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Computational aspects of modular parametrizations of elliptic curves

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Abstract

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Abstract goes here.

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GLOSSARY

ARGUMENT: replacement text which customizes a \LaTeX macro for each particular usage.

DEDICATION

to all of you

Chapter 1

INTRODUCTION

Chapter 2

COMPUTING THE MAZUR SWINNERTON-DYER CRITICAL SUBGROUP OF ELLIPTIC CURVES

2.0.1 Preliminaries

Let E be an elliptic curve over \mathbb{Q} and let $L(E, s)$ be the L -function of E . The rank part of the Birch and Swinnerton-Dyer (BSD) conjecture states that

$$\text{rank}(E(\mathbb{Q})) = \text{ord}_{s=1} L(E, s).$$

The right hand side is called the *analytic rank* of E , and is denoted by $r_{\text{an}}(E)$. The left hand side is called the *algebraic rank* of E . The rank part of the BSD conjecture is still open when $r_{\text{an}}(E) > 1$, and its proof for the case $r_{\text{an}}(E) = 1$ uses the *Gross-Zagier formula*, which relates the value of certain L -functions to heights of Heegner points.

Let N denote the conductor of E . The modular curve $X_0(N)$ is a nonsingular projective curve defined over \mathbb{Q} . Since E is modular (Breuil, Conrad, Diamond, and Taylor [BCDT01]), there is a surjective morphism $\varphi : X_0(N) \rightarrow E$ defined over \mathbb{Q} . Let ω_E be the invariant differential on E and let $\omega = \varphi^*(\omega_E)$. Then ω is a holomorphic differential on $X_0(N)$ and we have $\omega = cf(z)dz$, where f is the normalized newform attached to E and c is a nonzero constant. In the rest of the paper, we fix the following notations: the elliptic curve E , the conductor N , the morphism φ , and the differential ω . Let R_φ be the ramification divisor of φ .

Definition 2.0.1 (Mazur and Swinnerton-Dyer [MSD74]). The *critical subgroup* of E is

$$E_{\text{crit}}(\mathbb{Q}) = \langle \text{tr}(\varphi([z])) : [z] \in \text{supp } R_\varphi \rangle,$$

where $\text{tr}(P) = \sum_{\sigma: \mathbb{Q}(P) \rightarrow \mathbb{Q}} P^\sigma$.

Since the divisor R_φ is defined over \mathbb{Q} , every point $[z]$ in its support is in $X_0(N)(\overline{\mathbb{Q}})$, hence $\varphi([z]) \in E(\overline{\mathbb{Q}})$, justifying the trace operation. The group $E_{\text{crit}}(\mathbb{Q})$ is a subgroup of $E(\mathbb{Q})$. Observe that $R_\varphi = \text{div}(\omega)$, thus $\deg R_\varphi = 2g(X_0(N)) - 2$. In the rest of the paper, we use the notation $\text{div}(\omega)$ in place of the ramification divisor R_φ . In addition, we will assume E is an optimal elliptic curve, so φ is unique up to sign. This justifies the absence of φ in the notation $E_{\text{crit}}(\mathbb{Q})$.

Recall the construction of *Heegner points*: for an imaginary quadratic order $\mathcal{O} = \mathcal{O}_d$ of discriminant $d < 0$, let $H_d(x)$ denote its *Hilbert class polynomial*.

Definition 2.0.2. A point $[z] \in X_0(N)$ is a “*generalized Heegner point*” if there exists a negative discriminant d s.t. $H_d(j(z)) = H_d(j(Nz)) = 0$. If in addition we have $(d, 2N) = 1$, then $[z]$ is a *Heegner point*.

For any discriminant d , let E_d denote the quadratic twist of E by d . Then the Gross-Zagier formula in [GZ86] together with a non-vanishing theorem for $L(E_d, 1)$ (see, for example, Bump, Friedberg, and Hoffstein [BFH90]) implies the following

Theorem 2.0.3. (1) If $r_{\text{an}}(E) = 1$, then there exists a Heegner point $[z]$ on $X_0(N)$ such that $\text{tr}(\varphi([z]))$ has infinite order in $E(\mathbb{Q})$.

(2) If $r_{\text{an}}(E) \geq 2$, then $\text{tr}(\varphi([z])) \in E(\mathbb{Q})_{\text{tors}}$ for every “*generalized Heegner point*” $[z]$ on $X_0(N)$.

The first case in the above theorem is essential to the proof of rank BSD conjecture for $r_{\text{an}}(E) = 1$. We observe that the defining generators of the critical subgroup also take the form $\text{tr}(\varphi([z]))$. Then a natural question is:

Question 2.0.4. Does there exist an elliptic curve E defined over \mathbb{Q} such that $r_{\text{an}}(E) \geq 2$ and $\text{rank}(E_{\text{crit}}(\mathbb{Q})) > 0$?

We will show that the answer is negative for all elliptic curves with conductor $N < 1000$, using *critical polynomials* attached to elliptic curves.

2.0.2 Main results

Let E, N, φ , and ω be as defined previously, and write $\text{div}(\omega) = \sum_{[z] \in X_0(N)} n_z [z]$. Let j denote the j -invariant function.

Definition 2.0.5. The *critical j -polynomial* of E is

$$F_{E,j}(x) = \prod_{z \in \text{supp div}(\omega), j(z) \neq \infty} (x - j(z))^{n_z}.$$

Because $\text{div}(\omega)$ is defined over \mathbb{Q} and has degree $2g(X_0(N)) - 2$, we have $F_{E,j}(x) \in \mathbb{Q}[x]$ and $\deg F_{E,j} \leq 2g(X_0(N)) - 2$, where equality holds if $\text{div}(\omega)$ does not contain cusps. For any non-constant modular function $h \in \mathbb{Q}(X_0(N))$, the *critical h -polynomial* of E is defined similarly, by replacing j with h .

In this paper we give two algorithms *Poly Relation* and *Poly Relation-YP* to compute critical polynomials. The algorithm *Poly Relation* computes the critical j -polynomial $F_{E,j}$, and the algorithm *Poly Relation* computes the critical h -polynomial $F_{E,h}$ for some modular function h chosen within the algorithm. We then relate the critical polynomials to the critical subgroup via the following theorem. Recall that $H_d(x)$ denotes the Hilbert class polynomial associated to a negative discriminant d . We prove the following theorem.

Theorem 2.0.6. Suppose $r_{\text{an}}(E) \geq 2$, and assume at least one of the following holds:

- (1) $F_{E,h}$ is irreducible for some non-constant function $h \in \mathbb{Q}(X_0(N))$.
- (2) There exist negative discriminants D_k and positive integers s_k for $1 \leq k \leq m$ with $\mathbb{Q}(\sqrt{D_k}) \neq \mathbb{Q}(\sqrt{D_{k'}})$ for all $k \neq k'$, and an irreducible polynomial $F_0 \in \mathbb{Q}[x]$, such that

$$F_{E,j} = \prod_{k=1}^m H_{D_k}^{s_k} \cdot F_0.$$

Then $\text{rank}(E_{\text{crit}}(\mathbb{Q})) = 0$.

Combining Theorem 2.0.6 with our computation of critical polynomials, we verified the main result of this paper stated in the following corollary.

Corollary 2.0.7. *For all elliptic curves E of analytic rank 2 and conductor N smaller than 1000, the rank of $E_{\text{crit}}(\mathbb{Q})$ is zero.*

This paper is organized as follows: in Sections 2.1 and ??, we describe the algorithms *Poly Relation* and *Poly Relation-YP*. In Section ??, we prove Theorem 2.0.6. Last, in Section ??, we show a table of critical polynomials for all elliptic curves with rank 2 and conductor smaller than 1000, and prove Corollary 2.0.7.

2.1 The algorithm Poly relation

Let C/\mathbb{Q} be a nonsingular projective curve. For a rational function $r \in \mathbb{Q}(C)$, let $\text{div}_0(r)$ denote its divisor of zeros, and define $\deg r = \deg(\text{div}_0(r))$.

Definition 2.1.1. Let C/\mathbb{Q} be a nonsingular projective curve, and let r, u be two non-constant rational functions on C . A *minimal polynomial relation between r and u* is an irreducible polynomial $P(x, y) \in \mathbb{Q}[x, y]$ such that $P(r, u) = 0$ and $\deg_x(P) \leq \deg u, \deg_y(P) \leq \deg r$.

Minimal polynomial relation always exists and is unique up to scalar multiplication. Write $\text{div}(r) = \sum_{[z] \in X_0(N)} n_z[z]$ and $P(x, y) = f_n(y)x^n + \cdots + f_1(y)x + f_0(y)$. We will prove that

Proposition 2.1.2. *If $\mathbb{Q}(C) = \mathbb{Q}(r, u)$ and $\gcd(f_0(y), f_n(y)) = 1$, then there is a constant $c \neq 0$ s.t.*

$$f_0(y) = c \prod_{z \in \text{div}_0(r) \setminus \text{div}_\infty(u)} (y - u(z))^{n_z}.$$

Proof. Dividing $P(x, y)$ by $f_n(y)$, we get $x^n + \cdots + \frac{f_0(x)}{f_n(y)}$, which is a minimal polynomial of r over $\mathbb{Q}(u)$. So $\text{Norm}_{\mathbb{Q}(r, u)/\mathbb{Q}(u)}(r) = \frac{f_0(u)}{f_n(u)}$. The rest of the proof uses a fact on extensions of valuations (see, for example, [Stea, Theorem 17.2.2]), which we now quote.

Lemma 2.1.3. *Suppose v is a nontrivial valuation on a field K and let L be a finite extension of K . Then for any $a \in L$,*

$$\sum_{1 \leq j \leq J} w_j(a) = v(\text{Norm}_{L/K}(a)),$$

where the w_j are normalized valuations equivalent to extensions of v to L .

We continue with the proof. For any $z_0 \in C$ such that $u(z_0) \neq \infty$, consider the valuation $v = \text{ord}_{(u-u(z_0))}$ on $\mathbb{Q}(u)$. The set of extensions of v to $\mathbb{Q}(C) = \mathbb{Q}(r, u)$ is in bijection with $\{z \in C : u(z) = u(z_0)\}$. Take $a = r$ and apply Lemma 2.1.3, we obtain

$$\sum_{z: u(z)=u(z_0)} \text{ord}_z(r) = \text{ord}_{u-u(z_0)} \frac{f_0(u)}{f_n(u)}.$$

Combining the identities for all $z_0 \in C \setminus \text{div}_\infty(u)$, we have for some constant c ,

$$\prod_{z \in \text{div}(r): u(z) \neq \infty} (y - u(z))^{n_z} = c \cdot \frac{f_0(y)}{f_n(y)}.$$

If $r(z) = 0$, then the condition $\gcd(f_0(y), f_n(y)) = 1$ implies that $f_0(u(z)) = 0$ and $f_n(u(z)) \neq 0$. Therefore, since $\gcd(f_0, f_n) = 1$, we must have

$$f_0(y) = c \prod_{z \in \text{div}_0(r) \setminus \text{div}_\infty(u)} (y - u(z))^{n_z}.$$

This completes the proof. □

For completeness we also deal with the case where $u(z) = \infty$, which was left out in the above proof. The corresponding valuation on $\mathbb{Q}(u)$ is ord_∞ defined by $\text{ord}_\infty(g/h) = \deg g - \deg h$ for $0 \neq g, h \in \mathbb{Q}[u]$. We derive that

$$\sum_{z: u(z)=\infty} \text{ord}_z(r) = \deg f_n - \deg f_0.$$

Next we apply Proposition 2.1.2 to the computation of $F_{E,j}$. In the rest of the paper, $dj = j'(z)dz$ is viewed as a differential on $X_0(N)$. Fix the following two modular functions on $X_0(N)$:

$$r = j(j - 1728) \frac{\omega}{dj}, \quad u = \frac{1}{j}. \tag{2.1.1}$$

First we compute the divisor of r . Let $\mathcal{E}_2(N)$ and $\mathcal{E}_3(N)$ denote the set of elliptic points of order 2 and 3 on $X_0(N)$, respectively. Then

$$\operatorname{div}(dj) = -j^*(\infty) - \sum_{c=\text{cusp}} c + \frac{1}{2} \left(j^*(1728) - \sum_{z \in \mathcal{E}_2(N)} z \right) + \frac{2}{3} \left(j^*(0) - \sum_{z \in \mathcal{E}_3(N)} z \right). \quad (2.1.2)$$

Writing $j^*(\infty) = \sum_{c=\text{cusp}} e_c[c]$, we obtain

$$\operatorname{div}(r) = \operatorname{div}(\omega) + \frac{1}{2} \left(j^*(1728) + \sum_{z \in \mathcal{E}_2(N)} z \right) + \frac{1}{3} \left(j^*(0) + 2 \sum_{z \in \mathcal{E}_3(N)} z \right) - \sum_{c=\text{cusp}} (e_c - 1)[c]. \quad (2.1.3)$$

Note that (2.1.3) may not be the simplified form of $\operatorname{div}(r)$, due to possible cancellations when $\operatorname{supp} \operatorname{div}(\omega)$ contains cusps. But since the definition of $F_{E,j}$ only involves critical points that are not cusps, the form of $\operatorname{div}(r)$ in (2.1.3) works fine for our purpose.

Next we show $\mathbb{Q}(r, u) = \mathbb{Q}(X_0(N))$ for the functions r, u in (2.1.1). First we prove a lemma.

Lemma 2.1.4. *Let $N > 1$ be an integer and $f \in S_2(\Gamma_0(N))$ be a normalized newform. Suppose $\alpha \in \operatorname{SL}_2(\mathbb{Z})$ and $f|[\alpha] = f$, then $\alpha \in \Gamma_0(N)$.*

Proof. Write $\alpha = \begin{pmatrix} a & b \\ M & d \end{pmatrix}$. First we show that it suffices to consider the case where $d = 1$. Since $(M, d) = 1$, there exists $y, w \in \mathbb{Z}$ such that $My + dw = 1$. By replacing (y, w) with $(y + kd, w - kM)$ if necessary, we may assume $(y, N) = 1$. Now we can find $x, z \in \mathbb{Z}$ such that $\gamma = \begin{pmatrix} x & y \\ Nz & w \end{pmatrix} \in \Gamma_0(N)$, and $\alpha\gamma = \begin{pmatrix} * & * \\ M & 1 \end{pmatrix} \in \operatorname{SL}_2(\mathbb{Z})$ with $f|[\alpha\gamma] = f[\gamma] = f$. We then further reduce to the case where $\alpha = \begin{pmatrix} 1 & 0 \\ M & 1 \end{pmatrix}$, by noticing that $\begin{pmatrix} 1 & -b \\ 0 & 1 \end{pmatrix} \in \Gamma_0(N)$ and

$$\begin{pmatrix} 1 & -b \\ 0 & 1 \end{pmatrix} \cdot \begin{pmatrix} a & b \\ M & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ M & 1 \end{pmatrix}.$$

Let $w_N = \begin{pmatrix} 0 & -1 \\ N & 0 \end{pmatrix}$ be the Fricke involution on $X_0(N)$. Then $f|[w_N] = \pm f$, hence $f|[w_N \alpha w_N] = f$. We compute that $w_N \alpha w_N = \begin{pmatrix} -N & M \\ 0 & -N \end{pmatrix}$, thus $f(q) = f|[\begin{pmatrix} -N & M \\ 0 & -N \end{pmatrix}](q) = f(q\zeta_N^{-M})$, where $\zeta_N = e^{2\pi i/N}$. The leading term of $f(q)$ is q , while the leading term of $f(q\zeta_N^{-M})$ is $\zeta_N^{-M}q$. So we must have $\zeta_N^{-M} = 1$, i.e., $N \mid M$. Hence $\alpha \in \Gamma_0(N)$ and the proof is complete. \square

Proposition 2.1.5. *Let r, u be the two functions on $X_0(N)$ defined in (2.1.1), then $\mathbb{Q}(r, u) = \mathbb{Q}(X_0(N))$.*

Lemma 2.1.6. *Let g be the genus of $X_0(N)$. If $T \geq 2g - 2$ is a positive integer, then rj^T and u satisfy the second condition of Proposition 2.1.2.*

Proof. Let $r_1 = rj^T$. When $T \geq 2g - 2$, the support of $\text{div}_\infty(r_1)$ is the set of all cusps. Suppose $\gcd(f_n, f_0) > 1$. Let $p(y)$ be an irreducible factor of $\gcd(f_0, f_n)$. Consider the valuation ord_p on the field $K(y)$. Since $P(x, y)$ is irreducible, there exists an integer i with $0 < i < n$ such that $p(y) \nmid f_i$. Thus the Newton polygon of P with respect to the valuation ord_p has at least one edge with negative slope and one edge with positive slope. Therefore, for any Galois extension of L of $K(u)$ containing $K(r, u)$ and a valuation $\text{ord}_{\mathfrak{p}}$ on L extending ord_p , where \mathfrak{p} is an irreducible polynomial in $L[y]$ dividing $p(y)$, there exists two conjugates r', r'' of r such that $\text{ord}_{\mathfrak{p}}(r') < 0$ and $\text{ord}_{\mathfrak{p}}(r'') > 0$. This implies that $\text{div}_0(r') \cap \text{div}_\infty(r'') \neq \emptyset$. Fix $L = K(X(N))$, then all conjugates of r_1 in $K(X(N))/K(u)$ are of the form $r_1(\alpha z)$ for some $\alpha \in \text{SL}_2(\mathbb{Z})$. Hence the set of poles of any conjugate of r_1 is the set of all cusps on $X(N)$, a contradiction. \square

Note that for any $T \in \mathbb{Z}$, we have $\mathbb{Q}(rj^T, u) = \mathbb{Q}(r, u) = \mathbb{Q}(X_0(N))$. Hence when $T \geq 2g - 2$, the pair (rj^T, u) satisfies both assumptions of Proposition 2.1.2. We thus obtain

Theorem 2.1.7. *Let $T \geq 2g - 2$ be a positive integer and let*

$$P(x, y) = f_n(y)x^n + \cdots + f_1(y)x + f_0(y)$$

be a minimal polynomial relation of rj^T and u . Then there exist integers A, B and a nonzero constant c such that

$$F_{E,j}(y) = cf_0(1/y) \cdot y^A(y - 1728)^B.$$

The integers A and B are defined as follows. Let $\epsilon_i(N) = |\mathcal{E}_i(N)|$ for $i = 2$ or 3 and let $d_N = [SL_2(\mathbb{Z}) : \Gamma_0(N)]$, then $A = \deg f_n - T \cdot d_N - \frac{1}{3}(d_N + 2\epsilon_3(N))$ and $B = -\frac{1}{2}(d_N + \epsilon_2(N))$.

Proof. Write $\text{div}(\omega) = \sum_{[z] \in X_0(N)} n_z[z]$. Applying Proposition 2.1.2 to raj^T and u , we get

$$\prod_{z:u(z) \neq 0, \infty} (y - u(z))^{n_z} \cdot (y - 1/1728)^{\frac{1}{2}(d_N + \epsilon_2(N))} = cf_0(y) \quad (\text{a})$$

and

$$\sum_{z:u(z)=\infty} \text{ord}_z(\omega) + T \cdot d_N + \frac{1}{3}(d_N + 2\epsilon_3(N)) = \deg f_n - \deg f_0. \quad (\text{b})$$

To change from $u = \frac{1}{j}$ to j , we replace y by $1/y$ in (a) and multiply both sides by $y^{\deg f_0}$ to obtain

$$\prod_{z:j(z) \neq 0, \infty} (y - j(z))^{n_z} \cdot (y - 1728)^{\frac{1}{2}(d_N + \epsilon_2(N))} = cf_0(1/y)y^{\deg f_0}.$$

The contribution of $\{z \in \text{div}(\omega) : j(z) = 0\}$ to $F_{E,j}$ can be computed from (b), so

$$\begin{aligned} F_{E,j}(y) &= c \cdot y^{\deg f_n - \deg f_0 - T \cdot d_N - \frac{1}{3}(d_N + 2\epsilon_3(N))} y^{\deg f_0} \cdot (y - 1728)^{-\frac{1}{2}(d_N + \epsilon_2(N))} f_0(1/y) \\ &= c \cdot y^{\deg f_n - T \cdot d_N - \frac{1}{3}(d_N + 2\epsilon_3(N))} (y - 1728)^{-\frac{1}{2}(d_N + \epsilon_2(N))} f_0(1/y). \end{aligned}$$

□

Now we describe the algorithm *Poly Relation*.

Algorithm 1 *Poly relation*

Input: E = Elliptic Curve over \mathbb{Q} ; N = conductor of E ; f = the newform attached to E .

Values of $g = g(X_0(N))$, $d_N, \epsilon_2(N), \epsilon_3(N)$, and c_N = number of cusps of $X_0(N)$.

Output: The critical j -polynomial $F_{E,j}(x)$.

- 1: Fix a large integer M . $T := 2g - 2$.
 - 2: $r_1 := j^{2g-1}(j - 1728) \frac{f}{j}$, $u := \frac{1}{j}$.
 - 3: $\deg r_1 := (2g - 1)d_N - c_N$, $\deg u := d_N$.
 - 4: Compute the q -expansions of r_1 and u to q^M .
 - 5: Let $\{c_{a,b}\}_{0 \leq a \leq \deg u, 0 \leq b \leq \deg r_1}$ be unknowns, compute a vector that spans the one-dimensional vector space
 - 6: $K = \{(c_{a,b}) : \sum c_{a,b} r(q)^a u(q)^b \equiv 0 \pmod{q^M}\}$.
 - 7: $P(x, y) := \sum c_{a,b} x^a y^b$. Write $P(x, y) = f_n(y)x^n + \cdots + f_1(y)x + f_0(y)$.
 - 8: $A := \deg f_n - T \cdot d_N - \frac{1}{3}(d_N + 2\epsilon_3(N))$, $B := -\frac{1}{2}(d_N + \epsilon_2(N))$.
 - 9: Output $F_{E,j}(x) = cf_0(1/x) \cdot x^A(x - 1728)^B$.
-

Note that an upper bound on the number of terms M in the above algorithm can be taken to be $2 \deg r \deg u + 1$, by the following lemma.

Lemma 2.1.8. *Let $r, u \in \mathbb{Q}(X_0(N))$ be non-constant functions. If there is a polynomial $P \in \mathbb{Q}[x, y]$ such that $\deg_x P \leq \deg u$, $\deg_y P \leq \deg r$, and*

$$P(r, u) \equiv 0 \pmod{q^M}$$

for some $M > 2 \deg u \deg r$, then $P(r, u) = 0$.

Proof. Suppose $P(r, u)$ is non-constant as a rational function on $X_0(N)$, then $\deg P(r, u) \leq \deg r^{\deg u} u^{\deg r} = 2 \deg u \deg r$. It follows from $P(r, u) \equiv 0 \pmod{q^M}$ that $\text{ord}_{[\infty]} P(r, u) \geq M$. Since $M > 2 \deg u \deg r$, the number of zeros of $P(r, u)$ is greater than its number of poles, a contradiction. Thus $P(r, u)$ is a constant function. But then $P(r, u)$ must be 0 since it has a zero at $[\infty]$. This completes the proof. \square

Remark 2.1.9. When N is square free, there is a faster method that computes $F_{E,j}$ by computing the *Norm* of the modular form f , defined as $\text{Norm}(f) = \prod f|[A_i]$, where $\{A_i\}$ is a set of right coset representatives of $\Gamma_0(N)$ in $\text{SL}_2(\mathbb{Z})$. This approach is inspired by Ahrlgen and Ono [AO03], where j -polynomials of Weierstrass points on $X_0(p)$ are computed for p a prime.

Remark 2.1.10. In practice, in order to make the algorithm faster, we make different choices of r to make $\deg r$ small. Let η denote the Dedekind η -function and let $\Delta = \eta^{24}$ denote the discriminant modular form of level 1 and weight 12. When $4 \mid N$ we may take $r_4 = \frac{\omega j h_2}{d j (32 + h_4)}$, where $h_2 = \frac{\Delta(z) - 512\Delta(2z)}{\Delta(z) + 256\Delta(2z)}$ and $h_4 = (\eta(z)/\eta(4z))^8$. Then $\text{div}(r_4) = \text{div}(\omega) + D - D'$, where D and D' are supported on the cusps of $X_0(N)$, and $\deg D = c_N - \delta$, where δ is the number of cusps on $X_0(N)$ that are equivalent to $[\infty]$ modulo $\Gamma_0(4)$. Hence r_4 has a relatively small degree and is better suited for computation.

Remark 2.1.11. In order to speed up the computation, instead of taking $T = 2g - 2$ in the algorithm, we may take $T = 0$. First, if $\text{div}(\omega)$ does not contain cusps (for example, this happens if N is square free), then the functions r and u already satisfies the assumptions of Proposition 2.1.2. Second, if $\text{div}(\omega)$ does contain cusps, then $\deg(r)$ will be smaller than its set value in the algorithm, due to cancellation between zeros and poles. As a result, the vector space K will have dimension greater than 1. Nonetheless, using a basis of K , we could construct a set of polynomials $P_i(x, y)$ with $P_i(r, u) = 0$. Now $P(x, y)$ is the greatest common divisor of the $P_i(x, y)$.

We show a table of critical j -polynomials. Recall that $H_d(x)$ denotes the Hilbert class polynomial associated to a negative discriminant d . We use Cremona's labels [Cre] for elliptic curves in Table 2.1.1.

fixme: yang pair and the rest of ciritcal points paper

¹In this case $\text{div}(\omega) = [1/4] + [3/4] + [1/12] + [7/12]$ in supported on cusps.

Table 2.1.1: Critical polynomials for some elliptic curves with conductor smaller than 100

E	$g(X_0(N))$	Factorization of $F_{E,j}(x)$
37a	2	$H_{-148}(x)$
37b	2	$H_{-16}(x)^2$
44a	4	$H_{-44}(x)^2$
48a	3	1^1
67a	5	$x^8 + 1467499520383590415545083053760x^7 + \dots$
89a	7	$H_{-356}(x)$

Chapter 3

FOURIER EXPANSIONS OF MODULAR FORMS AT ALL CUSPS

Let k be a positive even integer and let $f \in S_k(\Gamma_0(N))$ be a nonzero cusp form. Then f has a Fourier expansion at the cusp infinity:

$$f = \sum_{n \geq 1} a_n q^n$$

where a_n are complex numbers and $q = e^{2\pi i \tau}$. We are concerned with the problem of computing the Fourier expansion of f at other cusps. When N is square-free, this problem is solved by Asai [Asa76]. The problem is studied in the Ph.D. thesis of Christophe Delaunay and in [Edixhoven], where a numerical algorithm is proposed. We will give a numerical algorithm to compute such expansions. Our approach is different from the one proposed in [Ed], for they require working at a higher level: to compute expansions at cusps of denominator Q , one needs to compute period matrices for forms of level NR^2 , where $R = \gcd(Q, \frac{N}{Q})$. As a contrast, our algorithm works at levels dividing N .

The main results of this chapter are Theorem 3.6.7 and Algorithm 4. The former gives a formula for the Fourier expansion of a newform $f \in S_k(\Gamma_0(N))$ at any cusp z of width one, and the latter describes how to use the formula to explicitly compute such expansion. Along the way, we will develop algorithms to compute the twists $f \otimes \chi$ and the pseudo-eigenvalue of newforms under the Fricke involution.

Section contains some examples.

3.1 Preliminaries

Let $N \geq 1$ be an integer and let $X_0(N)$ be the modular curve of level N .

Definition 3.1.1. Let z be a cusp on $X_0(N)$. If $z \neq \infty$, write $z = [a/c]$ with $\gcd(a, c) = 1$.

The *denominator* of z is

$$d_z = \gcd(c, N).$$

. If $z = \infty$, we set $d_\infty = N$. Choose $\alpha \in SL_2(\mathbb{Z})$ such that $\alpha(\infty) = z$. The *width* of z is

$$h_z = \left| \frac{SL_2(\mathbb{Z})_\infty}{(\alpha^{-1}\{\pm I\}\Gamma_0(N)\alpha)_\infty} \right|$$

where the subscript ∞ means taking the isotropy subgroup of ∞ in the corresponding group.

The width of a cusp can be computed in terms of its denominator. In fact, we have

Lemma 3.1.2. *If z is a cusp on $X_0(N)$, then*

$$h_z = \frac{N}{\gcd(d_z^2, N)}.$$

Proof. When $z = [\infty]$, we have $d_\infty = N$ and $h_\infty = 1$, so the formula holds trivially. Otherwise, write $z = [\frac{a}{c}]$ and find $\alpha = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in SL_2(\mathbb{Z})$. For $N' \in \mathbb{Z}$ we compute

$$\alpha \begin{pmatrix} 1 & N' \\ 0 & 1 \end{pmatrix} \alpha^{-1} = \begin{pmatrix} * & * \\ -c^2 N' & * \end{pmatrix}.$$

Hence $\begin{pmatrix} 1 & N' \\ 0 & 1 \end{pmatrix} \in (\alpha^{-1}\{\pm I\}\Gamma_0(N)\alpha)_\infty \iff N \mid c^2 N' \iff \frac{N}{\gcd(d_z^2, N)} \mid N'$. This completes the proof. \square

In particular, the width of a cusp z is one if and only if $N \mid d_z^2$.

Suppose f is a modular form on $\Gamma_0(N)$ of positive even weight k and $\alpha \in GL_2(\mathbb{Q})$. Recall the weight- k action is defined as

$$f|_\alpha(z) = (\det(\alpha))^{k/2} (cz + d)^{-k} f(\alpha z), \quad \alpha = \begin{pmatrix} a & b \\ c & d \end{pmatrix}.$$

In particular, if $\alpha \in SL_2(\mathbb{Z})$, then $f|_\alpha$ is a modular form on $\Gamma(N)$. So $f|_\alpha$ has a q -expansion, which is a power series in $q^{\frac{1}{N}}$. A natural thing to do is to define the expansion of f at the cusp z as the expansion of $f|_\alpha$. However, note that this may not be well-defined: in general the expansion depends on the choice of α . Nonetheless, when the cusp z has width one, the expansion is indeed well-defined as a power series in q .

Lemma 3.1.3. *Let z be a cusp on $X_0(N)$ with $h_z = 1$. Choose $\alpha \in SL_2(\mathbb{Z})$ such that $\alpha(\infty) = z$. Then $f|_\alpha$ is a cusp form on $\Gamma_1(N)$. Moreover, the function $f|_\alpha$ is independent of the choice of α .*

Proof. It is easy to verify that $\Gamma_1(N) \subseteq \alpha^{-1}\Gamma_0(N)\alpha$, hence the first claim holds. Now suppose $\beta \in SL_2(\mathbb{Z})$ is such that $\beta(\infty) = z$. Then $\alpha^{-1}\beta \in SL_2(\mathbb{Z})_\infty$. Since z has width one, we have $\alpha^{-1}\beta \in \alpha^{-1}\Gamma_0(N)\alpha$. Hence $\beta \in \Gamma_0(N)\alpha$, and it follows that $f|[\beta] = f|[\alpha]$. \square

In light of the lemma above, we define the q -expansion of f at a width one cusp z to be the q -expansion of $f|[\alpha]$, and denote it by f_z .

Assume further that f is an eigenform under the Atkin-Lehner operators. We will show that in order to compute the expansion of $f|[\alpha]$ for any $\alpha \in SL_2(\mathbb{Z})$, it suffices to do so for $\alpha = \begin{pmatrix} 1 & 0 \\ m & 1 \end{pmatrix}$, where $0 \leq m < N$ and $N \mid \gcd(m, N)^2$. In particular, it suffices to compute the expansions of f at a some cusps of width one.

Lemma 3.1.4. *For any $\alpha \in SL_2(\mathbb{Z})$, there exists a matrix $w_Q \in W_N$ and an upper triangular matrix $u \in GL_2(\mathbb{Q})$ such that $w_Q\alpha = \alpha'u$, where $\alpha' = \begin{pmatrix} a' & b' \\ c' & d' \end{pmatrix} \in SL_2(\mathbb{Z})$ satisfies $N \mid \gcd(N, c')^2$.*

Indeed, one may find Q using Lemma. Now $f|[\alpha] = f|[w_Q][w_Q\alpha] = f|[w_Q][\alpha'][u] = \lambda_Q(f)f[\alpha'][u] = \lambda_Q(f)f[\alpha'']|u|$, where α'' is of form $\begin{pmatrix} 1 & 0 \\ m & 1 \end{pmatrix}$. Note that for an upper triangular matrix $u = \begin{pmatrix} u_0 & u_1 \\ 0 & u_2 \end{pmatrix}$, we have $f[u](q) = f(q^{u_0/u_2}e^{2\pi i u_1/u_2})$.

3.2 Reducing to the case of newforms

The space $S_k(\Gamma_0(N))$ is spanned by elements of form $g(q^d)$, where g is newform of level $M \mid N$ and d is a divisor of $\frac{N}{M}$. Note that $g(q^d) = d^{-k/2}g\left(\frac{d}{0} \frac{0}{1}\right)$. For any $\alpha \in SL_2(\mathbb{Z})$, we can find $\alpha' \in SL_2(\mathbb{Z})$ and $u \in GL_2(\mathbb{Q})$ such that $\begin{pmatrix} d & 0 \\ 0 & 1 \end{pmatrix} \alpha = \alpha'u$. Hence to compute all expansions $f|[\alpha]$, it suffices to give an algorithm for newforms.

In the rest of this chapter, we will restrict ourselves to solving the following problem:

Problem 3.2.1. Let f be a normalized newform in $S_k(\Gamma_0(N))$ and z be a cusp on $X_0(N)$ of width one. Compute the q -expansion of f_z .

3.3 Twists of newforms

For $f \in S_k(\Gamma_1(N), \epsilon)$ a newform with expansion $f = \sum_n a_n(f)q^n$ and χ a Dirichlet character, the *twist* f_χ is a modular form with expansion $f_\chi(q) = \sum a_n(f)\chi(n)q^n$.

Lemma 3.3.1. [AWL78, Proposition 3.1] *Let $F \in S_k(\Gamma_1(N), \epsilon)$, where ϵ is a character of conductor N' . Let χ be a character modulo M . Put $\tilde{N} = \text{lcm}(N, N'M, M^2)$. Then $f_\chi \in S_k(\Gamma_1(\tilde{N}), \epsilon\chi^2)$.*

In particular, when ϵ is the trivial character and the conductor M of χ satisfies $M^2 \mid N$, we have $F_\chi \in S_k(\Gamma_1(N), \chi^2)$.

We write $f \otimes \chi$ for the unique newform such that $a_p(f \otimes \chi) = a_p(f_\chi)$ for all but finitely many primes p . From now, we refer to $f \otimes \chi$ as *the twist of f by χ* .

We quote two more results from [AWL78], which we will use extensively. First, we recall the definitions of U_d and B_d operators. For a modular form $f = \sum a_n q^n$ and a positive integer d , we put

$$f|U_d = \sum a_{nd} q^n, \quad f|B_d = \sum a_n q^{nd}.$$

It is easy to see that for any positive integers d, d' , we have U_d commutes with $B_{d'}$.

Lemma 3.3.2. [AWL78, Theorem 3.1] *Let $q \mid N$ and Q be the q -primary part of N . Write $N = QM$. Let F be a newform in $S_k(\Gamma_1(N), \epsilon)$ with $\text{cond}(\epsilon_Q) = q^\alpha, \alpha \geq 0$. Let χ be a character with conductor $q^\beta, \beta \geq 1$. Put $Q' = \max\{Q, q^{\alpha+\beta}, q^{2\beta}\}$. Then*

- (1) *For each prime $q' \mid M$, F_χ is not of level $Q'M/q$.*
- (2) *The exact level of F_χ is $Q'M$ provided (a) $\max\{q^{\alpha+\beta}, q^{2\beta}\} < Q$ if $Q' = Q$, or (b) $\text{cond}(\epsilon_Q \chi) = \max\{q^\alpha, q^\beta\}$ if $Q' > Q$.*

Lemma 3.3.3. [AWL78, Theorem 3.2] *Let $q \mid N$ and Q be the q -primary part of N . Write $N = QM$. Let χ be a character whose conductor equals a power of q . Let f be a newform in $S_k(\Gamma_1(N), \epsilon)$. Then $f \otimes \chi$ is a newform in $S_k(\Gamma_1(Q'M, \epsilon\chi^2))$, where Q' is a power of q . Moreover, we have*

$$f_\chi = f \otimes \chi - (f \otimes \chi)|U_q|B_q.$$

Since our goal is to compute expansions of newforms on $\Gamma_0(N)$, we will make the following assumptions: from now, unless otherwise noted, we assume f has trivial character, and that $\text{cond}(\chi)^2 \mid N$.

Next, we consider the problem of identifying the newform $f \otimes \chi$. This includes finding its level and its q -expansion to arbitrarily many terms. We will assume that we have an oracle which, given weight k and level N , computes the expansions of all newforms in $S_k(\Gamma_1(N))$ to arbitrarily many terms (for example, use the algorithm in [Steb]).

Now we proceed on how to recognise the level of $f \otimes \chi$ from the coefficients of f . One potential obstacle is that we do not know all Fourier coefficients of $f \otimes \chi$: we only know that $a_n(f \otimes \chi) = a_n(f)\chi(n)$ when $\gcd(n, N) = 1$. This can be overcome using a variant of Sturm's argument. First we prove a lemma.

Lemma 3.3.4. *Let $f \in S_k(N, \epsilon)$ be a normalized newform and q be any positive integer. Then $f|U_q|B_q \in S_k(Nq^2, \epsilon)$.*

Proof. It is a standard fact that for any integer $d \geq 1$, the map $f \mapsto f|B_d$ takes $S_k(N, \epsilon)$ to $S_k(Nd, \epsilon)$. To prove the lemma, we consider two separate cases. First, assume $q \nmid N$, then we have $T_q = U_q + q^{k-1}\epsilon(q)B_q$. By our assumption, we have $f|T_q = a_q(f)f$. Therefore, we have $f|U_q|B_q = f|(T_q - q^{k-1}\epsilon(q)B_q)|B_q = a_q(f)f|B_q - q^{k-1}\epsilon(q)f|B_q^2$. Hence $f|U_q|B_q \in S_k(Nq^2, \epsilon)$. Now assume $q \mid N$, so $U_q = T_q$. Hence $f|U_q|B_q = a_q(f)f|B_q \in S_k(Nq, \epsilon) \subseteq S_k(Nq^2, \epsilon)$. \square

The next proposition generalised the usual Sturm bound argument for modular forms.

Proposition 3.3.5. *Let g_1, g_2 be two normalised newforms of levels $N_1 \mid N_2$ and the same nebentypus character ϵ . Assume ϵ has prime power conductor $Q = q^\beta$ such that $Q^2 \mid N_1$. Let B be the Sturm bound for the congruence subgroup $\Gamma_1(Nq^2)$. Suppose*

$$a_n(g_1) = a_n(g_2), \text{ for all } 1 \leq n \leq B \text{ such that } \gcd(n, q) = 1.$$

Then $g_1 = g_2$.

Proof. Following [AWL78], we define the operator K_q on the space of modular forms by

$$g|K_q = g - g|U_q|B_q.$$

Then the assumption is equivalent to the statement that $\delta = (g_1 - g_2)|K_q$ has $a_n(\delta) = 0$ for all $1 \leq n \leq B$. Since $\delta \in S_k(Nq^2, \epsilon)$, Sturm's theorem implies $\delta = 0$. We then know from [DS06, Theorem 5.7.1] that $g_1 - g_2 \in S_k(N_2, \epsilon)^{old}$. Suppose $N_1 < N_2$, then g_1 is in the old subspace, hence so is g_2 , a contradiction. Therefore we must have $N_1 = N_2$. It follows that $g_1 - g_2 \in S_k(N_2, \epsilon)^{new}$, since g_1, g_2 are newforms. Since the new subspace and the old subspace intersect trivially, we must have $g_1 - g_2 = 0$. \square

Now we are ready to describe the algorithm.

Algorithm 2 Identifying $f \otimes \chi$

Input: k – a positive even integer; $f \in S_k(\Gamma_0(N))$ a normalized newform; χ a Dirichlet character of prime power conductor $Q = q^\beta$; $Q^2 \mid N$; B – a positive integer

Output: The level M_χ of $f \otimes \chi$ and the Fourier expansion of $f \otimes \chi$ up to q^B .

```

1: if  $Q = 1$  then
2:   return  $N$ .
3: end if
4:  $Q' := \text{cond}(\chi^2)$ ;  $N_0 := \frac{N}{q^{v_q(N)}}$ ;  $M_0 := Q'N_0$ ;  $t := \frac{N}{M_0} \in \mathbb{Z}$ .
5: for each positive divisor  $d$  of  $t$  do
6:   Set  $V_d := S_k(M_0d, \chi^2)$ .
7:   Compute a basis of newforms  $\{g_1^{(d)}, \dots, g_{s_d}^{(d)}\}$  of  $V_d$ .
8:   Set  $B_d :=$  the Sturm bound for  $\Gamma_1(M_0dq^2)$ .
9:   for  $1 \leq j \leq s_d$  do
10:    if  $a_n(g_j^{(d)}) = a_n(f)\chi(n)$  for all  $1 \leq n \leq B_d, \gcd(n, q) = 1$  then
11:      return  $M_0d$ .
12:    end if
13:   end for
14: end for
```

We give some sample computations applying the above algorithm.

Example 3.3.6. Let f be the normalised newform attached to the elliptic curve

$$E : y^2 + xy + y = x^3 - x - 2$$

of Cremona label **50a**. Then $f \otimes \chi$ is new of level 50 for all Dirichlet characters χ with modulus 5. In other words, f is 5-minimal.

As another example, we demonstrate a newform which is not p -minimal.

Example 3.3.7. Let f be the normalised newform attached to the elliptic curve

$$E : y^2 + xy = x^3 + x^2 - 25x - 111$$

of label **98a**. Let χ be the Dirichlet character modulo 7 defined by $\chi(3 \pmod{7}) = -1$. We found that $f \otimes \chi$ is a newform of level 14, with q -expansion

$$(f \otimes \chi)(q) = q - q^2 - 2q^3 + q^4 + 2q^6 + q^7 - q^8 + q^9 - 2q^{12} - 4q^{13} - q^{14} + O(q^{15}).$$

3.4 Pseudo-eigenvalues

Let ϵ be a Dirichlet character modulo N and let f be a newform in $S_k(N, \epsilon)$. For any divisor Q of N with $\gcd(Q, \frac{N}{Q}) = 1$, there is an algebraic number $w_Q(f)$ of absolute value one and a newform g in $S_k(N, \overline{\epsilon_Q} \epsilon_{N/Q})$ such that

$$W_Q(f) = w_Q(f)g,$$

Definition 3.4.1. The number $w_Q(f)$ is called the *pseudo-eigenvalue* of W_Q on f .

For ease of notations, we write $w(f) = w_N(f)$.

For a power series $f = \sum_{n \geq 0} a_n q^n$, its complex conjugate, denoted by f^* , is

$$f^*(q) = \sum \overline{a_n} q^n.$$

From [AWL78] we have $W_N(f) = w(f)f^*$. In the rest of this section, we describe an algorithm to efficiently compute $w(f)$ numerically. For a positive even integer k , let $\mathbb{M}(k)$ denote the space of weight- k modular symbols defined in [Steb]. The space $\mathbb{M}(k)$ is a quotient of $\mathbb{Z}[X, Y]_{k-2} \otimes \mathbb{P}^1(\mathbb{Q})^2$, and $GL_2(\mathbb{Q})$ acts on $\mathbb{M}(k)$ via the following rule

$$g(P(X, Y) \otimes \{\alpha, \beta\}) = P(g^{-1}(X, Y)^T) \{g(\alpha), g(\beta)\}.$$

Most importantly, there is a pairing between $\mathbb{M}(k)$ and the space of modular forms of weight k , defined as

$$\langle f, P(X, Y) \otimes \{\alpha, \beta\} \rangle_k = \int_{\alpha}^{\beta} f(z) P(z, 1) dz.$$

We will suppress the subscript k if its value is clear from context.

Lemma 3.4.2. *Let $M \in \mathbb{M}(k)$ and $f \in S_k(\Gamma_1(N))$. Then*

$$N^{\frac{k}{2}-1} \langle f | W_N, M \rangle = \langle f, W_N M \rangle.$$

Proof. See proof of [Steb, Proposition 8.17]. Note that the extra factor $N^{\frac{k}{2}-1}$ is due to the different constants involved in the definition of the weight- k action of $GL_2(\mathbb{Q})$ on modular forms. \square

The map

$$* : P(x, y) \{\alpha, \beta\} \mapsto P(-x, y) \{-\bar{\alpha}, -\bar{\beta}\}$$

defines the *star involution* on the space $\mathbb{M}(k)$. We have $\langle f^*, M \rangle = \overline{\langle f, M^* \rangle}$.

Lemma 3.4.3. *Let f be a normalised newform on $\Gamma_1(N)$ with positive even weight k and let $M \in \mathbb{M}(k)$ be such that $W_N(M) = N^{k/2-1} M^*$. Assume $\langle f, M \rangle \neq 0$. Then*

$$w(f) = \frac{\langle f, M \rangle}{\overline{\langle f, M \rangle}}.$$

Proof. Since $W_N^2(M) = N^{k-2}M$ for all $M \in \mathbb{M}(k)$, the assumption implies $W_N(M^*) = N^{k/2-1}M$. Now

$$\begin{aligned}
& N^{k/2-1}\langle f|W_N, M^*\rangle = \langle f, W_N(M^*)\rangle \\
& \implies N^{k/2-1}w(f)\langle f^*, M^*\rangle = N^{k/2-1}\langle f, M\rangle \\
& \implies w(f) = \frac{\langle f, M\rangle}{\langle f^*, M^*\rangle} \\
& \implies w(f) = \frac{\langle f, M\rangle}{\langle f, M\rangle}.
\end{aligned}$$

□

Suppose α, β are distinct points on the arc $\{z \in \mathbb{C} | \text{Im}(z) > 0, |z| = 1/\sqrt{N}\}$. Then it is easy to verify that $M = (xy)^{k/2-1} \otimes \{\alpha, \beta\}$ satisfies $W_N(M) = M^*$. Finally, we arrive at the algorithm to compute $w(f)$.

Algorithm 3 Computing the pseudo-eigenvalue of newforms.

Input: k – a positive even integer. $f \in S_k(\Gamma_1(N))$ a normalized newform.

Output: a numerical approximation of $w(f)$.

- 1: $n_0 := 10, z_0 := \frac{i}{\sqrt{N}}, \delta = 10^{-3}$.
 - 2: Randomly generate n_0 points $\{z_1, \dots, z_{N_0}\} \subseteq \{z | 0 < \text{Im}(z) < \frac{1}{2\sqrt{N}}, |z| = \frac{1}{\sqrt{N}}\}$.
 - 3: **for** $1 \leq i \leq n_0$ **do**
 - 4: compute the period integral $c_i = \int_{z_0}^{z_i} 2\pi i f(z) z^{\frac{k-2}{2}} dz$.
 - 5: $w_i \leftarrow c_i / \bar{c}_i$.
 - 6: **end for**
 - 7: **if** the standard deviation of w_1, \dots, w_{n_0} is less than δ **then**
 - 8: $w \leftarrow \frac{1}{n_0}(\sum_i w_i)$.
 - 9: **return** w .
 - 10: **else**
 - 11: **return** FAIL.
 - 12: **end if**
-

3.5 Formula for the Fourier expansion of f at width one cusps: Part 1

First we recall some notations from [AWL78].

Definition 3.5.1. For a positive integer c' , let $S'_c = \begin{pmatrix} 1 & \frac{1}{c'} \\ 0 & 1 \end{pmatrix}$. If χ is a character modulo c' , we define the operator on modular forms

$$f|R_\chi(c') = \sum_{u=0}^{c'-1} \bar{\chi}(u) f|S_{c'}^u.$$

Write R_χ in short for $R_\chi(\text{cond}(\chi))$. Note that $f|R_\chi = g(\bar{\chi})f_\chi$. Conversely, if $(a, M) = 1$, we have

$$\phi(c')S_{c'}^u = \sum_{\chi: \text{cond}(\chi)|c'} \chi(u) R_\chi(c'). \quad (3.5.1)$$

For our convenience, we define some operators, which are essentially the conjugates of S'_c and $R_\chi(c')$ by W_N . Let $A'_c = \begin{pmatrix} 1 & 0 \\ c' & 1 \end{pmatrix}$. Then it is easy to verify the following matrix identity.

Fact 3.5.2. $-N \cdot A_{N/c'}^{-1} = W_N S_{c'} W_N$.

From now on, we assume c is a divisor of N and $c' = \frac{N}{c}$. Then as operators on modular forms,

$$A_c^{-1} = W_N S_{c'} W_N.$$

Since $W_N^2 = id$ as operators, we have

$$A_c^{-u} = W_N S_{c'}^u W_N, \forall u \in \mathbb{Z}.$$

Parallel to the notion of $R_\chi(c')$, let $\Phi_\chi(c) = \sum_{u=0}^{c'-1} \bar{\chi}(u) A_c^{-u}$. Then $\Phi_\chi(c) = W_N R_\chi(c') W_N$. Similar to Formula 3.5.1, we have

$$\varphi(c') A_c^{-a} = \sum_{\text{cond}(\chi)|c'} \chi(a) \Phi_\chi(c) = \sum_{\text{cond}(\chi)|c'} \chi(a) W_N R_\chi(c') W_N. \quad (3.5.2)$$

Applying Formula 3.5.2 to f , we arrive at

$$f|_{[\frac{a}{c}]}(q) = \frac{1}{\varphi(c')} \sum_{\text{cond}(\chi)|c'} \chi(-a) f|W_N R_\chi(c') W_N. \quad (3.5.3)$$

$$= \frac{w(f)}{\varphi(c')} \sum_{\text{cond}(\chi)|c'} \chi(-a) f|R_\chi(c') W_N. \quad (3.5.4)$$

Now it left to compute the expansions of each $f|R_\chi(c') W_N$ in the sum.

3.6 Formula for the Fourier expansion of f at width one cusps: Part 2

In this section, we describe how to compute the expansion of $f|R_\chi(c')W_N$. First note that $T_p = U_p + \epsilon(p)p^{\frac{k}{2}}B_p$ as operators on $S_k(\Gamma_1(N), \epsilon)$. It follows that T_p commutes with B_d for any positive integer d .

We recall some notations and a result from [Del02].

Definition 3.6.1. [Del02, Definition III.2.4] For a Dirichlet character χ modulo $b = \prod_{j \in J} p_j^{\alpha_j}$. Let $r = |J|$. Decompose χ uniquely as $\chi = \chi_1 \cdots \chi_r$, where χ_i is a character modulo $p_j^{\alpha_j}$. We define $\text{cond}'(\chi)$ multiplicatively, by putting

$$\text{cond}'(\chi_j) = \begin{cases} \text{cond}(\chi_j) & \text{if } \text{cond}(\chi_j) > 1 \\ p_j & \text{else} \end{cases} \quad (3.6.1)$$

Also, if $I = \{j \in J : \chi_j \text{ is trivial character modulo } p_j^{\alpha_j}\}$, we put $tr = \prod_{j \in I} p_j^{\alpha_j}$ $nt = b/tr$, $\chi_{tr} = \prod_{j \in I} \chi_j$, and $\chi_{nt} = \chi/\chi_{tr}$. Then we set

$$g'(\chi) = (-1)^{|I|} \chi_{nt}(tr) g(\chi_{nt}). \quad (3.6.2)$$

Here $g(\chi)$ is the usual Gauss sum of χ : if χ is a character modulo d , then $g(\chi) = \sum_{a=1}^d e^{\frac{2\pi ia}{d}} \chi(a)$. If $\chi = \chi_0$ is the trivial character, we set $g(\chi_0) = 0$.

Lemma 3.6.2. [Del02, Prop 2.6] Let c' be an integer such that $c'^2 \mid N$. For a Dirichlet character χ mod c' , we have

$$f|R_\chi(c') = \begin{cases} g'(\bar{\chi}) f_{\chi_{nt}} & \text{if } \text{cond}'(\chi) = c' \\ 0 & \text{else.} \end{cases}$$

Using this lemma, we can simplify formula 3.5.3 to

$$f|_{\left[\frac{a}{c'}\right]} = \frac{w(f)}{\varphi(c')} \sum_{\text{cond}'(\chi)=c'} \chi(-a) g'(\bar{\chi}) f_{\chi_{nt}}|W_N. \quad (3.6.3)$$

Next, we compute $f_{\chi_{nt}}$ by the following: suppose $g = f \otimes \chi_{nt}$. Then

$$f_{\chi_{nt}} = g | \prod_{i=1}^r K_{p_i}. \quad (3.6.4)$$

Moreover, we have

$$K_p = 1 - U_p B_p = \begin{cases} 1 - (T_p - \chi_{nt}^2(p) p^{\frac{k}{2}} B_p) | B_p & p \nmid M \\ 1 - T_p | B_p & p \mid M \end{cases}. \quad (3.6.5)$$

Using the commutativity of T_* and B_* , we can write $f_{\chi_{nt}}$ in the form $\sum c_i(f \otimes \chi)(q^{d_i})$, where c_i and d_i are constants. To give a precise formula, we use the following notation. For a finite set S of integers, let $\pi(S) = \prod_{s \in S} s$ denote the product of all elements in S . For a Dirichlet character χ of conductor d , let S_χ be the set of prime divisors of d . For any positive integer M and any finite set of integers S , define

$$\mathcal{B}_{S,M} = \{(S_1, S_2) \in (2^\mathbb{Z})^2 \mid S_1, S_2 \subseteq S, S_1 \cap S_2 = \emptyset, \gcd(M, \pi(S_2)) = 1\} \quad (3.6.6)$$

Proposition 3.6.3. *Let $k \geq 2$ be an even integer and let f be a newform in $S_k(\Gamma_0(N))$.*

Then

$$f_{\chi_{nt}} = \sum_{(S_1, S_2) \in \mathcal{B}_{S_\chi, M}} (-1)^{|S_1|} a_{\pi(S_1)}(g_\chi) \pi(S_2)^{k/2} \chi_{nt}^2(\pi(S_2)) g_\chi | B_{\pi(S_1) \pi(S_2)^2}.$$

Here $g_\chi = f \otimes \chi$, M is the level of g_χ and $\mathcal{B}_{S_\chi, M}$ is as in 3.6.6.

Proof. This is a direct consequence of multiplying out 3.6.4 using 3.6.5, using the fact that T_p commutes with B_d , and noting that T_p acts as multiplication by $a_p(g_\chi)$ on g_χ . \square

Theorem 3.6.3 will be our starting point of computing the expansion of f at width one cusps. We will use it to compute $f_{\chi_{nt}} | W_N$. First we prove two lemmas.

Lemma 3.6.4. *Let f be a newform of even weight k on $\Gamma_1(M)$ and suppose d, N are positive integers such that $Md \mid N$. Then*

$$f | B_d | W_N = \left(\frac{N}{Md^2} \right)^{k/2} w(f) (f | B_{\frac{N}{Md}})^*.$$

Proof. Straightforward computation.

$$\begin{aligned}
f|B_d|W_N &= d^{-k/2} f| \begin{pmatrix} d & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 0 & -1 \\ N & 0 \end{pmatrix} \\
&= d^{-k/2} f| \begin{pmatrix} 0 & -1 \\ M & 0 \end{pmatrix} \begin{pmatrix} N/md & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} d & 0 \\ 0 & d \end{pmatrix} \\
&= \left(\frac{N}{Md^2} \right)^{k/2} f|W_M|B_{N/Md} \\
&= \left(\frac{N}{Md^2} \right)^{k/2} w(f)f^*|B_{N/Md} \\
&= \left(\frac{N}{Md^2} \right)^{k/2} w(f)(f|B_{N/Md})^*.
\end{aligned}$$

□

Before stating the second lemma, we quote another result in [Li75] on the coefficients of a newform at primes dividing the level.

Lemma 3.6.5. [Li75, Theorem 3 (iii)] *Let $f = \sum_{n \geq 1} a_n(f)q^n$ be a normalized newform in $S_k(\Gamma_1(N), \epsilon)$ and let p be a prime dividing N . Then*

- (1) *If ϵ is a character modulo N/p and $p^2 \mid N$, then $a_p(f) = 0$.*
- (2) *If ϵ is a character modulo N/p and $p^2 \nmid N$, then $a_p(f)^2 = \epsilon(p)p^{k-2}$.*
- (3) *If ϵ is not a character modulo N/p , then $|a_p(f)| = p^{\frac{k-1}{2}}$.*

Lemma 3.6.6. *Keep the notations in Proposition 3.6.3. If $(S_1, S_2) \in \mathcal{B}_{S_\chi, M_\chi}$ is such that $a_{\pi(S_1)}(g_\chi) \neq 0$. Then $M\pi(S_1)\pi(S_2)^2 \mid N$.*

Proof. Let p be a prime divisor of $N' := M\pi(S_1)\pi(S_2)^2$. If $p \nmid M$, then $\text{ord}_p(N') \leq \text{ord}_p(\text{cond}(\chi)^2) \leq \text{ord}_p(N)$. So we assume $p \mid M$, hence $p \nmid p(S_2)$. If $p \nmid p(S_1)$, then there's nothing to prove; if $p \mid \pi(S_1)$, we want to show that $\text{ord}_p(M) < \text{ord}_p(N)$. Suppose not, then $\text{ord}_p(M) = \text{ord}_p(N) \geq 2\text{ord}_p(\text{cond}(\chi))$. Since $\text{cond}(\chi^2) \leq \text{cond}(\chi)$, we know χ^2 is a character modulo M/p . Applying case (1) of Lemma 3.6.5 to the newform g_χ , we see that $a_p(g_\chi) = 0$, hence $a_{\pi(S_1)}(g_\chi) = 0$ by multiplicativity. □

Now we can state our main theorem from this chapter.

Theorem 3.6.7. *Let $k \geq 2$ be an even integer and let f be a normalized newform in $S_k(\Gamma_0(N))$. Let z be a cusp on $X_0(N)$ of width one. Write $z = [\frac{a}{d}]$ such that $\gcd(a, d) = 1$, $d \mid N$ and $N \mid d^2$. Let $d' = \frac{N}{d}$. Then the Fourier expansion of f at the cusp z is*

$$f_z(q) = \frac{w(f)}{\varphi(d')} \sum_{\chi: \text{cond}'(\chi)=d'} \chi(-a) g'(\bar{\chi}) w(f \otimes \chi) f_\chi^!(q).$$

Here

- $w(f)$ and $w(f \otimes \chi)$ are the pseudo-eigenvalues.
- $g'(\chi)$ is the modified Gauss sum defined in 3.6.2.
- cond' is the modified conductor of a Dirichlet character in 3.6.1.
- $f_\chi^!$ is as follows: let M_χ denote the level of $f \otimes \chi$. Then

$$f_\chi^! = \sum_{(S_1, S_2) \in \mathcal{B}_{S_{\chi_{nt}}, M_\chi}} (-1)^{|S_1|} a_{\pi(S_1)}(f \otimes \chi) \left(\frac{N}{M_\chi \pi(S_1)^2 \pi(S_2)^3} \right)^{k/2} \chi^2(\pi(S_2)) (f \otimes \chi | B_{\frac{N}{M_\chi \pi(S_1) \pi(S_2)^2}})^*$$

where the notations follow 3.6.3.

Proof. We start from formula 3.6.3:

$$f_{[\frac{a}{c}]} = \frac{w(f)}{\varphi(c')} \sum_{\text{cond}'(\chi)=c'} \chi(-a) g'(\bar{\chi}) f_{\chi_{nt}} | W_N.$$

From 3.6.3, we have

$$f_{\chi_{nt}} = \sum_{(S_1, S_2) \in \mathcal{B}_{S_\chi, M_\chi}} (-1)^{|S_1|} a_{\pi(S_1)}(f \otimes \chi) \pi(S_2)^{k/2} \chi_{nt}^2(\pi(S_2)) f \otimes \chi | B_{\pi(S_1) \pi(S_2)^2}.$$

To simplify notations, let $c(f, \chi, S_1, S_2) = (-1)^{|S_1|} a_{\pi(S_1)}(f \otimes \chi) \pi(S_2)^{k/2} \chi_{nt}^2(\pi(S_2))$. Then

$$\begin{aligned} f_{\chi_{nt}} | W_N &= \sum_{(S_1, S_2) \in \mathcal{B}_{S_\chi, M_\chi}} c(f, \chi, S_1, S_2) f \otimes \chi | B_{\pi(S_1) \pi(S_2)^2} W_N \\ &= \sum_{(S_1, S_2) \in \mathcal{B}_{S_\chi, M_\chi}} c(f, \chi, S_1, S_2) \left(\frac{N}{M_\chi (\pi(S_1) \pi(S_2)^2)^2} \right)^{k/2} w(f \otimes \chi) (f \otimes \chi | B_{\frac{N}{M_\chi \pi(S_1) \pi(S_2)^2}})^* \\ &= w(f \otimes \chi) f_\chi^!. \end{aligned}$$

Note that we applied Lemma 3.6.4 to obtain the penultimate equality, and we could do that because of Lemma 3.6.6. Now the result follows. \square

Theorem 3.6.7 gives us an algorithm to compute the expansion of f_z , which we will describe below. But first, we take a closer look at what ingredients goes into the expansion. Given a newform $f \in S_k(\Gamma_0(N))$ and a width one cusp z of denominator c . We need to consider the twist of f by all Dirichlet characters of conductor dividing c . For each such character χ , we then need to determine the level M_χ and q -expansion of the newform $f \otimes \chi$, the latter boils down to knowing $a_p(f \otimes \chi)$ for all primes $p \mid \text{cond}(\chi)$. Then we need to compute the pseudo-eigenvalues of $f \otimes \chi$. Finally, we combine these information together and apply Theorem 3.6.7 to compute f_z .

Algorithm 4 Computing Fourier coefficients of f at width one cusps

Input: $f \in S_k(\Gamma_0(N))$ a newform; a, c – coprime integers such that $N \mid c^2$; B – a positive integer.

Output: The first B Fourier coefficients of $f_{[\frac{a}{c}]}(q)$.

- 1: $c' \leftarrow N/c$. $X \leftarrow$ The set of all Dirichlet characters χ such that $\text{cond}'(\chi) = c'$.
 - 2: compute $w(f)$ using Algorithm 3.
 - 3: **for** χ in X **do**
 - 4: Using Algorithm 2, compute the level M_χ and the q -expansion of $g_\chi := f \otimes \chi$ to B terms.
 - 5: Compute $w(g_\chi)$ using Algorithm 3.
 - 6: **end for**
 - 7: Apply Theorem 3.6.7 to compute f_z to B terms.
-

3.7 A Converse Theorem

Given the work in previous sections, it is a natural question then to ask whether the information on twists of f is uniquely determined by the expansion of f at width one cusps. The answer is yes, and the precise statement is in the following theorem.

Theorem 3.7.1. *Let f be a normalized newform in $S_k(\Gamma_0(N))$. Assume the eigenvalue $w_N(f)$ is known. Suppose c is a positive divisor of N such that $N \mid c^2$. Then the expansions of f_z , where z runs through all cusps of denominator c , uniquely determines the following: for each Dirichlet character χ of such that $\text{cond}'(\chi) = c'$, the level M_χ , the pseudo-eigenvalue w_{M_χ} and the q -expansion of the newform $f \otimes \chi$.*

Proof. By plug in different a 's. We can solve for t_χ . Consider the first nonzero term of t_χ . Suppose

$$t_\chi = u_\chi q^{v_\chi} + O(q^{v_\chi+1}), \quad u_\chi \neq 0.$$

Assuming that χ has prime power conductor $p^\beta > 1$, we claim that

$$\left| \frac{v^{k/2}}{u} \right| = \begin{cases} p^{k/2} & \text{if } p \nmid M_\chi \\ p^{1/2} & \text{if } p \mid M_\chi \text{ and } a_p(g) \neq 0 \\ 1 & \text{else} \end{cases}$$

Proof of claim: the first and third case are easy to verify using Theorem 3.6.7. Now assume $p \mid M$ and $a_p(g_\chi) \neq 0$. By Lemma 3.6.5, we have $|a_p(g_\chi)| = p^{k/2-1/2}$ or $p^{k/2-1}$. However, $|a_p(g_\chi)| = p^{k/2-1}$ only if $p \parallel M_\chi$ and χ^2 is a character modulo M_χ/p . This means χ^2 is the trivial character. By Lemma 3.3.2, we compute the p -level of $f = g_\chi \otimes \bar{\chi}$: note that $\max p, p^{\alpha+\beta}, p^{2\beta} > p$, so (ii) applies and the p -level of f is equal to $\max(p^\alpha, p^\beta) = p^\beta$, i.e., $\text{ord}_p(N) = \beta$. This is impossible since we have $p^{2\beta} = \text{cond}(\chi)^2 \mid N$.

Therefore, we have $|a_p(g_\chi)| = p^{k/2-1/2}$ and the claim follows.

Since $k \geq 2$, we could determine which case we are in. Then we can read off M_χ and $w_M(g_\chi)$. For example, if we are in the second case, then the level can be computed via $M_\chi = \frac{N}{v_\chi p}$. Now the N/M_χ 's coefficient of t_χ is

$$\begin{aligned} a_{\frac{N}{M}}(t_\chi) &= w(g_\chi) \left(\frac{N}{M}\right)^{k/2} (1 - |a_p(g_\chi)|^2 \chi^2(p) p^{-k/2}) \\ &= w(g_\chi) \left(\frac{N}{M}\right)^{k/2} (1 - p^{k/2-1} \chi^2(p)). \end{aligned}$$

This allows us to solve $w(g_\chi)$. Finally, we compute $a_p(g_\chi)$ by $a_p(g) = \frac{-u_\chi}{w(g_\chi) \chi^2(p) \left(\frac{N}{Mp}\right)^{k/2}}$. The

value $a_p(g)$ determines the expansion of g_χ . Recursively, we could solve for all pn -coefficients of g_χ , from which we deduce its complete q -expansion.

In the general case, we consider the following subsets of S_χ . Let $S_1^* = \{p \in S_\chi : p \mid M\}$, $S_2^* = S_\chi \setminus S_1^*$, and $\widetilde{S}_1^* = \{p \in S_1^* : a_p(g_\chi) \neq 0\}$.

It follows that the leading term of t_χ belongs to the summand corresponding to $(\widetilde{S}_1^*, S_2^*)$ in Theorem 3.6.7. Still writing the leading term as $u_\chi q^{v_\chi}$, we have

$$u_\chi = w(g_\chi) \chi^2(p(S_2)) a_{p(\widetilde{S}_1^*)}(g_\chi) p(\widetilde{S}_1^*)^{-k} (p(S_2^*))^{-3k/2} \left(\frac{N}{M_\chi} \right)^{k/2}, \quad v_\chi = \frac{N}{M_\chi p(\widetilde{S}_1^*) p(S_2^*)^2}.$$

Similar to the prime power conductor case above, we have $|a_{p(\widetilde{S}_1^*)}(g_\chi)| = p(\widetilde{S}_1^*)^{k/2-1/2}$. So

$$|v_\chi^k u_\chi^{-2}| = p(\widetilde{S}_1^*) p(S_2^*)^2. \quad (3.7.1)$$

Hence we can factor $|v_\chi^k u_\chi^{-2}|$ and obtain $p(\widetilde{S}_1^*)$ and $p(S_2^*)$. Then M_χ can be solved using v_χ . Plug it back into u_χ , we obtain $a_{p(\widetilde{S}_1^*)} w(g_\chi)$. Finally, for each $p \in \widetilde{S}_1^*$, the $v_\chi p$'s coefficient of t_χ allows us to compute $a_{p(\widetilde{S}_1^*)/p}(g_\chi) w(g_\chi)$. These together determine $w(g_\chi)$ and $a_{p(\widetilde{S}_1^*)}$. The other Fourier coefficients of g_χ can then be computed recursively. \square

3.8 Fields of definitions

In the previous sections, we have described an algorithm to compute the Fourier coefficients of f_z . In fact, the Fourier coefficients are algebraic numbers. More precisely, if c is the denominator of z and $c' = N/c$, then $f_z(q) \in K_f(\zeta_{c'})[[q]]$. Here K_f is the number field generated by the Fourier coefficients of f (at the cusp ∞). Although this result is well-known, we include a proof for the reader's convenience.

Lemma 3.8.1. *Let c be a cusp of denominator d and let $d' = N/d$. Then*

$$\mathbb{Q}(\{a_n(f, c)\}) \subseteq \mathbb{Q}(\{a_n(f)\}, \zeta_{d'}).$$

(fixme: add proof), maybe Diamond-Im [DI95].

3.9 Denominators

3.10 Examples

Let $E = \mathbf{50a}$ and consider the 4 cusps of denominator 10 on $X_0(50)$. The corresponding first terms of q -expansions at these cusps are

$$\begin{aligned} a_1(f, \frac{1}{10}) &= \frac{1}{5}\zeta_5^3 - \frac{3}{5}\zeta_5^2 + \frac{3}{5}\zeta_5 - \frac{1}{5} \\ a_1(f, \frac{3}{10}) &= \frac{3}{5}\zeta_5^3 + \frac{6}{5}\zeta_5^2 + \frac{4}{5}\zeta_5 + \frac{2}{5} \\ a_1(f, \frac{7}{10}) &= \frac{2}{5}\zeta_5^3 - \frac{1}{5}\zeta_5^2 - \frac{4}{5}\zeta_5 - \frac{2}{5} \\ a_1(f, \frac{9}{10}) &= -\frac{6}{5}\zeta_5^3 - \frac{2}{5}\zeta_5^2 - \frac{3}{5}\zeta_5 - \frac{4}{5} \end{aligned}$$

As another examples, let $E = \mathbf{98a}$ and $z = [\frac{1}{14}]$. We computed numerically that

$$\begin{aligned} f_z(q) &= (-0.755001687308946 - 0.172324208281817i)q + (0.441471704846525 - 0.916725441095080i)q^2 \\ &\quad + (1.39294678431094 + 1.11083799261729i)q^3 + (0.696473392155471 - 0.555418996308649i)q^4 \\ &\quad + (1.51000337461789 - 0.344648416563641i)q^6 + (-3.80647894157196 \times 10^{-16} - 3.02371578407382i)q^7 \\ &\quad + (0.755001687308946 + 0.172324208281817i)q^8 + (-0.441471704846525 + 0.916725441095080i)q^9 \\ &\quad + (-0.882943409693050 - 1.83345088219016i)q^{12} + (-3.02000674923578 + 0.689296833127282i)q^{13} \\ &\quad + (3.80647894157196 \times 10^{-16} + 3.02371578407382i)q^{14} + O(q^{15}) \end{aligned}$$

3.11 Applications

One applications of the computation done in this chapter is the norm method to the computation of j -polynomials introduced in Chapter . Recall that the issue with the norm method for non-square free level is computing the expansions of form $f|_\gamma$, where γ runs over the set of right coset representatives of $\Gamma_0(N)$ in $SL_2(\mathbb{Z})$. As we have seen, it suffices to compute the expansions of f at all width one cusps.

3.12 Automorphic representations; norm of first terms

References: [BH06], [LW10]. [Bru12]. [Kra90]. [JL72].

In this section, we will restrict ourselves to the case when the Fourier coefficients of f are rational numbers. Then f induces an admissible representation π_f of $GL_2(\mathbb{A}_{\mathbb{Q}})$. We will see that the expansion of f at all cusps can also be computed from the local component $\pi_{f,p}$. Loeffler and Weinstein gave an algorithm to compute such local components.

We will restrict ourselves to the simplest case when f is twist-minimal, which means that the conductor of π_f is the smallest among all twists $\pi_{f \otimes \chi}$.

We will follow the notations of David Loeffler and use the formula of [Bru12].

Okay, what is my heuristics for general k ? What is it for $\Gamma_1(N)$? What happens on the automorphic side?

Also there's the question about normalization, which was never specified.

Raw data?

Let z be a width one cusp of denominator c . Then the first coefficient $a_1(f_z)$ is an element in $K_f(\zeta_{c'})$. For simplicity, we assume that $c' = p^\alpha$ is a prime power. It can be proved using automorphic representations + local langlands correspondence that there exists β such that $p^\beta a_1(f_z) \in \bar{\mathbb{Z}}$. One question is: what prime ideals appears in the prime factorisation of $(a_1(f, z))$? It seems from our numerical data, that

$$\text{ord}_{\mathfrak{q}}(a_1(f_z)) > 0 \implies \mathfrak{q} \cap \mathbb{Z} \equiv \pm 1 \pmod{p}.$$

The following is a table of data.

(fix: add table)

3.12.1 Cuspidal local constants

We keep the assumptions that f is a newform attached to an elliptic curve E/\mathbb{Q} and f is twist-minimal. Assume p is a prime dividing the conductor N of E such that $v_p(N) = 2$. Then there exists a character $\varphi : \mathbb{F}_{p^2}^\times \rightarrow \mathbb{C}^\times$ which determines $\pi_{f,p}$. We will prove

Lemma 3.12.1. *Let $\psi : \mathbb{Q}_p \rightarrow \mathbb{C}^\times$ be a character of level one (e.g. $\psi(x) = e(\{\frac{x}{p}\}_p)$). Then*

$$\epsilon(\pi_{f,p}, 1/2, \psi) = \frac{-1}{p} \sum_{x \in \mathbb{F}_{p^2}^\times} \psi(x + x^p) \varphi(x).$$

If χ is a Dirichlet character such that the $f \otimes \chi$ has the same level as f . Then

$$\epsilon(\pi_{f \otimes \chi, p}, 1/2, \psi) = \frac{-1}{p} \sum_{x \in \mathbb{F}_{p^2}^\times} \psi(x + x^p) \varphi(x) \bar{\chi}(x^{p+1}).$$

Proof. By [BH06], taking $n = r = 1$, we have

$$p^2 \epsilon(\pi_{f, p}, 1/2, \psi) \cdot \text{id} = \sum_{x \in GL_2(\mathbb{F}_p)} \psi(\text{tr}(x)) \pi_{f, p}^\vee(x). \quad (3.12.1)$$

where $\pi_{f, p}^\vee$ denotes the contragredient representation. The representation $\pi_{f, p}$ has dimension $(p - 1)$. Taking traces, we obtain

$$p^2(p - 1) \epsilon(\pi_{f, p}, 1/2, \psi) \cdot \text{id} = \sum_{x \in GL_2(\mathbb{F}_p)} \psi(\text{tr}(x)) \text{Tr}(\pi_{f, p}^\vee(x)). \quad (3.12.2)$$

By assumption, $\pi_{f, p}$ arises from a cupsidal representation of the finite group $GL_2(\mathbb{F}_p)$, which is in turn induced from φ . (See Fulton-Harris), we have formulae for $\text{Tr}(\pi_{f, p}^\vee(x))$. Splitting the sum corresponding to four types of conjugacy classes, we computed $S_1 = (p - 1) \sum_{x \in \mathbb{F}_p^\times} \psi(2x)$, $S_2 = (p^2 - 1) \sum_{x \in \mathbb{F}_p^\times} \psi(2x)(-1)$, $S_3 = 0$, and $S_4 = (p^2 - p)/2 \sum_{x \in \mathbb{F}_{p^2} \setminus \mathbb{F}_p} \psi(\text{tr}(x)) (\overline{\varphi(x) + \varphi(x^p)})$. So the sum on the right hand side of 3.12.2 equals $(p - p^2) \sum_{x \in \mathbb{F}_{p^2}^\times} \psi(\text{tr}(x)) \overline{\varphi(x)}$. Dividing by $p^2(p - 1)$ gives the formula.

□

Moreover, since E is defined over \mathbb{Q} , the character of $\pi_{f, p}$ takes rational values. Hence the order of φ is 3, 4 or 6. The local Langlands correspondence claims that the order of φ is equal to the order of the inertia subgroup of $\text{Gal}(L/\mathbb{Q})$, where L is the smallest number field over which E acquires good reduction (to-do: check this). For $p = 2$ or 3 , the order of φ can be determined using results of [Kra90]. The case $p \geq 5$ is easy, as we have the following lemma:

Lemma 3.12.2. [Kra90, Proposition 1] *Let Δ denote the minimal discriminant of E . Then for $p \geq 5$, the order of φ is equal to $\frac{12}{\gcd(12, v_p(\Delta))}$.*

We remark that for elliptic curves, $v_2(N)$ is at most 8 and $v_3(N)$ is at most 5. For the sake of simplicity, we do not treat the case when $v_p(N) > 2$ here, but we point out the local constants can be also computed from formula in [BH06], once the local component is determined using [LW10].

Example 3.12.3. An example with trivial central character. Let f be the newform attached to $E = \mathbf{121a}$. Using Sage, we computed $w(f) = -1$. Since the weight of f is 2, we know $\epsilon_\infty = -1$ (since the central character of π_f is trivial, the level of the additive character ψ_∞ does not matter). The discriminant of E is $\Delta = -121$, so φ has order 6. Using Lemma 3.12.1, we computed that $\epsilon_{11}(\pi_{f,11}, 1/2) = -1$. This verifies $w(f) = -\prod_{p \leq \infty} \epsilon_p$.

Example 3.12.4. We give an example with nontrivial central character. Let f be as in the previous example, and let χ be the Dirichlet character of \mathbb{F}_{11}^\times defined by $\chi(2) = e^{2\pi i/10}$. Lemma 3.12.1 gives

$$\epsilon_{11}(\pi_{f \otimes \chi, 11}, 1/2) = 0.64.. + 0.76..i$$

an algebraic number with minimal polynomial $x^{20} + 109/121x^{15} + 2861/1331x^{10} + 109/121x^5 + 1$. So $w = -\epsilon_{11}\epsilon_\infty = \epsilon_{11}$. Using the numerical algorithm 3, we compute $w(f \otimes \chi) = 0.642573377564283 + 0.766224154177894i$. This confirms the computation.

3.13 Norm of first terms computations

We keep the assumptions from the previous section, that f is a newform in $S_2(\Gamma_0(N))$, attached to an elliptic curve E/\mathbb{Q} . We assume f is twist-minimal and $p \geq 5$ is a prime dividing the conductor N such that $v_p(N) = 2$. In this case, the cusp $z_p = [\frac{-p}{N}]$ is of width one, and the q -expansion of f at z_p takes an especially simple form. We summarize this in the lemma below.

Lemma 3.13.1. *With the assumptions above, there exists a Galois-invariant set of numbers $\{b_1, \dots, b_{p-1}\} \subseteq \mathbb{Q}(\zeta_p)$, such that*

$$f_{z_p}(q) = \sum_{n \geq 1} a_n(f) b_n \pmod{pq^n}.$$

More precisely, the b_j are given by

$$b_j = w(f) \sum_{\chi: \text{cond}(\chi)=p} g(\bar{\chi}) w(f \otimes \chi) \chi(n)$$

Proof. First, the assumptions imply that $a_n(f) = 0$ if $p \mid n$. So the right hand side of the formula is well-defined. The formulae then follow directly from Theorem 3.6.7. We have $b_j \in \mathbb{Q}(\zeta_p)$ since the cusp z_p is defined over $\mathbb{Q}(\zeta_p)$. (fixme: check this). Moreover, the cusps $\{z_p^{(j)} = \frac{-jp}{N} : 1 \leq j \leq p-1\}$ form a Galois orbit on $X_0(N)$, and one has

$$a_n(f_{z_p^{(j)}}) = a_{jn}(f_{z_p}), \forall n \geq 1, 1 \leq j \leq p-1.$$

In particular, we have $\{b_j\} = \{a_1(f_{z_p^{(j)}})\}$. Since the latter set is Galois-invariant, so is the former. \square

We remark that it is clear from the formula of b_j that they are algebraic number. However, the formula does not imply directly that they lie in $\mathbb{Q}(\zeta_p)$.

We give another formula of $a_1(f_{z_p})$ in light of the previous section.

Lemma 3.13.2. *Keeping the assumptions in the previous two sections, we have*

$$a_1(f_{z_p}) = \frac{\sum_{x \in \mathbb{F}_{p^2}^\times} \psi(x + x^p + x^{p+1}) \varphi(x)}{\sum_{x \in \mathbb{F}_{p^2}^\times} \psi(x + x^p) \varphi(x)}.$$

Example 3.13.3. Let f be the newform attached to $E = \mathbf{49a}$. One checks that f is twist-minimal and φ has order 4. Using Lemma 3.13.2, we computed

$$a_1(f_{-1/7}) = -\frac{5}{7}\zeta_7^5 - \frac{3}{7}\zeta_7^4 - \frac{1}{7}\zeta_7^3 + \frac{1}{7}\zeta_7^2 + \frac{3}{7}\zeta_7 - \frac{2}{7} = 0.623489... + 1.29468...i.$$

The numerical algorithm gives $a_1(f_{-1/7}) = 0.623489801858733... + 1.29468991410431...i$. Hence our formulae are consistent for this example.

It is of interest to determine the factorization of $a_1(f_{z_p})$ as a principal fractional ideal in $\mathbb{Q}(\zeta_p)$. For example, they relate to critical points of modular parametrization of E in the following way:

Lemma 3.13.4. *Let \mathfrak{q} be a prime ideal in $\mathbb{Q}(\zeta_p)$ lying above a prime $q \neq p$, such that $\mathfrak{q} \mid a_1(f_{z_p})$ and $q \nmid a_1(f_{z_p})$. Then $F_{E,j}(x)$ is not integral at q .*

Proof. fixme: add proof. □

From the above discussion, we see that there are at most three possibilities for each p , corresponding to the order of φ being 3, 4 or 6.

Consider $N_{f,p} = \text{Norm}(a_1(f_{z_p})) \in \mathbb{Z}$. It is easy to show that we always have $p \mid N_{f,p}$. The following is a table of the prime divisors of the norm, when such primes exist.

As an observation, we found that the primes l in the third column of the above table all satisfy a congruence relation

$$l \equiv \pm 1 \pmod{p}.$$

It would be interesting to prove or disprove this in general.

Table 3.13.1: table of prime divisors $l \neq p$ of $N_{f,p}$

p	order of φ	primes $l \neq p, l \mid N_{f,p}$
17	3	509
19	4	37
23	3	1103
23	4	47
29	3	173
31	4	557
41	3	1209, 9103
41	6	163
43	4	4129
47	3	13034039
47	4	2819
53	3	107, 317, 8161
53	6	107
59	3	1061, 537173407
59	4	827, 42953
67	4	2143, 10853
71	3	634532719903
71	4	6613947917
71	6	3407

Chapter 4

INDEX OF CHOW-HEEGNER POINTS

We consider a special case of the Chow-Heegner points that has a simple description due to Shouwu Zhang. Let E, F be nonisogenous elliptic curves defined over \mathbb{Q} of the same conductor N . The Chow-Heegner point $P_{E,F} \in E(\mathbb{Q})$ is constructed by the following procedure: take any point on $F(\mathbb{C})$, take its inverse image on $X_0(N)$, then map that image down to E and take the sum the resulting points. In [?], Darmon, Daub, Lichtenstein and Rotger developed an algorithm to compute Chow-Heegner points via iterated integrals. In [?], Stein developed a fast and conceptually easy algorithm to numerically compute Chow-Heegner points. The following theorem is proved by Yuan-Zhang-Zhang in [?]:

Theorem 4.0.1 (Yuan-Zhang-Zhang). *Let $L(E, F, F, s) = L(E, s)L(E, \text{Sym}^2(F), s)$. Assume that the local root numbers of $L(E, F, F, s)$ at every prime of bad reduction is $+1$ and that the root number at infinity is -1 . Then*

$$\hat{h}(P_{E,F}) = (\star) \cdot L'(E, F, F, \frac{1}{2}),$$

where (\star) is nonzero.

In particular, when the analytic rank of E is at least two, the Chow-Heegner point $P_{E,F}$ is torsion. When the rank of $E(\mathbb{Q})$ is one, we consider the index $i_{E,F} = [E(\mathbb{Q})/\text{tors} : \mathbb{Z}P_{E,F}]$. Theorem 4.0.1 combined with the Bloch-Kato conjecture on critical values of motivic L -functions suggests that this index might be linked to interesting arithmetic invariants related to E and F .

Numerical evidence in [?] suggests that the index $i_{E,F}$ is always divisible by 2, when it is finite. I proved the following theorem.

Theorem 4.0.2. *Let $\sigma_0(N)$ denote the number of distinct prime factors of N . If*

$$\sigma_0(N) > \log_2(\#E[2](\mathbb{Q})) + \log_2(\#F[2](\mathbb{Q})) + 2,$$

then $P_{E,F} \in 2E(\mathbb{Q})$. Hence the index $i_{E,F}$ is divisible by 2, if it is finite.

I prove the theorem in Section 5.2. In Section 5.3, I develop an exact algorithm to compute the Chow-Heegner point, using the methods in Chapter 2.

fixme: add all possible references: [CE], [Darmon], [Daub].

4.1 Definitions

We recall the definition from [?]. Consider a pair E, F of nonisogenous optimal elliptic curves over \mathbb{Q} of the same conductor N and fix modular parametrizations from $X_0(N)$ to both curves.

Let $(\varphi_E)_*$ and $(\varphi_F)^*$ denote the pushforward and pullback map on divisors. Let $Q \in F(\mathbb{C})$ be any point, we define

$$P_{E,F,Q} = \sum (\varphi_E)_*(\varphi_F)^*(Q),$$

where \sum means the sum of the points in the divisor, using the group law on E . By [?, Proposition 1.1], $P_{E,F,Q}$ is independent of the choice of Q . Let $P_{E,F} = P_{E,F,Q}$ for any choice of Q . Since we may choose $Q = \mathcal{O} \in \mathbb{F}(\mathbb{Q})$, it follows that $P_{E,F} \in E(\mathbb{Q})$.

4.2 The index

We make the additional assumption $r_{an}(E) = 1$. We consider the index

$$i_{E,F} = [E(\mathbb{Q})/tors : \mathbb{Z}P_{E,F}].$$

Lemma 4.2.1. *Let $F : y^2 = x^3 + Ax + B$ be an elliptic curve with $A, B \in \mathbb{Z}$. Let $f_t(x) = x^3 + Ax + B - t^2$. Then the set*

$$\{K_t = \mathbb{Q}[x]/(f_t(x)) : t \in \mathbb{Z}\}$$

contains infinitely many cubic fields.

Proof. Note there are finitely many t such that $f_t(x)$ is reducible. So we may forget about this possibility. Now suppose for contradiction that the claim is false, then there exists a cubic field K such that $K = K_t$ for infinitely many $t \in \mathbb{Z}$. However, $K = K_t$ implies $E(K)$ contains the point $P_t = (\alpha_t, t)$, where α_t is a root of $f_t(x)$. Moreover, the point P_t lies in $E(\mathcal{O}_K)$. Hence we must have that $E(\mathcal{O}_K)$ is an infinite set. This is a contradiction, since $E(\mathcal{O}_K)$ is finite by Siegel's theorem. \square

We quote a lemma of Calegari and Emerton [?].

Lemma 4.2.2. *Let E/k be an elliptic curve and let A be the group of automorphisms of E as a curve over k . Suppose W is a finite elementary abelian 2-subgroup of A . Then the order of W divides twice the order of $E[2](k)$.*

Theorem 4.2.3. *Let E, F be elliptic curves defined over \mathbb{Q} , with the same conductor N . Let $\sigma_0(N)$ denote the number of distinct prime factors of N . If*

$$\sigma_0(N) > \log_2(|E(\mathbb{Q})[2]|) + \log_2(|F(\mathbb{Q})[2]|) + 2,$$

then $P_{E,F} \in 2E(\mathbb{Q})$. In particular, if $\sigma_0(N) \geq 7$, then the condition holds automatically, and $P_{E,F} \in 2E(\mathbb{Q})$.

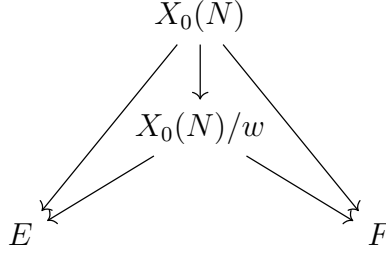
Proof. Consider the group \mathcal{W} of Atkin-Lehner involutions on $X_0(N)$. This group is elementary 2-abelian, and it descends to automorphisms on F and automorphisms on E , as curves. So we have a map

$$\pi : \mathcal{W} \rightarrow \text{Aut}(E) \times \text{Aut}(F)$$

By the Lemma above, we have $\text{im}(p_1 \circ \pi) \leq 2|E[2](\mathbb{Q})|$ and $\text{im}(p_2 \circ \pi) \leq 2|F[2](\mathbb{Q})|$. Hence the size of the image of π is bounded above by $4|E(\mathbb{Q})[2]| \cdot |F(\mathbb{Q})[2]|$. But we also know that

$$|\mathcal{W}| = 2^{\sigma_0(N)}.$$

Hence our assumption implies that $\ker(\pi)$ is nontrivial. Equivalently, there exists $w \in \mathcal{W}$ that acts as identity on both E and F . Now we consider the following diagram:



Let $\mathcal{O} \in F(\mathbb{Q})$ be the identity element. We have $P_{E,F} = P_{E,F,\mathcal{O}} = \sum(\varphi_E)_*(\varphi_F)^*(O) = \sum(\tilde{\varphi}_E)_*\pi_*\pi^*\tilde{\varphi}_F^*(O) = 2\sum(\tilde{\varphi}_E)_*\tilde{\varphi}_F^*(O) \in 2E(\mathbb{Q})$. \square

4.3 Applying the idea of IPR to the computation of Chow-Heegner points

Take E, F to be two non-isogeneous elliptic curves of the same conductor N and let φ, ψ denote the modular parametrisations of E, F , respectively. The *Chow-Heegner point* associated to the pair (E, F) is defined by

$$P_{E,F} = \sum \varphi(\psi^*(Q)), \forall Q \in F(\mathbb{C})$$

We state two facts: (1) $P_{E,F}$ is independent of the choice of Q (hence the notation); (2) $P_{E,F} \in E(\mathbb{Q})$.

We present an algorithm that either computes the Chow-Heegner point $P_{E,F}$ or outputs fail. Let x_E, y_E, x_F, y_F be the compositions of φ, ψ with the x and y coordinate functions on E and F , respectively. Note that there exists an algorithm to compute the q -expansions of x_E, x_F, y_E and y_F .

Example 4.3.1. Consider $E = 89a$ and $F = 89b$. Here $\deg(\varphi_E) = 2$ and $\deg(\varphi_F) = 5$. Let $D = \varphi(\psi^*(\infty)) \in \text{div}(E)$. Define $G_1(x) = \prod_{P \in D} (x - x(P))$ and $G_2(y) = \prod_{P \in D} (y - y(P))$. Using **IPR-CH**, we computed

$$G_1(x) = x^4 + \frac{13}{4}x^3 + \frac{17}{4}x^2 + \frac{21}{4}x + \frac{9}{2}, \quad G_2(y) = y^4 + \frac{1}{8}y^3 + \frac{21}{4}y^2 + \frac{7}{2}y + 3.$$

It turns out that $G_1(x)$ is irreducible over \mathbb{Q} . Let K be its splitting field, and write $G_1(x) = \prod (x - a_i)$ with $a_i \in K$. For each a_i , we found that $b_i = -\frac{8}{9}a_i^3 - \frac{20}{9}a_i^2 - \frac{28}{9}a_i - \frac{10}{3}$ is

Algorithm 5 Using polynomial relation to compute the Chow-Heegner point $P_{E,F}$

Input: E, F = non-isogeneous elliptic curves of conductor N ; q -expansions of x_E, y_E, x_F, y_F .

Output: the Chow-Heegner point $P_{E,F}$.

- 1: $u_E \leftarrow (x_F)^{-1}$ and $u_F \leftarrow (x_E)^{-1}$.
 - 2: Use the algorithm **IPR** to compute an irreducible polynomial $F(x, y)$ such that $F_{E,F}(u_E, u_F) = 0$.
 - 3: $f_{ch,x}(x) \leftarrow F_{E,F}(x, 0)$.
 - 4: Repeat steps 2-5 for $v_E = (y_E)^{-1}$ and u_F , get $f_{ch,y}(x)$.
 - 5: $K \leftarrow$ the splitting field of $f_{ch,x}$. Write $f_{ch,x}(x) = \prod (x - a_i), a_i \in K$.
 - 6: **for** each a_i **do**
 - 7: Find a point $p_i = (a_i, b_i)$ on $E(\bar{\mathbb{Q}})$.
 - 8: **if** $f_{ch,y}(b_i) = 0$ **then**
 - 9: $P_i = p_i$.
 - 10: **else**
 - 11: $P_i = -p_i$.
 - 12: **end if**
 - 13: **end for**
 - 14: Output $P_{E,F} = \sum_i P_i$.
-

the corresponding root of G_2 such that $(a_i, b_i) \in E$. Hence

$$P_{E,F} = \sum_{i=1}^4 P_i, \text{ where } P_i = (a_i, b_i).$$

Carrying out the summation in Sage, we obtain $P_{E,F} = (\frac{3}{4}, -\frac{15}{8})$. This agrees with Stein's result for the pair **(89a,89b)** in [?].

Chapter 5

**THINGS I TRIED TO DO BUT DID NOT END UP GIVING A
NICE RESULT**

generalizing the “congruence number” definition using other cusps.

Prove the “ $\pm 1 \bmod p$ ” guess.

Generalize another paper by William on computing order of component groups. (The original paper uses a trick which William fails to remember).

Computing the critical subgroup for 5077a (multimodular is not practical).

Critical points of reduction of modular parametrization.

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