

# Practical Stuff on Linux

## Disk encryption

— `dm-crypt` (`dm` = device mapper)

`/dev/sda2`  $\equiv$  your real hard drive.  
everything is encrypted.

`dm-crypt` gives an abstraction layer on top of the real disk:

\$ `cryptsetup -- /dev/sda2 cryptroot`

now `/dev/mapper/cryptroot` works like a hard drive. Setup f.s. on `cryptroot`!

Note: if you have a magnetic HD, first overwrite w/ random bytes.

LUKS does a good job of using passwords to access HD.

- Strong KDF by default.
- Downside: lose plausible deniability
- upside: deleting the header wipes out the disk (back up LUKS header?)
- Multiple passwords can be used!

Header has something like this:

$E_{\text{pwd}_0}(k)$

$E_{\text{pwd}_1}(k)$

$\vdots$

$E_{\text{pwd}_n}(k)$

$k \equiv$  actual AES key for the disk.

Note: can also encrypt ~~boot~~ if you have grub 2  
as your boot loader.  
kernel images  
live here!

Bad news: cold boot attacks!

DRAM decay slows with colder temperatures...

Most bits might be unchanged after 5-10 minutes  
even with a crude cooling approach!

Forcefully power off computer & reboot  
(from attacker's usb)

Countermeasure: lock down BIOS - don't boot from  
usb w/o another pwd.

Counter-counter measure: remove DRAM, place in another  
computer

Counter-counter-counter measure: solder it??

Other notes: how to find keys? How to recover  
if they have a few bits toggled?

Bad news:  $\exists$  redundant info that can be used  
as a sort of error correction

(e.g. key schedules for symmetric enc,  
& public keys)

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Authentication, etc.

Recall best practices for handling passwords:

- Store hash
- make hash expensive to compute  $\leftarrow$  slows down dictionary attack
- salt  $\leftarrow$  prevents "offline" attack

On linux: add "rounds = XXX..." in /etc/passwd/passwd

Documentation: man crypt  
man unix-passwd  
man shadow

What about non-system passwords?

password safes like keepassX aren't bad.

gpg can also be used for something basic.

SSH public key authentication FTW!

In your sshd\_config look for these:

PasswordAuthentication no  
ChallengeResponseAuthentication no  
AllowUsers yourusername  
PermitRootLogin no

Protecting your ssh private key: use ssh-keygen -a rounds

Protecting sshd from randos/port scanners:

cool trick with iptables: "port knocking"

1337  
12345  
9999  
22

attempting connections  
on these ports in order  
opens sshd on 22

Biometric authentication? How to store w/ similar protections to passwords? Biometric scans might be slightly different each time.

ECC?

.	.	↘.	.	.
.	.	.	.	.
.	.	.	.	.
.	.	.	↘.	.

Attapt 1:  $s \equiv \text{scan}$ .

store salt,  $H(\text{salt} \parallel \text{ECC}(s))$

to check new scan,  $s'$ ,

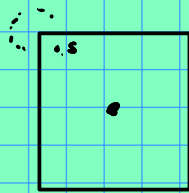
see if  $H(\text{salt} \parallel \text{ECC}(s'))$  equals the above.



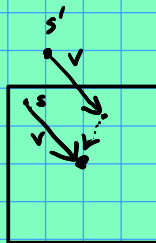
Problem: ideally, any "close" scan should work to authenticate.

This would be true if original looked to be a code word.

But what if original looks like this:



Solution: "fuzzy sketch": store the ECC vector  $v$  in addition to the hash + salt.



Modified scheme: store salt,  $H(\text{salt} \parallel \text{ECC}(s))$ ,  $v$

Now to check a new scan  $s'$ :

see if  $H(\text{salt} \parallel \text{ECC}(v + s'))$  matches.

$s'$  close to  $s \Rightarrow$  authentication succeeds.

Leaves small amount of info (v), but not unreasonable.

for other items:

run command as root: use sudo.

Note: can explicitly list commands in the sudoers file...

edit file as root? (e.g. /etc/ssh/sshd-config)

~~Sudo via file~~

sudoedit file



(The latter runs the editor w/ normal privilege to edit a temp file.

Gives temp file after editor closes)

run command with lesser privilege?

sudo runuser -u nobody [command]

Note on setuid programs (e.g. sudo!)

These run w/ the permissions of the owner, ← usually root!  
no matter what user started them!

Useful for things like passwd command:

normal user runs it... and /etc/shadow is modified!

Such programs must be written very carefully...

e.g. the access/open pattern:

```
if (access("file", —) )  
    exit(1);  
open("file", ...)
```

access checks if  
real user would  
have access to file.

Possible abuse:

touch /tmp/file

# run setuid program referencing /tmp/file

# between access + open:

In -sf /etc/shadow /tmp/file ⊗

"Time of check to Time of Use" bug.

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Big picture items... "Reflections on Trusting Trust" by K. Thompson

What if someone modified the C compiler to miscompile certain functions... e.g. the 'login' command.

if (looks like compiling 'login()')...

// add extra password that is  
// always accepted.

else if (looks like I'm compiling myself...)

// add the above  
// (as well as this)

Now if we revert the C compiler source code to its original form, but still compile w/ a bad binary, the login command will be backdoored.

Other attacks could be even more subtle:

make binary for encryption produce output that's  
— indistinguishable from that of the  
uncompromised also

— leaks the secret key used to encrypt !!

Idea: use IV to encrypt key?

use a PRF  $F: \{0,1\}^* \rightarrow \{0,1\} \times \{0,1\}^N$

Compromised algorithm could rejection-sample ciphertexts (using different randomness each time) until  $F(c) = (b, i)$  is such that

$$b = K[i]$$

$K$  = victim's encryption key.

Defense: Reproducible Builds

- Establish connection between source code & binary.
- gives instructions by which any one could produce an identical binary package.

(Debian is doing pretty well, btw...)