Today

- Learning Outcomes
 - Define
 - Supervisor mode
 - User mode
 - Processor privilege level
 - System call
 - File descriptor
 - File permissions (mode)
 - Use the system calls open, close, read, write.
- Reading
 - Chapter 10: 10.1-10.4

Processor Privilege levels

- Modern hardware has multiple privilege levels.
- Different software can run with different privileges.
- Processor hardware typically provides (at least) two different modes of operation:
 - User mode: how all "regular programs" run.
 - Kernel mode or supervisor mode: how the OS runs.
 - Most processors have only two modes; x86 has four; some older machines had 8!
- The mode in which a piece of software is running determines:
 - What instructions may be executed.
 - What memory locations may be accessed (enforced through translation).
 - Whether the program can directly interface with devices.

System Calls

- When a program wants to do something that it cannot do in user mode (e.g., interact with a disk), it needs to ask the kernel to perform the activity on its behalf.
- System calls are the mechanism that accomplishes this.
 - They are a way to for a program to explicitly request the operating system to do something.
 - System calls appear in section 2 of the man pages.

The POSIX File System API

- The file system APIs are implemented as library functions (in libc) that invoke system calls.
- open: takes a pathname (the name of a file) and returns a file descriptor (fd), which you will use to access that file.
- close: takes an fd and frees up all the OS and library resource that have been allocated to it.
- read: reads data from the file referenced by fd into a user-specified buffer.
- write: takes data from a buffer and writes it to the file referenced by fd.

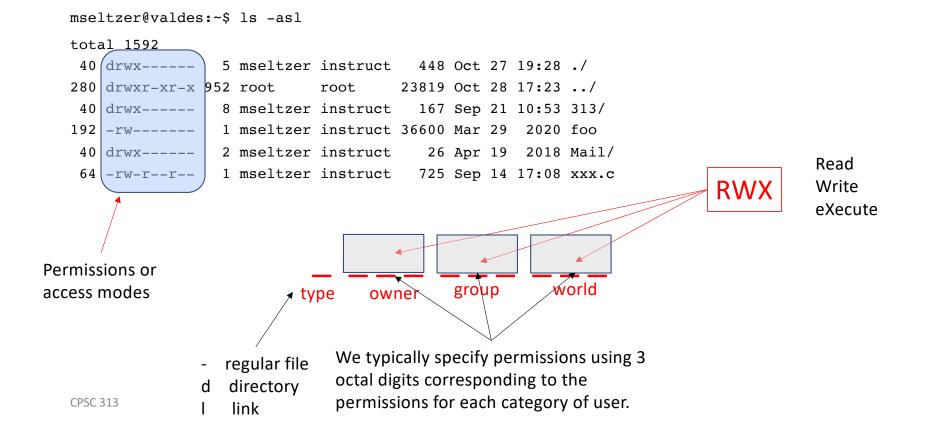
A Copy Program in Pseudo code

```
copy(infile, outfile)
fd_in = open infile for reading
fd_out = open outfile for writing
while (there is data left to read)
  read data from fd_in
  write data to fd_out
close fd_in
close fd_out
```

open

- int open(const char *path, int oflag, int mode)
- Parameters:
 - path: the path name of the file you wish to open
 - oflag: flag values that tell you things such as:
 - How you want to access the file: O RDONLY, O WRONLY, O RDWR
 - What to do if the file does/does not exist (O CREATE, O EXCL
 - How the file should behave with respect to (OS) caching (O SYNC, O DIRECT).
 - mode: if the file does not exist and you create it, what should its access mode be set to (e.g., who is allowed to read/write it). You can ignore this for now; we'll come back to it.
- Return value: a file descriptor
 - non-negative integer on success; -1 on error).
- Example
 - open("myfile", O_CREAT, 0600)

A Brief Digression on File System Modes



close

- int close(int fd)
- Parameters:
 - fd: the file descriptor to be closed
- Return value: 0 on success; -1 on error
- Example
 - int ret = close(fd);

read

Note: There was some voiceover in the video that might be confusing: A return value of 0 is the **only return that always indicates end of file**. A short read *may* indicate end of file but does not necessarily do so. Therefore, after a short read, you should try to read again.

- ssize_t read(int fd, void *buf, size_t nbyte)
- Parameters:
 - fd: file descriptor for the file from which to read
 - buf: a buffer into which the bytes read are placed
 - nbyte: the number of bytes to read
- Return value: a ssize t (a signed long)
 - On success: the number of bytes read
 - On end-of-file: 0
 - On failure: -1
- Example
 - char buf[4096];
 - ssize_t bytes_read = read(fd, buf, 4096);

write

- ssize_t write(int fd, const void *buf, size_t nbyte)
- Parameters:
 - fd: file descriptor for the file to which to write
 - buf: a buffer containing the data to write
 - nbyte: the number of bytes to write
- Return value: a ssize t (a signed long)
 - On success: the number of bytes written
 - On failure: -1
- Example
 - ssize_t bytes_written = write(fd, buf, 4096);

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