Y86 Stack Frames

Topics

 How do procedures organize data such as local variables, return addresses, etc.?

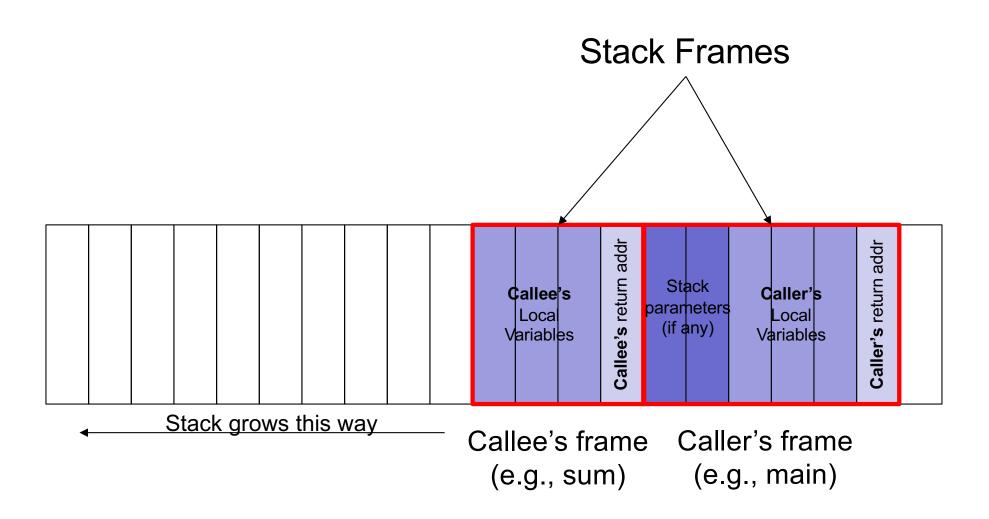
Learning Objectives

- Draw a stack frame illustrating how stack frames are established after a function call.
- Implement construction and teardown of a call frame.

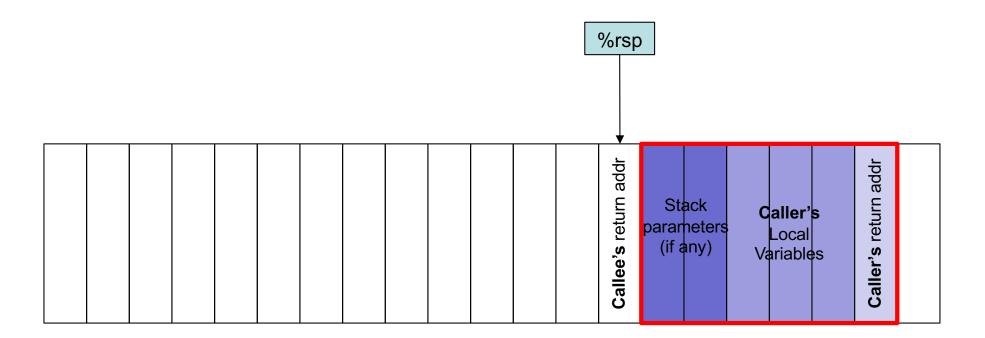
Using the Stack

- CALL: Uses the stack to store a return address
- We also use the stack to:
 - Transmit parameters (if we cannot use registers)
 - Store local variables
- The structure we use to store this information is called a stack frame.

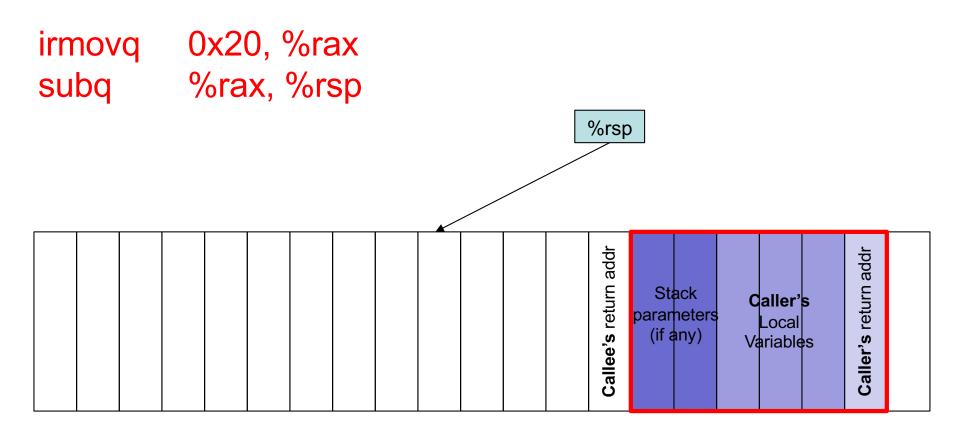
Stack Frames



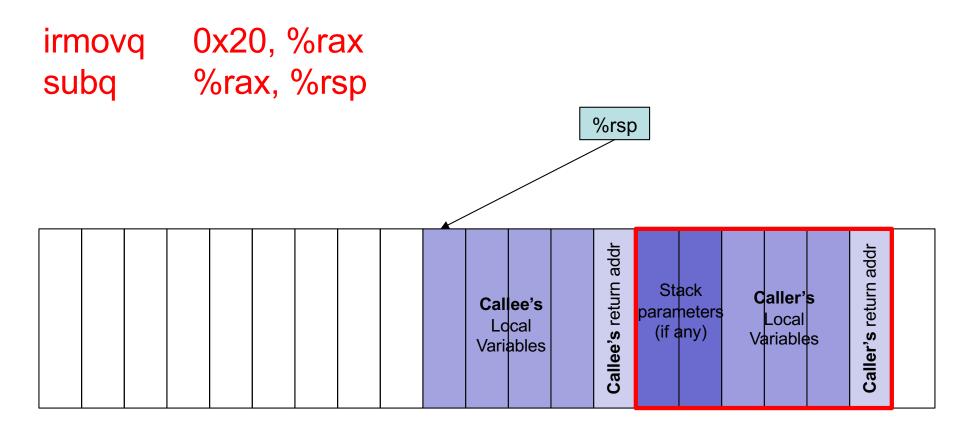
Stack Frames: no Base Pointer (1)



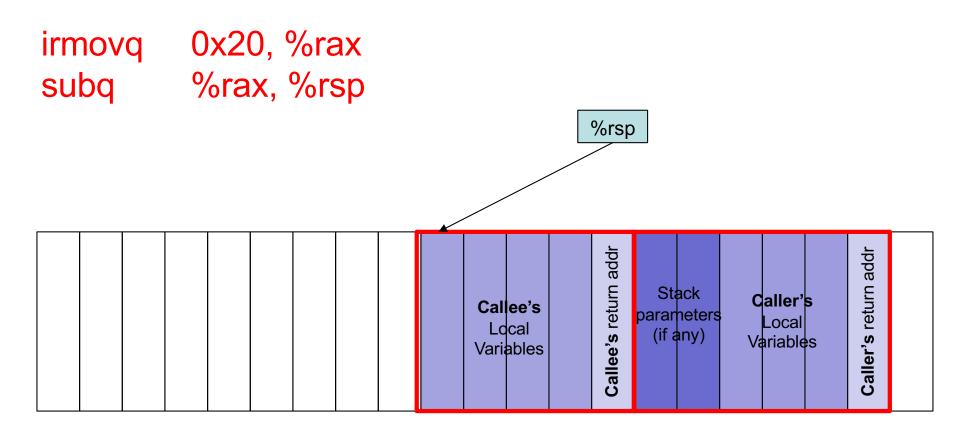
Stack Frame Setup: no Base Pointer (2)



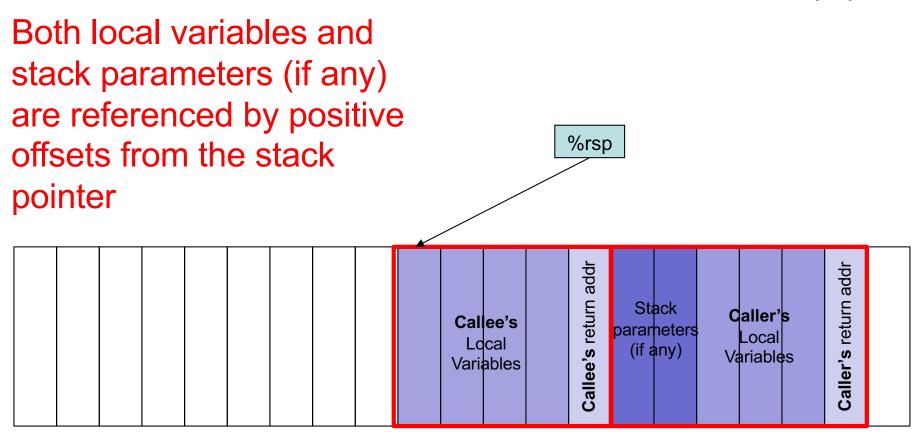
Stack Frame Setup: no Base Pointer (3)



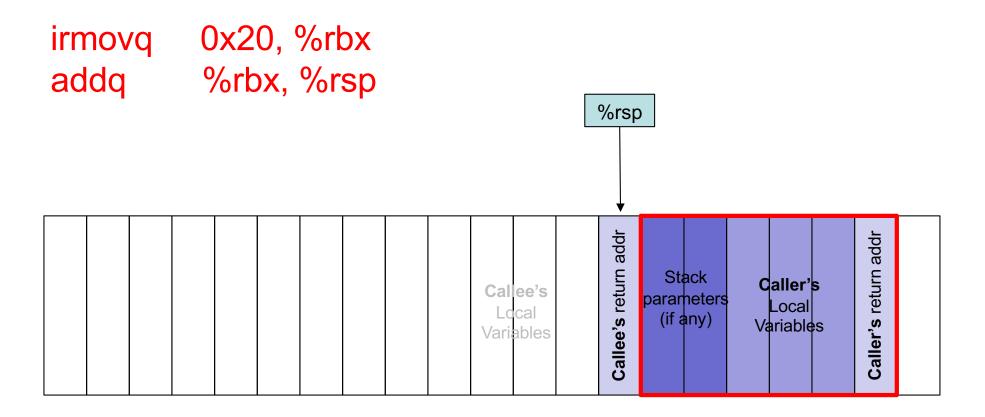
Stack Frame Setup: no Base Pointer (4)



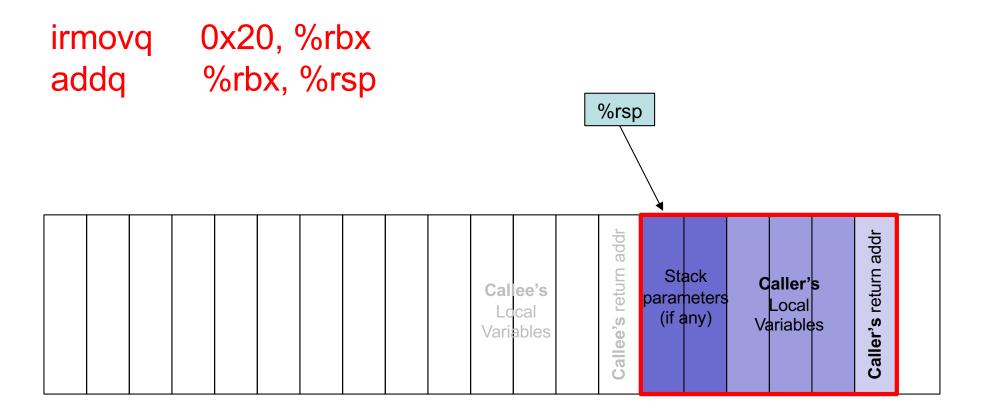
Stack Frame Setup: no Base Pointer (5)



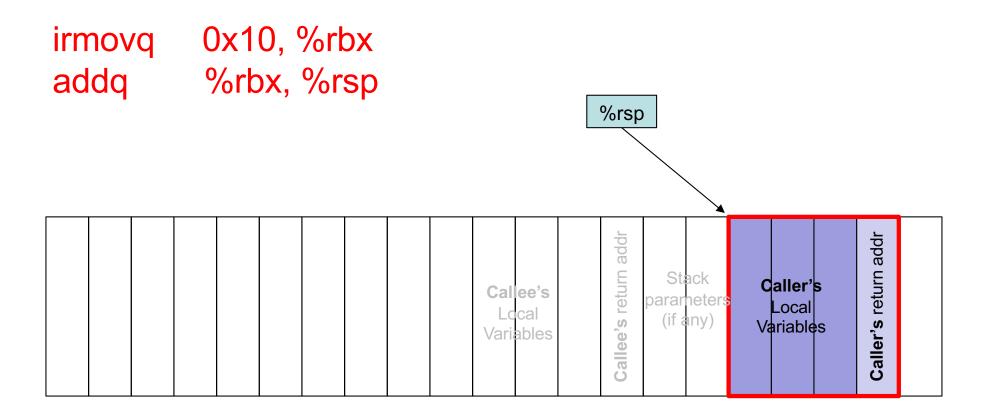
Stack Frame Teardown: no Base Pointer (6)



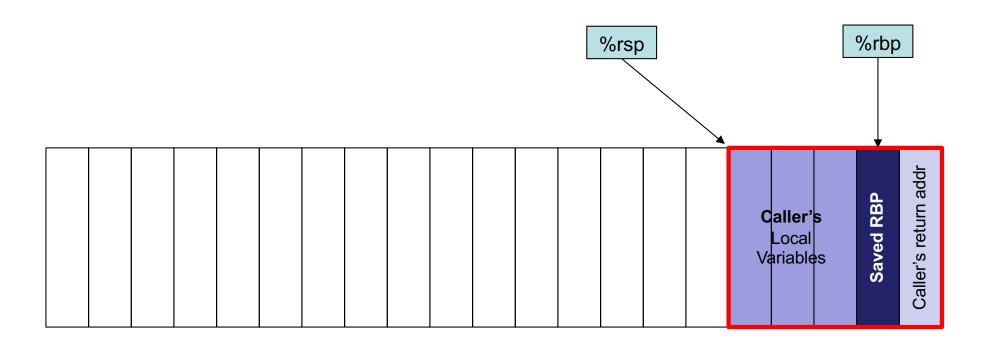
Stack Frame Teardown: no Base Pointer (7)



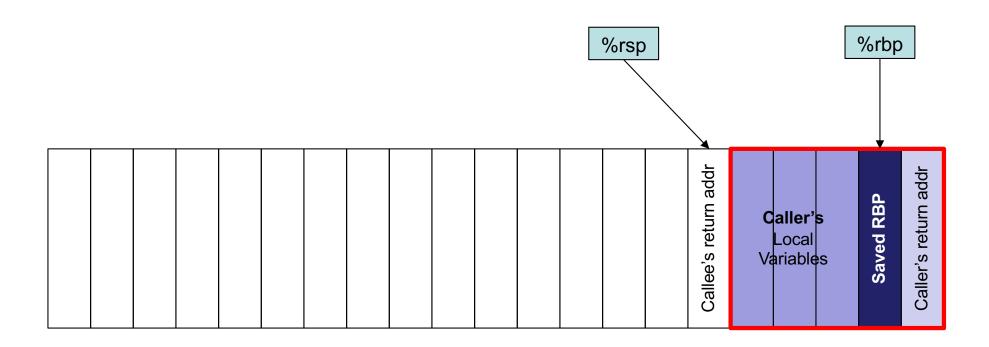
Stack Frame: no Base Pointer (8)



Stack Frames: with Base Pointer (1)

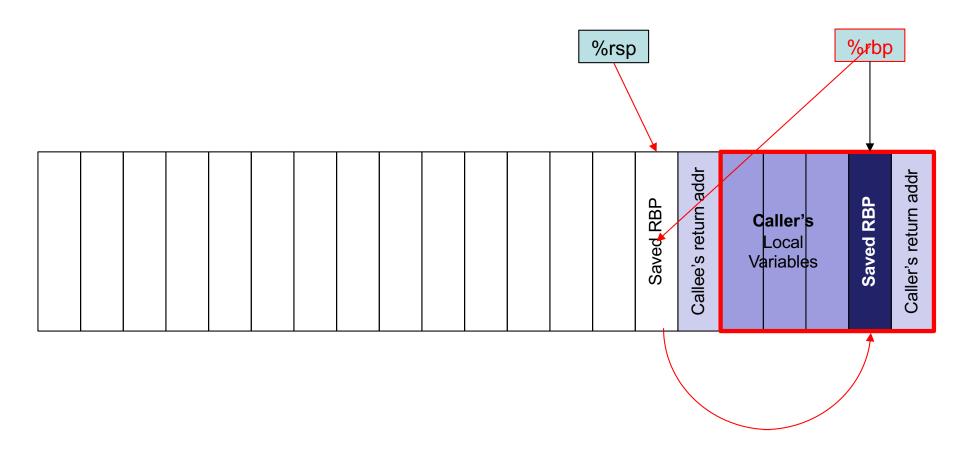


Stack Frame Setup: with Base Pointer (2)



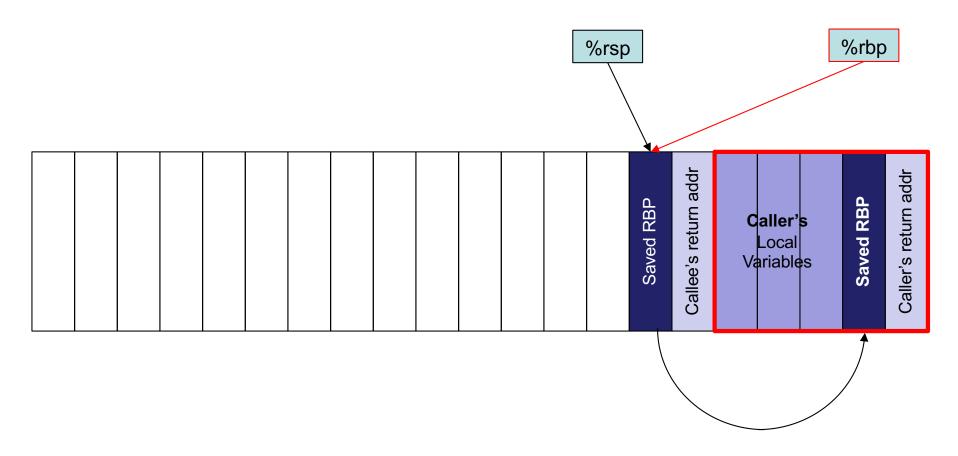
Stack Frame Setup: with Base Pointer (3)

pushq %rbp

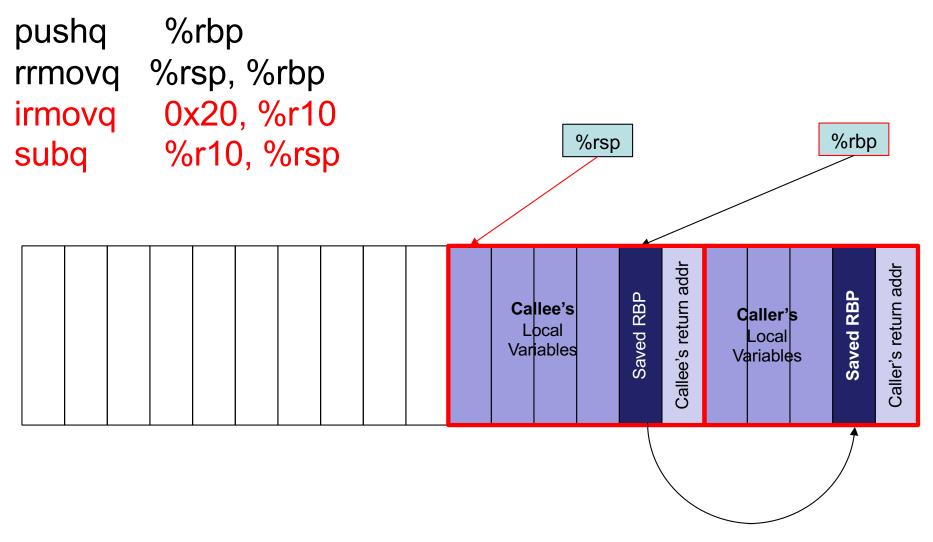


Stack Frame Setup: with Base Pointer (4)

pushq %rbp
rrmovq %rsp, %rbp

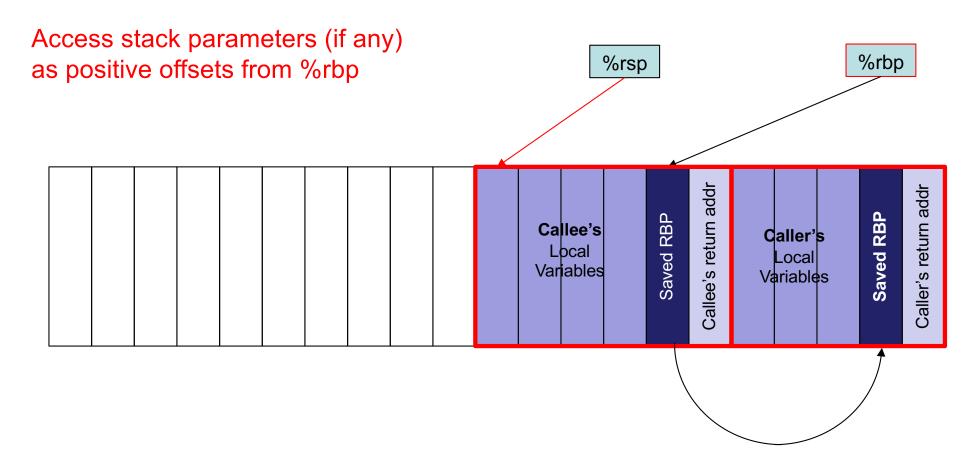


Stack Frame Setup: with Base Pointer (5)



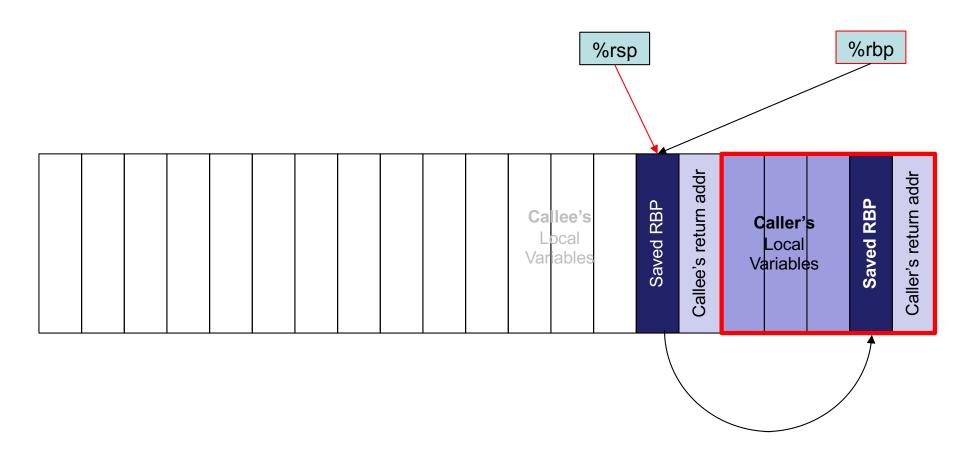
Stack Frames: with Base Pointer (6)

Access locals as negative offsets from %rbp



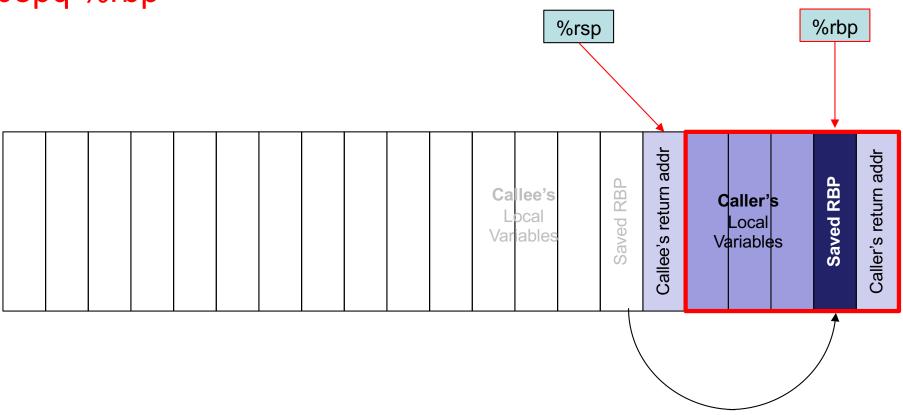
Stack Frame Teardown: with Base Pointer (7)

irmovq 0x20, %r10 addq %r10, %rsp



Stack Frame Teardown: with Base Pointer (8)

irmovq 0x20, %r10 addq %r10, %rsp popq %rbp



Stack Frame Teardown: with Base Pointer (9)

irmovq 0x20, %r10 addq %r10, %rsp popq %rbp

