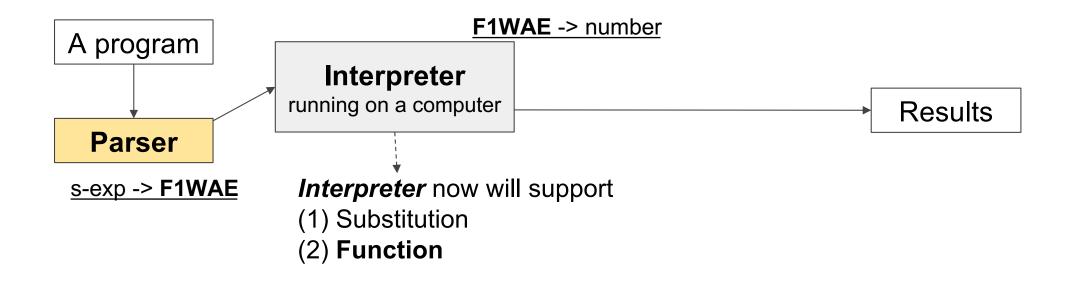
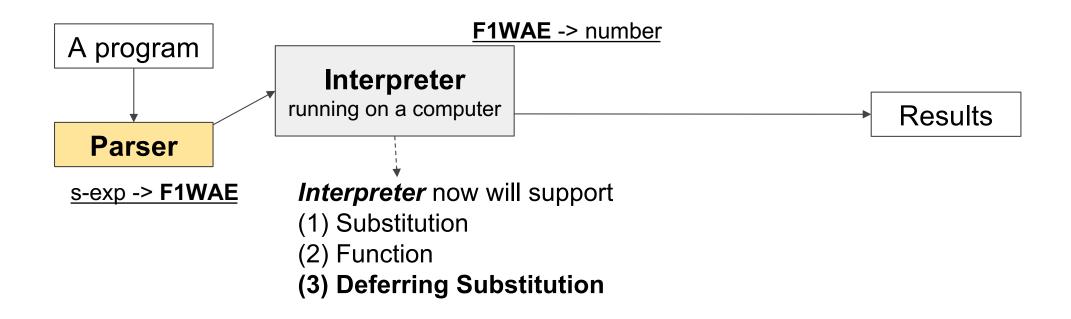
## ITP30011 L9 Deferring Substitution

Lecture09 JC

#### Big Picture (modeling languages: substitution)



#### Big Picture (modeling languages: substitution)



```
{with {x 3}
    {with {y 4}
        {with {z 5}
            {+ x {+ y z}}}}
= \{ with \{ y 4 \} \}
    {with {z 5}
        {+ 3 {+ y z}}}}
= {with z 5}
    \{+3\{+4z\}\}\}
```

```
{with {x 3}
    {with {y 4}
        {with {z 5}
             {+ x {+ y z}}}}
= \{ with \{ y 4 \} \}
    {with {z 5}
        {+ 3 {+ y z}}}
= \{ with z 5 \}
    \{+3\{+4z\}\}\}
= \{+ 3 \{+ 4 5\}\}
```

```
{with {x 3}
    {with {y 4}
        {with {z 5}
            {+ x {+ y z}}}}
= \{ with \{ y 4 \} \}
    {with {z 5}
        \{+3 \{+yz\}\}\}
= {with z 5}
    \{+3\{+4z\}\}\}
= \{+ 3 \{+ 4 5\}\}
```

Complexity of the interpreter for this program:

 $O(n^2)$ 

```
{with {x 3}
    {with {y 4}
        {with {z 5}
            {+ x {+ y z}}}}
= \{ with \{ y 4 \} \}
    {with {z 5}
        {+ 3 {+ y z}}}
= \{ with z 5 \}
    \{+3\{+4z\}\}\}
= \{+ 3 \{+ 4 5\}\}
```

WE CAN do BETTER??

### Let's improve Substitution!

 $\Rightarrow$  Our new language AE  $\rightarrow$  WAE  $\rightarrow$  F1WAE with better substitution

#### Substitution

```
; [contract] subst: WAE symbol number -> WAE
(define (subst wae idtf val)
    (type-case WAE wae
         num
                    (n)
                            wae]
         [add (lr) (add (subst lidtf val) (subst ridtf val))]
         [sub (lr) (sub (subst lidtf val) (subst ridtf val))]
         [with (i v e) (with i (subst v idtf val)
                                               (if (symbol=? i idtf)
                                                                 (subst e idtf val)))]
         [id
                    (s)
                                 (if (symbol=? s idtf) (num val) wae)]))
; {with \{x \ 10\} \ \{... \{with \ \{y \ 17\} \ x\}\} \Rightarrow 10 \text{ for } x \text{ in } \{with \ \{y \ 17\} \ x\}
                                                                                  \Rightarrow {with {y 17} 10}
(test (subst (with 'y (num 17) (id 'x)) 'x 10) (with 'y (num 17) (num 10)))
; \{\text{with } \{x \ 10\} \ \{... \{\text{with } \{x \ x\} \ x\}\}\} \ \Rightarrow 10 \text{ for } x \text{ in } \{\text{with } \{x \ x\} \ x\} \ \Rightarrow \{\text{with } \{x \ 10\} \ x\}
(test (subst (with 'x (id 'x) (id 'x)) 'x 10) (with 'x (num 10) (id 'x)))
```

#### Cost of Substitution

```
(interp (parse '{with {x 1}}
{with {y 2}}
{+ 100 {+ 99 {+ 98 ... {+ y x}...}))
```

#### Cost of Substitution

#### Cost of Substitution

\* n = the number of abstract syntax tree nodes of a program

With identifiers, evaluation will take  $O(n^2)$  time!

# How can we do better? Do you have any idea?

```
(interp (parse '{with {x 1} [] {with {y 2} {+ 100 {+ 99 {+ 98 ... {+ y x}...}))}
```

```
(interp (parse '{with {x 1} } {with {y 2} } {+ 100 {+ 99 {+ 98 ... {+ y x}...}})) \Rightarrow (interp (parse '{with {y 2} [x=1] } {+ 100 {+ 99 {+ 98 ... {+ y x} ...}})) \Rightarrow (interp (parse '{+ 100 {+ 99 {+ 98 ... {+ y x} ...}} [y=2 x=1] )) \Rightarrow ... \Rightarrow (interp (parse 'y [y=2 x=1]))
```

Always add to start, then always check from start

```
(interp (parse '{with {x 1}}
                     \{+ \{with \{x 2\} x\}\}
                         x}}))
\Rightarrow
(interp (parse '{+ {with {x 2} x} [x=1]
                      x}))
\Rightarrow
(+
       (interp (parse '\{\text{with } \{x \ 2\} \ x\} [x=1] ))
    (interp (parse 'x [x=1])))
\Rightarrow
(+ (interp (parse 'x [x=2 x=1])) (interp (parse 'x [x=1])))
```

```
(interp (parse '{with {x 1}}
                      \{+ \{with \{x 2\} x\}\}
                         x}}))
\Rightarrow
(interp (parse '\{+\} {with \{x 2\} x\} [x=1]
                      x}))
\Rightarrow
(+
       (interp (parse '\{with \{x 2\} x\} [x=1] ))
    (interp (parse 'x [x=1])))
\Rightarrow
(+ (interp (parse 'x [x=2 x=1])) (interp (parse 'x [x=1])))
\Rightarrow (+ 2 1)
```

#### Representing Deferred Substitution

```
Change
; interp : WAE -> number
to
; interp : WAE DefrdSub -> number
   (define-type DefrdSub
      [mtSub]
      [aSub (name symbol?)
               (value number?)
               (saved DefrdSub?)])
; example instance
(aSub 'x 1 (aSub 'y 4 (aSub 'x 2 (mtSub))))
```

<sup>\*</sup> mtSub: mt stands for 'empty' cache (repository)

<sup>\*</sup> aSub: non-empty cache, a pair of an identifier and a value for substitution and the next pair.

```
(interp (parse '{with {x 1}}
                     {with {y 2}
                        {+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (mtSub))
\Rightarrow
(interp (parse '{with {y 2}}
                {+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (aSub 'x 1 (mtSub)))
\Rightarrow
(interp (parse '{+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (...)
```

```
(interp (parse '{with {x 1}}
                    {with {y 2}
                       {+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (mtSub))
\Rightarrow
(interp (parse '{with {y 2}}
                {+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (aSub 'x 1 (mtSub)))
\Rightarrow
(interp (parse '{+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (aSub 'y 2 (aSub 'x 1 (mtSub))))
```

```
(interp (parse '{with {x 1}}
                     {with {y 2}
                         {+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})(mtSub))
\Rightarrow
(interp (parse '{with {y 2}}
                 {+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (aSub 'x 1 (mtSub)))
\Rightarrow
(interp (parse '{+ 100 {+ 99 {+ 98 ... {+ y x} ... }}})
        (aSub 'y 2 (aSub 'x 1 (mtSub))))
\Rightarrow \dots
\Rightarrow
(interp (parse 'y) (aSub 'y 2 (aSub 'x 1 (mtSub))))
```

#### WAE Interpreter with DefrdSub

#### WAE Interpreter without DefrdSub

```
; interp : WAE DefrdSub -> number
(define (interp wae ds)
   (type-case WAE wae
      [num (n)
                       n]
      [add (I r) (+ (interp I ds) (interp r ds))]
      [sub (lr) (- (interp l ds) (interp r ds))]
      [with (i v e) (interp (subst e i (interp v)))]
         (s) (error 'interp "free identifier")]))
      [id
```

For deferred substitution, we need a helper function to lookup a value of the id, s from ds!!

#### WAE Interpreter with DefrdSub

#### WAE Interpreter with DefrdSub

```
; lookup: symbol DefrdSub -> number
(define (lookup name ds)
   (type-case DefrdSub ds
                     (error 'lookup "free identifier")]
      [mtSub
      [aSub (i v saved) (if (symbol=? i name)
                                      (lookup name saved))]))
(test (lookup 'x (aSub 'x 1 (mtSub))) 1)
(test (lookup 'y (aSub 'x 1 (aSub 'y 4 (mtSub)))) 4)
```

```
{deffun {f x} {+ 1 x}}
(interp (parse '{with {y 2} {f 10}} []))
```

```
{deffun {f x} {+ 1 x}}

(interp (parse '{with {y 2} {f 10}} []))

\Rightarrow

(interp (parse '{f 10} [y=2]))
```

```
{deffun {f x} {+ 1 x}}

(interp (parse '{with {y 2} {f 10}})

\Rightarrow

(interp (parse '{f 10} [y=2]))

\Rightarrow

(interp (parse '{+ 1 x} [...]))
```

Interpreting function body starts with only one substitution

What goes wrong if you extend the old substitution?

```
{deffun {f x} {+ y x}}
(interp (parse '{with {y 2} {f 10}} []))
```

What goes wrong if you extend the old substitution?

```
{deffun {f x} {+ y x}}

(interp (parse '{with {y 2} {f 10}} []))

\Rightarrow

(interp (parse '{f 10} [y=2]))
```

What goes wrong if you extend the old substitution?

```
{deffun {f x} {+ y x}}
(interp (parse '{with {y 2} {f 10}} []))
⇒
(interp (parse '{f 10} [y=2]))
⇒
(interp (parse '{+ y x} [x=10 y=2]))
⇒
12 Wrong!
```

What goes wrong if you extend the old substitution?

Interpreting function body starts with only one substitution

#### Scope

- Static scope
   In a language with static scope, the scope of an identifier's binding is a syntactically delimited region.
- Dynamic scope
   In a language with dynamic scope, the scope of an identifier's binding is the entire remainder of the execution during which that binding is in effect.

## Scope

```
{deffun {f p} n}
{with {n 5} {f 10}
```

- Static scope
   In a language with static scope, the scope of an identifier's binding is a syntactically delimited region.
- Dynamic scope
   In a language with dynamic scope, the scope of an identifier's binding is the entire remainder of the execution during which that binding is in effect.

## Scope

```
{deffun {f p} n}
{with {n 5} {f 10}
```

- Static scope
   In a language with static scope, the scope of an identifier's binding is a syntactically delimited region.
   The code signals an error.
- Dynamic scope
   In a language with dynamic scope, the scope of an identifier's binding is the entire remainder of the execution during which that binding is in effect.
   The code evaluates to 5.

# Static and dynamic scope in Perl

```
x = 50;
sub fun2 {
 return $x;
sub fun1 {
 my $x = 10;
 my \$y = fun2();
 return $y;
print fun1();
```

```
$x = 50:
sub fun2 {
 return $x;
sub fun1 {
 local $x = 10;
 my \$y = fun2();
 return $y;
print fun1();
```

Static scope: Output 50

dynamic scope: Output 10

# Static scope only in C

```
int x = 50;
int fun2() {
  return x;
                                                    Static scope: Output 50
int fun1() {
  int x = 10;
  int y = fun2();
  return y;
<del>-int main() {</del>
  printf("%d", fun1());
```

Which one is more reasonable? Why??

```
; interp : F1WAE list-of-FucDef DefrdSub -> number
(define (interp f1wae fundefs ds
   (type-case F1WAE f1wae
   [app (ftn arg)
        (local [(define a-fundef (lookup-fundef ftn fundefs))]
          (interp (fundef-body a-fundef)
                 fundefs
                 (interp arg fundefs ds)
                 ... ))]))
```

```
; interp : F1WAE list-of-FucDef DefrdSub -> number
(define (interp f1wae fundefs ds
   (type-case F1WAE f1wae
   [app (ftn arg)
       (local [(define a-fundef (lookup-fundef ftn fundefs))]
          (interp (fundef-body a-fundef)
                fundefs
                (aSub (fundef-arg-name a-fundef)
                    (interp arg fundefs ds)
```

```
; interp : F1WAE list-of-FucDef DefrdSub -> number
(define (interp f1wae fundefs ds)
   (type-case F1WAE f1wae
   [app (ftn arg)
        (local [(define a-fundef (lookup-fundef ftn fundefs))]
                  (fundef-body a-fundef)
          (interp
                 fundefs
                 (aSub (fundef-arg-name a-fundef)
                                                         Add substitution info
                     (interp arg fundefs ds)
                                                         for the function parameter
                     (mtSub))))]))
```

(test (interp (parse '{f 1}) (list (parse-fd '{deffun (f x) {+ x 3}})) (mtSub)) 4)

```
; interp : F1WAE list-of-FucDef DefrdSub -> number
(define (interp f1wae fundefs ds)
(type-case F1WAE f1wae
 [num (n)
             n]
 [add (I r) (+ (interp I fundefs ds) (interp r fundefs ds))]
 [sub (lr) (- (interp l fundefs ds) (interp r fundefs ds))]
 [with (i v e) (interp e fundefs (aSub i (interp v fundefs ds) ds))]
      (s) (lookup s ds)]
 ſid
 [app (f a) (local
                              [(define a-fundef (lookup-fundef f fundefs))]
                              (interp (fundef-body a-fundef)
                                 fundefs
                            (aSub (fundef-arg-name a-fundef)
                                       (interp a fundefs ds)
                                             (mtSub))
                     ))]))
```

(test (interp (parse '{f 1}) (list (parse-fd '{deffun (f x) {+ x 3}})) (mtSub)) 4)

## Compare with F1WAE Interpreter in L08

```
; [contract] subst: F1WAE symbol number -> F1WAE
(define (subst f1wae idtf val)
         (type-case F1WAE f1wae
                  [num
                                             f1wae]
                           (n)
                  add
                           (Ir)
                                              (add (subst | idtf val) (subst r idtf
val))]
                                                       (sub (subst | idtf val) (subst r
                  [sub
                                    (lr)
idtf val))]
                           (i v e) (with i (subst v idtf val) (if (symbol=? i idtf) e
                  with
         (subst e idtf val)))]
                  [id
                                                       (if (symbol=? s idtf) (num val)
                                    (s)
f1wae)]
                                                      (subst a idtf val))]))
                           (f a)
                                             (app f
                  lapp
```

# Topics we cover and schedule (tentative)

- Racket tutorials (L2,3)
- Modeling languages (L4,5)
- Interpreting arithmetic (L5)
- Language principles
  - Substitution (L6-7)
  - Function (L8)
  - Deferring Substitution (L9)
  - First-class Functions (L10-L12)
  - Laziness (L13,14)
  - Recursion (L15,16)

- Mutable data structures
   (L17,18,19,20)
- Variables (L21,22)
- Continuations (L23-26)
- Guest Video Lecture (L27)

JC

#### **TODO**

Read Chapter 6. First-Class Functions

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<sup>\*</sup> Slides are from Prof. Sukyoung Ryu's PL class in 2018 Spring or created by JC based on the main text book.