Computer Security Capstone

Project 4: Capture The Flag (CTF)

Chi-Yu Li (2024 Spring)
Computer Science Department
National Yang Ming Chiao Tung University

Goal

 Understand the exploitation of basic programming bugs, Linux system knowledge, and reverse-engineering

- You will learn about
 - □ Solving basic CTF problems
 - □ Investigating C/Linux functions deeply instead of simply using them
 - □ What buggy codes are and how they can be exploited

What is CTF?

- A traditional outdoor game
 - ☐ Two teams each have a flag
 - □ Objective: to capture the other team's flag



From Wikipedia

- In computer security, it is a type of cryptosport: a computer security competition
 - ☐ Giving participants experience in securing a machine
 - Required skills: reverse-engineering, network sniffing, protocol analysis, system administration, programming, etc.
 - □ How?
 - A set of challenges is given to competitors
 - Each challenge is designed to give a "Flag" when it is countered

A CTF Example

A toy CTF

\$ python -c 'v = input(); print("flag:foobar") if v == "1" else print("failed")'

- ☐ You should enter "1" to pass the *if* statement and get the flag (flag:foobar)
- □ Otherwise, "failed" is obtained

Requirements

- Linux/Unix environment is required
 - ☐ Connecting to our CTF servers for all the tasks

- You are NOT allowed to team up: one student one team
 - □ Discussions are allowed between teams, but any collaboration is prohibited
- TA: Cheng-I Hu

How to Proceed?

- Connecting to each CTF server: nc <ip> <port>
 - □ ip: 140.113.207.245
 - □ port is given at each problem
 - ☐ The program of each problem runs as a service at the server
 - ☐ You can do whatever you are allowed to do
- You can use python with pwntools, too

How to Proceed? (Cont.)

- For each CTF problem, you should
 - □ analyze its given executable files or source code files
 - □ interact with the server to get a flag
 - ☐ The flag format: CSC2025{[a-zA-Z0-9_]+}
- You will need to submit the programs
 - ☐ run the programs when you demo

What If Get Stuck?

- Learn to use "man" in UNIX-like systems
 - ☐ If you don't know something, ask "man"
 - □ e.g., what is man?
 - \$ man man
- Learn to find answers with FIRST-HAND INFORMATION/REFERENCE
 - □ Google is your best friend (Using ENGLISH KEYWORDS!!)
 - ☐ First-hand information: Wikipedia, cppreference.com, devel mailing-list, etc.
 - ☐ First-hand reference: papers, standards, spec, man, source codes, etc.
 - □ Second-hand information: blog, medium, ptt, reddit, stackoverflow post, etc.

Two Tasks

- Task I: Basic CTF problems (70%)
- Task II: Advance CTF problems (30%)
- Download all given executable and source files from e3
 - □ CTF Server using ubuntu 24.04 (for some problem to calculate address)

Task I: Basic CTF Problems

■ Task I-1: Password Checker (20%)

■ Task I-2: Secure Random (20%)

● Task I-3: Simple Shell (15%)

■ Task I-4: Simple ROP (15%)

Task I-1: Password Checker

- Goal: Learn how type conversion works in C/C++
- Server port: 30170

- Hints
 - □ <u>Implicit conversions of type</u>

Task I-2: Secure Random

- Goal: learn about the glibc PRNG
- Server port: 30171

- Hints
 - □ Is the random function really random?
 - Make sure you have time synchronization in your environment!!

Task I-3: Simple Shell

- Goal: learn to identify basic logic flaw and buffer overflow in source codes
- Server port: 30172

- Hints
 - ☐ Inspect the code, where buffer overflow can occur?
 - What can you modify?
 - □ Inspect the impact of overflow by using gdb
 - You may want to install gdb extensions like gef

Task I-4: Simple ROP

- Goal: Given buffer overflow, try to find a way to open up a shell for remote command execution!!
- Server port: 30173

- Hints
 - ☐ Inspect the code, where buffer overflow can occur?
 - □ With NX enabled, you cannot write shell code for buffer overflow
 - □ Stack buffer overflow
 - Return-oriented programming
 - You may want to use tools like ROPgadget to find gadget for ROP

Task II: Advance CTF Problems

• Task II-1: ret2Flag (10%)

● Task II-2: Simple RTOS (10%)

Task II-3: Hard ROP(10%)

Task II-1: Ret2Flag

- Goal: Learn exploit buffer overflow to control program flow
- Server port: 30174

objdump -d ./

Hints

- □ Inspect the code, where buffer overflow can occur?
- ☐ How can you bypass <u>canary</u> protection?
- How can you find the function address?
 - You may want to find address that related with putFlag
- ☐ Try to leak the information you need!!!

Task II-2: Simple RTOS

- Goal: learn to identify dangerous function usage
- Server port: 30175

- Hints
 - ☐ How do you use printf normally?
 - Which conversion specifier can modify variable?
 - How can you return to the function you want?

Task II-3: Hard ROP

- Goal: Try to ROP with libc gadget!!
- Server port: 30176
- Hints
 - ☐ First, try to leaking every thing you need
 - ☐ Try to use libc gadget

Important: How to Prepare Your Program?

- Must provide a Makefile which compiles your source codes into senven executable file
- You can use any language and library you want
 - ☐ Use your environment to demo
 - □ Do not hardcode the flag in your program
- Test requirements for your program
 - ☐ Do not need user interaction to get flag
 - For online tasks, you can only input server IP and port
 - For local tasks, you can only input file path
 - Must print flag to stdout

Important: How to Demo Your Program?

- Download your code from e3
- Run make if needed
- Run your executables
- Ask some questions about your code
- Binary file for all task will not change
 - You can hardcode some symbol address if you need
 - ☐ FLAG during demo will change to avoid hardcode the flag

Project Submission

- Due date: 5/28 11:55 p.m.
- Submission rules
 - □ Put all your files into a directory and name it using your student ID(s)
 - □ Zip the directory and upload the zip file to New e3
 - ☐ A sample of the zip file: 1234567.zip 1234567
 - Makefile (if needed)

 - L ...
 - (Please have a studentID folder in your zip)
 - ☐ If files are not in a directory after unzip, 10 points will be deducted.

Questions?

Useful Info

- command
 - □ checksec
 - □ readelf
- pwntools
 - □ connect to server and control what content will be sent to it
 - ☐ generate shellcode, attach gdb ...etc.
- gdb
 - □ normal plugins: pwngdb / gef /peda
 - dynamic analysis of the program

```
rbp →
func:
         push rbp
         mov rbp, rsp
                                   Call fun = push next_rip
         sub rsp, 0x30
                                              jmp func
         move eax, 0x0
                                                                   rsp \rightarrow
         leave
         ret
main:
         call func
rip \rightarrow
         mov eax, 0x0 // address 0x4005a0
          ...
```

high address

Stack frame of main

rbp → func: push rbp mov rbp, rsp Call fun = push next_rip sub rsp, 0x30 jmp func move eax, 0x0 leave rsp \rightarrow ret main: call func rip \rightarrow mov eax, 0x0 // address 0x4005a0 ...

high address Stack frame of main 0x4005a0 (return address)

rbp → func: rip 👈 push rbp mov rbp, rsp sub rsp, 0x30 move eax, 0x0 leave rsp \rightarrow ret main: call func mov eax, 0x0 // address 0x4005a0 ...

high address

0x4005a0 (return address)

Stack frame of main

rbp → func: push rbp mov rbp, rsp rip 👈 sub rsp, 0x30 move eax, 0x0 leave ret rsp → main: call func mov eax, 0x0 // address 0x4005a0 ...

high address Stack frame of main 0x4005a0 (return address) old rbp

```
func:
          push rbp
         mov rbp, rsp
         sub rsp, 0x30
rip 👈
          move eax, 0x0
          leave
         ret
                                                           rbp \rightarrow rsp \rightarrow
main:
         call func
          mov eax, 0x0 // address 0x4005a0
          ...
```

high address

Stack frame of main

0x4005a0 (return address)

old rbp

```
func:
        push rbp
        mov rbp, rsp
        sub rsp, 0x30
rip 👈
        move eax, 0x0
        leave
        ret
                                                            rbp →
main:
        call func
        mov eax, 0x0 // address 0x4005a0
                                                            rsp ->
         ...
```

high address Stack frame of main 0x4005a0 (return address) old rbp Local variables of func()

...

Example: Stack frame during a function call

```
func:
        push rbp
                                leave = mov rsp, rbp
        mov rbp, rsp
                                       pop rbp
        sub rsp, 0x30
        move eax, 0x0
rip 👈
        leave
        ret
main:
        call func
        mov eax, 0x0 // address 0x4005a0
```

high address Stack frame of main 0x4005a0 (return address) old rbp Local variables of func()

rbp →

rsp -

...

Example: Stack frame during a function call

```
func:
         push rbp
                                   leave = mov rsp, rbp
         mov rbp, rsp
                                           pop rbp
         sub rsp, 0x30
         move eax, 0x0
rip 👈
         leave
         ret
                                                           rbp \rightarrow rsp \rightarrow
main:
         call func
         mov eax, 0x0 // address 0x4005a0
```

high address Stack frame of main

0x4005a0 (return address)

old rbp

Local variables of func()

```
rbp →
func:
        push rbp
        mov rbp, rsp
        sub rsp, 0x30
                                   ret = pop rip
        move eax, 0x0
        leave
                                                            rsp →
rip 👈
        ret
main:
        call func
        mov eax, 0x0 // address 0x4005a0
         ...
```

high address Stack frame of main 0x4005a0 (return address) old rbp Local variables of func()

rbp → func: push rbp mov rbp, rsp sub rsp, 0x30 move eax, 0x0 rsp → leave ret main: call func mov eax, 0x0 // address 0x4005a0 rip 💙 ...

high address

Stack frame of main

0x4005a0 (return address)

old rbp

Local variables of func()

Common Security Protection in Binary

- Canary
 - Put canary value before old rbp and return address

- PIE/ALSR
 - ☐ Randomize the address space of a process
 - The offset between different symbol still the same!!

Common Security Protection in Binary

- Relro
 - □ Lazy binding option for program
 - Full Relro will have GOT table read only before calling main function
 - Partial Relro make GOT table writable and resolve symbol after calling main function

- NX
 - Making stack not executable.
 - Make shellcode not able to run on stack.

Common Security Protection in Binary

- You may check the protection mechanism in binary using checksec
 - □ slimm609/checksec: Checksec

```
(env) huroy@build-server:~/csc2025-project4/hard_rop$ checksec --file hard rop
  '/home/huroy/csc2025-project4/hard_rop/hard_rop'
   Arch:
               amd64-64-little
               Full RELRO
   RELRO:
   Stack:
              Canary found
              NX enabled
   NX:
              PIE enabled
  PIE:
               Enabled
   SHSTK:
               Enabled
   IBT:
   Stripped:
               No
   Debuginfo: Yes
```

Example: Stack frame with canary

func:

push rbp mov rbp, rsp sub rsp, 0x30 rax,QWORD PTR fs:0x28 QWORD PTR [rbp-0x8],rax



rax,QWORD PTR [rbp-0x8] rax,QWORD PTR fs:0x28 call <__stack_chk_fail@plt> leave

ret

rbp - 8 → **→**

rbp

rsp

high address

Stack frame of main

0x4005a0 (return address)

old rbp

Canary, first bytes is \x00

Local variables of func()