

# Bare Demo of IEEEtran.cls for Journals

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**Abstract**—The abstract goes here.

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## I. MODELING

### A. Bearing Constraints

Suppose we have two sensors at positions  $S_1 \in \mathbb{R}^2$  and  $s_2 \in \mathbb{R}^2$ . When the communication signal is sent from  $s_1$  to  $s_2$ , the bearing  $\theta$  of the vector from  $s_1$  to  $s_2$ , can be measured. In addition, we assume there exists an error bound  $\delta$  for the angle measurement, or in inequalities,  $\theta - \delta \leq \angle(s_2 - s_1) \leq \theta + \delta$ . This angular condition can be characterized by two separate linear constraints. The first one is a semi-plane defined by the straight-line corresponding to the angle  $\theta + \delta$ , or equivalently,

$$[-\sin(\theta + \delta), \cos(\theta + \delta)]x \leq 0 \quad (1)$$

where  $x \in \mathbb{R}^2$ . The condition that the vector  $s_2 - s_1$  is within this semi-plane can thus expressed as:

$$[-\sin(\theta + \delta), \cos(\theta + \delta)](s_2 - s_1) \leq 0 \quad (2)$$

Similarly, the other semi-planes defined by the line with angle  $\theta - \delta$  can be obtained as:

$$[-\sin(\theta - \delta), \cos(\theta - \delta)]x \geq 0 \quad (3)$$

and the constraint for  $s_1$  and  $s_2$  becomes

$$[-\sin(\theta - \delta), \cos(\theta - \delta)](s_2 - s_1) \geq 0 \quad (4)$$

Note that the two semi-planes have different sign in their inequality expression due to the fact that (1) models the lower half of the plane while (3) models the upper half of it.

### B. Ranging Constraints

Due to the ranging error, the distance from  $s_1$  to  $s_2$  usually is not identical to the measurement  $d$ , but in a range between  $d - \eta^- > 0$  and  $d + \eta^+$ , where  $\eta^- \in \mathbb{R}$ ,  $\eta^- > 0$ ,  $\eta^+ \in \mathbb{R}$ ,  $\eta^+ > 0$ . The ranging constraint thus becomes,

$$d - \eta^- \leq \|s_2 - s_1\| \leq d + \eta^+ \quad (5)$$

Note that the range constraint  $d - \eta^- \leq \|s_2 - s_1\|$  is not convex. Direct solution of it may end up with local optima, or even violation of this constraint. For this consideration, we relax it to a semi-plane constraint formed by the straight-line passing through the two intersection points of the arc and the

two lines associated with angles  $\theta \pm \delta$ . This semi-plane is described as follows (**draw a figure to show it**),

$$-[\cos \theta, \sin \theta]x + (d - \eta^-) \cos \delta \leq 0 \quad (6)$$

This equation can be obtained directly by noticing that the separation line is normal to the vector  $[\cos \theta, \sin \theta]^T$ , and has a distance of  $(d - \eta^-) \cos \delta$  from the origin. Therefore, we have the following equivalently for  $s_1$  and  $s_2$ ,

$$-[\cos \theta, \sin \theta](s_2 - s_1) + (d - \eta^-) \cos \delta \leq 0 \quad (7)$$

This convex constraint approaches the non-convex constraint  $d - \eta^- \leq \|s_2 - s_1\|$  when  $\delta$  goes to zero. Generally, the set formed by (7) includes  $\{(s_1, s_2) \in \mathbb{R}^2 \times \mathbb{R}^2, d - \eta^- \leq \|s_2 - s_1\|\}$  as a subset. In practice,  $\delta$  is very small and (7) provides a tight relaxation of  $d - \eta^- \leq \|s_2 - s_1\|$ .

As to the range condition  $\|s_2 - s_1\| \leq d + \eta^+$ , we equivalently write it as follows,

$$(s_2 - s_1)^T(s_2 - s_1) \leq (d + \eta^+)^2 \quad (8)$$

### C. Restriction on the Variable Space

For beacon-free localization, the translation or the rotation of all coordinates has no impact to the relative ranging and bearing of all pair of sensors. In this paper, we aim to find one feasible solution satisfying the constraints formulated in previous sections. Although a set of coordinates with a large absolute value can also satisfy all the constraints, the accuracy may reduce in this situation due to fixed point operations of large values when an algorithm is realized on a computationally restrictive sensor network. For this consideration, we impose additional constraints as follows,

$$-\epsilon \leq s_i \leq \epsilon \quad \text{for } i = 1, 2 \quad (9)$$

where  $\epsilon \in \mathbb{R}^2$  defines the allowed variable space of  $s_i$ . In implementation, the value of  $\epsilon$  can be assigned according to prior knowledge on the size of the space occupied by the wireless sensor network.

### D. Problem Formulation

We are now ready to present the beacon-free localization problem in a mathematical formulation. Taking into account (2), (4), (7), (8), and (10), this problem is defined to find a feasible solution of the following inequality set,

$$[-\sin(\theta_{ij} + \delta_{ij}), \cos(\theta_{ij} + \delta_{ij})](x_j - x_i) \leq 0 \quad (10a)$$

$$[-\sin(\theta_{ij} - \delta_{ij}), \cos(\theta_{ij} - \delta_{ij})](x_j - x_i) \geq 0 \quad (10b)$$

$$-[\cos \theta_{ij}, \sin \theta_{ij}](x_j - x_i) + (d_{ij} - \eta_{ij}^-) \cos \delta_{ij} \leq 0 \quad (10c)$$

$$(x_j - x_i)^T(x_j - x_i) \leq (d_{ij} + \eta_{ij}^+)^2 \quad (10d)$$

$$-\epsilon \leq x_i \leq \epsilon \quad (10e)$$

$$\forall j \in \mathcal{N}(i), \forall i = 1, 2, \dots, |V|$$

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where  $|V|$  is the total number of sensors in the network  $G = (V, E)$  with  $V$  and  $E$  denoting the vertex set and the edge set respectively. The constraint (10) is composed of 4 scalar inequalities for each communication link  $i - j$ , and 4 for each sensor  $i$  (note that (10e) involves 4 inequalities as  $x_i$  is a two-dimensional vector). Overall, there are  $4|E| + 4|V|$  inequalities for the whole network. With the increase of  $|V|$ ,  $|E|$  increases at a speed of  $O(|V|^2)$  for some graph topologies, which makes it prohibitive to solve (10) in a centralized manner. In this paper, we therefore explore a distributed computational algorithm for an efficient solution of (10).

## II. ITERATIVE ALGORITHM DESIGN

### A. Problem Reformulation

Constraint (10a), (10b) and (10c) can be written into the following compact form,

$$A_{ij}(x_j - x_i) + b_{ij} \leq 0 \quad (11)$$

where

$$A_{ij} = \begin{bmatrix} -\sin(\theta_{ij} + \delta_{ij}) & \cos(\theta_{ij} + \delta_{ij}) \\ \sin(\theta_{ij} - \delta_{ij}) & -\cos(\theta_{ij} - \delta_{ij}) \\ -\cos \theta_{ij} & -\sin \theta_{ij} \end{bmatrix} \quad (12)$$

$$b_{ij} = \begin{bmatrix} 0 \\ 0 \\ (d_{ij} - \eta_{ij}^+) \cos \delta_{ij} \end{bmatrix} \quad (13)$$

To derive theoretical conditions on the solution, we regard this problem as the following optimization with a virtual objective function:

$$\min f(x_1, x_2, \dots, x_{|V|}) = 0 \quad (14a)$$

$$\text{s.t. } A_{ij}(x_j - x_i) + b_{ij} \leq 0 \quad (14b)$$

$$(x_j - x_i)^T(x_j - x_i) \leq (d_{ij} + \eta_{ij}^+)^2 \quad (14c)$$

$$-\epsilon \leq x_i \leq \epsilon \quad (14d)$$

$$\forall j \in \mathcal{N}(i), \forall i = 1, 2, \dots, |V|$$

### B. Dual Space Analysis

Define  $\lambda_{ij} \in \mathbb{R}^3$ ,  $\lambda_{ij} \geq 0$  as the dual variable associated with the constraint (14b), and  $\mu_{ij} \in \mathbb{R}$ ,  $\mu_{ij} \geq 0$  as the dual variable associated with the constraint (14c). Accordingly, define a Lagrangian function as

$$L = \frac{1}{2} \sum_{i=1}^{|V|} \sum_{j \in \mathcal{N}(i)} \left( k_0 \lambda_{ij}^T (A_{ij}(x_j - x_i) + b_{ij}) + \frac{1}{2} \mu_{ij} ((x_j - x_i)^T(x_j - x_i) - (d_{ij} + \eta^+)^2) \right) \quad (15)$$

where  $k_0 > 0$ . Using KKT condition, we get the following equivalent form of the problem,

$$\sum_{i=1}^{|V|} (y_i - x_i)^T \frac{\partial L}{\partial x_i} \geq 0 \quad \forall y_i \in \Omega \quad (16a)$$

$$\begin{cases} (\lambda_{ij})_k > 0 & \text{if } (A_{ij}(x_j - x_i) + b_{ij})_k = 0 \\ (\lambda_{ij})_k = 0 & \text{else if } (A_{ij}(x_j - x_i) + b_{ij})_k < 0 \end{cases} \quad (16b)$$

$$\forall k = 1, 2, 3, 4 \quad (16c)$$

$$\begin{cases} \mu_{ij} > 0 & \text{if } (x_j - x_i)^T(x_j - x_i) - (d_{ij} + \eta^+)^2 = 0 \\ \mu_{ij} = 0 & \text{else if } (x_j - x_i)^T(x_j - x_i) - (d_{ij} + \eta^+)^2 < 0 \end{cases} \quad (16d)$$

where the set  $\Omega$  is defined as  $\Omega = \{x \in \mathbb{R}^2, -\epsilon \leq x \leq \epsilon\}$ ,  $(\lambda_{ij})_k$  denotes the  $k$ -th element of the vector  $\lambda_{ij}$ , and  $(A_{ij}(x_j - x_i) + b_{ij})_k$  denotes the  $k$ -th element of the vector  $A_{ij}(x_j - x_i) + b_{ij}$ .  $\frac{\partial L}{\partial x_i}$  is computed as

$$\frac{\partial L}{\partial x_i} = - \sum_{j \in \mathcal{N}(i)} (k_0 A_{ij}^T \lambda_{ij} + \mu_{ij} (x_j - x_i)) \quad (17)$$

According to the properties of projection to convex set, (16) is equivalently converted to,

$$x_i = P_\Omega(x_i - k_1 \frac{\partial L}{\partial x_i}) \quad (18a)$$

$$\lambda_{ij} = (\lambda_{ij} + k_2 A_{ij}(x_j - x_i) + k_2 b_{ij})^+ \quad (18b)$$

$$\mu_{ij} = (\mu_{ij} + k_3 (x_j - x_i)^T(x_j - x_i) - k_3 (d_{ij} + \eta^+)^2)^+ \quad (18c)$$

where  $k_1 > 0$ ,  $k_2 > 0$ ,  $k_3 > 0$ ,  $P_\Omega(x)$  is defined as the projection of  $x \in \mathbb{R}^2$  to the set  $\Omega$ . In concrete,

$$y = P_\Omega(x) \Leftrightarrow y_i = \begin{cases} \epsilon_i & \text{if } x_i \geq \epsilon_i \\ -\epsilon_i & \text{if } x_i \leq -\epsilon_i \\ x_i & \text{otherwise} \end{cases} \quad \text{for } i = 1, 2 \quad (19)$$

where  $y = [y_1, y_2]^T \in \mathbb{R}^2$ ,  $\epsilon = [\epsilon_1, \epsilon_2]^T$ . For  $x = [x_i] \in \mathbb{R}^n$ ,  $y = [y_i] \in \mathbb{R}^n$  with  $n$  being  $n = 1, 2, 3, \dots$ ,  $y = x^+$  is defined in entry-wise as  $y_i = \max\{x_i, 0\}$ . The equation set (18) is derived using KKT condition and projection theorems. It is an equivalent solution of (10). The solution of (10) cannot be directly solved due to the nonlinearity in it. Instead, we use the following procedure, the equilibrium point of which

exactly follows (18), to solve (10) iteratively,

$$\begin{aligned} x_i(t) = & (1 - c_0)x_i(t-1) + c_0 \mathbf{P}_\Omega \left( x_i(t-1) \right. \\ & + k_1 \sum_{j \in \mathcal{N}(i)} [k_0 A_{ij}^T \lambda_{ij}(t-1) + \mu_{ij} x_j(t-1) \\ & \left. - \mu_{ij} x_i(t-1)] \right) \end{aligned} \quad (20a)$$

$$\begin{aligned} \lambda_{ij}(t) = & (1 - c_0)\lambda_{ij}(t-1) + c_0 [\lambda_{ij}(t-1) \\ & + k_2 A_{ij} x_j(t-1) - k_2 A_{ij} x_i(t-1) + k_2 b_{ij}]^+ \end{aligned} \quad (20b)$$

$$\begin{aligned} \mu_{ij}(t) = & (1 - c_0)\mu_{ij}(t-1) + c_0 [\mu_{ij}(t-1) \\ & + k_3 \|x_j(t-1) - x_i(t-1)\|^2 - k_3 (d_{ij} + \eta^+)^2]^+ \end{aligned} \quad (20c)$$

where the expression of  $\frac{\partial L}{\partial x_i(t-1)}$  in (17) is substituted into (20a).

### III. CONVERGENCE ANALYSIS

*Theorem 1:* For any small value  $\epsilon_0 > 0$ , there always exist  $\delta_0 > 0$ ,  $n_0 > 0$  and  $x_i^*$  (for  $i = 1, 2, \dots, |V|$ ), which is a solution of (10), such that for all  $0 < c_0 < \delta_0$  and  $t > n_0$ ,  $\|x_i(t) - x_i^*\| \leq \epsilon_0$  for  $i = 1, 2, \dots, |V|$ , where  $x_i(t)$  is the solution generated by the dynamic iteration (20) at time step  $t$  with  $k_0, k_1, k_2, k_3 > 0$  and any feasible initial values  $x_i(0) \in \Omega$ ,  $\lambda_{ij}(0) \geq 0$ ,  $\mu_{ij}(0) \geq 0$ ,  $\forall i = 1, 2, \dots, |V|$ ,  $j \in \mathcal{N}(i)$ .

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### IV. CONCLUSION

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#### APPENDIX A

##### PROOF OF THE FIRST ZONKLAR EQUATION

Appendix one text goes here.

#### APPENDIX B

Appendix two text goes here.

#### ACKNOWLEDGMENT

The authors would like to thank...

#### REFERENCES

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**Michael Shell** Biography text here.

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**John Doe** Biography text here.

**Jane Doe** Biography text here.