



CoderCareer: Discussing Coding Interview Questions from Google, Amazon, Facebook, Microsoft, etc

Thursday, October 27, 2011

No. 15 - Fibonacci Sequences

Problem: Please implement a function which returns the n^{th} number in Fibonacci sequences with an input n . Fibonacci sequence is defined as:

$$f(n) = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ f(n-1) + f(n-2) & n > 1 \end{cases}$$

Analysis: It is a classic interview questions to get numbers in Fibonacci sequences. We have different solutions for it, and their performance varies a lot.

Solution 1: Inefficient recursive solution

Fibonacci sequences are taken as examples to lecture recursive functions in many C/C++ textbooks, so most of candidates are familiar with the recursive solution. They feel confident and delighted when they meet this problem during interviews, because they can write the following code in short time:

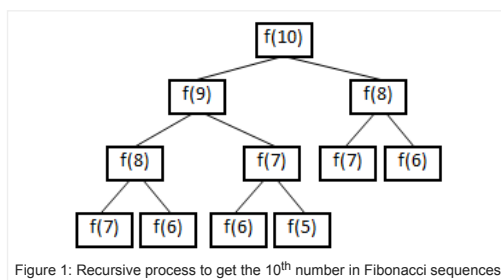
```
long long Fibonacci(unsigned int n)
{
    if(n <= 0)
        return 0;

    if(n == 1)
        return 1;

    return Fibonacci(n - 1) + Fibonacci(n - 2);
}
```

Our textbooks take Fibonacci sequences as examples for recursive functions does not necessarily mean that recursion is a good solution for Fibonacci sequences. Interviewers may tell candidates that the performance of this recursive solution is quite bad, and ask them to analyze root causes.

Let us take $f(10)$ as an example to analyze the recursive process. We have to get $f(9)$ and $f(8)$ before we get $f(10)$. Meanwhile, $f(8)$ and $f(7)$ are needed before we get $f(9)$. The dependency can be visualized in a tree as shown in Figure 1:



```

        return result[n];

    long long    fibNMinusOne = 1;
    long long    fibNMinusTwo = 0;
    long long    fibN = 0;
    for(unsigned int i = 2; i <= n; ++ i)
    {
        fibN = fibNMinusOne + fibNMinusTwo;

        fibNMinusTwo = fibNMinusOne;
        fibNMinusOne = fibN;
    }

    return fibN;
}

```

Solution 3: $O(\log n)$ solution

Usually interviewers expect the $O(n)$ solution above. However, there is an $O(\log n)$ solution available, which is based on an uncommon equation as shown below:

$$\begin{bmatrix} f(n) & f(n-1) \\ f(n-1) & f(n-2) \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1}$$

It is not difficult to prove this equation with mathematical induction. Interested readers may have try.

Now the only problem is how to calculate power of a matrix. We can calculate power with exponent n in $O(\log n)$ time with the following equation:

$$a^n = \begin{cases} a^{n/2} * a^{n/2} & n \text{ is even} \\ a^{(n-1)/2} * a^{(n-1)/2} * a & n \text{ is odd} \end{cases}$$

The source code to get power of a matrix looks complicated, which is listed below:

```

#include <cassert>

struct Matrix2By2
{
    Matrix2By2
    (
        long long m00 = 0,
        long long m01 = 0,
        long long m10 = 0,
        long long m11 = 0
    )
    :m_00(m00), m_01(m01), m_10(m10), m_11(m11)
    {
    }

    long long m_00;
    long long m_01;
    long long m_10;
    long long m_11;
};

Matrix2By2 MatrixMultiply
(
    const Matrix2By2& matrix1,
    const Matrix2By2& matrix2
)
{
    return Matrix2By2(
        matrix1.m_00 * matrix2.m_00 + matrix1.m_01 * matrix2.m_10,
        matrix1.m_00 * matrix2.m_01 + matrix1.m_01 * matrix2.m_11,
        matrix1.m_10 * matrix2.m_00 + matrix1.m_11 * matrix2.m_10,
        matrix1.m_10 * matrix2.m_01 + matrix1.m_11 * matrix2.m_11);
}

Matrix2By2 MatrixPower(unsigned int n)
{
    assert(n > 0);

    Matrix2By2 matrix;
    if(n == 1)
    {
        matrix = Matrix2By2(1, 1, 1, 0);
    }
    else if(n % 2 == 0)
    {
        matrix = MatrixPower(n / 2);
        matrix = MatrixMultiply(matrix, matrix);
    }
    else if(n % 2 == 1)
    {
        matrix = MatrixPower((n - 1) / 2);
        matrix = MatrixMultiply(matrix, matrix);
        matrix = MatrixMultiply(matrix, Matrix2By2(1, 1, 1, 0));
    }
}

```

[No. 11 - Print Binary Trees from Top to Bottom](#)

[No. 10 - K-th Node from End](#)

[No. 09 - Numbers with a Given Sum](#)

[No. 08 - Calculate 1+2+...+n](#)

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Harry He

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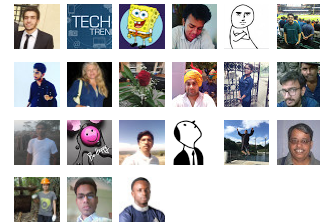
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```

        return matrix;
    }

    long long Fibonacci(unsigned int n)
    {
        int result[2] = {0, 1};
        if(n < 2)
            return result[n];

        Matrix2By2 PowerNMinus2 = MatrixPower(n - 1);
        return PowerNMinus2.m_00;
    }

```

Even though it cost only $O(\log n)$ time in theory, its hidden constant parameter is quite big, so it is not treated as a practical solution in real software development. Additionally, it is not a recommended solution during interviews since its implementation code is very complex.

The discussion about this problem is included in my book <Coding Interviews: Questions, Analysis & Solutions>, with some revisions. You may find the details of this book on Amazon.com, or [Apress](http://Apress.com).

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Posted by [Harry He](#) at 11:48 AM



Labels: [Algorithm](#), [Amazon](#), [Interview Question](#), [Microsoft](#)

10 comments:



buddystripes February 1, 2013 at 1:53 AM

Hey! Thanks for posting all of these cool problems. I wanted to suggest an ever faster solution. It's constant time if you solve out the closed form solution for the nth fibonacci number.

<http://math.stackexchange.com/questions/55922/fibonacci-recurrence-relations>

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Harry He February 4, 2013 at 3:40 PM

Thanks for your reply. The equation listed in the URL you provided is cool, but is not better than the solution in this post: (1) Precision issue. The equation is based on the value of $\sqrt{5}$. The calculation on double values is not as accurate as calculation on integer values. (2) The time efficiency is not constant. We can only calculate $[(1 + \sqrt{5}) / 2]^n$ and $[(1 - \sqrt{5}) / 2]^n$ in $O(\log n)$ time.



Sahil Rally October 16, 2013 at 4:55 PM

Hi Harry

Why to consider Precision Issue iff it is giving accurate answer ?

[Reply](#)



Satish Dixit September 25, 2015 at 12:29 AM

very nice explanation. One of the different fibonacci problem is mentioned in this link <http://techno-terminal.blogspot.in/2015/09/fibonacci-problem-calculate-sum-fx-1.html>

[Reply](#)



djax February 8, 2016 at 6:26 AM

Binet's formula solves this instantaneously.

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djax February 8, 2016 at 6:27 AM

Accuracy is maintained, all the floating values cancel in the numerator and denominator



djax February 8, 2016 at 6:28 AM

"The calculation on double values not as accurate..." Who would ever ask for the 2.2th fibonacci number? That is meaningless.



djax February 8, 2016 at 6:35 AM

Accuracy is maintained, all the floating values cancel in the numerator and denominator

[Reply](#)



djax February 8, 2016 at 6:37 AM

Binet's formula solves this instantaneously.

[Reply](#)



Unknown May 19, 2017 at 4:48 AM

Solution one is not only inefficient it is also incorrect.

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