Tree sort

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A **tree sort** is a sort algorithm that builds a binary search tree from the elements to be sorted, and then traverses the tree (inorder) so that the elements come out in sorted order. Its typical use is sorting elements adaptively: after each insertion, the set of elements seen so far is available in sorted order.

Contents

- 1 Efficiency
- 2 Example
- 3 See also
- 4 External links

Efficiency

Adding one item to a binary search tree is on average an $O(\log n)$ process (in big O notation), so adding n items is an $O(n \log n)$ process, making tree sort a 'fast sort'. But adding an item to an unbalanced binary tree needs O(n) time in the worst-case, when the tree resembles a linked list (degenerate tree), causing a worst case of $O(n^2)$ for this sorting algorithm. This worst case occurs when the algorithm operates on an already sorted set, or one that is nearly sorted. Expected $O(n \log n)$ time can however be achieved in this case by shuffling the array.

The worst-case behaviour can be improved upon by using a self-balancing binary search tree. Using such a tree, the algorithm has an $O(n \log n)$ worst-case performance, thus being degree-optimal for a comparison sort. When using a splay tree as the

Tree sort Class Sorting algorithm Data structure Array **Worst-case performance** $O(n^2)$ (unbalanced) $O(n \log n)$ (balanced) $O(n \log n)$ **Best-case performance** $O(n \log n)$ Average performance Worst-case space $\Theta(n)$ complexity

binary search tree, the resulting algorithm (called splaysort) has the additional property that it is an adaptive sort, meaning that its running time is faster than $O(n \log n)$ for inputs that are nearly sorted.

Example

The following tree sort algorithm in pseudocode accepts an array of comparable items and outputs the items in ascending order:

```
STRUCTURE BinaryTree

BinaryTree:LeftSubTree

Object:Node

BinaryTree:RightSubTree

PROCEDURE Insert(BinaryTree:searchTree, Object:item)

IF searchTree.Node IS NULL THEN

SET searchTree.Node TO item
```

```
ELSE
       IF item IS LESS THAN searchTree.Node THEN
           Insert(searchTree.LeftSubTree, item)
            Insert(searchTree.RightSubTree, item)
PROCEDURE InOrder(BinaryTree:searchTree)
   IF searchTree.Node IS NULL THEN
       EXIT PROCEDURE
   ELSE
        InOrder(searchTree.LeftSubTree)
        EMIT searchTree.Node
        InOrder(searchTree.RightSubTree)
PROCEDURE TreeSort(Array:items)
   BinaryTree:searchTree
   FOR EACH individualItem IN items
        Insert(searchTree, individualItem)
   InOrder(searchTree)
```

In a simple functional programming form, the algorithm (in Haskell) would look something like this:

In the above implementation, both the insertion algorithm and the retrieval algorithm have $O(n^2)$ worst-case scenarios.

See also

 Heapsort: builds a binary heap out of its input instead of a binary search tree, and can be used to sort in-place (but not adaptively).

External links

- Binary Tree Java Applet and Explanation (http://www.qmatica.com/D ataStructures/Trees/BST.html)
- Tree Sort of a Linked List (http://www.martinbroadhurst.com/articles/sorting-a-linked-list-by-turning-it-into-a-binary-tree.html)
- Tree Sort in C++ (http://www.martinbroadhurst.com/cpp-sorting.html #tree-sort)



Categories: Sorting algorithms

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