

Languages

Language: a set of strings

String: a sequence of symbols
from some alphabet

Example:

Strings: cat, dog, house

Language: {cat, dog, house}

Alphabet: $\Sigma = \{a, b, c, \square, z\}$

Languages are used to describe
computation problems:

$$PRIMES = \{2, 3, 5, 7, 11, 13, 17, \square \}$$

$$EVEN = \{0, 2, 4, 6, \square \}$$

Alphabet: $\Sigma = \{0, 1, 2, \square, 9\}$

Computation is translated to set membership

Example computation problem:

Is number x prime?

Equivalent set membership problem:

$$x \in PRIMES = \{2, 3, 5, 7, 11, 13, 17, \dots\}?$$

Alphabets and Strings

An alphabet is a set of symbols

Example Alphabet: $\Sigma = \{a, b\}$

A string is a sequence of symbols from the alphabet

Example Strings

a

ab

abba

aaabbbbaabab

String variables

u = ab

v = bbbbaaa

w = abba

Decimal numbers alphabet $\Sigma = \{0,1,2,\dots,9\}$

102345

567463386

Binary numbers alphabet $\Sigma = \{0,1\}$

100010001

101101111

Unary numbers alphabet $\Sigma = \{1\}$

Unary number: 1 11 111 1111 11111

Decimal number: 1 2 3 4 5

String Operations

$$w = a_1 a_2 \boxtimes \dots \boxtimes a_n$$

aaa

$$v = b_1 b_2 \boxtimes \dots \boxtimes b_m$$

bbb

Concatenation

$$wv = a_1 a_2 \boxtimes \dots \boxtimes a_n b_1 b_2 \boxtimes \dots \boxtimes b_m$$

aaabbbb

$$v = a_1 a_2 \boxtimes a_n \qquad w = ababaaabbb$$

Reverse

$$w^R = bbbbaaababa$$

String Length

$$w = a_1 a_2 \boxtimes a_n$$

Length: $|w| = n$

Examples: $|abba| = 4$

$$|aa| = 2$$

$$|a| = 1$$

Length of Concatenation

$$|uv| = |u| + |v|$$

Example: $u = aab, \quad |u| = 3$
 $v = abaab, \quad |v| = 5$

$$|uv| = |aababaab| = 8$$

$$|uv| = |u| + |v| = 3 + 5 = 8$$

Empty String

A string with no letters is denoted: λ or ε

Acts as a neutral element

Observations: $|\lambda| = 0$

$$\lambda w = w\lambda = w$$

$$\lambda abba = abba\lambda = ab\lambda ba = abba$$

Substring

Substring of string:

a subsequence of consecutive characters

String

Substring

abbab

ab

abbab

abba

abbab

b

abbab

bbab

Prefix and Suffix

string *abbab*

Prefixes

Suffixes

λ

abbab

a

bbab

ab

bab

abb

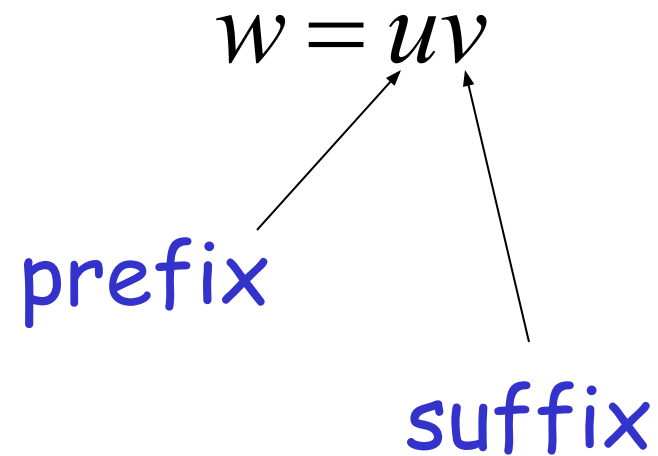
ab

abba

b

abbab

λ



Exponent Operation

$$w^n = \underbrace{w w w \dots w}_n$$

Example: $(abba)^2 = abbaabba$

Definition: $w^0 = \lambda$

$$(abba)^0 = \lambda$$

The * Operation

Σ^* : the set of all possible strings from
alphabet Σ

$$\Sigma = \{a, b\}$$

$$\Sigma^* = \{\lambda, a, b, aa, ab, ba, bb, aaa, aab, \dots\}$$

The + Operation

Σ^+ : the set of all possible strings from alphabet Σ except λ

$$\Sigma = \{a, b\}$$

$$\Sigma^* = \{\lambda, a, b, aa, ab, ba, bb, aaa, aab, \Box \}$$

$$\Sigma^+ = \Sigma^* - \lambda$$

$$\Sigma^+ = \{a, b, aa, ab, ba, bb, aaa, aab, \Box \}$$

Powers of an alphabet

Let Σ be an alphabet.

- Σ^k = the set of all strings of length k
- $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$
- $\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \dots$

Languages

A language over alphabet Σ
is any subset of Σ^*

Example: $\Sigma = \{a, b\}$

$$\Sigma^* = \{\lambda, a, b, aa, ab, ba, bb, aaa, \dots\}$$

Languages: $\{\}$

$$\{\lambda\}$$

$$\{a, aa, aab\}$$

$$\{\lambda, abba, baba, aa, ab, aaaaaa\}$$

More Language Examples

Alphabet $\Sigma = \{a, b\}$

An infinite language $L = \{a^n b^n : n \geq 0\}$

λ
 ab
 $aabb$
 $aaaaabbbbbb$

} $\in L$

$bbabb \notin L$

$abb \notin L$

Prime numbers

Rule: Numbers divisible by 1 and itself

Alphabet $\Sigma = \{0,1,2,\square,9\}$

Language:

$PRIMES = \{x : x \in \Sigma^* \text{ and } x \text{ is prime}\}$

$PRIMES = \{2,3,5,7,11,13,17,\square\}$

Even and odd numbers

Alphabet $\Sigma = \{0,1,2,\dots,9\}$

Languages:

$EVEN = \{x : x \in \Sigma^* \text{ and } x \text{ is even}\}$

$EVEN = \{0,2,4,6,\dots\}$

$ODD = \{x : x \in \Sigma^* \text{ and } x \text{ is odd}\}$

$ODD = \{1,3,5,7,\dots\}$

Addition (of unary numbers)

Alphabet: $\Sigma = \{1, +, =\}$

Language:

$$ADDITION = \{x + y = z : x = 1^n, y = 1^m, z = 1^k, \\ n + m = k, n \geq 1, m \geq 1\}$$

$$11 + 111 = 11111 \in ADDITION$$

$$111 + 111 = 111 \notin ADDITION$$

$$ADDITION = \{1 + 1 = 11, 1 + 11 = 111, 11 + 1 = 111, 11 + 11 = 1111, \dots\}$$

Squares (of unary numbers)

Alphabet: $\Sigma = \{1, \#\}$

Language:

$$SQUARES = \{x\#y : x = 1^n, y = 1^m, m = n^2\}$$

$$11\#1111 \in SQUARES$$

$$111\#1111 \notin SQUARES$$

$$SQUARES = \{1\#1, 11\#1111, 111\#1111111111, \dots\}$$

Two special languages

Empty language

$\{\}$ or \emptyset

Language with
empty string

$\{\lambda\}$

Size of a language (number of elements):

$$|\{\}| = 0$$

$$|\{\lambda\}| = 1$$

$$|\{a, aa, ab\}| = 3$$

$$|\{\lambda, aa, bb, abba, baba\}| = 5$$

Note that:

Sets

$$\emptyset = \{\} \neq \{\lambda\}$$

Set size

$$|\{\}| = |\emptyset| = 0$$

Set size

$$|\{\lambda\}| = 1$$

String length

$$|\lambda| = 0$$

Operations on Languages

The usual set operations:

$$\begin{aligned}\{a, ab, aaaa\} \cup \{bb, ab\} &= \{a, ab, bb, aaaa\} && \text{union} \\ \{a, ab, aaaa\} \cap \{bb, ab\} &= \{ab\} && \text{intersection} \\ \{a, ab, aaaa\} - \{bb, ab\} &= \{a, aaaa\} && \text{difference}\end{aligned}$$

Complement: $\bar{L} = \Sigma^* - L$

$$\overline{\{a, ba\}} = \{\lambda, b, aa, ab, bb, aaaa, \boxtimes\}$$

Reverse

Definition: $L^R = \{w^R : w \in L\}$

Examples: $\{ab, aab, baba\}^R = \{ba, baa, abab\}$

$$L = \{a^n b^n : n \geq 0\}$$

$$L^R = \{b^n a^n : n \geq 0\}$$

Concatenation

Definition: $L_1L_2 = \{xy : x \in L_1, y \in L_2\}$

Example: $\{a, ab, ba\}\{b, aa\}$

$$= \{ab, aaa, abb, abaa, bab, baaa\}$$

Another Operation

Definition: $L^n = \underbrace{L \boxtimes L \boxtimes \dots \boxtimes L}_n$

$$\{a, b\}^3 = \{a, b\}\{a, b\}\{a, b\} = \{aaaa, aab, aba, abb, baa, bab, bba, bbb\}$$

Special case: $L^0 = \{\lambda\}$

$$\{a, bba, aaa\}^0 = \{\lambda\}$$

Example

$$L = \{a^n b^n : n \geq 0\}$$

$$L^2 = \{a^n b^n a^m b^m : n, m \geq 0\}$$

$$aabbbaaabb \in L^2$$

$$??????? \notin L^2$$

Star-Closure (Kleene *)

All strings that can be constructed from L

Definition: $L^* = L^0 \cup L^1 \cup L^2 \cup \dots$

Example:

$$\{a, bb\}^* = \left\{ \begin{array}{l} \lambda, \\ a, bb, \\ aa, abb, bba, bbbb, \\ aaa, aabb, abba, abbbb, \dots \end{array} \right\} \begin{array}{l} L^0 \\ L^1 \\ L^2 \\ L^3 \end{array}$$

Kleene's Closure/ Kleene's Star

- Given Σ , then the Kleene Star Closure of the alphabet Σ , denoted by Σ^* , is the collection of all strings defined over Σ , including Λ .
- It is to be noted that Kleene Star Closure can be defined over any set of strings.

Examples

If $\Sigma = \{x\}$

$$\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$$

Then $\Sigma^* = \{\Lambda, x, xx, xxx, xxxx, \dots\}$

If $\Sigma = \{0,1\}$

Then $\Sigma^* = \{\Lambda, 0, 1, 00, 01, 10, 11, \dots\}$

If $\Sigma = \{aaB, c\}$

Then $\Sigma^* = \{\Lambda, aaB, c, aaBaaB, aaBc, caaB, cc, \dots\}$

Note

Languages generated by Kleene Star Closure of set of strings, are infinite languages. (By infinite language, it is supposed that the language contains infinite many words, each of finite length).

Positive Closure

Definition: $L^+ = L^1 \cup L^2 \cup L^3 \cup \Box$

Note that: $L^* = L^0 \cup L^+$

$$\{a, bb\}^+ = \left\{ \begin{array}{l} a, bb, \\ aa, abb, bba, bbbb, \\ aaa, aabb, abba, abbbb, \Box \end{array} \right\} \begin{array}{l} L^1 \\ L^2 \\ L^3 \end{array}$$

Valid/In-valid alphabets

While defining an alphabet, an alphabet may contain letters consisting of group of symbols for

Example

$\Sigma_1 = \{B, aB, bab, d\}$.

Now consider an alphabet

$\Sigma_2 = \{B, Ba, bab, d\}$ and a string **BababB**.

This string can be tokenized in two different ways

(Ba), (bab), (B)

(B), (abab), (B)

Which shows that the second group cannot be identified as a string, defined over $\Sigma = \{a, b\}$.

As when this string is scanned by the compiler (Lexical Analyzer), first symbol B is identified as a letter belonging to Σ , while for the second letter the lexical analyzer would not be able to identify, so while defining an alphabet it should be kept in mind that ambiguity should not be created.

Remarks

- While defining an alphabet of letters consisting of more than one symbols, no letter should be started with the letter of the same alphabet *i.e. one letter should not be the prefix of another. However, a letter may be ended in a letter of same alphabet.*

Conclusion

$\Sigma_1 = \{B, aB, bab, d\}$

$\Sigma_2 = \{B, Ba, bab, d\}$

Σ_1 is a valid alphabet while Σ_2 is an in-valid alphabet.

Descriptive definition of language

The language is defined, describing the conditions imposed on its words.

Example

The language L of strings of odd length, defined over $\Sigma = \{a\}$, can be written as

$$L = \{a, aaa, aaaaa, \dots\}$$

Example

The language L of strings that does not start with a , defined over $\Sigma = \{a, b, c\}$, can be written as

$$L = \{b, c, ba, bb, bc, ca, cb, cc, \dots\}$$

Example

The language L of strings of length 2, defined over $\Sigma = \{0,1,2\}$, can be written as

$$L = \{00, 01, 02, 10, 11, 12, 20, 21, 22\}$$

Example

The language L of strings ending in 0, defined over $\Sigma = \{0,1\}$, can be written as

$$L = \{00, 10, 000, 010, 100, 110, \dots\}$$

Example

The language **EQUAL**, of strings with number of a's equal to number of b's, defined over $\Sigma = \{a,b\}$, can be written as

$$L = \{\Lambda, ab, ba, aabb, abab, baba, abba, \dots\}$$

Example

The language **EVEN-EVEN**, of strings with even number of a's and even number of b's, defined over $\Sigma = \{a,b\}$, can be written as

$$L = \{\Lambda, aa, bb, aaaa, aabb, abab, abba, baab, baba, bbaa, bbbb, \dots\}$$

Example

The language **INTEGER**, of strings defined over $\Sigma = \{-, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, can be written as
$$\text{INTEGER} = \{\dots, -2, -1, 0, 1, 2, \dots\}$$

Example

The language **EVEN**, of strings defined over $\Sigma = \{-, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, can be written as
$$\text{EVEN} = \{\dots, -4, -2, 0, 2, 4, \dots\}$$

Example

The language $\{a^n b^n\}$, of strings defined over $\Sigma = \{a, b\}$, as $\{a^n b^n : n=1, 2, 3, \dots\}$, can be written as
$$L = \{ab, aabb, aaabbb, aaaabbbb, \dots\}$$

Example

The language $\{a^n b^n a^n\}$, of strings defined over $\Sigma = \{a, b\}$, as $\{a^n b^n a^n : n=1, 2, 3, \dots\}$, can be written as
$$L = \{aba, aabbbaa, aaabbbbaaa, aaaabbbbbaaaaa, \dots\}$$

Example

The language **factorial**, of strings defined over $\Sigma=\{0,1,2,3,4,5,6,7,8,9\}$ i.e.

$\{1,2,6,24,120,\dots\}$

Example

The language **FACTORIAL**, of strings defined over $\Sigma=\{a\}$, as $\{a^n! : n=1,2,3,\dots\}$, can be written as $\{a,aa,aaaaaa,\dots\}$.

It is to be noted that the language **FACTORIAL** can be defined over any single letter alphabet.

Example

The language **DOUBLE-FACTORIAL**, of strings defined over $\Sigma=\{a, b\}$, as

$\{a^n!b^n! : n=1,2,3,\dots\}$, can be written as $\{ab, aabb, aaaaaabbbbbbb,\dots\}$

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