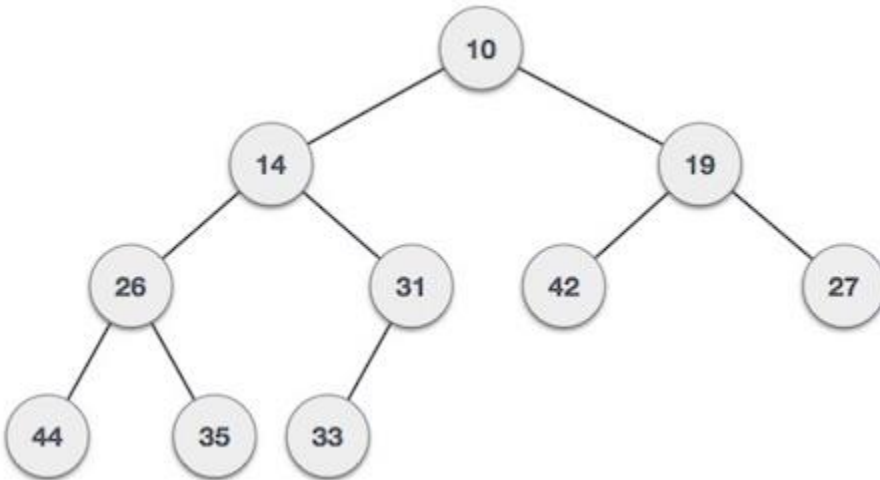


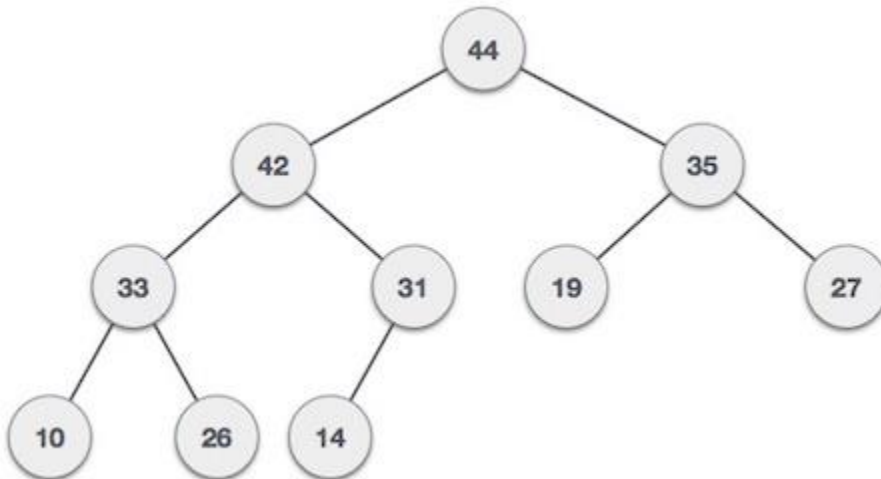
A heap is a specialized tree-based data structure that satisfied the heap property:

- If B is a child node of A, then  $\text{key}(A) \geq \text{key}(B)$ . This implies that an element with the greatest key is always in the root node, and so such a heap is sometimes called a max-heap. Of course, there's also a min-heap.

**Min-Heap** – where the value of root node is less than or equal to either of its children.

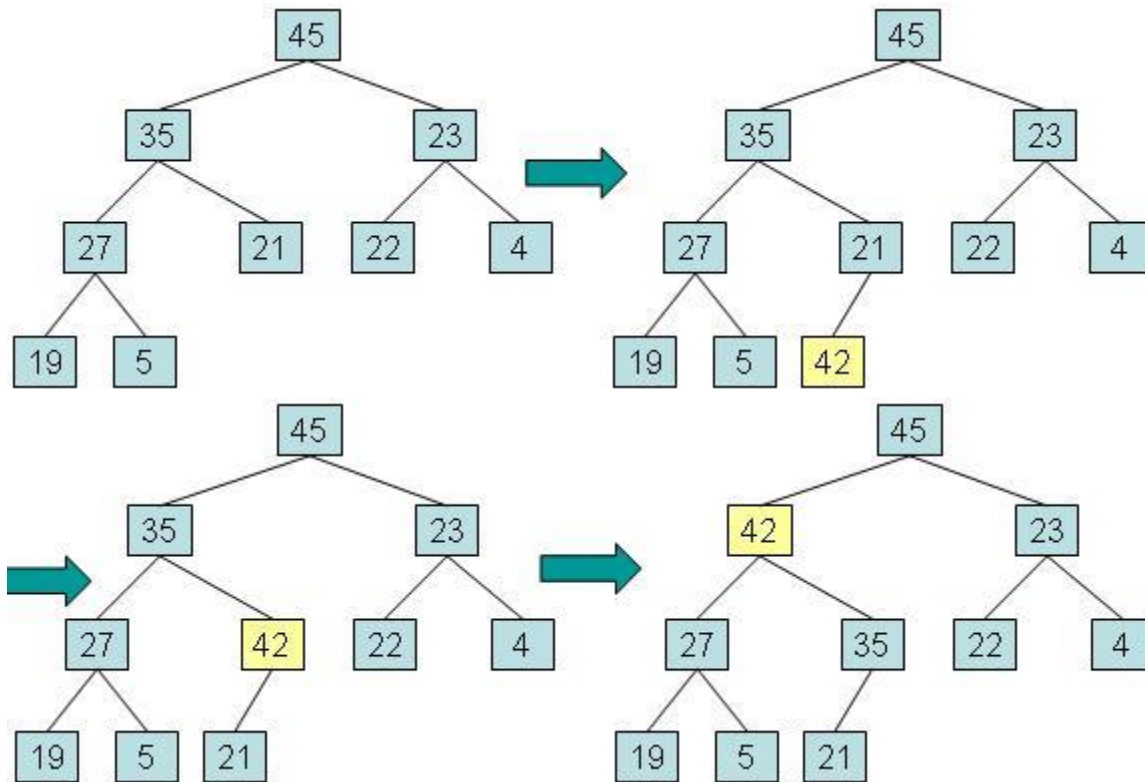


**Max-Heap** – where the value of root node is greater than or equal to either of its children.



## Adding an Element to a Heap

Example: We want to insert a node with value 42 to the heap on the left.

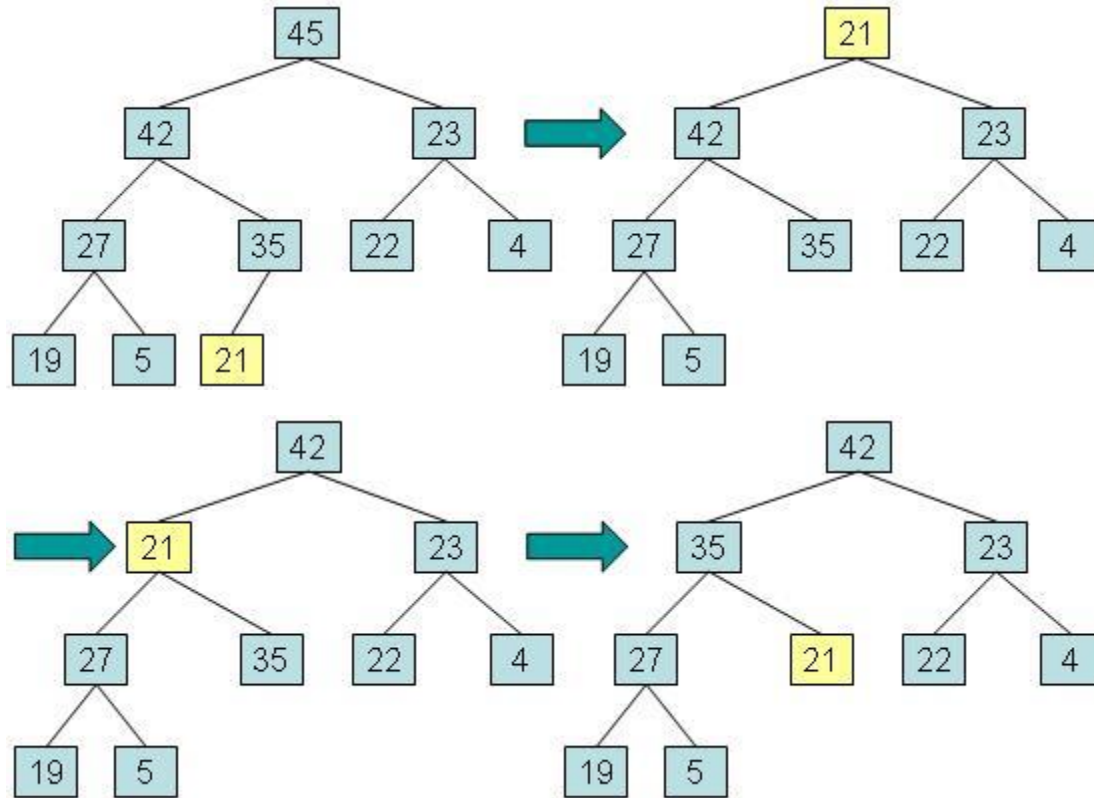


The above process is called **reheapification upward**.

### Pseudocode for Adding an Element:

**Step 1** - Create a new node at the end of heap.  
**Step 2** - Assign new value to the node.  
**Step 3** - Compare the value of this child node with its parent.  
**Step 4** - If value of parent is less than child, then swap them.  
**Step 5** - Repeat step 3 & 4 until Heap property holds.

## Removing the Root of a Heap



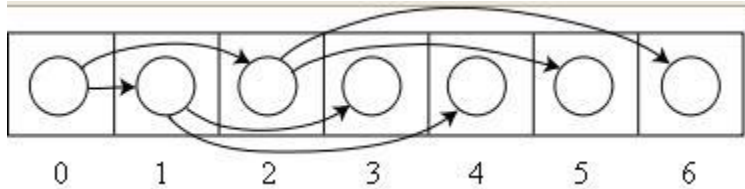
The above process is called **reheapification downward**.

### Pseudocode for Removing the Root:

- Step 1 - Remove root node.
- Step 2 - Move the last element of last level to root.
- Step 3 - Compare the value of this child node with its parent.
- Step 4 - If value of parent is less than child, then swap them.
- Step 5 - Repeat step 3 & 4 until Heap property holds.

## Heap Implementation:

A more common approach is to store the heap in an array. Since heap is always a complete binary tree, it can be stored compactly. No space is required for pointers; instead, the parent and children of each node can be found by simple arithmetic on array indices.



The rules (assume the root is stored in  $arr[0]$ ):

- For each index  $i$ , element  $arr[i]$  has children at  $arr[2i + 1]$  and  $arr[2i + 2]$ , and the parent at  $arr[\text{floor}((i - 1)/2)]$ .

## Time Complexity:

Function  $\text{reheapUp}()$  has time complexity  $O(\log n)$ , when you add  $n$  elements to the heap, each takes  $O(\log n)$ , so total heap sort complexity  $O(n \log n)$

Similarly,  $\text{reheapDown}()$  has time complexity  $O(\log n)$ , when you delete  $n$  elements to the heap, each takes  $O(\log n)$ , so total heap sort complexity  $O(n \log n)$

In general, **heap sort is  $O(n \log n)$**  for Best, Average and worst case.

## Heap/PriorityQueue (min/max):

- **Find Min/Find Max:**  $O(1)$
- **Insert:**  $O(\log n)$
- **Delete Min/Delete Max:**  $O(\log n)$
- **Lookup, Delete** (if at all provided):  $O(n)$ , we will have to scan all the elements as they are not ordered like BST