#### CS525: Advanced Database Organization

#### Notes 1: Introduction to DBMS Implementation

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Slides: adapted from a course taught by Hector Garcia-Molina, Stanford

#### Core Terminology Review

- Data
  - any information worth preserving, most likely in electronic form
- Database
  - organized collection of interrelated data that models some aspect of the real-world.
- Query
  - an operation that extracts specified data from the database.
- Relation
  - an organization of data into a two-dimensional table, where rows (tuples) represent basic entities or facts of some sort, and columns (attributes) represent properties of those entities.
- Schema
  - a description of the structure of the data in a database, often called "metadata"

#### Database Management System (DBMS)

- A DBMS is software that allows applications to store and analyze information in a database.
- A general-purpose DBMS is designed to allow the definition, creation, querying, update, and administration of databases.

#### Advanced Database Organization?

- = Database Implementation
- = How to implement a database system
- and have fun doing it ;-)

#### What do you want from a DBMS?

- Keep data around (persistent)
- Answer questions (queries) about data
- Update data

## Isn't Implementing a Database System Simple?

■ Relation ⇒ Statements ⇒ Results

# Introduction the $M \, \mathrm{EGATRON} \, \, 3000$ Database Management System

- "Imaginary" database System
- The latest from MEGATRON Labs
- Incorporates latest relational technology
- UNIX compatible
- Lightweight & cheap!

#### MEGATRON 3000 Implementation Details

- ullet MEGATRON 3000 uses the file system to store its relations
- R elations stored in files (ASCII)
- Use a separate file per entity/relation.
- The application has to parse the files each time they want to read/update records.
  - e.g., relation Students(name,id,dept) is in /usr/db/Students
  - The file Students has one line for each tuple.
  - Values of components of a tuple are stored as a character string, separated by special marker character #

Smith	#	123	#	CS
Jonson	#	522	#	ΕE
		•		

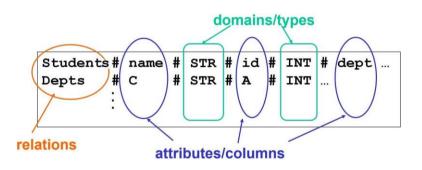
## MEGATRON 3000 Implementation Details

- The database schema is stored in a special file
- Schema file (ASCII) in /usr/db/schema
  - For each relation, the file schema has a line beginning with that relation name, in which attribute names alternate with types.
  - The character # separates elements of these lines.

```
Students # name # STR # id # INT # dept ...

Depts # C # STR # A # INT ...
```

## $MEGATRON\ 3000$ Implementation Details

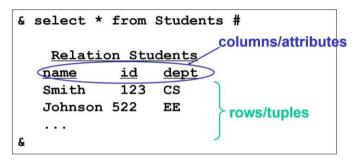


## $MEGATRON\ 3000\ Sample\ Sessions$

```
% MEGATRON3000
    Welcome to MEGATRON 3000!
&
    :
    :
    quit
%
```

 $\,$  We are now talking to the  $\rm MEGATRON~3000$  user interface, to which we can type SQL queries in response to the Megatron prompt (&).

## MEGATRON 3000 Sample Sessions



A # ends a query

#### MEGATRON 3000 Sample Sessions

Execute a query and send the result to printer

```
& select *
from Students | LPR #
&
```

Result sent to LPR (printer).

## $MEGATRON\ 3000\ Sample\ Sessions$

Execute a query and store the result in a new file

```
& select *
  from Students
  where id < 100 | LowId #
&</pre>
```

New relation LowId created.

#### How MEGATRON 3000 Executes Queries

To execute

```
SELECT * FROM R WHERE <condition>
```

- Read schema to get attributes of R
- Check validity of condition
- Display attributes of R as the header
- Read file R; for each line:
  - Check condition
  - If TRUE, display the line as tuple

## MEGATRON 3000 Query Execution

To execute

```
SELECT * FROM R WHERE <condition> | T
```

- Process select as before but omit Step 3
- Write results to new file T
  - Append new line to dictionary usr/db/schema

#### MEGATRON 3000 Query Execution

- Consider a more complicated query, one involving a join of two relations
   R, S
- To execute

```
SELECT A, B FROM R, S WHERE <condition>
```

- Read schema to get R,S attributes
- Read R file, for each line r:
  - Read S file, for each line s:
    - f a Create join tuple f r & f s
    - Check condition
    - If TRUE, Display r, s[A, B]

- DBMS is not implemented like our imaginary MEGATRON 3000
- Described implementation is inadequate for applications involving significant amount of data or multiple users of data
- Partial list of problems follows

- Tuple layout on disk is inadequate with no flexibility when the database is modified
- e.g., change String from CS to CSDept in one Students tuple, we have to rewrite the entire file
  - ASCII storage is expensive
  - Deletions are expensive

- Search expensive; no indexes
  - e.g., cannot find tuple with given key quickly
  - Always have to read full relation

- Brute force query processing
- e.g.,

```
SELECT * FROM R, S WHERE R.A = S.A and S.B > 1000
```

- Much better if use index to select tuples that satisfy condition (Do select using S.B >1000 first)
- More efficient join (sort both relations on A and merge)

- No buffer manager
  - There is no way for useful data to be buffered in main memory; all data comes off the disk, all the time
  - e.g., need caching.

- No concurrency control
  - Several users can modify a file at the same time with unpredictable results.

- No reliability
- e.g., in case of error/crash, say, power failure or leave operations half done
  - Can lose data

- No security
- e.g., file system security is coarse
  - Unable to restrict access, say, to some fields of a relation and not others

- No application program interface (API)
  - e.g., how can a payroll program get at the data?

Cannot interact with other DBMSs.

No GUI

#### This Course

 Introduce students to better way of building a database management systems.

#### Reading assignment

- Refresh your memory about basics of the relational model and SQL
  - from your earlier course notes
  - from some textbook
  - http://cs.iit.edu/~cs425/schedule.html



#### Next

Notes 2: Hardware