CSE 100 Advanced Data Structures

Homework 3

Fall 2018

Due on: Thursday 11/08 (*40 points*)

Instructors: Leo Porter (Sec. A)

& Debashis Sahoo (Sec. B)

Name: ______ PID: ______ *Date*: 11 /____ / 2018

Instructions

- 1. Answer each problem in the boxes or circles provided. Any writing outside of the boxes *will NOT be graded*. Do not turn in responses recorded on separate sheets.
- 2. Handwritten or typed responses are accepted. In either case, make sure all answers are in the appropriate boxes.
- 3. All responses *must* be neat and legible. Illegible answers will result in zero points.
- 4. Make sure to scan in portrait mode and to select the corresponding pages on Gradescope for each question.
- 5. You may use code from any of the class resources, including Stepik. You may not use code from other sources.
- 1. (8 points) *Linear Probing*: The following hash table was created using linear probing:

Hash function: h(x) = (5 * x)%8

О	1	2	3	4	5	6	7
11	3	2	7	12	23		19

a.	(3 points - Correctness) List all of the elements that could have been inserted into the hash table <i>first</i>
b.	(3 points - Correctness) List all of the elements that could have been inserted into the hash table <i>last</i> .
c.	(2 points - Completeness) Provide a possible insertion order of the elements.

2. (6 points - **Correctness**) *Bloom Filter*: Given the following bloom filter:

• hash function 1: $h_1(x) = (x^3)\%13$ • hash function 2: $h_2(x) = (x * 5)\%13$

O	1	-	2	3	4	5	6	7	8	9	10	11	12
	1		1			1			1		1		1

a.	What are 4 possible distinct values between 1 and 13 (inclusive) which could be inserted to create the bloc filter above?	om
b.	List all the false positives for integers between 1 and 13 (inclusive) in your bloom filter given inserted elem you chose above.	nents

3.	(7 points - Correctness) Cuckoo Hashing: The following hash tables and their respective hash functions are used
	for cuckoo hashing.

• Hash function 1: $h_1(x) = (x\%7)\%5$

0	1	2	3	4
7			17	46

• Hash function 2: $h_2(x) = (x+3)\%5$

O	1	2	3	4
47	33			

Draw the 2 tables	listed above after inse	rting 10. If you thir	K It leads to all III	mine cycle write	IIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIII
Provide an insert	value that will cause ar	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause a	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause ar	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause ar	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause a	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause a	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause a	n infinite cycle. Sho	w the cycle that is	s created.	
Provide an insert	value that will cause a	n infinite cycle. Sho	w the cycle that is	s created.	

a.	What	is the average-case	time complexit	y for insert in a Ha	sh Table.	
O <i>O</i> (1)		$O(\log n)$	0	O(n)	$ O(n \log n) $	$O(n^2)$
b	What	is the average-case	time complexit	y for find in a Hasl	n Table.	
O <i>O</i> (1)		$O(\log n)$	0	O(n)	$ O(n \log n) $	$O(n^2)$
c.		is the worst-case tied List is used in the			ash Table using Separate	Chaining (assume a
O <i>O</i> (1)		$O(\log n)$	O	O(n)	$ O(n \log n) $	$O(n^2)$
d		is the worst-case ti s used in the bucket			Րable using Separate Chai	ning (assume a Linked
O <i>O</i> (1)		$O(\log n)$	O	O(n)	$ O(n \log n) $	$O(n^2)$
e.		is the worst-case ti			ash Table using a Hash Tation)	able using Separate
O <i>O</i> (1)		$O(\log n)$	0	O(n)	$ O(n \log n) $	$ O(n^2) $
f.		is the worst-case ti in the bucket imple		for find in a Hash T	Гable using Separate Chai	ning (assume a BST is
O <i>O</i> (1)		$O(\log n)$	0	O(n)	$ O(n \log n) $	$O(n^2)$
g.	What	is the load factor o	f a Hash Table.			
	0	The minimum n	umber of eleme	nts you need to ins	ert before the Hash Table	performs optimally
	Ō			rrent size of hash t		- · ·
	Ö				ou can insert in your hash	table
	0	The ratio #inser	ted elements/#c	collisions		
	\cap	The ratio #occur	oied hash values	/current size of ha	sh table	

4. (6 points - Correctness) Know Your Facts:

h.	At whice	ch load factor do you general	ly want to increase the si	ze of your hash table?	
	0	0.9			
	Ο	0.7			
	Ο	151			
	0	It depends on your collision	n resolving strategy		
	Ο	It depends on the current s	ize of your hash table		
i.		s the worst-case time comple used in the bucket implemen		sh Table using Separate	Chaining (assume an AVL
O (1)		$ O(\log n) $	O(n)	$ O(n \log n) $	$O(n^2)$
j.		s the worst-case time comple the bucket implementation		able using Separate Chair	ning (assume an AVL tree
O (1)		$ O(\log n) $	O(n)	$ O(n \log n) $	$O(n^2)$
k.	What is	s the worst-case time comple	exity for insertion using C	Cuckoo Hashing?	
O (1)		$ O(\log n) $	O(n)	$ O(n \log n) $	$O(n^2)$
l.	What is	s the worst-case time comple	exity for find using Cucko	o Hashing?	
O (1)		$ O(\log n) $	O(n)	$ O(n \log n) $	$O(n^2)$

5.	(7 points - Correctness)	Load Factor:	You implemented	a hash table u	sing linear p	robing and the	e following l	hash
	function:							

$$h(x) = (x+4)\%5$$

You defined 0.7 to be the maximum load factor for your hash table and when exceeded you want to approximately double the size of your hash table.

This is the current state of your hash table:

O	1	2	3	4
6		13	24	

a. What is the current load factor of your hash table?

b. Show your hash function and hash table after inserting 5.

П	
ı	

c. What is the current load factor of your new hash table?



6. (6 points - Correctness) Collision Handling: Consider two hash functions:

- $h_1(x) = x\%7$
- $h_2(x) = 5 (x\%3)$

Insert the following keys into a hash table of size 7:

13, 501, 128, 37, 2

a. Insert the keys using linear probing (using h_1)

0 1 2 3 4 5 6			
1 2 3 4 5 5 S			
1 2 3 4 5 5 S		\neg	
1 2 3 4 5 5 S	0		
2 3 4 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5			
2 3 4 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5			
3 4 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5	1		
3 4 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5			
4 5	2		
4 5			
4 5			
5	3		
5			
5	4		
	5		
6	J		
6			
	6		

b. Insert the keys using separate chaining (using h_1)

	<u> </u>
o	
1	
2	
3	
4	
5	
6	
6	

c. Ir	c. Insert the keys using double hashing (use h_1 as primary and h_2 as secondary hash function)			
	T	1		
o				
1				
2				
3				
4				
5				
6				