Follow up1: how to optimize space

```
// LastMax is an optimization to S[i], using O(1) space.
00 public int getMax(int[] s) {
     if (s == null || s.length < 1) {
01
02
          return 0;
03
04
     int max = Integer.MIN VALUE;
     int lastMax = s[0];
0.5
06
     for (int i = 1; i < s.length; i++) {
           lastMax = Math.max(s[i], s[i] + lastMax);
07
08
           max = Math.max(max, lastMax);
09
10
     return max;
11 }
```

FollowUp 2

What if you are **requested** to return the left and right index of the subarray that has the largest sum.

```
index 0 1 2 3 4 5 6
input {1, 2, 4, -1, -12, 10 -1},

R
solu_L
solu_R
Question1: when to update solu_L = L (when max is updated)
when to update solu_R = R (when max is updated)
when to update L (L = i when M[i-1] < 0)
when to update R (update all the time each round)
```

Solution2: DP

String1 = n letters xxxxxxxxx n String2 = m letters yyyyyyyyyy m

M[n][m] = ???????? M[n-1][m] or M[n][m-1] or M[n-1][m-1]

XXXXXXXX

XXXXXXXX

xxxx xxx Y = [i][j] represents the minimum number of actions needed to convert

xxxxxxxxx the first i-letters of s1 to the first j letters of s2

XXXXXXXX

XXXXXXXX

xxxxxxxX [n][m]

Solution:

```
1. Base case: M[0][0] = 0, and M[0][m] = m AND M[n][0] = n;
0 1 2 3 4 5 6 7
1 x x x xx
2
3
```

- Induction rule:
 - M[i][j] represents the minimum number of actions to transform substring (the first i letters of s1 to the first j letters of s2.
 - b. M[i][j] = case 1(do nothing) M[i-1][j-1] if the i-th letter of s1 == the j-th letter of s2

```
case 2 (replace) 1 + M(i-1, j-1)
case 3 (delete) 1 + M(i-1, j)
case 4 (insert) 1 + M(i, j-1)
```

= min (case 1, case 2, case3, case 4)

```
subproblem: M[i][j]: represents the max size of square with the
coordinate[i][j] as its bottom right corner.
          M[i][j] = 0 if A[i][j] == 0;
                   = 1 + min(M[i-1][j-1], M[i][j-1], M[i-1][j]) else
0 0 0 0 0
1 1 1 1 0
1 1 1 1 0
1 1 1 0 0
1 1 1 1 0 0
M[i][j] =
0 0 0 0 0
1 1 x x x
1 x x x x
1 x x x x
1 x x x x
```