



Course:  
Program:  
Date:  
Section:  
Roll No:  
Quiz:

Advance Database Concepts  
BS (Computer Science)  
Tue 21-Feb-2023  
BCS-6A  
1 (CCT)

Course Code:  
Semester:  
Total Marks:

CS4064  
Spring 2023  
15

Q. Consider the following schedule of actions, listed in the order they are submitted to the DBMS:

S:  $r_1(X)$ ,  $w_2(X)$ ,  $w_1(X)$ ,  $r_3(X)$ ,  $c_1$ ,  $c_2$ ,  $c_3$ .

For each of the following concurrency control mechanisms, describe how the concurrency control mechanism handles the schedule. Assume that the timestamp of transaction  $T_i$  is  $i$ . For lock-based concurrency control mechanisms, add lock and unlock requests to the above schedule of actions as per the locking protocol. The DBMS processes actions in the order shown. If a transaction is blocked, assume that all its actions are queued until it is resumed; the DBMS continues with the next action (according to the listed schedule) of an unblocked transaction.

1. Rigorous 2PL with timestamps used for deadlock detection (Use wait-for-graph to deal with deadlock)
2. Strict Timestamp Ordering (Assume  $T_1 < T_2 < T_3$ )
3. Validation Concurrency Control Technique (Use defer the validation until a later time when the conflicting transactions have finished)

1 -  $T_1$        $T_2$        $T_3$        $R/H$

$r_1(X)$   
 $r_1(X)$

$w_2(X)$   
wait for  $T_1$

$w_1(X)$   
upgrade lock  
 $w_1(X)$

$r_3(X)$  wait for  $T_1$

$c_1$   
 $u_1(X) \rightarrow$

$w_2(X)$   
 $c_2$  wakeup call  
 $u_2(X)$

$r_3(X)$  wakeup call  
 $c_3$

2 - (STRICT TO)

$T_1$

$T_2$

$T_3$

$\boxed{XWR}$

$\boxed{XRD}$

$r_1(x)$

$\{T_1\}$

$w_2(x)$

$\{T_2\}$

$w_1(x)$

abort

restarts

later

$r_3(x)$

waits for  $T_2$

↓

wakeup call

$r_3(x)$

$C_3$

$C_2$

$\{T_3\}$

3-  $T_1$ :

BV: True (As there are no overlapping transactions)

FV:  $\{x\} \cap \{x\} \neq \emptyset$  fails (delay)  
 $T_1(WS)$   $T_2, T_3(RS)$

$T_2$ : BV: True

FV:  $\{x\} \cap \{x\} \neq \emptyset$  Fails  
 $T_2(WS)$   $T_1/T_3(RS)$

$T_3$ : BV: True

FV:  $\{ \} \cap \{x\} = \emptyset$  True  
 $T_3(WS)$   $T_1/T_2(RS)$

$\therefore T_3$  successfully completed.

But revalidations for  $T_1 \wedge T_2$ :

$T_1$ : BV:  $\{x\} \cap \{ \} = \emptyset$  True  
 $T_1(RS)$   $T_3(WS)$

FV:  $\{x\} \cap \{ \} = \emptyset$  True  
 $T_1(WS)$   $T_2(RS)$

$T_1$  completed successfully

$T_2$ :  $T_2(RS)$   $T_1/T_2(WS)$

BV:  $\{ \} \cap \{x\} = \emptyset$  True

FV:  $\{x\} \cap \{ \} = \emptyset$  True  
 No active are present

$T_2$  completed successfully