perp$ , denote:
  - an unknown value (e.g. home phone no.)
  - a missing value (office phone no.)
  - an attribute that is not applicable (e.g., visa status for foreign status).

Name	Ssn	Home_phone	Address	Office_phone	Age	Gpa
Benjamin Bayer	305-61-2435	(817)373-1616	2918 Bluebonnet Lane	NULL	19	3.21
Chung-cha Kim	381-62-1245	(817)375-4409	125 Kirby Road	NULL	18	2.89
Dick Davidson	422-11-2320	NULL	3452 Elgin Road	(817)749-1253	25	3.53
Rohan Panchal	489-22-1100	(817)376-9821	265 Lark Lane	(817)749-6492	28	3.93
Barbara Benson	533-69-1238	(817)839-8461	7384 Fontana Lane	NULL	19	3.25

# Constraints: NULL values

- **Constraints** are restrictions on actual values in a database state.
- Derived from rules in miniworld that database represents.

```
CREATE TABLE Employee (  
    EmpID          int,  
    SSN            int,  
    Name           varchar NOT NULL,  
    Surname        varchar NOT NULL,  
    Salary         int,  
    Department     int,  
    Supervisor     int)
```

# Insertion of new data

```
CREATE TABLE Employee (  
    EmpID      int,  
    SSN        int,  
    Name       varchar NOT NULL,  
    Surname    varchar NOT NULL,  
    Salary     int,  
    Department int,  
    Supervisor int)
```

```
INSERT INTO Employee  
VALUES (1,1985123,'John','McEnroe',43000,4,13)
```

```
INSERT INTO Employee(EmpID,Name,Surname,SSN)  
VALUES (2,'Roger','Federer',198511155482)
```



## Feature 2: Tuple identification

- **Superkey:** specifies *uniqueness constraint* that no two distinct tuples in any state  $r$  of  $R$  can have same value for set of attributes.
- A set  $K$  of the attributes in a relation schema  $R$ , i.e.  $K \subseteq R$ , such that:
  - $t_1, t_2 \in r(R), t_1 \neq t_2 \rightarrow t_1[K] \neq t_2[K]$

Examples of superkeys:

- EmpID, SSN, Name, Surname, Salary
- EmpID, Surname
- SSN
- EmpID, SSN, Name
- EmpID
- ...

EmpID	SSN	Name	Surname	Salary
001	18160421	Charlotte	Brontë	13 000
002	18180730	Emily	Brontë	12 500
003	18200117	Anne	Brontë	11 000

# Question

Given a relation with 4 attributes, how many superkeys will be present *for sure*?

A: 0

B: 1

C: 4

D: 16

# Basic concepts: Candidate keys

- A superkey  $k$  is minimal if we cannot remove any attributes and still have the uniqueness constraint i.e. no other superkey  $K'$  such that  $K' \subset K$ .
- Every minimal superkey is called a **candidate key** for  $R$ .

Examples of candidate keys:

- EmpID, SSN, Name, Surname, Salary
- EmpID, Surname
- **SSN**
- EmpID, SSN, Name
- **EmpID**
- ...

<u>EmpID</u>	<u>SSN</u>	Name	Surname	Salary
001	18160421	Charlotte	Brontë	13 000
002	18180730	Emily	Brontë	12 500
003	18200117	Anne	Brontë	11 000

# Basic concepts: Primary key

- Used to uniquely identify a tuple in R.
- Database designer chooses one of R's candidate keys as **primary key** (or just **key**).
- Preferable to choose primary key with single attribute.

Examples of primary key:

- EmpID, SSN, Name, Surname, Salary
- EmpID, Surname
- **SSN**
- EmpID, SSN, Name
- **EmpID**
- ...

<u>EmpID</u>	SSN	Name	Surname	Salary
001	18160421	Charlotte	Brontë	13 000
002	18180730	Emily	Brontë	12 500
003	18200117	Anne	Brontë	11 000

# Relational keys: Summary

We need to identify specific tuples in a relation.

- **Superkey:** a set of attributes that allows us to do it.
  - A relation can have many superkeys.
- **Candidate key:** a minimal superkey.
  - A relation can have many candidate keys – normally, less than the superkeys.
  - UNIQUE in SQL.
- **Primary key:** the single candidate key we choose to identify records in a table.
  - PRIMARY KEY in SQL.
  - No NULLs allowed.

# Constraints: Candidate keys

```
CREATE TABLE Employee (  
    EmpID          int UNIQUE,  
    SSN            int UNIQUE,  
    Name           varchar,  
    Surname        varchar,  
    Salary         int,  
    Department     int,  
    Supervisor     int,  
)
```

# Constraints: Primary key

```
CREATE TABLE Employee (  
    EmpID          int PRIMARY KEY,  
    SSN            int UNIQUE,  
    Name           varchar(100) NOT NULL,  
    Surname        varchar(100) NOT NULL,  
    Salary         int,  
    Department     int,  
    Supervisor     int,  
    )
```

# Exercise: Key constraints

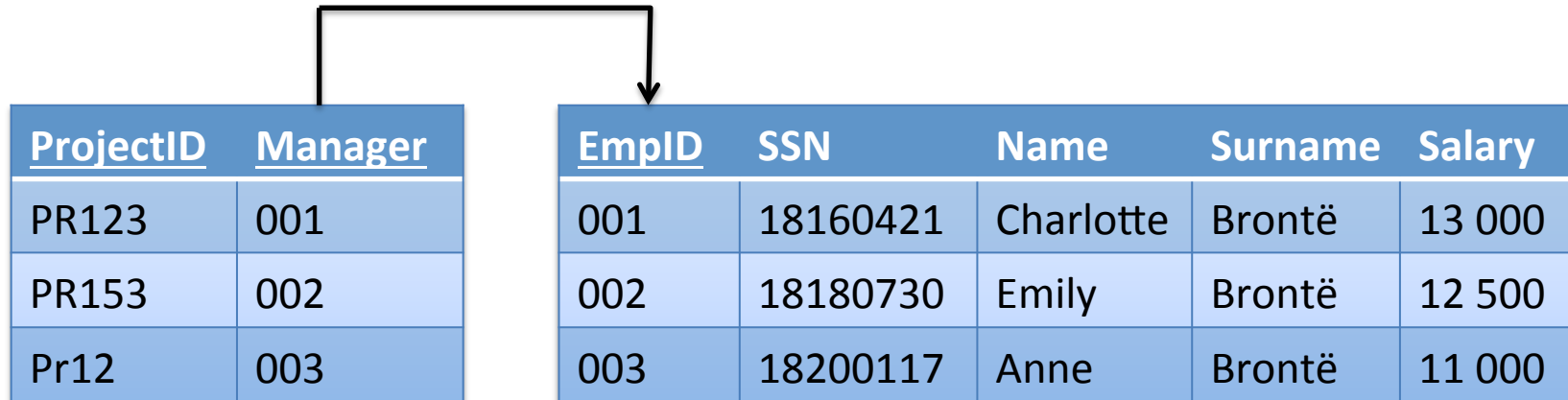
A1	A2	A3	A4
a	1	alpha	0
a	2	beta	0
b	3	gamma	0
b	4	beta	0
b	5	beta	0
c	1	alpha	0
d	2	gamma	0

- A) A2 is a candidate key
- B) A4 is the primary key
- C) A1, A2 is a candidate key
- D) A2, A3 is superkey
- E) None of the previous answers
- F) No answer



# Basic concepts: Foreign key constraints

- References between relations are made *by value* using *foreign keys*.
- “Manager” is a foreign key referencing the EmpID key in the Employee relation.



# Constraints: Referential integrity

- Specified between **two relations** to maintain consistency.
- A tuple in one relation that refers to another relation must refer to an **existing** tuple in that relation.

```
CREATE TABLE Employee (  
    EmpID          int PRIMARY KEY,  
    SSN            int UNIQUE,  
    Name           varchar NOT NULL,  
    Surname        varchar NOT NULL,  
    Salary         int,  
    Department     int,  
    Supervisor     int references Employee(EmpID),  
    )
```

# Exercise: Identify the Foreign Keys

## EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	------------	-------	---------	-----	--------	-----------	-----

## DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
-------	----------------	---------	----------------

## DEPT\_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
----------------	------------------

## PROJECT

Pname	<u>Pnumber</u>	Plocation	Dnum
-------	----------------	-----------	------

## WORKS\_ON

<u>Essn</u>	<u>Pno</u>	Hours
-------------	------------	-------

## DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
-------------	-----------------------	-----	-------	--------------

# Exercise: Identify Referential integrity constraints

## EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	------------	-------	---------	-----	--------	-----------	-----

## DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
-------	----------------	---------	----------------

## DEPT\_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
----------------	------------------

## PROJECT

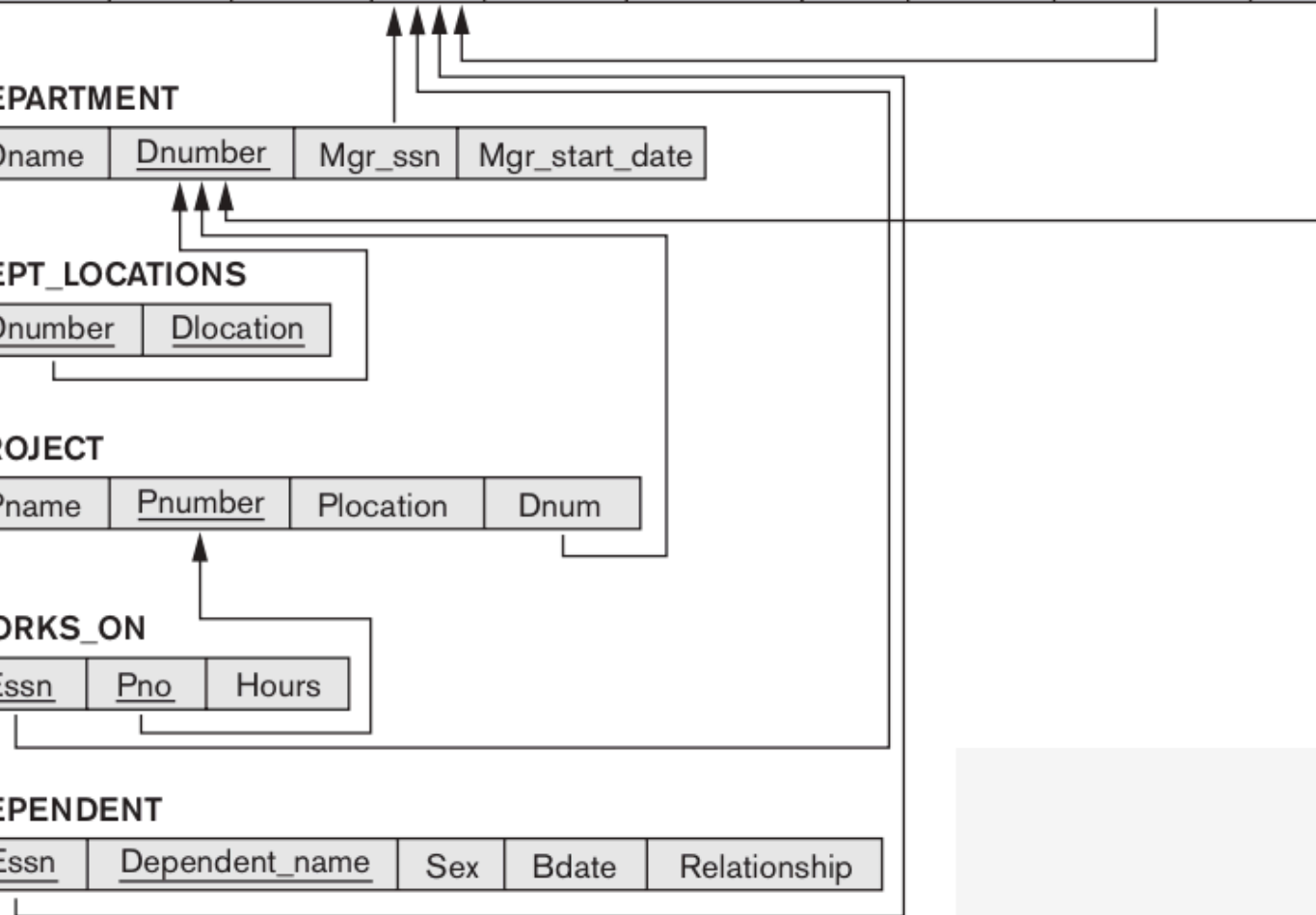
Pname	<u>Pnumber</u>	Plocation	Dnum
-------	----------------	-----------	------

## WORKS\_ON

<u>Essn</u>	<u>Pno</u>	Hours
-------------	------------	-------

## DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
-------------	-----------------------	-----	-------	--------------



# A summary

```
CREATE TABLE Employee (  
    EmpID          int PRIMARY KEY,  
    SSN            int UNIQUE,  
    Name           varchar NOT NULL,  
    Surname       varchar NOT NULL,  
    Salary        int,  
    Department int  
        references Department(Number),  
    Supervisor   int  
        references Employee(EmpID),  
)
```

```
CREATE TABLE Department (  
    Number int PRIMARY KEY,  
    Name   varchar NOT NULL  
)
```

```
CREATE TABLE Manager (  
    EmpID      int  
        references Employee(EmpID),  
    Department int  
        references Department(Number),  
    StartDate  date,  
    PRIMARY KEY(EmpID,Department)  
)
```

# Insertion, deletion and update

**INSERT INTO** Employee  
VALUES (1, 'John', ...

**INSERT INTO** Manager  
<SQL query>

**DELETE FROM** Employee  
WHERE EmpID = 1

**UPDATE** Employee  
SET EmpID = 3  
WHERE EmpID = 1

# Integrity constraints: Reactions

- Deleting tuples or modifying values in a relation may make references inconsistent.
- We can instruct the database to automatically ensure that referential constraints are kept consistent.
- Possible instructions are:

**on delete**

**< cascade | set null | set default | no action >**

**on update**

**< cascade | set null | set default | no action >**

# Delete reactions

- **cascade**: deletes are propagated.
- **set null**: the referring attribute values are set to null.
- **set default**: the referring attribute values are set to their default.
- **no action**: deletion is not authorized.



# Update reactions

- **cascade**: the value is updated also on the referring attributes.
- **set null**: the referring attribute values are set to null.
- **set default**: the referring attribute values are set to their default.
- **no action**: deletion is not authorized.

# Reactions to deletion / updates

```
CREATE TABLE Employee (  
    EmpID        int PRIMARY KEY,  
    SSN          int UNIQUE,  
    Name         varchar NOT NULL,  
    Surname      varchar NOT NULL,  
    Salary       int,  
    Department   int  
        references Department(Number)  
        on delete set null,  
    Supervisor   int  
        references Employee(EmpID)  
        on update cascade,  
    unique (Name, Surname)  
)
```

# Schema updates

- It is possible to modify the design of a relational database using two main statements:

- **ALTER**

- **DROP**

Table name

Only if the  
table is empty

DROP TABLE Employee

||

DROP TABLE Employee **RESTRICT**

DROP TABLE Employee **CASCADE**

# Details about the syntax

MySQL 5.7 Reference Manual :: 13 SQL Statement Syntax

## Chapter 13 SQL Statement Syntax

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This chapter describes the syntax for the [SQL](#) statements supported by MySQL.

# Relational keys: Summary

We need to identify specific tuples in a relation.

- **Superkey:** a set of attributes that allows us to do it.
  - A relation can have many superkeys.
- **Candidate key:** a minimal superkey.
  - A relation can have many candidate keys – normally, less than the superkeys.
  - UNIQUE in SQL.
- **Primary key:** the single candidate key we choose to identify records in a table.
  - PRIMARY KEY in SQL.
  - No NULLs allowed.
- **Foreign key:** used to “connect” relations and reference tuples in other relations.
  - FOREIGN KEY ... REFERENCES in SQL.