# Data Wrangling and Data Analysis Heterogeneous Data Integration (Part B) Entity Linkage

**Hakim Qahtan** 

Slides prepared by

**Prof. Yannis Velegrakis** 

Department of Information and Computing Sciences
Utrecht University



### I won a trip to L.A.

That is Lekanopedio Attikis

Not the famous **L**os **A**ngeles!!











capital of UK, host city of the IV Olympic Games, host city of the XIV Olympic Games, future host of the XXX Olympic Games, city of the Westminster Abbey, city of the London Eye, the city described by Charles Dickens in his novels, ...





London 런던 ๑ฉ๗ लंडन लंदन લંડન ለንደን ロンドン ศษ ลอนดอน இலண்டன் ლონდონი Llundain Londain Londe Londen Londen Londen Londinium London Londona Londonas Londoni Londono Londra Londres Londrez Londyn Lontoo Loundres Luân Đôn Lunden Lundúnir Lunnainn Lunnon וויני ווויני וויני וו

capital of UK, host city of the IV Olympic Games, host city of the XIV Olympic Games, future host of the XXX Olympic Games, city of the Westminster Abbey, city of the London Eye, the city described by Charles Dickens in his novels, ...

http://sws.geonames.org/2643743/ http://en.wikipedia.org/wiki/London http://dbpedia.org/resource/Category:London

. . .



#### ... or ...

#### How many "entities" have the same name?

- London, KY
- London, Laurel, KY
- London, OH
- London, Madison, OH
- London, AR
- London, Pope, AR
- London, TX
- London, Kimble, TX
- London, MO
- London, MO

. . .

- London, London, MI
- London, London, Monroe, MI
- London, Uninc Conecuh County, AL
- London, Uninc Conecuh County, Conecuh, AL
- London, Uninc Shelby County, IN
- London, Uninc Shelby County, Shelby, IN
- London, Deerfield, WI
- London, Deerfield, Dane, WI
- London, Uninc Freeborn County, MN



#### ... or ...

#### How many "entities" have the same name?

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- London, Deerfield, WI
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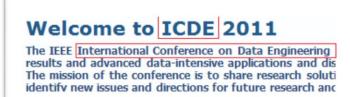
- London, Jack 2612 Almes Dr Montgomery, AL (334) 272-7005
- London, Jack R
   2511 Winchester Rd
   Montgomery, AL 36106-3327
   (334) 272-7005
- London, Jack
   1222 Whitetail Trl
   Van Buren, AR 72956-7368
   (479) 474-4136
- London, Jack
   7400 Vista Del Mar Ave
   La Jolla, CA 92037-4954
   (858) 456-1850

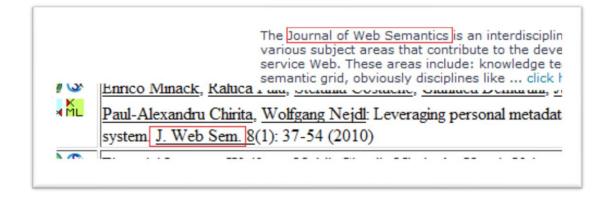
...

### Reasons of Different Descriptions

#### Text variations:

- Misspellings
- Acronyms
- Transformations
- Abbreviations
- etc.







### Reasons for Different Descriptions

- Text variations
- Local knowledge:
  - Each source uses different formats
     e.g., person from publication vs. person from email
  - Lack of global coordination for identifier assignment



### On-the-Fly Entity-Aware Query Processing in the Presence of Linkage

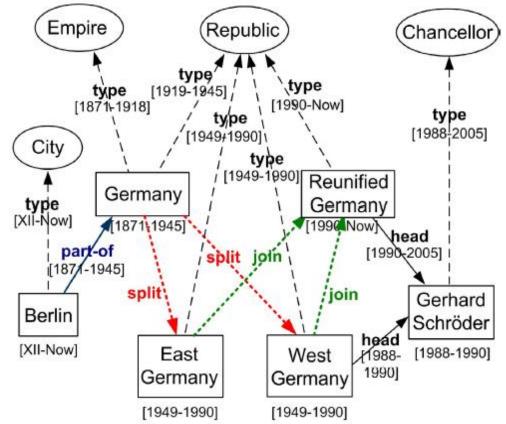




Claudia Niederée L3S Research Center Hannover, Germany niederee@L3S.de Yannis Velegrakis University of Trento Trento, Italy velgias@disi.unitn.eu

### Reasons for Different Descriptions [Velegrakis BM09]

- Text variations
- Local knowledge
- Evolving nature of data:
  - Entity alternative names
  - appearing in time
  - Updates in entity data



Jacqueline Lee Bouvier







### Reasons for Different Descriptions

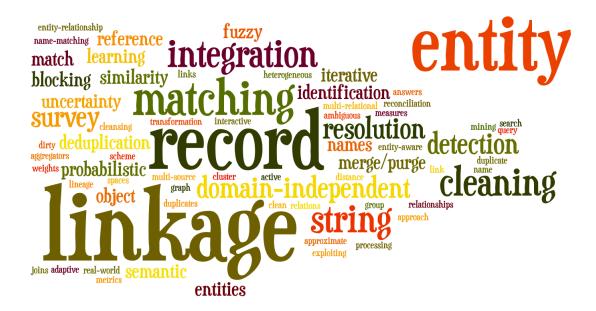
- Text variations
- Local knowledge
- Evolving nature of data
- New functionality:
  - Import data collections from various applications
    - e.g., Wikipedia data used in Freebase



### **Entity Resolution**

[Dong et al., Book 2015] [Elmagarmid et al., TKDE 2007]:

identify the different structures/records that model the same real-world object.





### Why it is useful

- Improves data quality and integrity
- Fosters re-use of existing data sources
- Optimize space

#### Application areas:

Linked Data, Social Networks, census data, price comparison portals



### Challenges for ER

- Variety Semantic
  - Semi-structured data → unprecedented levels of heterogeneity
  - Numerous entity types & vocabularies
  - LOD Cloud\*: ~50,000 predicates, ~12,000 vocabularies



### Outline

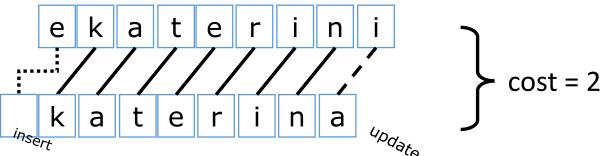
1. Atomic similarity methods



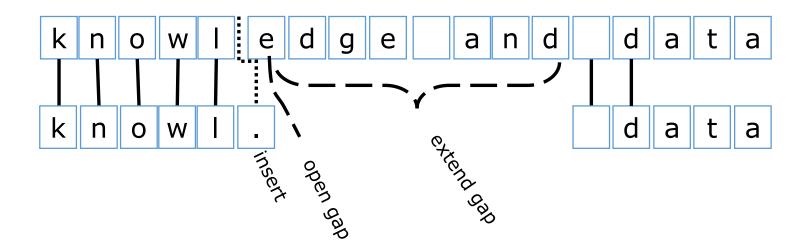
- Examples of targeting cases:
- Publication authors: "John D. Smith" vs. "J. D. Smith"
- Journal names: "Transactions on Knowledge and Data Engineering"
   vs. "Trans. Knowl. Data Eng."

#### **Edit Distance:**

- Number of operations to convert from 1<sup>st</sup> to 2<sup>nd</sup> string
- Operations in Levenstein distance [Lev66]
  - → delete, insert, and update a character with cost 1



- Gap Distance:
- Overcome limitation of edit distance with shortened strings
- Considers two extra operations [Nav01]
  - → open gap, and extend gap (with small cost)



$$cost = 1 + o + 8e$$

- Jaro similarity [Jar89]:
- Small strings, e.g., first and last names

$$JaroSim(S_1, S_2) = \frac{1}{3} \left( \frac{C}{|S_1|} + \frac{C}{|S_2|} + \frac{C - T}{C} \right)$$

 $C \rightarrow \text{common characters in } S_1 \text{ and } S_2$ 

Example: "DEIS"vs. "DESI"

$$C=4, T=2/2,$$

JaroSim= 
$$\frac{1}{3} \left( \frac{4}{4} + \frac{4}{4} + \frac{4-1}{4} \right) = 0.9167$$

- Jaro-Winkler similarity [Win99]:
- Extension that gives higher weight to matching prefix
- Increasing it's applicability to names
- P is a scaling factor (0.1 by default)
- L is the length of the common prefix up to maximum 4
- **Example:** Compute  $J_w(arnab, aranb)$ 
  - $JaroSim(arnab, aranb) = \frac{1}{3} \left( \frac{5}{5} + \frac{5}{5} + \frac{4}{5} \right) = 0.933$

- Soundex:
- Converts each word into a phonetic encoding by assigning the same code to the string parts that sound the same
- Similarity is computed between the corresponding phonetic encodings
- Examples: Meier and Mayer or Smith, Smyth, Smithe, Smeth, Smeeth
- Remarks:
- Surveys: [CRF03], [Win06]
- Existing API with these methods:
  - SecondString: <a href="http://secondstring.sourceforge.net/">http://secondstring.sourceforge.net/</a>
  - SimMetrics: <a href="http://www.dcs.shef.ac.uk/~sam/simmetrics.html">http://www.dcs.shef.ac.uk/~sam/simmetrics.html</a>

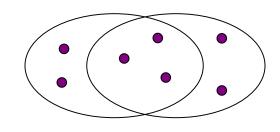


### Outline

- 1. Atomic similarity methods
- 2. Similarity methods for sets



### Similarity methods for sets



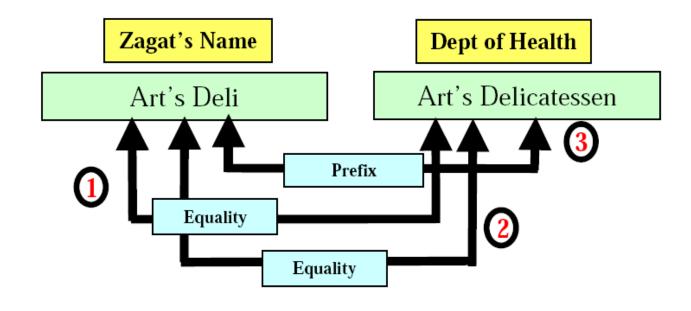
- Jaccard distance/similarity
  - The Jaccard similarity of two sets is the size of their intersection divided by the size of their union:
  - $sim(C_1, C_2) = |C_1 \cap C_2| / |C_1 \cup C_2|$

3 in intersection
7 in union
Jaccard similarity= 3/7
Jaccard distance = 4/7

- Jaccard distance:  $d(C_1, C_2) = 1 |C_1 \cap C_2| / |C_1 \cup C_2|$
- Similarity between {a, b, c, d} and {a, b, e, f} = 2/6 = 1/3
- Jaccard bag similarity counts the repetition of the elements
- The similarity between  $\{a,a,a,b\}$  and  $\{a,a,b,b,c\} = 3/9 = 1/3$

### Similarity methods for sets

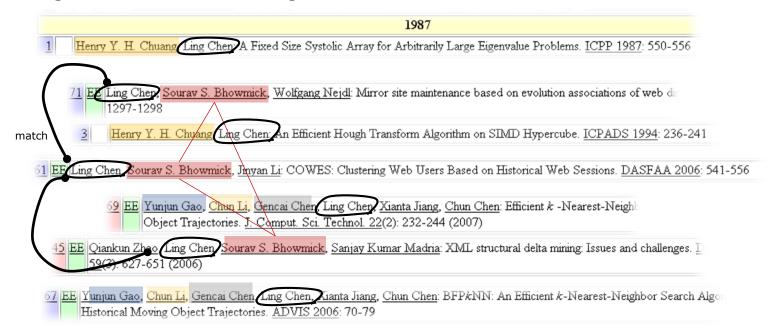
- Using transformations [TKM02]:
- 1. Analyze data to generate transformations
  - Unary transform:
    - Equality, Stemming, Soundex,
       Abbreviation (e.g., 3rd or third)
  - N-ary transformations:
    - Initial, Prefix, Suffix, Substring Acronym, Abbreviation
- 2. Calculate transformation weights
- 3. Apply on candidate mappings





### Similarity methods for sets

- Group Linkage [OKLS07] (Survey [EIV07]):
- Considers groups of relational records
  - o not individual relational records
- Groups match when:
  - 1. High similarity between data of individual records
  - 2. Large fraction of matching records, i.e., no. 1





### Outline

- 1. Motivation: Entity Resolution
- 2. Atomic similarity methods
- 3. Similarity methods for sets
- 4. Facilitating inner-relationships

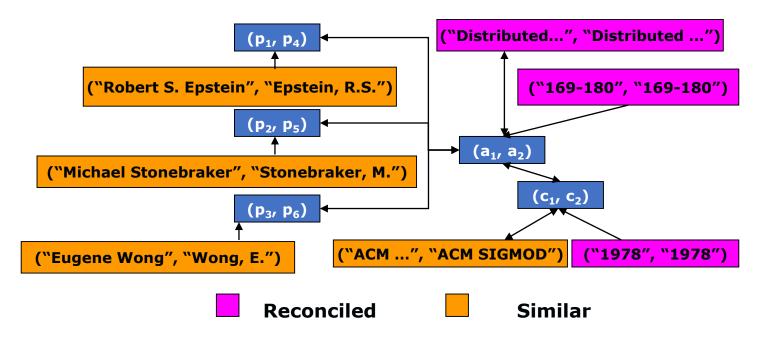


- General idea
- Heterogeneous data
  - Lack of schema information
  - Variations in entity descriptions
  - Incomplete or missing values
- Improve effectiveness by considering data semantics
- Example → Reference Reconciliation



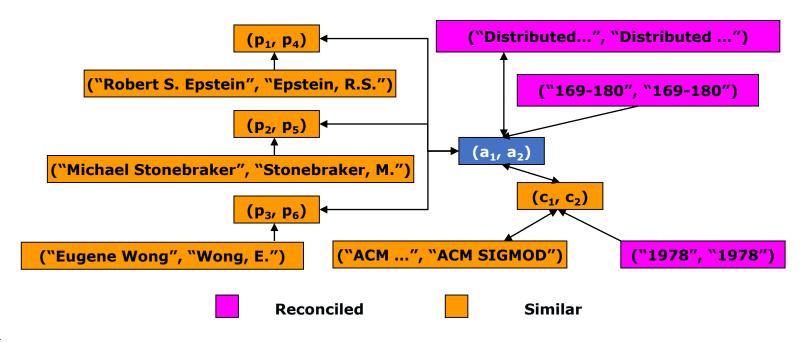
- Reference Reconciliation [DHM05]
- 1. Build a dependency graph

```
a1: ("Distributed...", "169-180, c1:(1978, "ACM..."), "Robert S. Epstein", "Michael Stonebraker", "Eugene Wong") a2: ("Distributed...", "169-180, c2:(1978, "ACM SIGMOD"), "Epstein, R.S.", "Stonebraker, M.", "Wong, E.")
```



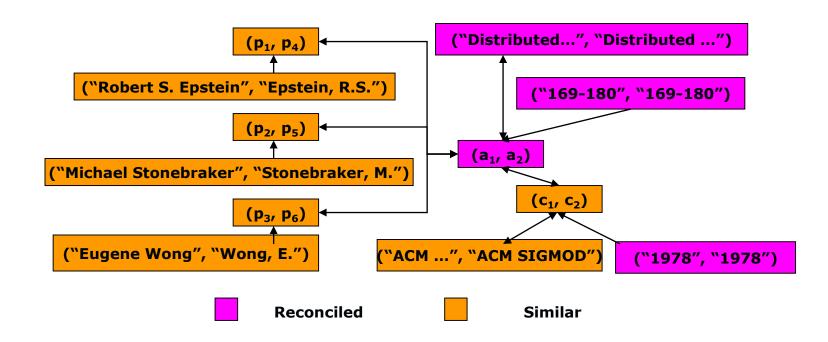


- Reference Reconciliation [DHM05]
- 1. Build a dependency graph
- 2. Exploit information and relationships



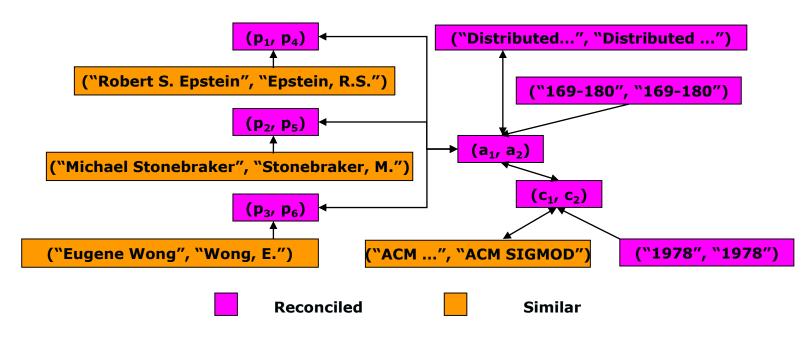


- Reference Reconciliation [DHM05]
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- Reference Reconciliation [DHM05]
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### Outline

- 1. Motivation: Entity Resolution
- 2. Atomic similarity methods
- 3. Similarity methods for sets
- 4. Facilitating inner-relationships
- 5. Methods in uncertain data



### Methods in uncertain data

- General idea:
- Keep conflicting relations, e.g., [AFM06], [RDS07], [DS07a], [DHY07]
  - Lack of resolution rules to correctly resolve and merge relations
  - No merging, but maintain results in the database
  - Relations are alternative representations of the same real world object
- Entity representation with probability indicates...
  - Reliability of the source
  - Output of the matching process
  - o Etc.

CU	sto	m	er
-u		,,,,,	

	<u>custId</u>	name	income	prob		
$s_1$	c1	John	\$120K	0.9		
$s_2$	c1	John	\$80K	0.1		
$s_3$	c2	Mary	\$140K	0.4		
$s_4$	c2	Marion	\$40K	0.6		



### Methods in uncertain data

 Entity-Aware querying over prob. linkages (different perspectives)

[loannou & Velegrakis 2010]

- Not merging the entities using threshold
- Keep probabilistic linkages alongside the original data
- Use them during query processing
- The idea of possible Worlds
- Query:
  - o "J. K. Rowling" movies in "2002"
- Assume no linkages:
  - o zero results
- Possible answer with linkages:
  - o merge(e1, e2)
  - o merge(e<sub>1</sub>, e<sub>2</sub>, e<sub>3</sub>)

```
title: Harry Potter and the Chamber of Secrets
        starring: Daniel Radcliffe
        starring: Emma Watson
                 J.K. Rowling
        writer:
         genre:
                  Fantasy
        title: Harry Potter and the Chamber of Secrets
                  2002
         date:
        starring: Daniel Radcliffe
         starring: Emma Watson
        title: Harry Potter and the Chamber of Secrets
        genre: Fantasy
         author: J.K. Rowling
        codename: The Big Blue
                    California
        location:
0.8
        name: International Business Machines
               New York
               2002
```



### **Similarity Methods for Sets**

The case of documents

#### **A Common Metaphor**

- Many problems can be expressed as finding "similar" sets
  - Find near-neighbors in <a href="high-dimensional">high-dimensional</a> space
- Examples:
  - Pages with similar words
    - For duplicate detection, classification by topic, plagiarism
  - Customers who purchased similar products (e.g. Movies)
  - Products with similar customer sets (e.g. fans)
  - Images with similar features
    - Users who visited similar websites

#### Documents as High-Dim Data

- Converting documents to set
  - Simple approaches:
    - Document = set of words appearing in document
    - Document = set of "important" words
- Need to account for ordering of words!
- A different way: Shingles!

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38

#### **Define: Shingles**

- A *k*-shingle (or *k*-gram) for a document is a sequence of *k* tokens that appears in the doc
  - Tokens can be characters, words or something else, depending on the application
  - Assume tokens = characters for examples
- Example: k=2; document  $D_1$  = abcab Set of 2-shingles:  $S(D_1)$  = {ab, bc, ca}
  - Option: Shingles as a bag (multiset), count ab twice: S'(D<sub>1</sub>)
     = {ab, bc, ca, ab}

# **Compressing Shingles**

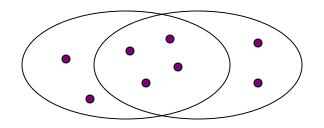
- Represent a document by the set of hash values of its kshingles
  - Idea: Two documents could (rarely) appear to have shingles in common, when in fact only the hash-values were shared
- Example: k=2; document  $D_1$ = abcab Set of 2-shingles:  $S(D_1)$  = {ab, bc, ca} Hash the shingles:  $h(D_1)$  = {1, 5, 7}

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# **Similarity Metric for Shingles**

- Represent document D<sub>1</sub> as a set of its k-shingles C<sub>1</sub>=S(D<sub>1</sub>)
- Equivalently, each document is a 0/1 vector in the space of k-shingles
  - Each unique shingle is a dimension
  - Vectors are very sparse
- A natural similarity measure is the Jaccard similarity:

$$sim(D_1, D_2) = |C_1 \cap C_2|/|C_1 \cup C_2|$$



# **Working Assumption**

- Documents that have lots of shingles in common have similar text, even if the text appears in different order
- Caveat: You must pick k large enough, or most documents will have most shingles
  - k = 5 is OK for short documents
  - k = 10 is better for long documents

# Challenges for ER

- Variety Semantic
  - Semi-structured data → unprecedented levels of heterogeneity
  - Numerous entity types & vocabularies
  - LOD Cloud\*: ~50,000 predicates, ~12,000 vocabularies
- Volume Performance
  - Millions of entities
  - Billions of name-value pairs describing them
  - LOD Cloud\*:  $>5,5\cdot10^7$  entities,  $\sim1,5\cdot10^{11}$  triples
  - Too many documents, Too few memory



#### **Motivation**

- Suppose we need to find near-duplicate documents among N=1 million documents
- Naïvely, we would have to compute pairwise
   Jaccard similarities for every pair of docs
  - $N(N-1)/2 \approx 5*10^{11}$  comparisons
  - At 10<sup>5</sup> secs/day and 10<sup>6</sup> comparisons/sec, it would take **5 days**
- For N = 10 million, it takes more than a year...

#### Find pairs of similar docs

Main idea: Candidates

Instead of keeping a count of each pair, only keep a count of candidate pairs!

- -- Pass 1: Take documents and hash them to buckets such that documents that are similar hash to the same bucket
- -- Pass 2: Only compare documents that are candidates (i.e., they hashed to a same bucket)

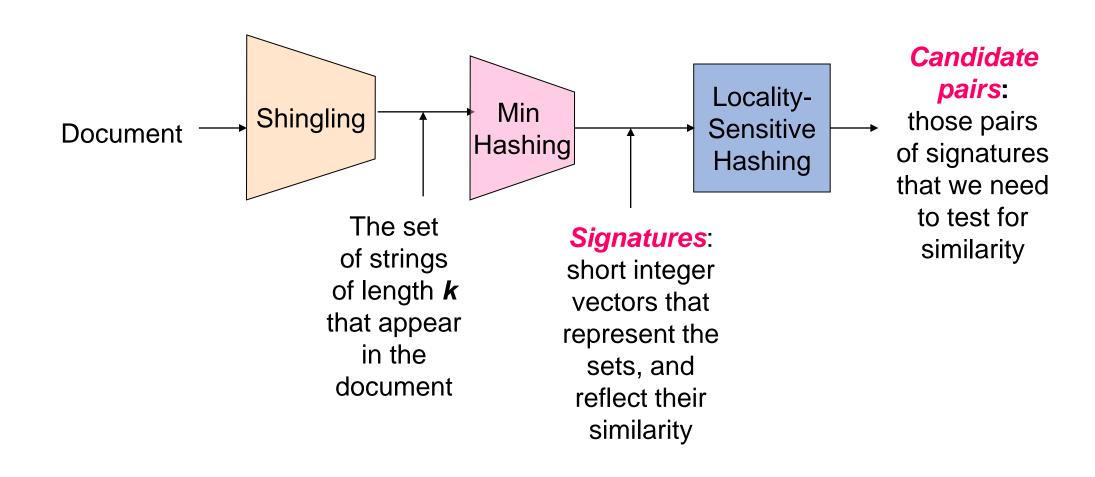
Benefits: Instead of O(N<sup>2</sup>) comparisons, we need O(N) comparisons to find similar documents



#### 3 Essential Steps for Similar Docs

- 1. Shingling: Convert documents to sets
- Min-Hashing: Convert large sets to short signatures, while preserving similarity
- 3. Locality-Sensitive Hashing: Focus on pairs of signatures likely to be from similar documents
  - Candidate pairs!

# The Big Picture



#### **Distance Measures**

- Goal: Find near-neighbors in high-dim. space
  - We formally define "near neighbors" as points that are a "small distance" apart
- For each application, we first need to define what "distance" means

#### From Sets to Boolean Matrices

- Rows = elements (shingles)
- Columns = sets (documents)
  - 1 in row **e** and column **s** if and only if **e** is a member of **s**
  - Column similarity is the Jaccard similarity of the corresponding sets (rows with value 1)
  - Typical matrix is sparse!
- Each document is a column:
  - Example:  $sim(C_1, C_2) = ?$ 
    - Size of intersection = 3; size of union = 6,
       Jaccard similarity (not distance) = 3/6
    - $d(C_1,C_2) = 1 (Jaccard similarity) = 3/6$

**Documents** 

Shingles	1	1	1	0
	1	1	0	1
	0	1	0	1
	0	0	0	1
	1	0	0	1
	1	1	1	0
	1	0	1	0

# **Finding Similar Columns**

- So far:
  - Documents → Sets of shingles
  - Represent sets as Boolean vectors in a matrix
- Next goal: Find similar columns while computing small signatures
  - Similarity of columns == similarity of signatures

# Hashing Columns (Signatures)

- Key idea: "hash" each column C to a small signature h(C), such that:
  - (1) h(C) is small enough that the signature fits in RAM
  - (2)  $sim(C_1, C_2)$  is the same as the "similarity" of signatures  $h(C_1)$  and  $h(C_2)$
- Goal: Find a hash function  $h(\cdot)$  such that:
  - If  $sim(C_1, C_2)$  is high, then with high prob.  $h(C_1) = h(C_2)$
  - If  $sim(C_1, C_2)$  is low, then with high prob.  $h(C_1) \neq h(C_2)$
- Hash docs into buckets. Expect that "most" pairs of near duplicate docs hash into the same bucket!
- Clearly, the hash function depends on the similarity metric:
  - Not all similarity metrics have a suitable hash function
- There is a suitable hash function for the Jaccard similarity: It is called Min-Hashing

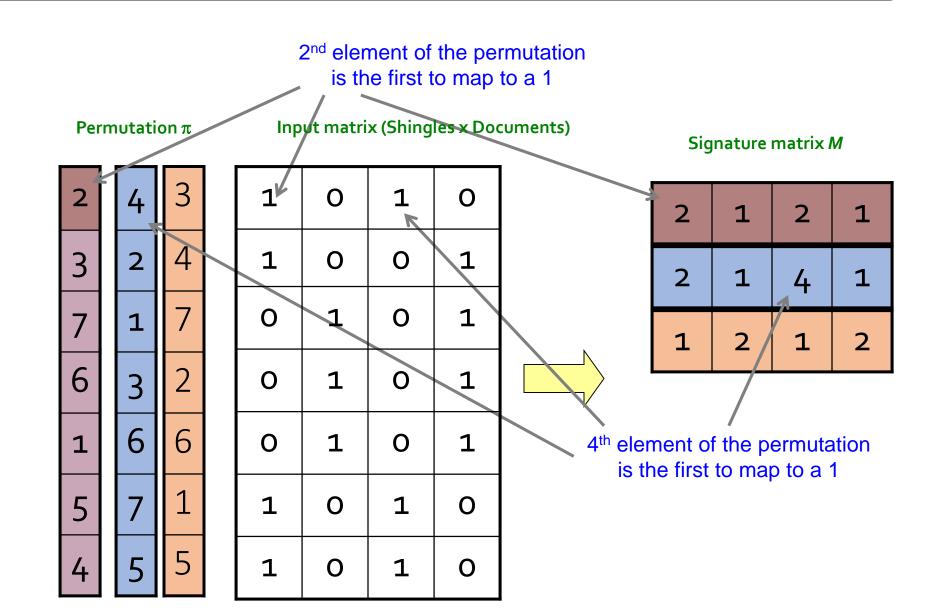
# Min-Hashing

- Imagine the rows of the Boolean matrix permuted under random permutation  $\pi$
- Define a "hash" function  $h_{\pi}(C)$  = the index of the first (in the permuted order  $\pi$ ) row in which column C has value 1:

$$h_{\pi}(\mathbf{C}) = \min_{\pi} \pi(\mathbf{C})$$

• Use several (e.g., 100) independent hash functions (that is, permutations) to create a signature of a column

# Min-Hashing Example



# Similarity for Signatures

- We know:  $Pr[h_{\pi}(C_1) = h_{\pi}(C_2)] = sim(C_1, C_2)$
- Now generalize to multiple hash functions
- The *similarity of two signatures* is the fraction of the hash functions in which they agree
- Note: Because of the Min-Hash property, the similarity of columns is the same as the expected similarity of their signatures

# Min-Hashing Example

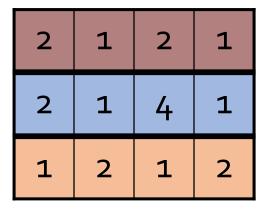
#### Permutation $\pi$

#### Input matrix (Shingles x Documents)

#### Signature matrix M

2	4	3
3	2	4
7	1	7
6	3	2
1	6	6
5	7	1
4	5	5

1	0	1	0
1 0		0	1
0	1	0	1
0	1	0	1
0	1	0	1
1 0		1	0
1	0	1	0

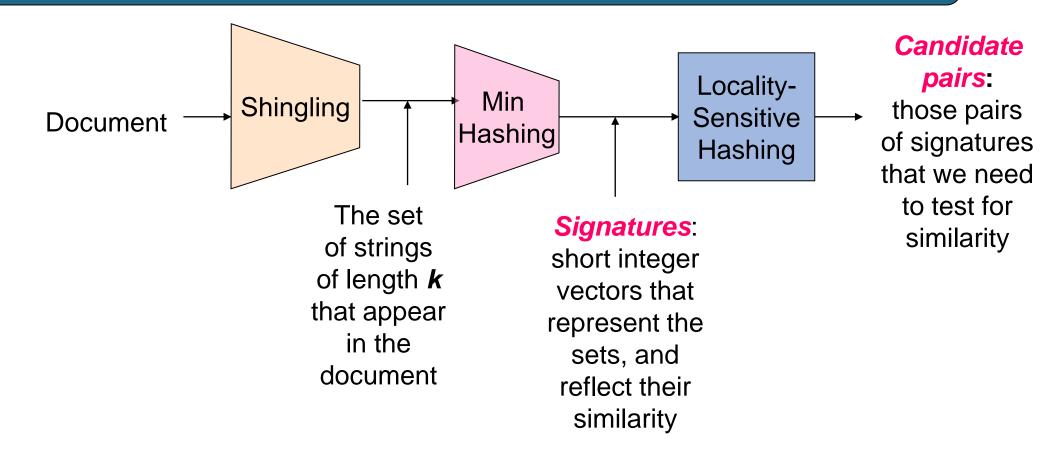


#### **Similarities:**

Col/Col 0.7
Sig/Sig 0.6

	1-3	2-4	1-2	3-4
	0.75			
ig	0.67	1.00	0	0

# Locality-Sensitive Hashing



Step 3: Locality-Sensitive Hashing:
 Focus on pairs of signatures likely to be from similar documents

#### LSH for Min-Hash

- Big idea: Hash columns of signature matrix M several times
- Arrange that (only) similar columns are likely to hash to the same bucket, with high probability
- Candidate pairs are those that hash to the same bucket

• (Blocking)

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#### **Performance**

- Merge-purge [HS95],[HS98]:
- Idea: same entities will share information
- Create a key for each relation (e.g., email)
- Sort relations according to key
- Compare only a limited set of records in each iteration

# **Standard Blocking**

Earliest, simplest form of blocking.

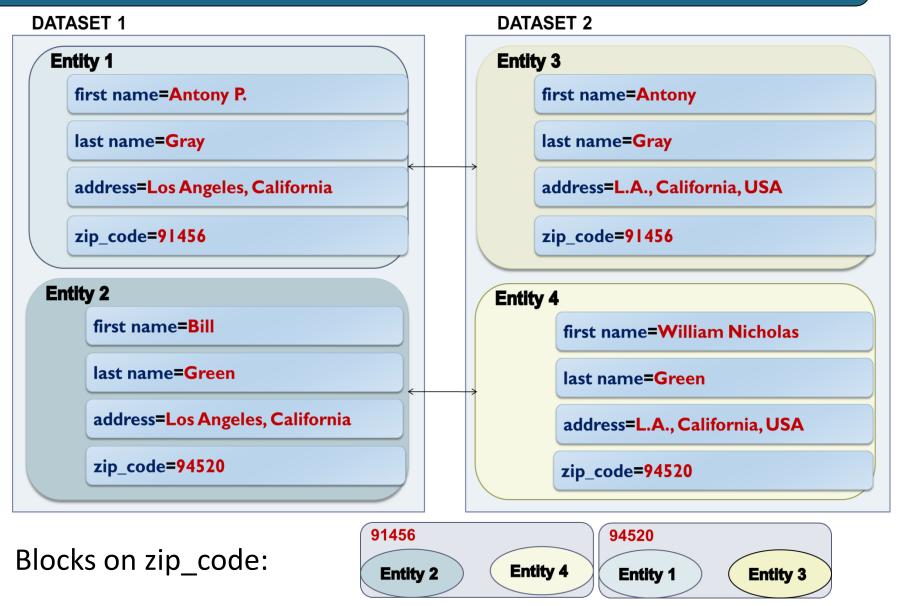
[Fellegi et. al., JASS 1969]

#### Algorithm:

- 1. Select the most appropriate attribute name(s) w.r.t. noise and distinctiveness.
- 2. Transform the corresponding value(s) into a Blocking Key (BK)
- For each BK, create one block that contains all entities having this BK in their transformation.

Works as a hash function! → Blocks on the **equality** of BKs

# **Example of Standard Blocking**





# Thank you for your attention! Questions?

Disclaimer: Much of the material presented originates from a number of different presentations and courses of the following people: Yannis Velegrakis (Utrecht University), Jeff Ullman (Stanford University), Bill Howe (U of Washington), Martin Fouler (Thought Works), Ekaterini Ioannou (Tilburg University), Themis Palpanas (U of Paris-Descartes). Copyright stays with the authors. No distribution is allowed without prior permission by the authors.

