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# Computational aspects of modular parametrizations of elliptic curves

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## Abstract

Computational aspects of modular parametrizations of elliptic curves

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Abstract goes here.

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# ${\bf GLOSSARY}$

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# **DEDICATION**

to all of you

## Chapter 1

# FOURIER EXPANSIONS OF CUSPIDAL MODULAR FORMS FORMS AT CUSPS

Let k be a positive even integer and let  $f \in S_k(\Gamma_0(N))$  be a nonzero cusp form. We are concerned with the problem of computing the Fourier expansion of f at cusps of width one other than the cusp  $[\infty]$ . Note that such cusps exist if and only if N is not square-free. We will give two algorithms, one numerical and the other exact, to compute such expansions. The question is studied in the Ph.D. thesis of Christophe Delaunay. We draw insight from another preprint by F.Brunault. The question is also studied in [Edixhoven], where numerical algorithm is given. The algorithm in [Ed] for computing expansions requires working at a higher level: to compute expansions at cusps of denominator Q, one needs to compute period matrices for forms of level  $NR^2$ , where  $R = \gcd(Q, \frac{N}{Q})$ . As a contrast, our algorithm works at levels dividing N.

The main result of this chapter is Theorem 1.6.7 and Algorithm 3, which computes the Fourier expansion of a newform  $f \in S_k(\Gamma_0(N))$  at any width one cusp z. Along the way, we have developed algorithms to compute the twists  $f \otimes \chi$  and the pseudo-eigenvalue of newforms under the Fricke involution. Section contains some examples. In section, we describe an alternative approach and

#### 1.1 Preliminaries

Let  $N \geq 1$  be an integer and let  $X_0(N)$  be the modular curve of level N.

**Definition 1.1.1.** Let z be a cusp on  $X_0(N)$ . If  $z \neq \infty$ , write z = [a/c] with gcd(a, c) = 1. The *denominator* of z is

$$d_z = \gcd(c, N).$$

. If  $z = \infty$ , we set  $d_{\infty} = N$ . Choose  $\alpha \in SL_2(\mathbb{Z})$  such that  $\alpha(\infty) = z$ . The width of z is

$$h_z = \left| \frac{SL_2(\mathbb{Z})_{\infty}}{(\alpha^{-1} \{ \pm I \} \Gamma_0(N) \alpha)_{\infty}} \right|$$

where the subscript  $\infty$  means taking the isotropy subgroup of  $\infty$  in the corresponding group.

The width of a cusp can be computed in terms of its denominator. In fact, we have

**Lemma 1.1.2.** If z is a cusp on  $X_0(N)$ , then

$$h_z = \frac{N}{\gcd(d_z^2, N)}.$$

*Proof.* When  $z = [\infty]$ , we have  $d_{\infty} = N$  and  $h_{\infty} = 1$ , so the formula holds trivially. Otherwise, write  $z = \begin{bmatrix} a \\ c \end{bmatrix}$  and find  $\alpha = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in SL_2(\mathbb{Z})$ . For  $N' \in \mathbb{Z}$  we compute

$$\alpha \begin{pmatrix} 1 & N' \\ 0 & 1 \end{pmatrix} \alpha^{-1} = \begin{pmatrix} * & * \\ -c^2 N' & * \end{pmatrix}.$$

Hence  $\begin{pmatrix} 1 & N' \\ 0 & 1 \end{pmatrix} \in (\alpha^{-1}\{\pm I\}\Gamma_0(N)\alpha)_{\infty} \iff N \mid c^2N' \iff \frac{N}{\gcd(d_z^2,N)} \mid N'$ . This completes the proof.

In particular, the width of a cusp z is one if and only if  $N \mid d_z^2$ .

Suppose f be a cusp form on  $\Gamma_0(N)$  and  $\alpha \in SL_2(\mathbb{Z})$ . Then  $f[\alpha]$  is a cusp form on  $\Gamma(N)$ . So  $f[\alpha]$  has a q-expansion, which is a power series in  $q^{\frac{1}{N}}$ . A natural thing to do is to define the expansion of f at the cusp z as the expansion of  $f|[\alpha]$ . However, one must note that this may not be well-defined: the expansion depends on the choice of  $\alpha$ . Nonetheless, when the denominator of the cusp is sufficiently divisible by the prime divisors of N, the expansion is well-defined as a power series in q.

**Lemma 1.1.3.** Let z be a cusp on  $X_0(N)$  with width one. Choose  $\alpha \in SL_2(\mathbb{Z})$  such that  $\alpha(\infty) = z$ . Then  $f|[\alpha]$  is a cusp form on  $\Gamma_1(N)$ . Moreover, the function  $f|[\alpha]$  is independent of the choice of  $\alpha$ .

Proof. It is easy to verify that  $\Gamma_1(N) \subseteq \alpha^{-1}\Gamma_0(N)\alpha$ , hence the first claim holds. Now suppose  $\beta \in SL_2(\mathbb{Z})$  is such that  $\beta(\infty) = z$ . Then  $\alpha^{-1}\beta \in SL_2(\mathbb{Z})_{\infty}$ . Since z has width one, we have  $\alpha^{-1}\beta \in \alpha^{-1}\Gamma_0(N)\alpha$ . Hence  $\beta \in \Gamma_0(N)\alpha$ , and it follows that  $f|[\beta] = f|[\alpha]$ .

In light of the lemma above, we define the q-expansion of f at a width one cusp z to be the q-expansion of  $f|[\alpha]$ , and denote it by  $f_z$ .

Assume further that f is an eigenform under the Atkin-Lehner operators. We will show that in order to compute the expansion of  $f|[\alpha]$  for any  $\alpha \in SL_2(\mathbb{Z})$ , it suffices to do so for  $\alpha = \begin{pmatrix} 1 & 0 \\ m & 1 \end{pmatrix}$ , where  $0 \le m < N$  and  $N \mid \gcd(m, N)^2$ . In particular, it suffices to compute the expansions of f at a some cusps of width one.

**Lemma 1.1.4.** For any  $\alpha \in SL_2(\mathbb{Z})$ , there exists a matrix  $w_Q \in W_N$  and an upper triangular matrix  $u \in GL_2(\mathbb{Q})$  such that  $w\alpha = \alpha' u$ , where  $\alpha' = \begin{pmatrix} a' & b \\ c' & d' \end{pmatrix} \in SL_2(\mathbb{Z})$  satisfies  $N \mid \gcd(N, c')^2$ .

Indeed, one may find Q using Lemma. Now  $f|[\alpha] = f|[w_Q][w_Q\alpha] = f|[w_Q][\alpha'][u] = \lambda_Q(f)f[\alpha'][u] = \lambda_Q(f)f[\alpha''][u]$ , where  $\alpha''$  is of form  $\begin{pmatrix} 1 & 0 \\ m & 1 \end{pmatrix}$ . Note that for an upper triangular matrix  $u = \begin{pmatrix} u_0 & u_1 \\ 0 & u_2 \end{pmatrix}$ , we have  $f[u](q) = f(q^{u_0/u_2}e^{2\pi i u_1/u_2})$ .

## 1.2 Reducing to the case of newforms

The space  $S_k(\Gamma_0(N))$  is spanned by elements of form  $g(q^d)$ , where g is newform of level  $M \mid N$  and d is a divisor of  $\frac{N}{M}$ . Note that  $g(q^d) = d^{-k/2}g \mid \begin{pmatrix} d & 0 \\ 0 & 1 \end{pmatrix}$ . For any  $\alpha \in SL_2(\mathbb{Z})$ , we can find  $\alpha' \in SL_2(\mathbb{Z})$  and  $u \in GL_2(\mathbb{Q})$  such that  $\begin{pmatrix} d & 0 \\ 0 & 1 \end{pmatrix} \alpha = \alpha'u$ . Hence to compute all expansions  $f \mid [\alpha]$ , it suffices to give an algorithm for newforms.

In the rest of this chapter, we will restrict ourselves to solving the following problem:

**Problem 1.2.1.** Let f be a normalized newform in  $S_2(\Gamma_0(N))$  and z be a cusp on  $X_0(N)$  of width one. Compute the q-expansion of  $f_z$ .

#### 1.3 Twists of newforms

For  $f \in S_k(\Gamma_1(N), \epsilon)$  a newform with expansion  $f = \sum_n a_n(f)q^n$  and  $\chi$  a Dirichlet character, the twist  $f_{\chi}$  is a modular form with expansion  $f_{\chi}(q) = \sum_n a_n(f)\chi(n)q^n$ .

**Lemma 1.3.1.** [AWL78, Proposition 3.1] Let  $F \in S_k(\Gamma_1(N), \epsilon)$ , where  $\epsilon$  is a character of conductor N'. Let  $\chi$  be a character modulo M. Put  $\tilde{N} = lcmN, N'M, M^2$ . Then  $F_{\chi} \in S_k(\Gamma_1(\tilde{N}), \epsilon \chi^2)$ .

In particular, when  $\epsilon$  is the trivial character and the conductor M of  $\chi$  satisfies  $M^2 \mid N$ , we have  $F_{\chi} \in S_k(\Gamma_1(N), \chi^2)$ .

We write  $f \otimes \chi$  for the unique newform such that  $a_p(f \otimes \chi) = a_p(f_{\chi})$  for all but finitely many primes p. From now, we refer to  $f \otimes \chi$  as the twist of f by  $\chi$ .

We quote two results from [Winnie], which we will use extensively. First, we recall the definitions of  $U_d$  and  $B_d$  operators. For a modular form  $f = \sum a_n q^n$  and a positive integer d, we put

$$f|U_d = \sum a_{nd}q^n, f|B_d = \sum a_n q^{nd}.$$

**Lemma 1.3.2.** [AWL78, Theorem 3.1] Let  $q \mid N$  and Q be the q-primary part of N. Write N = QM. Let F be a newform in  $S_k(\Gamma_1(N), \epsilon)$  with  $\operatorname{cond}(\epsilon_Q) = q^{\alpha}, \alpha \geq 0$ . Let  $\chi$  be a character with conductor  $q^{\beta}$ ,  $\beta \geq 1$ . Put  $Q' = \max\{Q, q^{\alpha+\beta}, q^{2\beta}\}$ . Then

- (1) For each prime  $q' \mid M, F_{\chi}$  is not of level Q'M/q.
- (2) The exact level of  $F_{\chi}$  is Q'M provided (a)  $\max\{q^{\alpha+\beta}, q^{2\beta}\} < Q$  if Q' = Q, or (b)  $\operatorname{cond}(\epsilon_Q \chi) = \max\{q^{\alpha}, q^{\beta}\}$  if Q' > Q.

**Lemma 1.3.3.** [AWL78, Theorem 3.2] Let  $q \mid N$  and Q be the q-primary part of N. Write N = QM. Let  $\chi$  be a character with conductor equal a power of q. Let F be a newform in  $S_k(\Gamma_1(N), \epsilon)$ . Then  $f \otimes \chi$  is a newform in  $S_k(\Gamma_1(Q'M, \epsilon \chi^2))$ , where Q' is a power of q, such that

$$F_{\chi} = f \otimes \chi - (f \otimes \chi)|U_q|B_q.$$

Since our goal is to compute expansions of newforms on  $\Gamma_0(N)$ , we will make the following assumptions: from now, unless otherwise noted, we assume f has trivial character, and that  $\operatorname{cond}(\chi)^2 \mid N$ .

Next, we consider the problem of identifying the newform  $f \otimes \chi$ . This includes finding its level, which we denote by  $M_{\chi}$ , and its q-expansion to arbitrarily many terms. We will

assume that we have an oracle which, given weight k and level N, computes the expansions of all newforms in  $S_k(\Gamma_1(N))$  to arbitrarily many terms (for example, use the algorithm in [Steb]).

Now we proceed on how to recognise the level of  $f \otimes \chi$  from the coefficients of f. One potential obstacle is that we do not know all Fourier coefficients of  $f \otimes \chi$ . We only know that  $a_n(f \otimes \chi) = a_n(f)\chi(n)$  when  $\gcd(n, N) = 1$ . This can be overcome using a variant of Sturm's argument. First we prove a lemma.

**Lemma 1.3.4.** Let  $f \in S_k(N, \epsilon)$  be a normalized newform. Then  $f|U_q|B_q \in S_k(Nq^2, \epsilon)$ .

Proof. We use a standard fact that for any integer  $d \geq 1$ , the map  $f \mapsto f|B_d$  takes  $S_k(N,\epsilon)$  to  $S_k(Nd,\epsilon)$ . To prove the lemma, we consider two separate cases. First, assume  $q \nmid N$ , then we have  $T_q = U_q + q^{k-1}\epsilon(q)B_q$ . By our assumption, we have  $f|T_q = a_q(f)f$ . Therefore, we have  $f|U_q|B_q = f|(T_q - q^{k-1}\epsilon(q)B_q)|B_q = a_q(f)f|B_q - q^{k-1}\epsilon(q)f|B_q^2$ . Hence  $f|U_q|B_q \in S_k(Nq^2,\epsilon)$ . Now assume  $q \mid N$ , so  $U_q = T_q$ . Hence  $f|U_q|B_q = a_q(f)f|B_q \in S_k(Nq,\epsilon) \subseteq S_k(Nq^2,\epsilon)$ .

The next proposition generalised the usual Sturm bound argument for modular forms.

**Proposition 1.3.5.** Let  $g_1$ ,  $g_2$  be two normalised newforms of levels  $N_1 \mid N_2$  and the same nybentypus character  $\epsilon$ . Assume  $\epsilon$  has prime power conductor  $Q = q^{\beta}$  such that  $Q^2 \mid N$ . Let B be the Sturm bound for the congruence subgroup  $\Gamma_1(Nq^2)$ . Suppose

$$a_n(g_1) = a_n(g_2)$$
, for all  $1 \le n \le B$  such that  $gcd(n,q) = 1$ .

Then  $g_1 = g_2$ .

*Proof.* Following [Winnie], we define the operator  $K_q$  on the space of modular forms by

$$g|K_q = g - g|U_q|B_q.$$

Then the assumption is equivalent to the statement that  $\delta = (g_1 - g_2)|K_q$  has  $a_n(\delta) = 0$  for all  $1 \le n \le B$ . Since  $\delta \in S_k(Nq^2, \epsilon)$ , Sturm's theorem implies  $\delta = 0$ .

We know from [DS06, Theorem 5.7.1] that  $g_1 - g_2 \in S_k(N_2, \epsilon)^{old}$ . Suppose  $N_1 < N_2$ , then  $g_1$  is in the old subspace, hence so is  $g_2$ , a contradiction. Therefore we must have  $N_1 = N_2$ . It follows that  $g_1 - g_2 \in S_k(N_2, \epsilon)^{new}$ , since  $g_1, g_2$  are newforms. This forces  $g_1 - g_2 = 0$ .  $\square$ 

Now we are ready to describe the algorithm.

# **Algorithm 1** Identifying $f \otimes \chi$

**Input:** k – a positive even integer;  $f \in S_k(\Gamma_0(N))$  a normalized newform;  $\chi$  a Dirichlet character of prime power conductor  $Q = q^{\beta}$ ;  $Q^2 \mid N$ ; B – a positive integer

**Output:** The level  $M_{\chi}$  of  $f \otimes \chi$  and the Fourier expansion of  $f \otimes \chi$  up to  $q^{B}$ .

```
1: if Q = 1 then
```

2: return N.

3: end if

4: 
$$Q' := \operatorname{cond}(\chi^2)$$
;  $N_0 := \frac{N}{q^{v_q(N)}}$ ;  $M_0 := Q'N_0$ ;  $t := \frac{N}{M_0} \in \mathbb{Z}$ .

5: **for** each positive divisor d of t **do** 

6: Set 
$$V_d := S_k(M_0 d, \chi^2)$$
.

7: Compute a basis of newforms  $\{g_1^{(d)}, \cdots g_{s_d}^{(d)}\}$  of  $V_d$ .

8: Set  $B_d :=$  the Sturm bound for  $\Gamma_1(M_0dq^2)$ .

9: **for**  $1 \le j \le s_d$  **do** 

10: **if** 
$$a_n(g_i^{(d)}) = a_n(f)\chi(n)$$
 for all  $1 \le n \le B_d, \gcd(n, q) = 1$  **then**

11: return  $M_0d$ .

12: end if

13: end for

14: end for

We give some sample computations applying the above algorithm.

**Example 1.3.6.** Let f be the normalised newform attached to the elliptic curve

$$E: y^2 + xy + y = x^3 - x - 2$$

of Cremona label **50a**. Then  $f \otimes \chi$  is new of level 50 for all Dirichlet characters  $\chi$  with modulus 5. In other words, f is 5-minimal.

As another example, we demonstrate a newform which is not p-minimal.

**Example 1.3.7.** Let f be the normalised newform attached to the elliptic curve

$$E: y^2 + xy = x^3 + x^2 - 25x - 111$$

of label **98a**. Let  $\chi$  be the Dirichlet character modulo 7 defined by  $\chi(3 \pmod{7}) = -1$ . We found that  $f \otimes \chi$  is a newform of level 14, with q-expansion (fixme: add expansion).

#### 1.4 Pseudo-eigenvalues

Let  $\epsilon$  be a Dirichlet character modulo N and let f be a newform in  $S_k(N, \epsilon)$ . For any divisor Q of N with  $gcd(Q, \frac{N}{Q}) = 1$ , there is an algebraic number  $w_Q(f)$  of absolute value one and a newform g in  $S_k(N, \overline{\epsilon_Q} \epsilon_{N/Q})$  such that

$$W_O(f) = w_O(f)q$$

**Definition 1.4.1.** The number  $w_Q(f)$  is called the *pseudo-eigenvalue* of  $W_Q$  on f. We write  $w(f) = w_N(f)$ .

For a power series  $f = \sum_{n \geq 0} a_n q^n$ , its complex conjugate, denoted by  $f^*$ , is

$$f^*(q) = \sum \overline{a_n} q^n.$$

From [AWL78] we have  $W_N(f) = w(f)f^*$ . In the rest of this section, we describe an algorithm to efficiently compute w(f) numerically. For a positive even integer k, let  $\mathbb{M}(k)$  denote the space of weight-k modular symbols defined in [Steb]. The space  $\mathbb{M}(k)$  is a quotient of  $\mathbb{Z}[X,Y]_{k-2} \otimes \mathbb{P}^1(\mathbb{Q})^2$ , and  $GL_2(\mathbb{Q})$  acts on  $\mathbb{M}(k)$  via the following rule

$$g(P(X,Y) \otimes \{\alpha,\beta\}) = P(g^{-1}(X,Y)^T)\{g(\alpha),g(\beta)\}.$$

Most importantly, there is a pairing between M(k) and the space of modular forms of weight k, defined as

$$\langle f, P(X, Y) \otimes \{\alpha, \beta\} \rangle_k = \int_{\alpha}^{\beta} f(z) P(z, 1) dz.$$

We will suppress the subscript k if its value is clear from context.

**Lemma 1.4.2.** Let  $M \in \mathbb{M}(k)$  and  $f \in S_k(\Gamma_1(N))$ . Then

$$N^{\frac{k}{2}-1}\langle f|W_N, M\rangle = \langle f, W_N M\rangle.$$

*Proof.* See proof of [Steb, Proposition 8.17]. Note that the extra factor  $N^{\frac{k}{2}-1}$  is due to the different constants involved in the definition of the weight-k action of  $GL_2(\mathbb{Q})$  on modular forms.

The map

$$*: P(x,y)\{\alpha,\beta\} \mapsto P(-x,y)\{-\bar{\alpha},-\bar{\beta}\}$$

defines the star involution on the space  $\mathbb{M}(k)$ . We have  $\langle f^*, M \rangle = \overline{\langle f, M^* \rangle}$ .

Now we are ready to prove the main theorem of this section.

**Theorem 1.4.3.** Let f be a normalised newform on  $\Gamma_1(N)$  with positive even weight k and let  $M \in \mathbb{M}(k)$  be such that  $W_N(M) = N^{k/2-1}M^*$ . Assume  $\langle f, M \rangle \neq 0$ . Then

$$w(f) = \frac{\langle f, M \rangle}{\langle f, M \rangle}.$$

*Proof.* Since  $W_N^2(M) = N^{k-2}M$  for all  $M \in M(k)$ , the assumption implies  $W_N(M^*) = N^{k/2-1}M$ . Now

$$N^{k/2-1}\langle f|W_N, M^*\rangle = \langle f, W_N(M^*)\rangle \text{ (Lemma ??)}$$

$$\Longrightarrow N^{k/2-1}w(f)\langle f^*, M^*\rangle = N^{k/2-1}\langle f, M\rangle$$

$$\Longrightarrow w(f) = \frac{\langle f, M\rangle}{\langle f^*, M^*\rangle}$$

$$\Longrightarrow w(f) = \frac{\langle f, M\rangle}{\langle f, M\rangle}.$$

Suppose  $\alpha, \beta \in \{z \in \mathbb{C} | Im(z) > 0, |z| = 1/\sqrt{N}\}$ . Then it is easy to verify that  $M = (xy)^{k/2-1} \otimes \{\alpha, \beta\}$  satisfies  $W_N(M) = M^*$ . Finally, we arrive at the algorithm to compute w(f).

## Algorithm 2 Computing the pseudo-eigenvalue of newforms.

**Input:** k – a positive even integer.  $f \in S_k(\Gamma_1(N))$  a normalized newform.

**Output:** a numerical approximation of w(f).

- 1:  $n_0 := 10, z_0 := \frac{i}{\sqrt{N}}. \delta = 10^{-3}.$
- 2: Randomly generate  $n_0$  points  $\{z_1, \dots, z_{N_0}\} \subseteq \{z | 0 < Im(z) < \frac{1}{2\sqrt{N}}, |z| = \frac{1}{\sqrt{N}}\}.$
- 3: **for**  $1 \le i \le n_0$  **do**
- 4: compute the period integral  $c_i = \int_{z_0}^{z_i} 2\pi i f(z) z^{\frac{k-2}{2}} dz$ .
- 5:  $w_i \leftarrow c_i/\bar{c}_i$ .
- 6: end for
- 7: if the standard deviation of  $w_1, \dots, w_{n_0}$  is less than  $\delta$  then
- 8:  $w \leftarrow \frac{1}{n_0} (\sum_i w_i)$ .
- 9: return w.
- 10: **else**
- 11: return FAIL.
- 12: **end if**

#### 1.5 Formula for the Fourier expansion of f at width one cusps: Part 1

First we recall some notations from [AWL78].

**Definition 1.5.1.** For a positive integer c', let  $S'_c = \begin{pmatrix} 1 & \frac{1}{c'} \\ 0 & 1 \end{pmatrix}$ . If  $\chi$  is a character modulo c', we define the operator on modular forms

$$f|R_{\chi}(c') = \sum_{u=0}^{c'-1} \bar{\chi}(u)f|S_{c'}^{u}.$$

Write  $R_{\chi}$  in short for  $R_{\chi}(\operatorname{cond}(\chi))$ . Note that  $f|R_{\chi}=g(\bar{\chi})f_{\chi}$ . Conversely, if (a,M)=1,

we have

$$\phi(c')[S_{c'}^u] = \sum_{\chi:cond(\chi)|c'} \chi(u)R_{\chi}(c'). \tag{1.5.1}$$

For our convenience, we define a new set of operators, which are basically the conjugates of  $S'_c$  and  $R_{\chi}(c')$  by  $W_N$ . Let  $A'_c = \begin{pmatrix} 1 & 0 \\ c' & 1 \end{pmatrix}$ . Then we have

Fact 1.5.2. 
$$-N \cdot A_{N/c'}^{-1} = W_N S_{c'} W_N$$

From now on, we assume c is a divisor of N and  $c' = \frac{N}{c}$ . Then

$$A_c^{-1} = W_N S_{c'} W_N.$$

Since  $W_N^2 = id$  as operators, we have

$$A_c^{-u} = W_N S_{c'}^u W_N, \, \forall u \in \mathbb{Z}.$$

Parallel to the notion of  $R_{\chi}(c')$ , let  $\Phi_{\chi}(c) = \sum_{u=0}^{c'-1} \bar{\chi}(u) A_c^{-u}$ . Then  $\Phi_{\chi}(c) = W_N R_{\chi}(c') W_N$ . Similar to Formula 1.5.1, we have

$$\varphi(c')A_c^{-a} = \sum_{\operatorname{cond}(\chi)|c'} \chi(a)\Phi_{\chi}(c) = \sum_{\operatorname{cond}(\chi)|c'} \chi(a)W_N R_{\chi}(c')W_N. \tag{1.5.2}$$

Finally, applying Formula 1.5.2 to f, we arrive at

$$f_{\left[\frac{a}{c}\right]}(q) = \frac{1}{\varphi(c')} \sum_{\operatorname{cond}(\chi)|c'} \chi(-a) f|W_N R_{\chi}(c') W_N. \tag{1.5.3}$$

$$= \frac{w(f)}{\varphi(c')} \sum_{\operatorname{cond}(\chi)|c'} \chi(-a) f | R_{\chi}(c') W_N.$$
 (1.5.4)

Now it left to compute the expansions of each  $f|R_{\chi}(c')|W_N$  in the sum.

# 1.6 Formula for the Fourier expansion of f at width one cusps: Part 2

In this section, we describe how to compute the expansion of  $f|R_{\chi}(c')|W_N$ . First note the following identity between operators on  $S_k(\Gamma_1(N), \epsilon)$ :

$$T_p = U_p + \epsilon(p) p^{\frac{k}{2}} B_p.$$

We recall some notations and a result from [Del02].

**Definition 1.6.1.** [Del02, Definition III.2.4] Put cond'( $\chi$ ) multiplicatively. If  $\chi_j$  is trivial character modulo  $p_j^{a_j}$ , set cond'( $\chi_j$ ) =  $p_j$ .  $\chi = \chi_{nt}\chi_{tr}$ . Put

$$g'(\chi) = (-1)^{|I|} \chi_{nt}(tr) g(\chi_{nt}).$$

**Lemma 1.6.2.** [Del02, Prop 2.6] Let c' be such that  $c'^2 \mid N$ . For a Dirichlet character  $\chi$  mod c', we have

$$f|R_{\chi}(c') = \begin{cases} g'(\bar{\chi})f_{\chi_{nt}} & if \operatorname{cond}'(\chi) = c\\ 0 & else \end{cases}$$

Then we compute  $f_{\chi_{nt}}$  by the following: suppose  $g = f \otimes \chi_{nt}$ . Then

$$f_{\chi_{nt}} = \prod_{i=1}^{r} g | K_{p_i}.$$

Moreover, we have

$$K_p = 1 - U_p B_p = \begin{cases} 1 - (T_p - \chi_{nt}^2(p) p^{\frac{k}{2}} B_p) | B_p & p \nmid M \\ 1 - T_p | B_p & p \mid M \end{cases}.$$

Using the commutativity of T and B, we can write  $f_{\chi_{nt}}$  in the form  $\sum c_i g(q_i^d)$ , where  $c_i$  and  $d_i$  are constants. To give a precise formula, we use the following notation. For a finite set S of integers, let  $p(S) = \prod_{s \in S} s$  denote the product of all elements in S.

**Theorem 1.6.3.** Let  $S_{\chi} = \{p_1, \dots, p_r\}$  be the set of prime divisors of  $\operatorname{cond}(\chi)$ . Let  $\mathcal{B}_{\chi,M} = \{(S_1, S_2) | S_1, S_2 \subseteq S_{\chi}, S_1 \cap S_2 = \emptyset, \gcd(M, p(S_2)) = 1.\}$ . Write  $g_{\chi} = f \otimes \chi$ . Then

$$f_{\chi_{nt}} = \sum_{(S_1, S_2) \in B_{S,M}} a_{p(S_1)}(g_{\chi}) p(S_2)^{k/2} \chi_{nt}^2(p(S_2)) g_{\chi} |B_{p(S_1)p(S_2)^2}.$$

Theorem 1.6.3 will be our starting point of computing the expansion of f at width one cusps. We will use it to compute  $f_{\chi_{nt}}|W_N$ . First we prove two lemmas.

**Lemma 1.6.4.** Let f be a newform of even weight k on  $\Gamma_1(M)$  and suppose d, N are positive integers such that  $Md \mid N$ . Then

$$f|B_d|W_N = \left(\frac{N}{Md^2}\right)^{k/2} w_M(f)\overline{f|B_{\frac{N}{Md}}}.$$

where  $w_M(f)$  is the pseudo-eigenvalue of f defined in previous sections.

*Proof.* Straightforward computation.

$$f|B_{d}|W_{N} = d^{-k/2}f|\begin{pmatrix} d & 0 \\ 0 & 1 \end{pmatrix}\begin{pmatrix} 0 & -1 \\ N & 0 \end{pmatrix}$$

$$= d^{-k/2}f|\begin{pmatrix} 0 & -1 \\ M & 0 \end{pmatrix}\begin{pmatrix} N/md & 0 \\ 0 & 1 \end{pmatrix}\begin{pmatrix} d & 0 \\ 0 & d \end{pmatrix}$$

$$= \left(\frac{N}{Md^{2}}\right)^{k/2}f|W_{M}|B_{N/Md}$$

$$= \left(\frac{N}{Md^{2}}\right)^{k/2}w_{M}(f)\bar{f}|B_{N/Md}$$

$$= \left(\frac{N}{Md^{2}}\right)^{k/2}w_{M}(f)\bar{f}|B_{N/Md}.$$

Before stating the second lemma, we quote another result of Winnie Li on the coefficients of a newform at primes dividing the level.

**Lemma 1.6.5.** [Li75, Theorem 3 (iii)] Let  $f = \sum_{n\geq 1} a_n(f)q^n$  be a normalized newform in  $S_k(\Gamma_1(N), \epsilon)$ , p a prime dividing N. Then

- (1) If  $\epsilon$  is a character modulo N/p and  $p^2 \mid N$ , then  $a_p(f) = 0$ .
- (2) If  $\epsilon$  is a character modulo N/p and  $p^2 \nmid N$ , then  $a_p(f)^2 = \epsilon(p)p^{k-2}$ .
- (3) If  $\epsilon$  is not a character modulo N/p, then  $|a_p(f)| = p^{\frac{k-1}{2}}$ .

**Lemma 1.6.6.** Using notations in Theorem 1.6.3, and assume  $(S_1, S_2) \in \mathcal{B}_{S,M}$  is such that  $a_{p(S_1)}(g_\chi) \neq 0$ . Then  $Mp(S_1)p(S_2)^2 \mid N$ .

Proof. Let p be a prime divisor of  $N' := Mp(S_1)p(S_2)^2$ . If  $p \nmid M$ , then  $\operatorname{ord}_p(N') \leq \operatorname{ord}_p(\operatorname{cond}(\chi)^2) \leq \operatorname{ord}_p(N)$ . So we assume  $p \mid M$ , hence  $p \nmid p(S_2)$ . If  $p \nmid p(S_1)$ , then it follows from  $M \mid N$ ; if  $p \mid p(S_1)$ , we want to show that  $\operatorname{ord}_p(M) < \operatorname{ord}_p(N)$ . Suppose not, then  $\operatorname{ord}_p(M) = \operatorname{ord}_p(N) \geq 2\operatorname{ord}_p(\operatorname{cond}(\chi))$ . Since  $\operatorname{cond}(\chi^2) \leq \operatorname{cond}(\chi)$ , we know  $\chi^2$  is a character modulo M/p. Applying Lemma 1.6.5 to the newform  $g_{\chi}$  on level M, we see that  $a_p(g_{\chi}) = 0$ , hence  $a_{p(S_1)}(g_{\chi}) = 0$  by multiplicativity.

Applying Lemma 1.6.4 to Theorem 1.6.3, we finally arrive at

**Theorem 1.6.7.** Let f be a normalized newform in  $S_2(\Gamma_0(N))$  and z = [a/c] be a cusp on  $X_0(N)$  of width one. Then  $f_z$  is

$$f_z = \frac{w(f)}{\varphi(c')} \sum_{\chi: \text{cond}'(\chi) = c'} \chi(-a) g'(\bar{\chi}) w(g_\chi) t_\chi.$$

Here  $t_{\chi}$  is as follows: let  $M_{\chi}$  denote the level of  $g_{\chi} := f \otimes \chi$ . Then

$$t_{\chi} = \sum_{(S_1, S_2) \in B_{S_{\chi_{nt}}, M_{\chi}}} (-1)^{|S_1|} a_{p(S_1)}(g_{\chi}) \left( \frac{N}{M_{\chi} p(S_1)^2 p(S_2)^3} \right)^{k/2} \chi(p(S_2)) \overline{g_{\chi} | B_{\frac{N}{M_{\chi} p(S_1) p(S_2)^2}}}.$$

This theorem gives us an algorithm to compute the expansion of  $f_z$ , which we describe below. But first, we take a closer look at what ingredients goes into the expansion. Given a newform  $f \in S_k(\Gamma_0(N))$  and a width one cusp z of denominator c. We need to consider the twist of f by all Dirichlet characters of conductor dividing c. For each such character  $\chi$ , we then need to determine the level  $M_{\chi}$  and q-expansion of the newform  $f \otimes \chi$ , the latter boils down to knowing  $a_p(f \otimes \chi)$  for all  $p \mid \text{cond}(\chi)$ . Then we need to compute the. Finally, we combine these information together and apply Throem 1.6.7 to compute  $f_z$ .

# **Algorithm 3** Computing Fourier coefficients of f at width one cusps

**Input:**  $f \in S_k(\Gamma_0(N))$  a newform; a, c – coprime integers such that  $N \mid c^2$ ; B – a positive integer.

**Output:** The first B Fourier coefficients of  $f_z(q)$ .

- 1:  $c' \leftarrow N/c$ .  $X \leftarrow$  The set of all Dirichlet characters  $\chi$  such that  $\operatorname{cond}'(\chi) = c'$ .
- 2: compute w(f) using Algorithm 2.
- 3: for  $\chi$  in X do
- 4: Compute  $M_{\chi}$  and the expansion of  $g_{\chi} := f \otimes \chi$  to B terms, using Algorithm 1
- 5: Compute  $g'(\bar{\chi})$ .
- 6: Compute  $w(g_{\chi})$  using Algorithm 2.
- 7: end for
- 8: Apply Theorem 1.6.7 to compute  $f_z$  to B terms.

#### 1.7 A Converse Theorem

Given the work in previous sections, it is a natural question then to ask whether the information on twists of f is uniquely determined by the expansion of f at width one cusps. The answer is yes, and the precise statement is in the following theorem.

**Theorem 1.7.1.** Let f be a normalized newform in  $S_k(\Gamma_0(N))$ . Assume the eigenvalue  $w_N(f)$  is known. Suppose c is a positive divisor of N such that  $N \mid c^2$ . Then the expansions of  $f_z$ , where z runs through all cusps of denominator c, uniquely determines the following: for each Dirichlet character  $\chi$  of such that  $\operatorname{cond}'(\chi) = c'$ , the level  $M_\chi$ , the pseudo-eigenvalue  $w_{M_\chi}$  and the q-expansion of the newform  $f \otimes \chi$ .

*Proof.* By plug in different a's. We can solve for  $t_{\chi}$ . Consider the first nonzero term of  $t_{\chi}$ . Suppose

$$t_{\chi} = u_{\chi} q^{v_{\chi}} + O(q^{v_{\chi}+1}), \ u_{\chi} \neq 0.$$

Assuming that  $\chi$  has prime power conductor  $p^{\beta} > 1$ , we claim that

$$\left| \frac{v^{k/2}}{u} \right| = \begin{cases} p^{k/2} & if p \nmid M_{\chi} \\ p^{1/2} & if p \mid M_{\chi} \text{ and } a_{p}(g) \neq 0 \\ 1 & else \end{cases}$$

Proof of claim: the first and third case are easy to verify using Theorem 1.6.7. Now assume  $p \mid M$  and  $a_p(g_\chi) \neq 0$ . By Lemma 1.6.5, we have  $|a_p(g_\chi)| = p^{k/2-1/2}$  or  $p^{k/2-1}$ . However,  $|a_p(g_\chi)| = p^{k/2-1}$  only if  $p \mid M_\chi$  and  $\chi^2$  is a character modulo  $M_\chi/p$ . This means  $\chi^2$  is the trivial character. By another theorem of Winnie, we compute the p-level of  $f = g_\chi \otimes \bar{\chi}$ : note that  $\max p, p^{\alpha+\beta}, p^{2\beta} > p$ , so (ii) applies and the p-level of f is equal to  $\max(p^\alpha, p^\beta) = p^\beta$ , i.e.,  $\operatorname{ord}_p(N) = \beta$ . This is impossible since we have  $p^{2\beta} = \operatorname{cond}(\chi)^2 \mid N$ .

Therefore, we have  $|a_p(g_\chi)| = p^{k/2-1/2}$  and the claim follows.

Since  $k \geq 2$ , we could determine which case we are in. Then we can read off  $M_{\chi}$  and  $w_M(g_{\chi})$ . For example, if we are in the second case, then the level can be computed via

 $M_{\chi} = \frac{N}{v_{\chi}p}$ . Now the  $N/M_{\chi}$ 's coefficient of  $t_{\chi}$  is

$$a_{\frac{N}{M}}(t_{\chi}) = w(g_{\chi})(\frac{N}{M})^{k/2}(1 - |a_{p}(g_{\chi})|^{2}\chi^{2}(p)p^{-k/2})$$
$$= w(g_{\chi})(\frac{N}{M})^{k/2}(1 - p^{k/2-1}\chi^{2}(p)).$$

This allows us to solve  $w(g_{\chi})$ . Finally, we compute  $a_p(g_{\chi})$  by  $a_p(g) = \frac{-u_{\chi}}{w(g_{\chi})\chi^2(p)(\frac{N}{Mp})^{k/2}}$ . The value  $a_p(g)$  determines the expansion of  $g_{\chi}$ . Recursively, we could solve for all pn-coefficients of  $g_{\chi}$ , from which we deduce it complete q-expansion.

In the general case, we consider the following subsets of  $S_{\chi}$ . Let  $S_1^* = \{p \in S_{\chi} : p \mid M\}$ ,  $S_2^* = S_{\chi} \setminus S_1^*$ , and  $\widetilde{S_1^*} = \{p \in S_1^* : a_p(g_{\chi}) \neq 0\}$ .

It follows that the leading term of  $t_{\chi}$  belongs to the summand corresponding to  $(\widetilde{S_1^*}, S_2^*)$  in Theorem 1.6.7. Still writing the leading term as  $u_{\chi}q^{v_{\chi}}$ , we have

$$u_{\chi} = w(g_{\chi})\chi^{2}(p(S_{2}))a_{p(\widetilde{S_{1}^{*}})}(g_{\chi})p(\widetilde{S_{1}^{*}})^{-k}(p(S_{2}^{*})^{-3k/2}\left(\frac{N}{M_{\chi}}\right)^{k/2}, v_{\chi} = \frac{N}{M_{\chi}p(\widetilde{S_{1}^{*}})p(S_{2}^{*})^{2}}.$$

Similar to the prime power conductor case above, we have  $|a_{p(\widetilde{S_1^*})}(g_\chi)| = p(\widetilde{S_1^*})^{k/2-1/2}$ . So

$$|v_{\chi}^{k}u_{\chi}^{-2}| = p(\widetilde{S_{1}^{*}})p(S_{2}^{*})^{2}. \tag{1.7.1}$$

Hence we can factor  $|v_\chi^k u_\chi^{-2}|$  and obtain  $p(\widetilde{S_1^*})$  and  $p(S_2^*)$ . Then  $M_\chi$  can be solved using  $v_\chi$ . Plug it back into  $u_\chi$ , we obtain  $a_{p(\widetilde{S_1^*})}w(g_\chi)$ . Finally, for each  $p\in\widetilde{S_1^*}$ , the  $v_\chi p$ 's coefficient of  $t_\chi$  allows us to compute  $a_{p(\widetilde{S_1^*})/p}(g_\chi)w(g_\chi)$ . These together determines  $w(g_\chi)$  and  $a_{p(\widetilde{S_1^*})}$ . The other Fourier coefficients of  $g_\chi$  can then be computed recursively.

#### 1.8 Fields of definitions

**Lemma 1.8.1.** Let c be a cusp of denominator d and let d' = N/d. Then

$$\mathbb{Q}(\{a_n(f,c)\}) \subseteq \mathbb{Q}(\{a_n(f)\},\zeta_{d'}).$$

## 1.9 Examples

Let  $E = \mathbf{50a}$  and consider the 4 cusps of denominator 10 on  $X_0(50)$ . The corresponding first terms of q-expansions at these cusps are

$$a_1(f, \frac{1}{10}) = \frac{1}{5}\zeta_5^3 - \frac{3}{5}\zeta_5^2 + \frac{3}{5}\zeta_5 - \frac{1}{5}$$

$$a_1(f, \frac{3}{10}) = \frac{3}{5}\zeta_5^3 + \frac{6}{5}\zeta_5^2 + \frac{4}{5}\zeta_5 + \frac{2}{5}$$

$$a_1(f, \frac{7}{10}) = \frac{2}{5}\zeta_5^3 - \frac{1}{5}\zeta_5^2 - \frac{4}{5}\zeta_5 - \frac{2}{5}$$

$$a_1(f, \frac{9}{10}) = -\frac{6}{5}\zeta_5^3 - \frac{2}{5}\zeta_5^2 - \frac{3}{5}\zeta_5 - \frac{4}{5}$$

## 1.10 Applications

One applications of the computation done in this chapter is the norm method to the computation of j-polynomials introduced in Chapter . Recall that the issue with the norm method for non-square free level is computing the expansions of form  $f|\gamma$ , where  $\gamma$  runs over the set of right coset representatives of  $\Gamma_0(N)$  in  $SL_2(\mathbb{Z})$ . As we have seen, it suffices to compute the expansions of f at all width one cusps.

#### 1.11 Norm guess and data

Fixme

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