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EXTENDS Integers, Sequences, FiniteSets, TLC CONSTANTS N, FAILNUM ASSUME N \leq 5 \land 0 \leq FAILNUM \land FAILNUM \leq 4 Nodes \stackrel{\triangle}{=} 1 \dots N
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--algorithm syncCon1
{ variables FailNum = FAILNUM;
   up = [n \in Nodes \mapsto TRUE];
   pt = [n \in Nodes \mapsto 0];
   t = [n \in Nodes \mapsto FALSE];
   d = [n \in Nodes \mapsto -1];
   mb = [n \in Nodes \mapsto \{\}];
   define {
   SetMin(S) \stackrel{\triangle}{=} CHOOSE \ i \in S : \forall j \in S : i \leq j
    }
   macro MaybeFail( ) {
     if ( FailNum > 0 \land up[self] )
        { either
            \{ up[self] := FALSE; FailNum := FailNum - 1; \}
              or skip; };
    }
   fair process ( n \in Nodes )
   variable v = 0, Q = \{\};
P: \mathbf{if} \ (\ up[self]) \ \{
   v := self;
    Q := Nodes;
PS: while ( up[self] \land Q \neq \{\} ) {
       with (p \in Q)
          MaybeFail(); In process n, each time we add v to mb[p], this process might be fail, and once the process
                           fail, we set up[self] fail and after that, all operations will be invalid.
         if (up[self]) mb[p] := mb[p] \cup \{v\}; In process n, add v to mb[1] to mb[N]
          Q := Q \setminus \{p\}; For each element p in Q, we have to broadcast value v to mb[p], after adding v to mb[p],
                               remove p from Q in case of adding v to mb[p] again.
        } ;
     if ( up[self] ) pt[self] := pt[self] + 1;
PR: await (up[self] \land (\forall i \in Nodes : pt[i] = pt[self])); If no operations are failed and all process enter in round 1
     d[self] := SetMin(mb[self]); Set the minimum of mb[p] to d[self], if success, all d[p] will be the same.
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t[self] := TRUE; The process is terminated

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};
 BEGIN TRANSLATION
Variables FailNum, up, pt, t, d, mb, pc
 define statement
SetMin(S) \stackrel{\triangle}{=} CHOOSE \ i \in S : \forall j \in S : i \leq j
Variables v, Q
vars \triangleq \langle FailNum, up, pt, t, d, mb, pc, v, Q \rangle
ProcSet \stackrel{\Delta}{=} (Nodes)
Init \stackrel{\triangle}{=} Global variables
            \wedge FailNum = FAILNUM
            \land up = [n \in Nodes \mapsto TRUE]
            \land pt = [n \in Nodes \mapsto 0]
            \land t = [n \in Nodes \mapsto FALSE]
            \land d = [n \in Nodes \mapsto -1]
            \land mb = [n \in Nodes \mapsto \{\}]
            Process n
            \wedge v = [self \in Nodes \mapsto 0]
            \land Q = [self \in Nodes \mapsto \{\}]
           \land pc = [self \in ProcSet \mapsto "P"]
P(self) \triangleq \land pc[self] = "P"
                \wedge IF up[self]
                        THEN \wedge v' = [v \text{ EXCEPT } ! [self] = self]
                                 \land Q' = [Q \text{ EXCEPT } ![self] = Nodes]
                                 \land pc' = [pc \text{ EXCEPT } ![self] = "PS"]
                        ELSE \land pc' = [pc \text{ EXCEPT } ! [self] = \text{"Done"}]
                                 \land UNCHANGED \langle v, Q \rangle
                \land UNCHANGED \langle FailNum, up, pt, t, d, mb \rangle
PS(self) \triangleq \land pc[self] = "PS"
                  \wedge IF up[self] \wedge Q[self] \neq \{\}
                          THEN \wedge \exists p \in Q[self]:
                                          \wedge IF FailNum > 0 \wedge up[self]
                                                 THEN \wedge \vee \wedge up' = [up \text{ EXCEPT } ![self] = \text{FALSE}]
                                                                  \wedge FailNum' = FailNum - 1
                                                               \vee \wedge \text{True}
                                                                  \wedge Unchanged \langle FailNum, up \rangle
                                                  ELSE \land TRUE
                                                           \land UNCHANGED \langle FailNum, up \rangle
                                          \wedge IF up'[self]
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THEN \wedge mb' = [mb \text{ EXCEPT } ![p] = mb[p] \cup \{v[self]\}]
                                                  ELSE ∧ TRUE
                                                           \wedge mb' = mb
                                          \land Q' = [Q \text{ EXCEPT } ![self] = Q[self] \setminus \{p\}]
                                    \land pc' = [pc \text{ EXCEPT } ! [self] = "PS"]
                                    \wedge pt' = pt
                          ELSE \wedge IF up[self]
                                           THEN \wedge pt' = [pt \text{ EXCEPT } ![self] = pt[self] + 1]
                                           ELSE \land TRUE
                                                     \wedge pt' = pt
                                    \land pc' = [pc \text{ EXCEPT } ! [self] = "PR"]
                                    \land UNCHANGED \langle FailNum, up, mb, Q \rangle
                  \wedge UNCHANGED \langle t, d, v \rangle
PR(self) \triangleq \land pc[self] = "PR"
                  \land (up[self] \land (\forall i \in Nodes : pt[i] = pt[self]))
                  \wedge d' = [d \text{ EXCEPT } ! [self] = SetMin(mb[self])]
                  \wedge t' = [t \text{ EXCEPT } ! [self] = \text{TRUE}]
                  \land pc' = [pc \text{ EXCEPT } ! [self] = \text{"Done"}]
                  \land UNCHANGED \langle FailNum, up, pt, mb, v, Q \rangle
n(self) \stackrel{\triangle}{=} P(self) \vee PS(self) \vee PR(self)
Next \stackrel{\triangle}{=} (\exists self \in Nodes : n(self))
               V Disjunct to prevent deadlock on termination
                  (\forall self \in ProcSet : pc[self] = "Done") \land UNCHANGED vars)
Spec \stackrel{\triangle}{=} \wedge Init \wedge \Box [Next]_{vars}
             \land \forall self \in Nodes : WF_{vars}(n(self))
Termination \triangleq \Diamond(\forall self \in ProcSet : pc[self] = "Done")
 END TRANSLATION
agreement \stackrel{\Delta}{=} (\forall i, j \in Nodes : t[i] \land t[j] \Rightarrow d[i] = d[j])
\*When no crash happened, the program will excute without any problems and all d[i]
\*will be the same. However, if some nodes crash, the crashed process will be in round 0
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- \\*and other uncrashed processes will be in round 1, so other uncrashed processes will
- \\*not execute SetMin(S), so all d[i] will be -1. Only if all processes are not crashed can \\*we get minimum value.
- **\\*** Modification History
- \* Last modified Tue Oct 24 21:46:27 EDT 2017 by lenovo
- \\* Created Wed Oct 11 00:01:22 EDT 2017 by lenovo