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1. χ^2 Test for uniform distribution

We briefly recall the property and usage of the χ^2 test for uniform distributions over a finite set S. Suppose M samples $y_1, \dots, y_M \in S$. We partition S into r subsets

$$S = \sqcup_{i=1}^r S_i,$$

For each $1 \leq j \leq r$, we compute the expected number of samples that would fall in the i-th subset: $c_j := |S_j|M/|S|$. Then we compute the actual number of samples, $t_j := |\{1 \leq i \leq r : y_i \in S_j\}|$. Finally, the χ^2 value is computed as

$$\chi^{2}(S, y) = \sum_{j=1}^{r} \frac{(t_{j} - c_{j})^{2}}{c_{j}}.$$

Note that degree of freedom in this test is d=r-1. To decide whether the samples are from a uniform distribution, we can either look up a table of χ^2 values, or use an approximation rule: when df is large, the χ^2 distribution can be well-approximated by a normal distribution N(d,2d); for example, if it turns out that $\chi^2 \notin (d-c\sqrt{2d},d+c\sqrt{2d})$, then the confidence we have that the samples are not taken from a uniform distribution is $2\Phi(c)-1$.

2. Attack

Now we describe our attack. It relies on the assumption that there exists a prime ideal \mathfrak{q} such that error $e \pmod{\mathfrak{q}}$ is not uniformly distributed in the finite field $F = R/\mathfrak{q}R$. The attack loop through all q^f possibilities of $\bar{s} = s \pmod{\mathfrak{q}}$. For each guess s', it computes the values $\bar{e}' = \bar{b} - \bar{a}s' \pmod{\mathfrak{q}}$. If the guess is wrong, or if the samples are taken from the uniform distribution instead of an RLWE instance, then \bar{e}' is uniformly distributed in F and it would pass the χ^2 test. On the other hand, if the guess is correct, then we expect the errors \bar{e}' to fail the χ^2 test.

Let $N=q^f$, the number of χ^2 tests we run in the attack. For the attack to be successful, we need the (N-1) tests corresponding to wrong guess of $s\pmod{\mathfrak{q}}$ to pass, and the one test corresponding to the correct guess to fail. Therefore, we need to choose the confidence interval of our χ^2 test so that it is unlikely for a set of samples coming from uniform distribution to fail the test. In practice, we choose the confidence level to be $\alpha=1-\frac{1}{10N}$. Let β denote the probability that the sample errors fails the uniform test with probability Then the probability that our algorithm will success is $p=(1-\frac{1}{10N})^{N-1}\beta$. Note that when N is large, $(1-\frac{1}{10N})^{N-1}$ is about $e^{-1/10}\approx 0.904$.

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Algorithm 1 χ^2 -test attack of RLWE(R, \mathfrak{q})

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Require: R = (K, q, \sigma) – an RLWE instance.
          R = \mathcal{O}_K – the ring of integers of K.
          n: the degree of K.
          \mathfrak{q}: a prime ideal in K above q.
          N: the cardinality of the finite field R/\mathfrak{q}.
          S: a collection of M (M = \Omega(N)) RLWE samples (a, b) \sim R.
Ensure: a guess of the value s \pmod{\mathfrak{q}}, or NON-RLWE, or INSUFFIICNET-SAMPLES
 1: \alpha \leftarrow 1 - \frac{1}{10N}.
 2: \omega \leftarrow \Phi^{-1}((1+\alpha)/2)
 3: G = \emptyset
 4: for s in F do
         for a, b in S do
 5:
             E \leftarrow \emptyset.
 6:
             \bar{a}, \bar{b} \leftarrow a \pmod{\mathfrak{q}}, b \pmod{\mathfrak{q}}.
 7:
             \bar{e} \leftarrow \bar{b} - \bar{a}s.
 8:
             add e to E.
 9:
10:
         end for
        Run \chi^2 test on E and obtain the value \chi^2(E).
11:
         if |\chi^{2}(E) - B - 1| > \omega \sqrt{2B - 2} then
12:
             add s to G
13:
         end if
14:
15: end for
16: if G = \emptyset then
          return NOT RLWE
17: else if G = \{g\} then
          return g
18: else
          return INSUFFIICNET-SAMPLES
19: end if
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Proof of correctness.

Proposition 2.1. The algorithm has time complexity $O(q^{2f})$. Let Δ be the statistical distance between $D_{R,\sigma} \pmod{\mathfrak{q}}$ and the uniform distribution on R/\mathfrak{q} . Then algorithm outputs with probability at least 0.904Δ (fixme: do the correct constant).