

Problem Set 4

Issued: Monday 16th November, 2020

Due: Monday 30th November, 2020

Notations: We use $\mathbf{x}, \mathbf{y}, \mathbf{w}$ and $\underline{\mathbf{x}}, \underline{\mathbf{y}}, \underline{\mathbf{w}}$ to denote random variables and random vectors.

4.1. Please review Chapter 12 in Cover's book, then you can get some ideas on how to find the K-L divergence in Sanov's Theorem. Let \mathbf{x}_i be i.i.d. $\sim \mathcal{N}(0, \sigma^2)$:

- (a) Find the behavior of $-\frac{1}{n} \log \mathbb{P} \left(\frac{1}{n} \sum_{i=1}^n \mathbf{x}_i^2 \geq \alpha^2 \right)$. This can be done from the first principles (since the normal distribution is nice) or by using Sanov's theorem.
- (b) What does the data look like if $\frac{1}{n} \sum_{i=1}^n \mathbf{x}_i^2 \geq \alpha^2$. That is, what is the distribution that minimizes the K-L divergence in the Sanov's theorem.

Solution: A simple conclusion in Chapter 12 tells that the maximum entropy distribution is of the form $f(x) \sim C e^{-\beta x^2}$ (Gaussian) and the constraint is $\mathbb{E}_f[\mathbf{x}^2] = \alpha^2$. Therefore it is $\mathcal{N}(0, \alpha^2)$.

$$\begin{aligned} \lim_{n \rightarrow \infty} -\frac{1}{n} \log \mathbb{P} \left(\frac{1}{n} \sum_{i=1}^n \mathbf{x}_i^2 \geq \alpha^2 \right) &= \arg \min_{P \in \{P: \mathbb{E}_P[\mathbf{x}^2] \geq \alpha^2\}} D(P \| Q) \\ &= D(\mathcal{N}(0, \alpha^2) \| \mathcal{N}(0, \sigma^2)) \\ &= \frac{1}{2} \frac{\alpha^2}{\sigma^2} + \frac{1}{2} \log \frac{\sigma^2}{\alpha^2} \end{aligned}$$

4.2. We hope to derive an asymptotic value of $\binom{n}{k}$.

- (a) Firstly, let's prove the lemma about Stirling's approximation of factorials, which we have used before.

$$\left(\frac{n}{e} \right)^n \leq n! \leq n \left(\frac{n}{e} \right)^n$$

Please justify the following steps:

$$\ln(n!) = \sum_{i=2}^{n-1} \ln i + \ln n \leq \dots$$

$$\ln(n!) = \sum_{i=1}^n \ln i \geq \dots$$

- (b) If $0 < p < 1$, and $k = \lfloor np \rfloor$, i.e., k is the largest integer less than or equal to np , then please find

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \binom{n}{k}$$

Could you explain it without Stirling's Approximation?

Now let p_i 's be a probability distribution on m symbols. Guess what is

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \binom{n}{\lfloor np_1 \rfloor \ \lfloor np_2 \rfloor \ \cdots \ \lfloor np_{m-1} \rfloor \ (n - \sum_{i=1}^{m-1} \lfloor np_i \rfloor)}$$

Solution:

(a)

$$\ln(n!) = \sum_{i=2}^{n-1} \ln i + \ln n \leq \int_2^{n-1} \ln x dx + \ln n$$

$$\ln(n!) = \sum_{i=1}^n \ln i \geq \int_0^n \ln x dx$$

Then you can get the approximation.

(b) By applying the approximation, we can easily get $\lim_{n \rightarrow \infty} \frac{1}{n} \log \binom{n}{k} = H(p) \triangleq -p \log p - (1-p) \log(1-p)$. Let's think about the conclusion, if we have a sequence with n i.i.d. rv from $\text{Bern}(p)$, when n goes to ∞ with high probability, the sequence will have k '0's. For coding the sequence, we need at least $\log \binom{n}{k}$ bits, consistent with the entropy of the sequence $nH(p)$.

That's what the result says.

The result is useful in the Sanov's theorem for Bernoulli rv's.

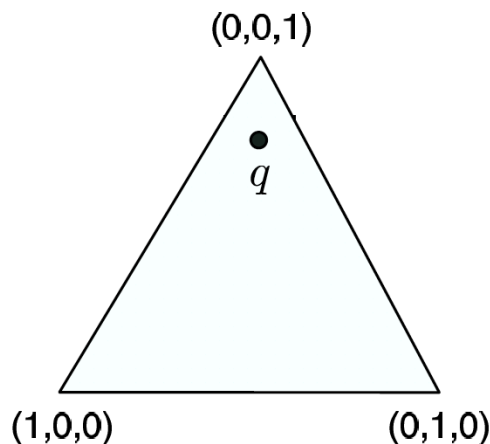
$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \binom{n}{\lfloor np_1 \rfloor \ \lfloor np_2 \rfloor \ \cdots \ \lfloor np_{m-1} \rfloor \ (n - \sum_{i=1}^{m-1} \lfloor np_i \rfloor)} = H(p_1, \dots, p_m)$$

4.3. Consider the set of distributions on $\Omega = \{0, 1, 2\}$ and note that they lie on the 2-simplex

$$\{p = (p_0, p_1, p_2) : p_0 + p_1 + p_2 = 1, p_0 \geq 0, p_1 \geq 0, p_2 \geq 0\}$$

represented by the triangular figure. Let y be a random variable such that $p_y(i) = p_i, i \in \{0, 1, 2\}$. Let $q = (1/6, 1/6, 2/3)$ be a particular probability mass function.

- Draw on the simplex the linear family corresponding to the expectation $\mathbb{E}[y] = 0$, i.e. draw $\mathcal{L}_0 = \{p : \mathbb{E}_p[y] = 0\}$.
- Draw $\mathcal{L}_{1/2} = \{p : \mathbb{E}_p[y] = 1/2\}$
- Specify the exponential family \mathcal{E} that passes through q and is orthogonal to $\mathcal{L}_{1/2}$, and draw the entire family on the 2-simplex.



Hint: Remember we introduced two versions of the exponential family, which are Lagrange-Multiplier induced one and parameterized one. You might be confused when you are facing cardinality-3 distributions, especially the Lagrange-Multiplier induced one. It is good if you can think about the equivalency of the two versions. Let's do the problem firstly under the parameterized version. That is $\mathcal{E} = \{\tilde{q} : \tilde{q} = qe^{sf(y)-\alpha(s)}\}$. Following the definition above, $f(y) = y$.

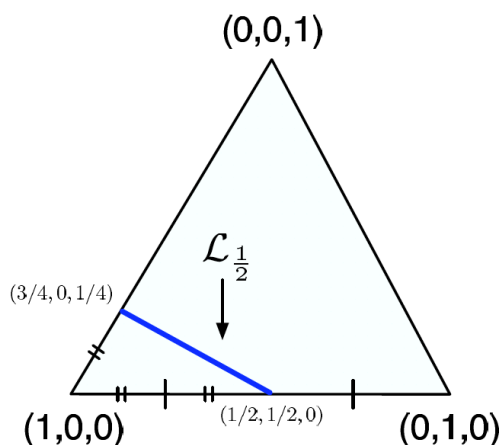
- (d) Calculate the I-projection p^* of q onto $\mathcal{L}_{1/2}$ and mark it on the simplex.
- (e) Draw $\mathcal{P} = \{p : \mathbb{E}_p[y] \leq 1/2\}$.
- (f) Calculate the I-projection p^* of q onto \mathcal{P} and mark it.

Hint: $D(\cdot \| q)$ is convex in its first argument.

Solution:

(a) Boundary point $(1, 0, 0)$. Figure omitted.

(b) We seek the set of points p such that $\mathbb{E}_p[y] = 0p_0 + 1p_1 + 2p_2 = 1/2$. Letting $\lambda = 2p_1$, it follows that $\{(3 - \lambda)/4, \lambda/2, (1 - \lambda)/4\}$.

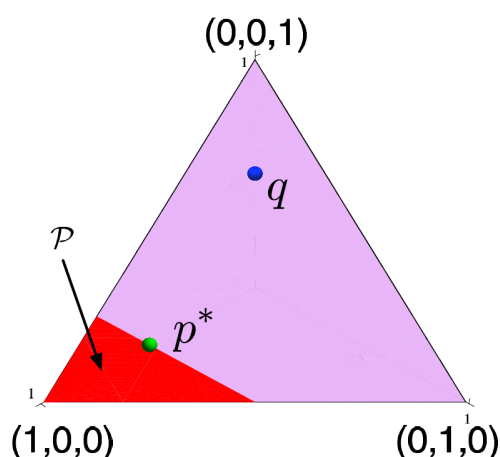
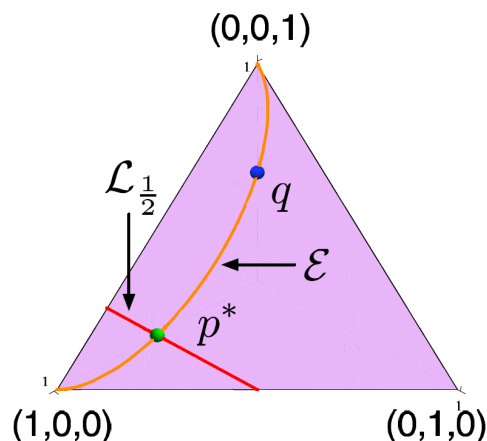


(c) The family is

$$\left(\frac{1}{1 + e^x + 4e^{2x}}, \frac{e^x}{1 + e^x + 4e^{2x}}, \frac{4e^{2x}}{1 + e^x + 4e^{2x}} \right).$$

Therefore, $p(0)p(2) = 4p^2(1) = 4(1p(0)p(2))^2$; the family is an ellipse.

(d) $p^* = (2/3, 1/6, 1/6)$.



(e)

(f) The set \mathcal{P} can be thought of as the collection of lines parallel to $\mathcal{L}_{1/2}$ of the form $L_\mu = p : \mathbb{E}_p[y] = \mu$ for $\mu \in [0, 1/2]$. Let

$$p_\mu^* \triangleq \arg \min_{p \in \mathcal{L}_\mu} D(p \| q)$$

$$p_\lambda = \lambda p_\mu^* + (1 - \lambda)q$$

where $\lambda \in [0, 1]$.

Then $\mathbb{E}_{p=0}[y] = \mathbb{E}_q[y] = 3/2$, and $\mathbb{E}_{p=1}[y] = E_{p_\mu^*}[y] = \mu \leq 1/2$. By continuity we see that there exists a λ^* such that $\mathbb{E}_{p_{\lambda^*}}[y] = 1/2$ and hence $p_{\lambda^*} \in \mathcal{L}_{1/2}$. Therefore,

$$D(p^* \| q) \leq D(p_{\lambda^*} \| q) \leq \lambda^1 \star D(p_\mu^* \| q) + (1 - \lambda^*)D(q \| q).$$

Hence for any given μ ,

$$D(p_\mu^* \| q) \geq \frac{1}{\lambda^*} D(p^* \| q) \geq D(p^* \| q).$$

Hence, $p^* = (2/3, 1/6, 1/6)$.

- 4.4. Let $q(y) > 0$ ($y = 0, 1, \dots$) be a probability mass function for a random variable y and let \mathcal{P} be the set of all PMFs defined over $\{0, \dots, M-1\}$ for a known constant M :

$$\mathcal{P} \triangleq \{p : p(y) = 0, \forall y \geq M\}.$$

We can represent each element p of \mathcal{P} as a M -dimensional vector $[p_0, \dots, p_{M-1}]^T$ that lies on a $(M-1)$ -dimensional simplex, i.e., $\sum_{m=0}^{M-1} p_m = 1$.

- Show that, for all $p \in \mathcal{P}$, $D(q \| p) = \infty$
- Show that, for all $p \in \mathcal{P}$, $D(p \| q) < \infty$
- Find the I-projection of q onto \mathcal{P} , $p^* = \arg \min_{p \in \mathcal{P}} D(p \| q)$, and the corresponding divergence $D(p^* \| q)$ in terms of $Q(y) \triangleq \mathbb{P}(y \leq y)$, the CDF of the random variable y .

Let \mathcal{P}_ϵ be the space of all PMFs with weight of ϵ on values M and above:

$$\mathcal{P}_\epsilon \triangleq \left\{ p : \sum_{y=M}^{\infty} p(y) = \epsilon \right\}$$

We can think of \mathcal{P}_ϵ as an extension of \mathcal{P} to the distributions defined for all integers that only allows limited weight to be allocated to the values outside $\{0, \dots, M-1\}$.

- Find the I-projection of q onto \mathcal{P}_ϵ , $p_\epsilon^* = \arg \min_{p \in \mathcal{P}_\epsilon} D(p \| q)$ and the corresponding divergence $D(p_\epsilon^* \| q)$ in terms of $Q(y)$. Show that $\lim_{\epsilon \rightarrow 0} D(p_\epsilon^* \| q) = D(p^* \| q)$.
- Show that \mathcal{P}_ϵ can be represented as a linear family of PMFs.
- Show that p_ϵ^* belongs to an exponential family through q and find the value of the parameter that corresponds to p_ϵ^* .

Solution:

(a) ■

(b) ■

(c) We form the Lagrangian,

$$L = D(p \| q) - \lambda \left(\sum_{y=0}^{M-1} p(y) - 1 \right).$$

Then,

$$p^* = q \exp(\lambda - 1).$$

According to the constraint, $\exp(\lambda - 1) = \frac{1}{Q(M-1)}$ and $D(p^*||q) = \log \frac{1}{Q(M-1)}$.

(d) Follow the similar procedures,

$$p_\epsilon^* = \begin{cases} q \frac{1-\epsilon}{Q(M-1)}, & y \leq M-1 \\ q \frac{\epsilon}{1-Q(M-1)}, & y \geq M \end{cases}.$$

Then,

$$D(p_\epsilon^*||q) = (1-\epsilon) \log \frac{1-\epsilon}{Q(M-1)} + \epsilon \log \frac{\epsilon}{1-Q(M-1)},$$

and the limitation is easy to verify.

(e) Let $f(y) = \mathbb{1}_{y \geq M}$, then $\mathcal{P}_\epsilon = \{p : \mathbb{E}_p[f(y)] = \epsilon\}$.

(f) $\alpha(s) = \log \left(\sum_{y=0}^{M-1} q(y) + \sum_{y=M}^{\infty} q(y)e^s \right)$

$$p_\epsilon^* = q \exp \left(\log \left(\frac{\epsilon}{1-\epsilon} \frac{Q(M-1)}{1-Q(M-1)} \right) \mathbb{1}_{y \geq M} - \log \frac{Q(M-1)}{1-\epsilon} \right)$$

4.5. *Joint Gaussian Distribution.* Suppose $\underline{\mathbf{x}} = (\mathbf{x}_1, \mathbf{x}_2)^T$ is a Gaussian random vector with $\mathbb{E}[\mathbf{x}_1] = \mathbb{E}[\mathbf{x}_2] = 0$, $\text{var}(\mathbf{x}_1) = \text{var}(\mathbf{x}_2) = \sigma^2$, and $\rho_{\mathbf{x}} \triangleq \rho(\mathbf{x}_1, \mathbf{x}_2)$ denoting the correlation coefficient between \mathbf{x}_1 and \mathbf{x}_2 . Let $\underline{\mathbf{y}} = (\mathbf{y}_1, \mathbf{y}_2)^T \triangleq \mathbf{A}\underline{\mathbf{x}}$, where

$$\mathbf{A} = \begin{bmatrix} 1 & -\rho_{\mathbf{x}} \\ 0 & 1 \end{bmatrix}.$$

Then, $\underline{\mathbf{y}}$ is also a Gaussian random vector, since it is a linear transformation of $\underline{\mathbf{x}}$.

(a) Calculate $\mathbf{K}_{\mathbf{x}} \triangleq \text{cov}(\underline{\mathbf{x}})$ and $\mathbf{K}_{\mathbf{y}} \triangleq \text{cov}(\underline{\mathbf{y}})$.

(b) Prove that $\rho(\mathbf{y}_1, g(\mathbf{y}_2)) = 0$, for all functions¹ $g(\cdot)$. *Hint:* First prove that $\mathbf{y}_1 \perp \mathbf{y}_2$.

(c) Prove that $\mathbb{E}[(\mathbf{x}_1 - \rho_{\mathbf{x}}\mathbf{x}_2)^2] \leq \mathbb{E}[(\mathbf{x}_1 - g(\mathbf{x}_2))^2]$, for all functions $g: \mathbb{R} \rightarrow \mathbb{R}$. *Hint:* Rewrite the inequality using \mathbf{y}_1 and \mathbf{y}_2 .

Solution:

(a)

$$\mathbf{K}_{\mathbf{x}} = \text{cov}(\underline{\mathbf{x}}) = \mathbb{E}[\underline{\mathbf{x}}\underline{\mathbf{x}}^T] - \mathbb{E}[\underline{\mathbf{x}}]\mathbb{E}[\underline{\mathbf{x}}]^T = \mathbb{E}[\underline{\mathbf{x}}\underline{\mathbf{x}}^T] = \sigma^2 \begin{bmatrix} 1 & \rho_{\mathbf{x}} \\ \rho_{\mathbf{x}} & 1 \end{bmatrix}.$$

¹Strictly speaking, $g(\cdot)$ is required to be measurable.

Then,

$$\mathbf{K}_y = \mathbf{A}\mathbf{K}_x\mathbf{A}^T = \sigma^2 \begin{bmatrix} 1 - \rho_x^2 & 0 \\ 0 & 1 \end{bmatrix}.$$

- (b) Since \mathbf{K}_y is a diagonal matrix, we know that $\rho(y_1, y_2) = 0$. Then, as \underline{y} is a Gaussian vector, we obtain $y_1 \perp y_2$, and thus the result follows.
- (c) This is equivalent to $\mathbb{E}[y_1^2] \leq \mathbb{E}[(y_1 + \rho_x y_2 - g(y_2))^2]$, which can be obtained from

$$\mathbb{E}[(y_1 + \rho_x y_2 - g(y_2))^2] = \mathbb{E}[y_1^2] + \mathbb{E}[(\rho_x y_2 - g(y_2))^2].$$

4.6. *Mathematical expectation and variance in estimation.* Suppose we want to estimate the value of y using an estimator \hat{y} , and using its MSE (Mean Square Error) to evaluate the goodness of estimate, defined as

$$\text{MSE}(\hat{y}) \triangleq \mathbb{E}[(y - \hat{y})^2].$$

The estimator \hat{y} could be chosen from a set \mathcal{A} , and our goal is to find the best estimator in \mathcal{A} which achieves the least MSE. Then the best estimator is called the MMSE (Minimum Mean Square Error) estimator.

- (a) Assume we want to use a real number to estimate y , i.e., $\mathcal{A} = \mathbb{R}$.
- Prove that $\mathbb{E}[y]$ is the MMSE estimator:

$$\mathbb{E}[y] = \arg \min_{\alpha \in \mathbb{R}} \mathbb{E}[(y - \alpha)^2].$$

- Evaluate this estimator's MSE.

- (b) Now you are allowed to use a function of x to estimate y , i.e., $\mathcal{A} = \{f(\cdot) : \mathcal{X} \mapsto \mathbb{R}\}$. Prove that:
- $\mathbb{E}[y|x]$ is the MMSE estimator:

$$\mathbb{E}[y|x] = \arg \min_{f: \mathcal{X} \mapsto \mathbb{R}} \mathbb{E}[(y - f(x))^2],$$

- The MSE of estimator $\mathbb{E}[y|x]$ is

$$\text{MSE}(\mathbb{E}[y|x]) = \mathbb{E}[\text{var}(y|x)].$$

- (c) Compare these two estimators. First, prove that

$$x \perp y \implies \text{MSE}(\mathbb{E}[y]) = \text{MSE}(\mathbb{E}[y|x]) \implies \forall f, \rho(f(x), y) = 0,$$

where $\rho(\cdot, \cdot)$ is the Pearson correlation coefficient. In general, which one of these two estimators would have less MSE than the other?

Solution: (a)(b) can be solved by direct computations. For (c), note that we have (cf. PS2.1 (a) iv)

$$\begin{aligned} \text{MSE}(\mathbb{E}[y]) &= \text{var}(y) = \mathbb{E}[\text{var}(y|x)] + \text{var}(\mathbb{E}[y|x]) = \text{MSE}(\mathbb{E}[y|x]) + \text{var}(\mathbb{E}[y|x]) \\ &\geq \text{MSE}(\mathbb{E}[y|x]). \end{aligned}$$

Therefore,

$$\text{MSE}(\mathbb{E}[y]) = \text{MSE}(\mathbb{E}[y|x]) \implies \text{var}(\mathbb{E}[y|x]) = 0 \implies \mathbb{E}[y|x] = \mathbb{E}[\mathbb{E}[y|x]] = \mathbb{E}[y].$$

As a result,

$$\mathbb{E}[f(x)y] = \mathbb{E}[f(x) \mathbb{E}[y|x]] = \mathbb{E}[f(x) \mathbb{E}[y]] = \mathbb{E}[f(x)] \mathbb{E}[y],$$

which implies $\rho(f(x), y) = 0$.

- 4.7. Consider the estimation of one-hot encoded vectors, where the settings are similar to those of Problem 3.3. In particular, suppose \mathbf{y} takes values from $\mathcal{Y} = \{1, 2, \dots, k\}$, then its one hot encoding is a k -dimensional vector defined as $\underline{\mathbf{y}} \triangleq (\mathbb{1}_{y=1}, \mathbb{1}_{y=2}, \dots, \mathbb{1}_{y=k})^T$, i.e., $\underline{\mathbf{y}}$ is the i -th vector of the standard basis if $y = i$.

Now, we would use $\hat{\underline{\mathbf{y}}}$ to estimate $\underline{\mathbf{y}}$, and use its MSE to evaluate the goodness of estimate. The MSE is defined similarly as the scalar case, except that the scalar quadratic operator is replaced by the ℓ_2 norm squared:

$$\text{MSE}(\hat{\underline{\mathbf{y}}}) \triangleq \mathbb{E}[\|\underline{\mathbf{y}} - \hat{\underline{\mathbf{y}}}\|_2^2].$$

Again, the estimator $\hat{\underline{\mathbf{y}}}$ could be chosen from a set \mathcal{A} .

- (a) Suppose we want to use a vector to estimate $\underline{\mathbf{y}}$, i.e., $\mathcal{A} = \mathbb{R}^k$. Prove that $\underline{P}_{\mathbf{y}}(\cdot)$ is the MMSE estimator:

$$\underline{P}_{\mathbf{y}}(\cdot) = \arg \min_{\underline{\alpha} \in \mathbb{R}^k} \mathbb{E}[\|\underline{\mathbf{y}} - \underline{\alpha}\|_2^2],$$

where $\underline{P}_{\mathbf{y}}(\cdot) \triangleq [P_{\mathbf{y}}(1), P_{\mathbf{y}}(2), \dots, P_{\mathbf{y}}(k)]^T$.

- (b) Now you are allowed to use a multivariate function of \mathbf{x} to estimate $\underline{\mathbf{y}}$, i.e., $\mathcal{A} = \{\underline{f} : \mathcal{X} \mapsto \mathbb{R}^k\}$. Prove that the MMSE estimator is $\underline{P}_{\mathbf{y}|\mathbf{x}}(\cdot|\mathbf{x})$:

$$\underline{P}_{\mathbf{y}|\mathbf{x}}(\cdot|\mathbf{x}) = \arg \min_{\underline{f}: \mathcal{X} \mapsto \mathbb{R}^k} \mathbb{E}[\|\underline{\mathbf{y}} - \underline{f}(\mathbf{x})\|_2^2],$$

where $\underline{P}_{\mathbf{y}|\mathbf{x}}(\cdot|\mathbf{x}) \triangleq [P_{\mathbf{y}|\mathbf{x}}(1|\mathbf{x}), P_{\mathbf{y}|\mathbf{x}}(2|\mathbf{x}), \dots, P_{\mathbf{y}|\mathbf{x}}(k|\mathbf{x})]^T$.

Solution: The MMSE estimators are the expectation and conditional expectation:

$$\underline{P}_{\mathbf{y}}(\cdot) = \mathbb{E}[\underline{\mathbf{y}}]$$

and

$$\underline{P}_{y|x}(\cdot|x) = \mathbb{E} [\underline{y}|x] .$$

- 4.8. The data $\mathbf{x}[n] = ar^n + \mathbf{w}[n]$ for $n = 0, \dots, N-1$ are observed. The random variables $\mathbf{w}[0], \dots, \mathbf{w}[N-1]$ are i.i.d. Gaussian random variables with zero mean and variance σ^2 . r is a non-zero constant. Find the Cramér-Rao bound for a . Does an efficient estimator exist? If so, what is it and what is its variance?

Solution: The Cramér-Rao bound for a is

$$\lambda_e(a) \geq \begin{cases} \sigma^2 \frac{1-r^2}{1-r^{2N}}, & |r| \neq 1 \\ \frac{\sigma^2}{N}, & |r| = 1 \end{cases} .$$

The efficient estimator exists, given by

$$a_{\text{eff}}(\mathbf{x}) = \frac{\underline{r}^T \underline{x}}{\underline{r}^T \underline{r}},$$

with $\underline{x} = [x[0], \dots, x[N-1]]^T$ and $\underline{r} = [r^0, \dots, r^{N-1}]^T$. Since it is efficient, the variance is given by the Cramér-Rao bound

$$\lambda_{\text{eff}}(a) = \begin{cases} \sigma^2 \frac{1-r^2}{1-r^{2N}}, & |r| \neq 1 \\ \frac{\sigma^2}{N}, & |r| = 1 \end{cases} .$$