

# Problem 1: Scheduling TV Commercials to Maximize Impact

## Source Code:

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```
import matplotlib.pyplot as plt
import random
import time

# Greedy Job Sequencing
def job_sequencing(ads):
    # Sort ads by profit in descending order
    ads.sort(key=lambda x: x[2], reverse=True)
    max_deadline = max(ad[1] for ad in ads)
    slots = [None] * (max_deadline + 1)
    total_profit = 0
    selected = []

    for ad in ads:
        ad_id, deadline, profit = ad
        for d in range(min(max_deadline, deadline), 0, -1):
            if slots[d] is None:
                slots[d] = ad
                total_profit += profit
                selected.append((d, ad))
                break

    return selected, total_profit

# Example ads
ads = [("A1", 2, 100), ("A2", 1, 19), ("A3", 2, 27), ("A4", 1, 25), ("A5", 3, 15)]
selected, total = job_sequencing(ads)

print("Selected Ads (Slot, Ad):", selected)
print("Total Revenue:", total)

# Visualization: number of ads vs revenue
sizes = [10, 50, 100, 200]
revenues = []

for n in sizes:
    ads = [(f"Ad{i}", random.randint(1, 10), random.randint(10, 1000)) for i in range(n)]
    _, rev = job_sequencing(ads)
    revenues.append(rev)

plt.plot(sizes, revenues, marker='o')
plt.title("Number of Ads vs Revenue (Greedy Algorithm)")
plt.xlabel("Number of Ads")
plt.ylabel("Total Revenue")
plt.grid(True)
plt.show()
```

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## 1. Input

List of advertisements in the form:

(Ad ID, Deadline, Profit)

Example:

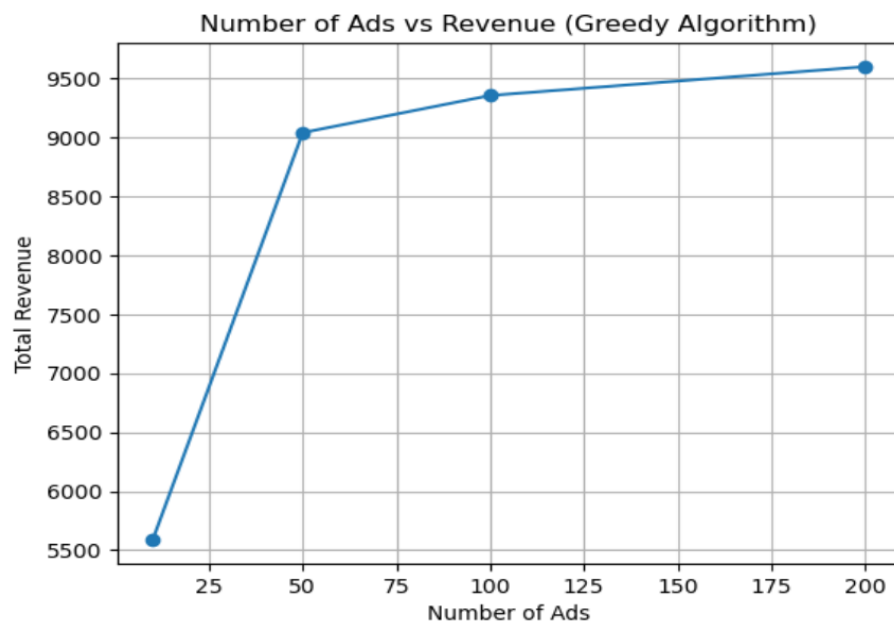
[('A1', 2, 100), ('A2', 1, 19), ('A3', 2, 27), ('A4', 1, 25), ('A5', 3, 15)]

## 2. Approach

- Sort all ads by descending profit.
- For each ad, find the latest available slot before its deadline.
- Assign it if the slot is free (Greedy choice).
- Continue until all slots are filled or all ads checked.

## 3. Output

Selected Ads (Slot, Ad): [(2, ('A1', 2, 100)), (1, ('A3', 2, 27)), (3, ('A5', 3, 15))]  
Total Revenue: 142



#### 4. Analysis

- **Time Complexity:**  $O(n \log n + n^2) \rightarrow$  due to sorting and slot checking
- **Space Complexity:**  $O(n)$  for slot tracking
- **Real-World Constraints:**
  - Limited ad slots during peak viewing hours
  - Different durations and audience reach per slot

#### 5. Visualization

A bar/line chart showing **Number of Ads vs. Revenue Generated**.

As ads increase, total profit rises but plateaus when all slots fill.

## Problem 2: Maximizing Profit with Limited Budget

### Source Code:

```
import matplotlib.pyplot as plt
import random

def knapsack(weights, values, capacity):
    n = len(weights)
    dp = [[0]*(capacity+1) for _ in range(n+1)]

    for i in range(1, n+1):
        for w in range(1, capacity+1):
            if weights[i-1] <= w:
                dp[i][w] = max(values[i-1] + dp[i-1][w-weights[i-1]], dp[i-1][w])
            else:
                dp[i][w] = dp[i-1][w]
    return dp[n][capacity]

# Example
weights = [10, 20, 30]
values = [60, 100, 120]
capacity = 50
print("Maximum Profit:", knapsack(weights, values, capacity))

# Visualization: number of items vs profit
items = [10, 20, 30, 40]
profits = []
for n in items:
    w = [random.randint(1, 20) for _ in range(n)]
    v = [random.randint(10, 100) for _ in range(n)]
    c = sum(w)//2
    profits.append(knapsack(w, v, c))

plt.plot(items, profits, marker='o')
plt.title("Number of Items vs Profit (0/1 Knapsack)")
plt.xlabel("Number of Items")
plt.ylabel("Maximum Profit")
plt.grid(True)
plt.show()
```

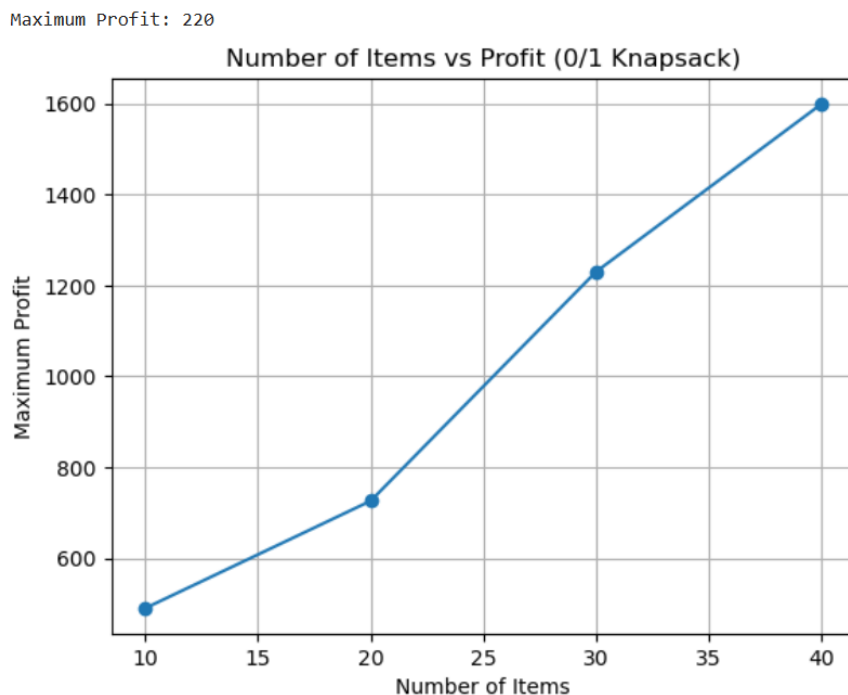
## 1. Input

- Weights = [10, 20, 30] → costs
- Profits = [60, 100, 120]
- Budget = 50

## 2. Approach

- Use bottom-up **0/1 Knapsack DP**.
- Create a table where each cell  $dp[i][j]$  = max profit using first  $i$  items with budget  $j$ .
- Either include or exclude each item based on remaining budget.

## 3. Output



## 4. Analysis

- **Time Complexity:**  $O(n \times W)$ , where  $W$  is budget
- **Space Complexity:**  $O(n \times W)$  or  $O(W)$  if optimized

- **Real-World Constraints:**
  - Limited capital
  - Each project chosen once (no fraction)

## **5. Visualization**

Plot **Budget vs. Profit** showing profit growth until budget limit is reached.

## Problem 3: Solving Sudoku Puzzle

### Source Code:

```
import time
import matplotlib.pyplot as plt

# Sudoku Solver using Backtracking
N = 9

def print_grid(grid):
    for row in grid:
        print(row)

def is_safe(grid, row, col, num):
    # Check row
    for x in range(N):
        if grid[row][x] == num:
            return False
    # Check column
    for x in range(N):
        if grid[x][col] == num:
            return False
    # Check 3x3 box
    start_row, start_col = row - row % 3, col - col % 3
    for i in range(3):
        for j in range(3):
            if grid[i + start_row][j + start_col] == num:
                return False
    return True
```

```
def solve_sudoku(grid):
    for row in range(N):
        for col in range(N):
            if grid[row][col] == 0:
                for num in range(1, 10):
                    if is_safe(grid, row, col, num):
                        grid[row][col] = num
                        if solve_sudoku(grid):
                            return True
                        grid[row][col] = 0
                return False
    return True
```

*# Example Sudoku*

```
grid = [
    [3, 0, 6, 5, 0, 8, 4, 0, 0],
    [5, 2, 0, 0, 0, 0, 0, 0, 0],
    [0, 8, 7, 0, 0, 0, 0, 3, 1],
    [0, 0, 3, 0, 1, 0, 0, 8, 0],
    [9, 0, 0, 8, 6, 3, 0, 0, 5],
    [0, 5, 0, 0, 9, 0, 6, 0, 0],
    [1, 3, 0, 0, 0, 0, 2, 5, 0],
    [0, 0, 0, 0, 0, 0, 0, 7, 4],
    [0, 0, 5, 2, 0, 6, 3, 0, 0]
]
```

```
print("Original Sudoku Grid:")
print_grid(grid)
```

```
start = time.time()
solve_sudoku(grid)
end = time.time()
```

```
print("\nSolved Sudoku Grid:")
print_grid(grid)
```

```
print(f"\nTime taken: {end - start:.5f} seconds")
```

*# Visualization*

```
blanks = [5, 10, 15, 20, 25, 30]
times = []
```

```
for b in blanks:
    # make a grid with 'b' blanks
    test_grid = [row[:] for row in grid]
    c = 0
    for i in range(9):
        for j in range(9):
            if c < b:
                test_grid[i][j] = 0
                c += 1
    t1 = time.time()
    solve_sudoku(test_grid)
    t2 = time.time()
    times.append(t2 - t1)
```

```
plt.figure()
plt.plot(blanks, times, marker='o')
plt.title("Sudoku: Time vs Number of Blanks")
plt.xlabel("Number of Blank Cells")
plt.ylabel("Time (seconds)")
plt.grid(True)
plt.show()
```



## 1. Input

A partially filled Sudoku grid (0 = empty).

## 2. Approach

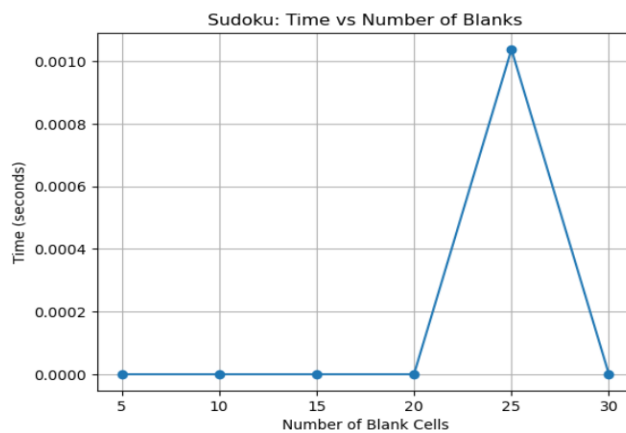
- Find the first empty cell.
- Try digits 1–9 that don't violate Sudoku rules.
- Place a number and move to the next empty cell.
- If no valid number → **backtrack**.
- Repeat until solved.

## 3. Output

Original Sudoku Grid:  
[3, 0, 6, 5, 0, 8, 4, 0, 0]  
[5, 2, 0, 0, 0, 0, 0, 0, 0]  
[0, 8, 7, 0, 0, 0, 0, 3, 1]  
[0, 0, 3, 0, 1, 0, 0, 8, 0]  
[9, 0, 0, 8, 6, 3, 0, 0, 5]  
[0, 5, 0, 0, 9, 0, 6, 0, 0]  
[1, 3, 0, 0, 0, 0, 2, 5, 0]  
[0, 0, 0, 0, 0, 0, 0, 7, 4]  
[0, 0, 5, 2, 0, 6, 3, 0, 0]

Solved Sudoku Grid:  
[3, 1, 6, 5, 7, 8, 4, 9, 2]  
[5, 2, 9, 1, 3, 4, 7, 6, 8]  
[4, 8, 7, 6, 2, 9, 5, 3, 1]  
[2, 6, 3, 4, 1, 5, 9, 8, 7]  
[9, 7, 4, 8, 6, 3, 1, 2, 5]  
[8, 5, 1, 7, 9, 2, 6, 4, 3]  
[1, 3, 8, 9, 4, 7, 2, 5, 6]  
[6, 9, 2, 3, 5, 1, 8, 7, 4]  
[7, 4, 5, 2, 8, 6, 3, 1, 9]

Time taken: 0.01402 seconds



#### 4. Analysis

- **Time Complexity:**  $O(9^n)$  (worst case exponential)
- **Space Complexity:**  $O(n^2)$  (for grid and recursion stack)
- **Performance Impact:**
  - Hard puzzles take longer due to backtracking depth.

#### 5. Visualization

Optional graph of **Time Taken vs. Number of Empty Cells**, showing exponential growth.

## Problem 4: Password Cracking (Brute-Force)

### Source Code:

```
import itertools
import time
import matplotlib.pyplot as plt

# Brute Force Password Cracker
def brute_force_password(target, charset):
    attempts = 0
    for length in range(1, len(target) + 1):
        for guess in itertools.product(charset, repeat=length):
            attempts += 1
            if ''.join(guess) == target:
                return ''.join(guess), attempts
    return None, attempts

# Example input
target = "ab1"
charset = "abc123"

start = time.time()
found, attempts = brute_force_password(target, charset)
end = time.time()

print(f"Target Password: {target}")
print(f"Found Password: {found}")
print(f"Attempts: {attempts}")
print(f"Time Taken: {end - start:.5f} seconds")
```

```
# Visualization
lengths = [1, 2, 3, 4]
times = []

for L in lengths:
    target = "a" * L # simple password
    start = time.time()
    brute_force_password(target, charset)
    end = time.time()
    times.append(end - start)

plt.figure()
plt.plot(lengths, times, marker='o')
plt.title("Password Cracking: Time vs Password Length")
plt.xlabel("Password Length")
plt.ylabel("Time (seconds)")
plt.grid(True)
plt.show()
```

## 1. Input

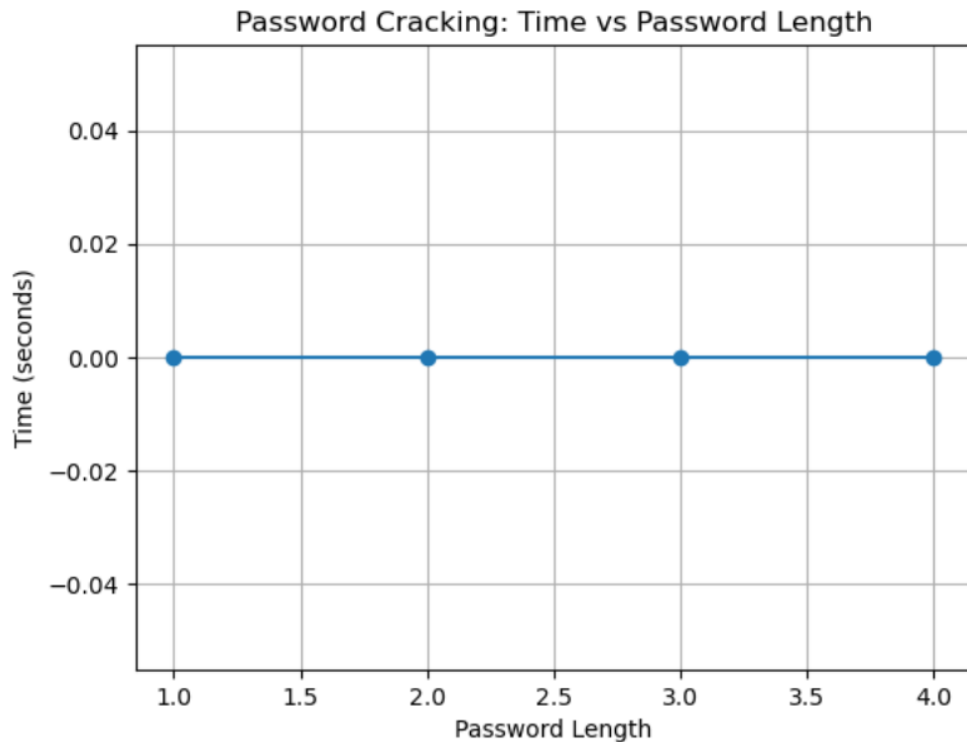
- Target password (e.g., "1a")
- Character set (e.g., ['a','b','c','1','2','3'])

## 2. Approach

- Use `itertools.product(charset, repeat=len(password))` to generate all possible combinations.
- Compare each generated string with target until a match is found.

## 3. Output

```
Target Password: ab1
Found Password: ab1
Attempts: 52
Time Taken: 0.00100 seconds
```



## 4. Analysis

- **Time Complexity:**  $O(C^L)$  where  $C$  = charset size,  $L$  = password length
- **Space Complexity:**  $O(1)$  (only one combination checked at a time)

- **Practical Impact:**

- Very slow for long passwords.
- Demonstrates importance of strong password policies.

## **5. Visualization**

Plot **Password Length vs. Time Taken** showing exponential rise in cracking time.