

THE TARSKI LAPLACIAN

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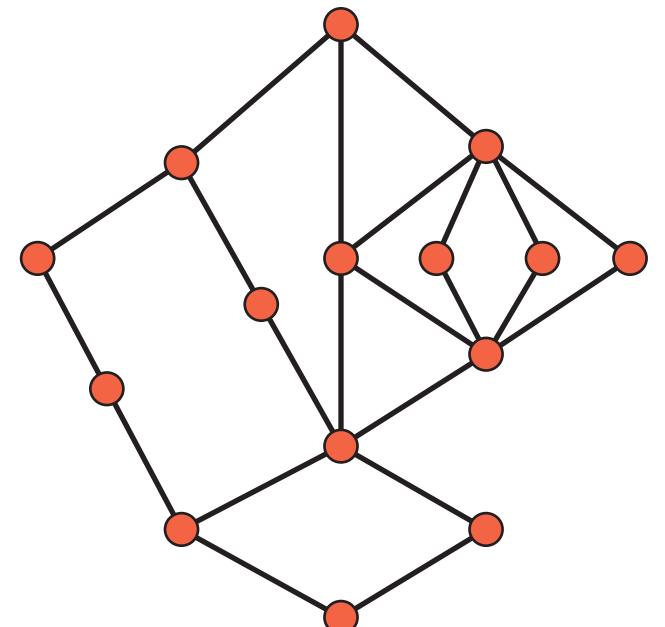
Joint work with Paige Randall-North, Robert Ghrist,
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Lattices and Galois Connections

Lattices are algebraic structures—much like vector spaces or symmetry groups—that distinguish themselves by their hierarchical nature. Given any two lattice elements, one may perform lattice operations “meet” and “join” denoted $x \wedge y$ and $x \vee y$. These operations behave like intersection and union, satisfying a number of axioms (symmetry, idempotence, absorption, identity), and inherit an order \leq . Lattices also are assumed to have a top element 1 and a bottom element 0. Examples of lattices include the powerset 2^X with intersection and union, the unit interval $[0, 1]$ with minimum and maximum, and the collection of subspaces of a vector space with intersection and sum.

Lattices are related by structure-preserving maps between them. These are called Galois connections. Galois connection are to lattices as linear maps and their adjoints (transpose) are to inner-product spaces. Formally, given a lattice \mathbf{P} and a lattice \mathbf{Q} , a **Galois connection** consists of a pair of maps $L: \mathbf{P} \leftrightarrows \mathbf{Q}: R$ such that $L(x) \geq y$ if and only if $R(y) \geq x$.

Example. Let X (e.g. users) and Y (e.g. movies) be sets, and $I \subseteq X \times Y$ a binary relation (e.g. user x watched movie y). Then, $(-)^\dagger: 2^X \leftrightarrows 2^Y: (-)^\ddagger$ is a Galois connection given by $\sigma^\dagger = \{y \in Y: (x, y) \in I \ \forall x \in X\}; \tau^\ddagger = \{x \in X: (x, y) \in I \ \forall y \in Y\}$



	SPACE WARS	LORD OF THE BANDS	HARRY CERAMICS
ALICE	Orange		Orange
BOB	Blue	Blue	
EVE			Red

Weighted Galois Connections

One challenge to designing information processing systems with lattices is that there are no “coefficients,” a necessity in order to employ traditional training techniques in machine learning. This leads to a notion of weighted lattices. A **weighted lattice** is a collection of maps $[X, \mathbb{L}]$ for which X is a set and \mathbb{L} is a lattice with additional structure called a **residuated lattice**. A residuated lattice is a tuple $\mathbb{L} = (L, \wedge, \vee, 0, 1, \otimes, e, \Rightarrow)$ such that $(L, \wedge, \vee, 0, 1)$ is a lattice and (L, \otimes, e) is a monoid (set with associative binary operation & identity) with a notion of implication \Rightarrow dual to \otimes . A notion of **subsethood** [Belohlavek 1999] quantifies the degree to which a two elements σ, τ of \mathbb{L} “are related”

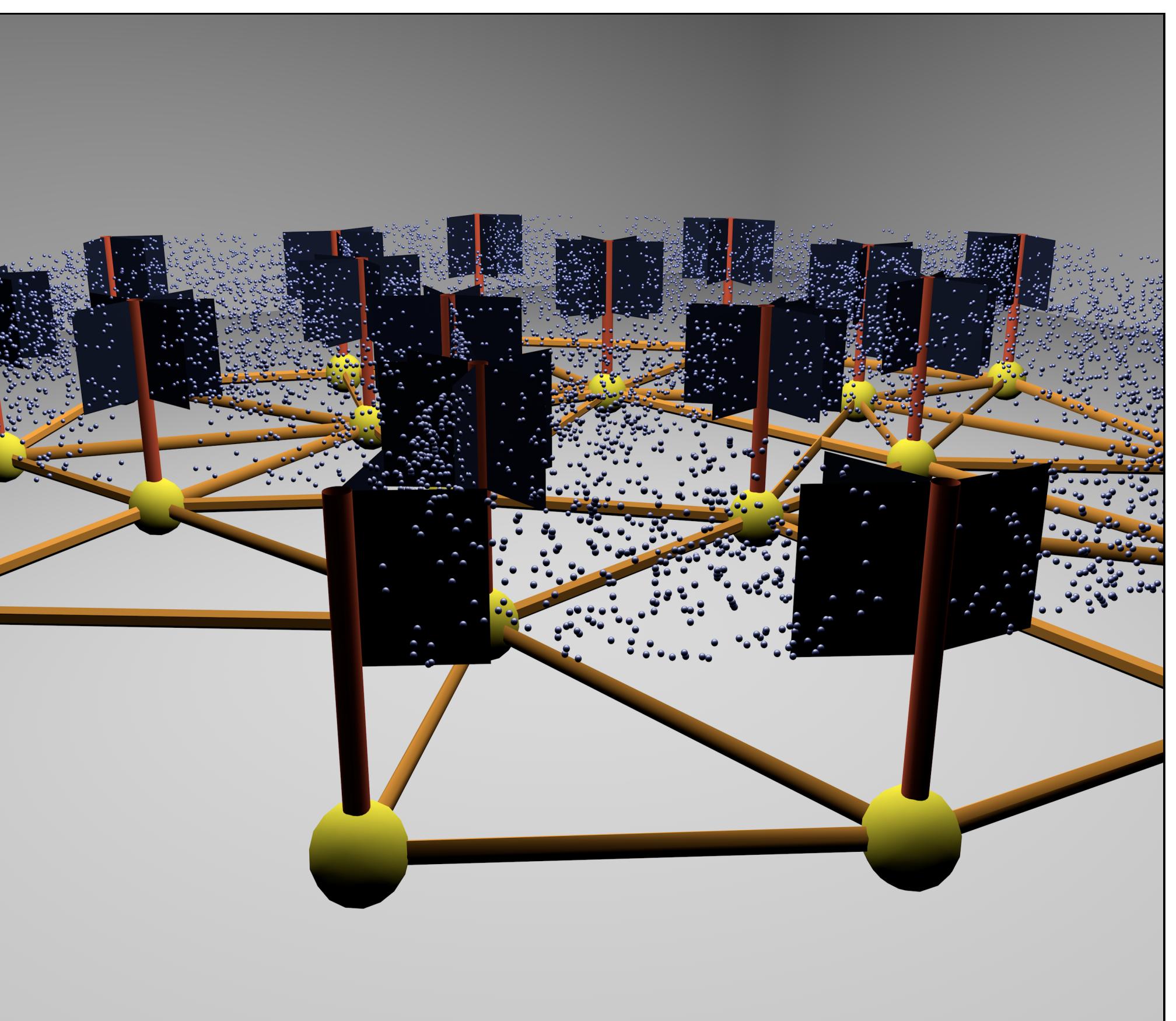
$$\text{Sub}(\sigma, \tau) = \bigwedge_{x \in X} \sigma(x) \Rightarrow \tau(x)$$

Given a weighted relation $I: X \times Y \rightarrow \mathbb{L}$, there is a **weighted Galois connection** between $[X, \mathbb{L}]$ and $[Y, \mathbb{L}]$ given by

$$\sigma^\dagger(y) = \bigwedge_{x \in X} \sigma(x) \Rightarrow I(x, y); \tau^\ddagger(y) = \bigwedge_{y \in Y} \sigma(y) \Rightarrow I(x, y).$$

Such a Galois connection satisfies the property

$$\text{Sub}(\tau^\dagger, \sigma) = \text{Sub}(\sigma^\dagger, \tau).$$



Sections vs. Concepts

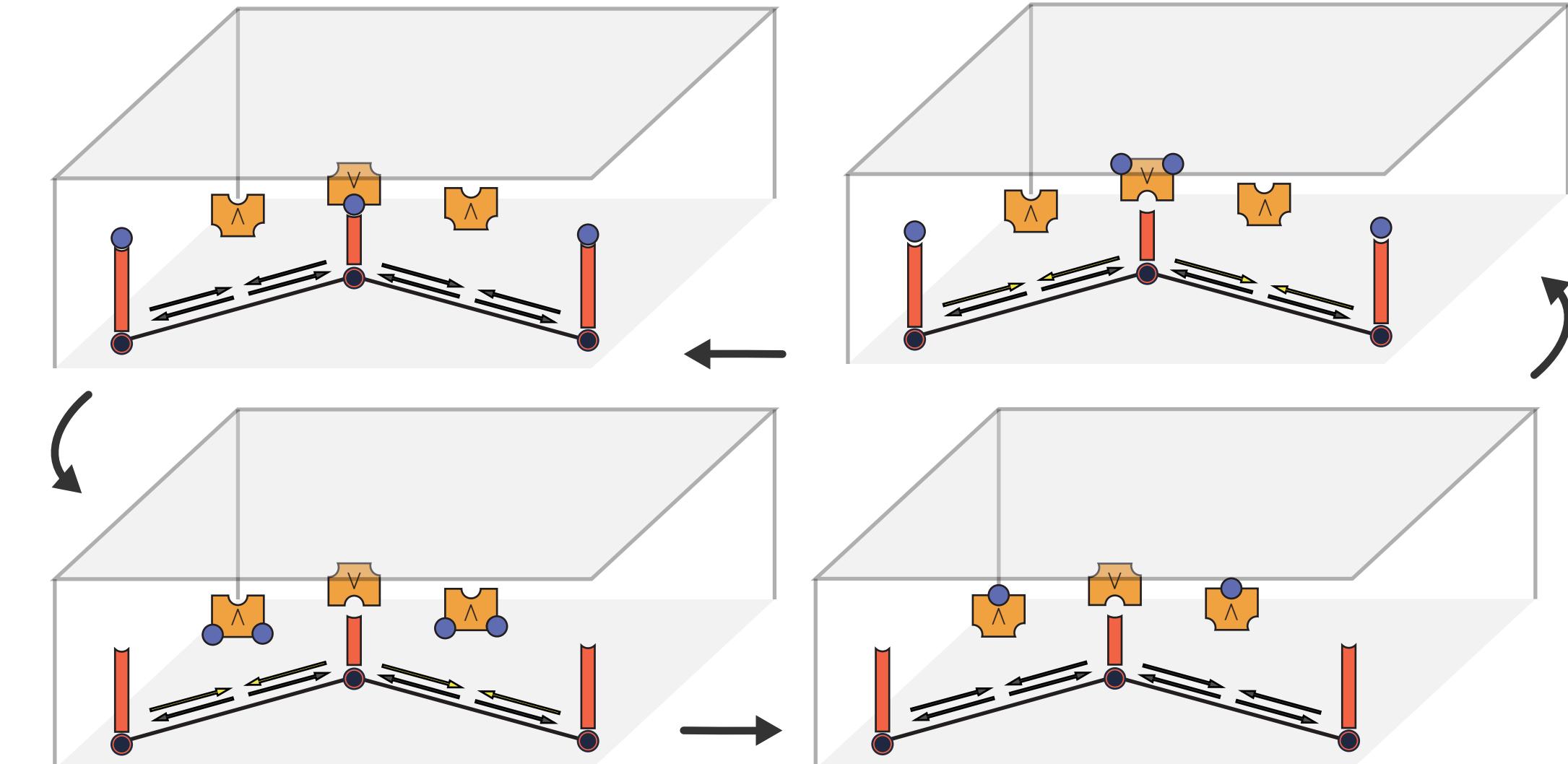
Both sections and concepts are consistent assignments of lattice-valued data to a network. The nomenclature of a section comes from sheaf theory, a deep area of mathematics studying local-to-global behavior, where the term (formal) concept was coined by Wille (1982) in deference to ontology. While there are multiple ways to associate lattices to a network, we will narrow our focus to **relation-weighted graphs**. Let $\mathcal{G} = (\mathcal{V}, \mathcal{E}, X, Y, \mathbb{L}, \mathcal{I})$ be a graph with a collection

$$\mathcal{I} = \left\{ I_{v,w}: X \times Y \rightarrow \mathbb{L} \right\}_{v \in \mathcal{V}, w \in \mathcal{W}}$$

of relations. A **section**, then, is a collection $(\sigma_v: X \rightarrow \mathbb{L})_{v \in \mathcal{V}}$ such that $\sigma_v^{\dagger v} = \sigma_w^{\dagger w}$ for all (v, e, w) .

A **concept** is a section with $(\tau_e: Y \rightarrow \mathbb{L})_{e \in \mathcal{E}}$ such that $\tau_e^{\dagger v} = \tau_{e'}^{\dagger v}$ for all (e, v, e') .

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Tarski Laplacian

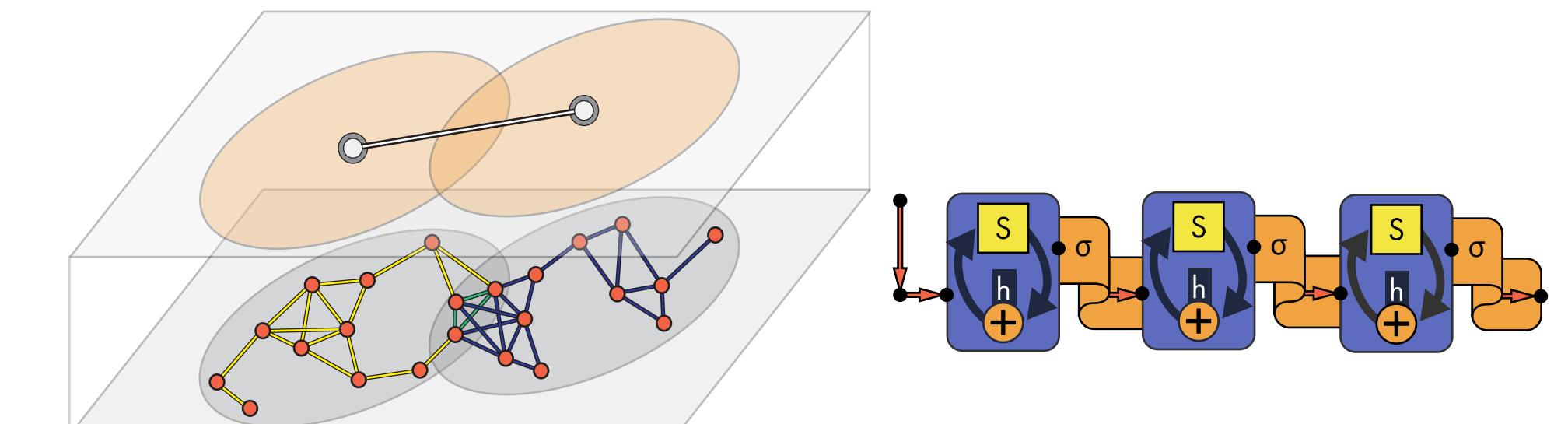
We are ready to define the Tarski Laplacian which acts on assignments of lattice-valued data over a network. The one direction of our research is using the Tarski Laplacian as a shift operator for GSP. Here's the setup: let $\mathcal{G} = (\mathcal{V}, \mathcal{E}, X, Y, \mathbb{L}, \mathcal{I})$ be a relation-weighted graph. Then the **Tarski Laplacian** is the operator

$$L: \prod_{v \in \mathcal{V}} [X, \mathbb{L}] \rightarrow \prod_{v \in \mathcal{V}} [X, \mathbb{L}]$$

given by

$$(L\sigma)_v = \bigvee_{(w,e) \in \mathcal{N}(v)} (\sigma_v^{\dagger e} \wedge \sigma_w^{\dagger e})^{\downarrow w}$$

Theorem. Let $\rho \in \mathbb{L}$, then $\bigwedge_{v \in \mathcal{V}} \text{Sub}(\sigma_v, (L\sigma)_v) \geq \rho$ if and only if $\text{Sub}(\sigma_v^{\dagger e}, \sigma_w^{\dagger e}) \geq \rho$ for all (v, e, w) . This theorem says the Tarski Laplacian, much like the graph Laplacian, “smoothes” out data in order to reach a section. (The figure above illustrates what is happening.)



Example. Let $X = \mathcal{V}_1, Y = \mathcal{V}_2$ be the vertex sets of two particular communication subgraphs $\mathcal{G}_1 = (\mathcal{V}_1, \mathcal{E}_1), \mathcal{G}_2 = (\mathcal{V}_2, \mathcal{E}_2)$ in a sensor network. Let $\mathbb{L} = ([0, 1], \min, \max, 0, 1, \cdot, \rightarrow)$, a residuated lattice with $s \rightarrow t = \max(t/s, 1)$. Assume $\mathcal{V}_1 \cap \mathcal{V}_2 \neq \emptyset$. Then, \mathcal{G}_1 and \mathcal{G}_2 are connected (though possibly “distant”). We build a relation weighted graph with $I(v_1, w) = e^{d(v_1, w)}$ and $I(v_1, w) = e^{d(v_1, w)}$ with $w \in \mathcal{V}_1 \cap \mathcal{V}_2$. With subgraph signals $s_1: \mathcal{V}_1 \rightarrow [0, 1], s_2: \mathcal{V}_2 \rightarrow [0, 1]$, we can “harmonize” them by applying the weighted Tarski Laplacian

$$L(s_1, s_2) = ((s_1^{\dagger 1} \wedge s_2^{\dagger 1})^{\downarrow 1}, (s_1^{\dagger 1} \wedge s_2^{\dagger 2})^{\downarrow 2}).$$

Explicitly, for $i \in \{1, 2\}$,

$$L(s_1, s_2)_i(v_i) = \min_{w \in \mathcal{V}_1 \cap \mathcal{V}_2} \frac{\min_{v_1 \in \mathcal{V}_1} \max_{v_2 \in \mathcal{V}_2} (\frac{s_1(v_1)}{s_1(v_1)}, 1), \min_{v_2 \in \mathcal{V}_2} \max_{v_1 \in \mathcal{V}_1} (\frac{s_2(v_2)}{s_2(v_2)}, 1)}{e^{-d(v_i, w)}}$$

We hypothesize that the weighted Tarski Laplacian prove useful in distributed training of GNNs.