## Centrality and Other Measures

CS 579 Online Social Network Analysis

Dr. Cindy Hood 9/23/25

### Homework Assignments

- HW #3 Network Metrics
  - Assigned soon
  - Good prep for Exam 1
- HW #4 Chicago Community Areas + Census Data
  - ▶ You may work in groups up to 4 students (no exceptions) on this hw
  - Assigned later this week
- Please contact TAs with questions on hw grading

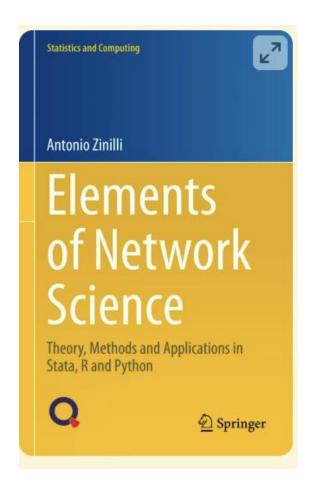
# Exams and Final Project Poster Presentation

- Exam 1 Oct 9 in class
- Exam 2 Dec 2 in class
- ► Final Project Poster Session Dec 4 in class
- Online students (sections 2 and 3) will have remote options

## **Teaching Assistants**

- Siva Krishna Golla
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  - Mondays 2-3pm on zoom
- Khush Dhiren Patel
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  - ► Wednesdays 11-12 online
- Aswith Sama
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  - ► Thursdays 3-4pm on zoom

#### Reference



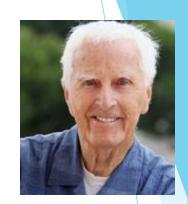
https://link.springer.com/book/10.1007/978-3-031-84712-7

#### Recall - Degree Centrality

- Total number of connections a vertex has
  - Degree of vertex
    - ► Total number of edges connected to a vertex
- Directed network
  - ▶ In-degree
  - Out-degree
- Can be considered a popularity measure
  - ► Is it a good popularity measure?

#### **Recall - Betweenness Centrality**

Another way of looking at centrality is by considering how important nodes are in connecting other nodes



Linton Freeman

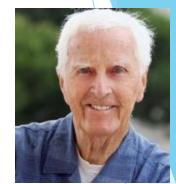
$$C_b(v_i) = \sum_{s \neq t \neq v_i} \frac{\sigma_{st}(v_i)}{\sigma_{st}}$$

 $\sigma_{st}$  The number of shortest paths from vertex s to t – a.k.a. information pathways

 $\sigma_{st}(v_i)$  The number of **shortest paths** from s to t that pass through  $v_i$ 

#### **Recall - Closeness Centrality**

The intuition is that influential/central nodes can quickly reach other nodes



These nodes should have a smaller average shortest path length to others

Linton Freeman

Closeness centrality: 
$$C_c(v_i) = \frac{1}{\overline{l}_{v_i}}$$

$$\bar{l}_{v_i} = \frac{1}{n-1} \sum_{v_j \neq v_i} l_{i,j}$$

#### Page Rank Centrality

- Measure of global centrality
  - Uses entire network to assess the significance of single node
- Assesses node's relevance based on importance of nodes that link to it
- Iterative process that estimates a node's importance in the network
  - Calculate centrality of a node based on importance of neighbors

#### Page Rank Calculation (Simple)

To begin with, each node is equally important:

$$PR_i^0 = \frac{1}{n}, \forall i$$
 n = number of nodes/pages

- Each node distributes its centrality equally to the nodes it links to (outgoing)
- Page Rank of a node is calculated by adding the Page Rank fractions of the vertices that have edges to it from the previous iteration

$$PR_i^{t} = \sum_{j=1}^{n} a_{ji} \frac{PR_j^{t-1}}{d_j^{out}}, \forall t \in \{1,k\}$$
 t = iteration

where  $d_j^{out}$  is the number of hyperlinks on the node/page j.

#### Page Rank Calculation

► The first iteration is

$$PR_i^1 = \frac{1}{n} \sum_{j=1}^n a_{ji} \frac{1}{d_j^{out}}$$

- Iterations continue until Page Rank values converge
  - ▶ When the difference between PR<sup>t+1</sup> and PR<sup>t</sup> is very small and is converging to 0

## Page Rank Example

Fig. 3.7 Graph example

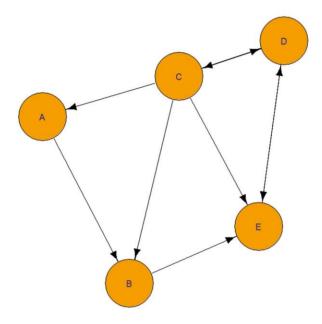
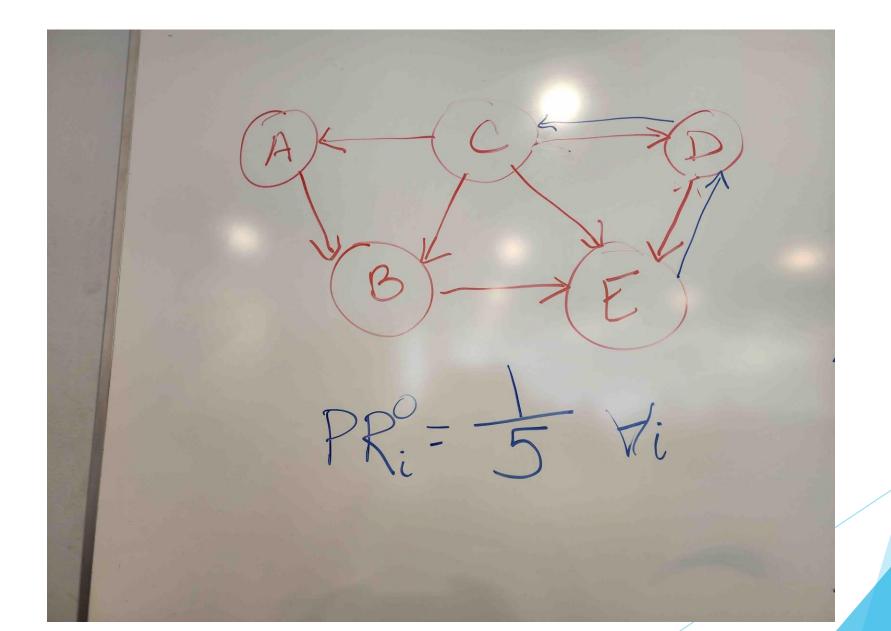
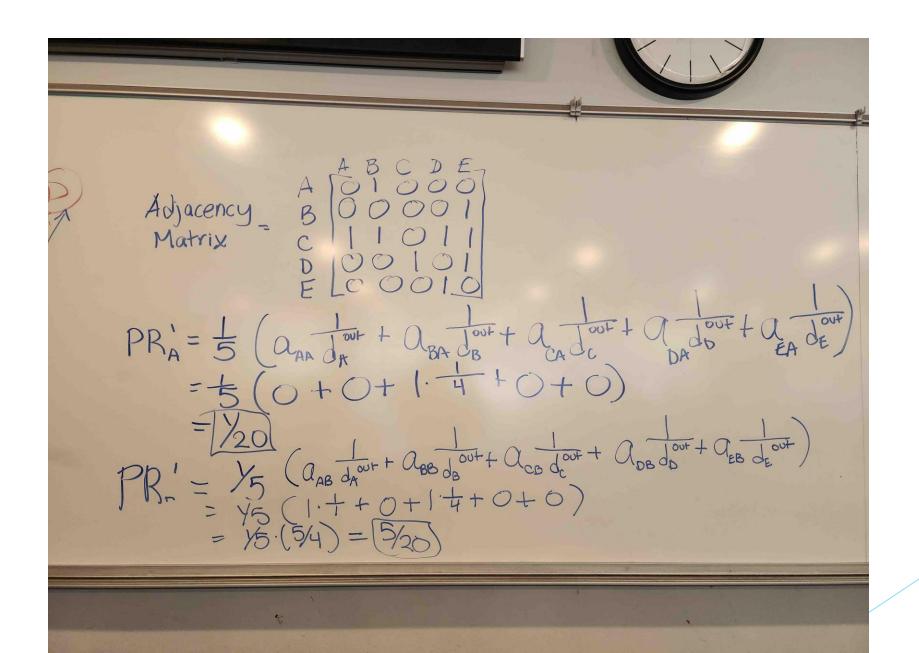


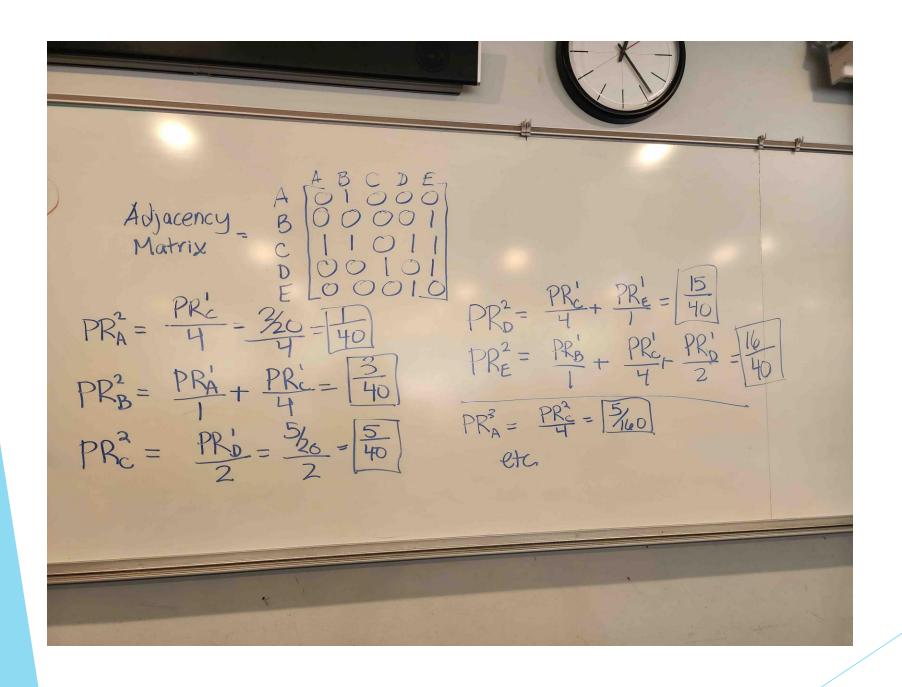
 Table 3.2 Example of PageRank computation (second step)

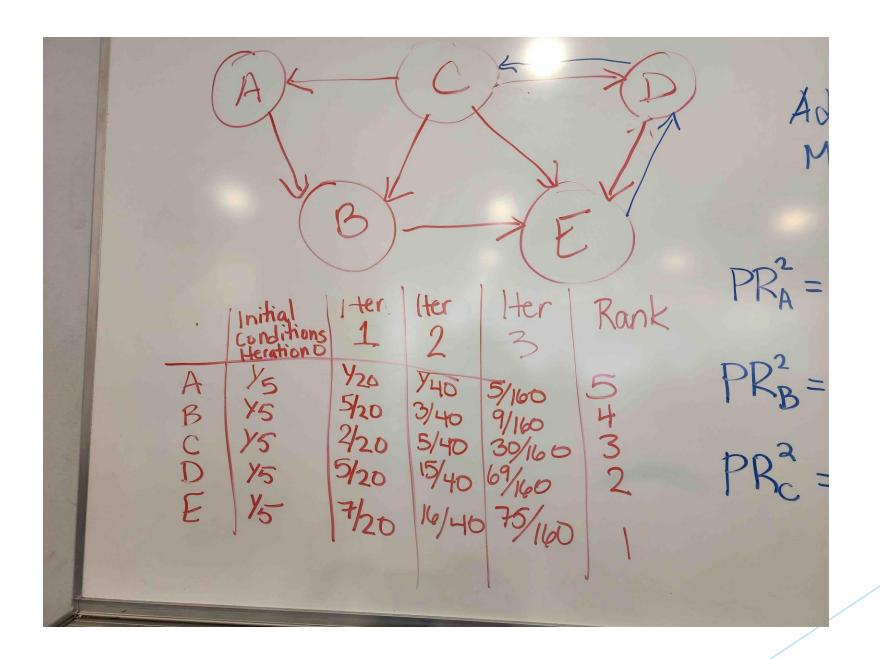
Nodes	Iteration 1	Iteration 2	Iteration 3	Final rank
A	1/5	1/20	1/40	5
В	1/5	5/20	3/40	4
С	1/5	1/10	5/40	3
D	1/5	5/20	15/40	2
E	1/5	7/20	16/40	1





PRo = 5 ( and down to alborate and down to a d = 15(0+1+1+1+1-2+0)= 1/20





### Node Strength (for weighted graphs)

For an undirected network, the W matrix defines the strength  $s_i$  of a node as the sum of the weights of the edges incident on a node as follows:

$$s_i = \sum_{j \in \mathbb{N}} w_{ij} \tag{3.12}$$

If the network is directed, there are two factors that determine the node strength: the quantity of incoming  $(s_i^{\text{in}})$  and outgoing  $(s_i^{\text{out}})$  weighted edges:

$$s_{i}^{in} = \sum_{j} w_{ji}$$

$$s_{i}^{out} = \sum_{j \in \mathbb{N}} w_{ij}$$

$$s_{i} = \sum_{j \in \mathbb{N}} w_{ij}$$
(3.13)

Total strength is finally defined as:

$$s_i = s_i^{in} + s_i^{out} \tag{3.14}$$

#### Node Strength

- Similar to degree, can obtain the strength distribution P(s)
- ▶ In case of directed networks, two distributions
  - $ightharpoonup P(s^{in})$
  - ► P(s<sup>out</sup>)
- Nodes with significantly greater strength than the rest of the network's nodes are referred to as hubs

Where might we use weighted graphs?

#### Recall Node's Neighborhood

A nodes neighborhood includes the group of nodes that it is connected to

$$N_i(G) = \left\{j : g_{ij} = 1\right\}$$

- Studying the neighborhood aids in identifying nodes or communities of nodes that share specific characteristics
- When individuals are motivated to emulate the behavior of their neighbors, cascading effects can occur
  - Can be observed when a new behavior initiates with a small group of early adopters and then spreads radially outward through the network

#### Neighborhood Degree Sequence

The neighborhood degree sequence for the node i,  $s_i$ , is derived from:

$$s_{i} = \left\{k_{1}^{i}, k_{2}^{i} k_{3}^{i} \dots, k_{n}^{i}\right\}$$
 (3.19)

with  $k_i$  equal to the degree of the nodes to which i is connected. Being a sequence of degrees:

$$k_1^i \le k_2^i \le k_3^i \dots \le k_n^i \tag{3.20}$$

#### Neighborhood Degree Sequence

- Analysis
  - May help in examining variance of neighborhood degree sequences
    - Context of hierarchical complexity
  - May analyze neighbors of same degree to determine how similar they are

#### Example Neighborhood Degree Sequence

Fig. 3.4 Graph example

Extra Credit #4: List degree sequences for each node

Section 1: in class only

Sections 2 and 3: submit to canvas

by 9/25

