

Q1 (10 points).

Answer: Use UDP (1 point) and build requirements in application layer (1 point).

Elements: Sliding Window (1 point) for flow control (1 point).

ACKs, Retransmits, timeouts (3 points) for reliable transfer (1 point).

Pipelining (1 point) for high performance (1 point).

Q2 (5 points).

Answer: Client contacts local DNS (1 point)

local DNS contacts root DNS (1 point)

root DNS contacts TLD DNS (1 point)

TLD DNS contacts authoritative DNS (1 point)

IP for www.cnn.com is propagated back on the same route (1 point)

Q3 (3 points).

Answer: There is not enough information to tell (1 points)

Because both GBN (1 point) and SR (1 point) will ACK each of the first two messages as they are received correctly, and there is not enough behaviour information for other scenarios.

Q4 (3 points).

*Answer: There 40 terabytes = $40 * 10^{12} * 8$ bits (1 point).*

*So, if using the dedicated link, it will take $40 * 10^{12} * 8 / (100 * 10^6) = 3200000$ seconds = 37 days (1 point).*

But with FedEx overnight delivery, you can guarantee the data arrives in one day (1 point).

Q5 (5 points).

Answer: The queueing delay is 0 for the first transmitted packet (1 point).

L/R for the second transmitted packet (1 point)

and generally, $(n-1)L/R$ for the n th transmitted packet (1 point).

Thus, the average delay for the N packets is (2 points):

$$\begin{aligned}
& (L/R + 2L/R + \dots + (N-1)L/R)/N \\
&= L/(RN) * (1 + 2 + \dots + (N-1)) \\
&= L/(RN) * N(N-1)/2 \\
&= LN(N-1)/(2RN) \\
&= (N-1)L/(2R)
\end{aligned}$$

Note that here we used the well-known fact:

$$1 + 2 + \dots + N = N(N+1)/2$$

Q6 (6 points).

Answer: Trudy can pretend to be Bob to Alice and vice versa (1 point)

Trudy can partially or completely modify the messages being sent from Bob to Alice and from Alice to Bob (1 point) Trudy can even drop the packets that are being sent by Bob to Alice (1 point). The latter can be done even in case of encryption. Attack based on the content of messages can be prevented by encryption (1 point)

Q7 (5 points).

Answer: An overlay network is a network that is using another network as a platform. The overlay network in a P2P system consists of the nodes participating in the (e.g. file sharing) system (1 point) and the logical links between the nodes (2 points).

There is edge in graph from node A to node B if there is a (TCP) connection between A and B (1 point).

An overlay network does not include routers (1 point).

Q8 (5 points).

Answer: An advantage of using the earlier values of `cwnd` and `ssthresh` at t_2 is that TCP would not have to go through slow start (1 point) and congestion avoidance to ramp up to the throughput value obtained at t_1 (1 point). A disadvantage of using these values is that they may be no longer accurate (1 point). In particular, if the path has become more congested between t_1 and t_2 , the sender will send a large window's worth of segments into an already more congested path (2 points).

Q9 (10 points). Please refer to the figure below, and answer the following questions.

(a) (1 point)

Answer: Slow start (2xMSS every RTT)

(b) (1 point)

Answer: Window size reached `ssthresh`

(c) (2 points)

Answer: 8 (before congestion) and 6 (after congestion)

(d) (1 point)

Answer: Congestion avoidance (MSS+1 every RTT)

(e) (1 point)

Answer: Reaction to Congestion

(f) (1 point)

Answer: TCP 2

(g) (1 point)

Answer: TCP 1

(h) (2 points)

Answer: Window size and window maintenance mechanism after congestion

Q10 (5 points).

Answer: Persistent HTTP connections can send multiple objects over one TCP connection.

1. Reduces the TCP handshaking overhead (2 points).

2. Reduces the repeated TCP slow starts (1 point).

Very effective if multiple objects come from the very same server (1 point), but not if they come from different servers (1 point).

Q11 (10 points).

(a) (5 points)

Answer: Two different ways of managing the transmission resources in communication network (1 point).

***Circuit switching:** explicit call setup phase; allocates an end-to-end path and resources; all data flows on same path (1 point).*

***Example:** classic telephony (1 point).*

Packet switching: no end-to-end setup; data divided into chunks called packets; each packet independently addressed and routed through network (1 point). Example: Internet (1 point).

(b) (5 points)

Answer: (transport) protocols (1 point)

***Connection-oriented:** a protocol with explicit phases for setup, data transfer, and teardown afterwards; involves state information at nodes (1 point). Example: TCP. (1 point)*

***Connection-less:** a protocol with no advance setup or release of resources, and minimal or no state information involved (1 point). Example: UDP (1 point).*

Q12 (5 points).

***Answer:** Since the link capacity is only 100 Mbps, so host A's sending rate can be at most 100Mbps (1 point). Still, host A sends data into the receive buffer faster than Host B can remove data from the buffer. The receive buffer fills up at a rate of roughly 50Mbps (1 point). When the buffer is full, Host B signals to Host A to stop sending data by setting RcvWindow = 0. Host A then stops sending until it receives a TCP segment with RcvWindow > 0 (1 point). Host A will thus repeatedly stop and start sending as a function of the RcvWindow values it receives from Host B (1 point). On average, the long-term rate at which Host A sends data to Host B as part of this connection is no more than 50Mbps (1 point).*

Q13 (10 points).

***Answer:** Electronic Mail Application (1 point)
POP3 (Mail Access) IMAP (Mail Access) SMTP (Mail Transfer) (6 points)
POP3 and IMAP are both mail access protocols. IMAP provides more functionalities (1 point), including folders and saved mailbox (2 points)*

Q14 (10 points).

Source port #							Dest port #						
Sequence number													
Acknowledgment number													
Header length		Unused		URG	ACK	PSH	RST	SYN	FIN	Receive window			
Internet checksum									Urgent data pointer				
Options													
Data													

Q15 (10 points).

***Answer:** Throughput = $\min \{R_s, R_c, R/M\}$*