Why the Candidates move further away? They're not crazy!

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#### Abstract

### 1 Introduction

## 2 Deterministic Model

We denote the policy space is  $\mathbb{R}^m$  and the set of voters  $N = \{v_1, v_2, \cdots, v_n\}$ , and their utility function  $u_i : \mathbb{R}^m \to \mathbb{R}$  which is smooth(or  $C^2$ , second differentiable) and concave, where  $v_i \in \mathbb{R}^m$  is the ideal position (maximum point of  $f_i$ ), for the *i*-th voter,  $1 \le i \le n$ . For the two candidates  $P_1$  and  $P_2$ ,  $P_1$  chooses a position first, and  $P_2$  chooses a position later. After both candidates complete their actions, the voters will vote for the candidate close to them separately. In general, for each position x such that  $P_1$  chooses, we could find a region S(x) as a subset of policy space  $\mathbb{R}^m$  such that  $P_2$  can choose any position in this region to beat  $P_1$ .

For given information of voters (i.e. all  $f_i$ ), we denote  $\alpha : \mathbb{R}^m \to \mathbb{R}^+, \alpha(x) = \mu(S(x))$ , where  $\mu(\cdot)$  is the Lebesgue measure.

**Theorem 2.1.**  $\alpha$  is a convex function. (expected)

We denote the collections of dominant voters as  $\mathcal{M} = \{M \subset N : \#M \geq \frac{1}{2} \#N\}$ . For a collection of voters M and a given position x  $P_1$  chooses, we denote the region for  $P_2$  could win all voters in M as  $V_M$ , denote the region for  $P_2$  could win the ith voter as  $B_i(x)$ .

$$B_{i}(x) = \{ y \in \mathbb{R}^{m} f_{i}(y) \geq f_{i}(x) \}, \quad V_{M}(x) = \bigcap_{i \in M} B_{i}(x)$$

$$\alpha(x) = \mu \left( \bigcup_{M \subset \mathcal{M}} \bigcap_{i \in M} B_{i}(x) \right)$$

$$= \mu \left( \bigcup_{M \subset \mathcal{M}} V_{M}(x) \right)$$

$$= \sum_{M \subset \mathcal{M}} \mu \left( V_{M}(x) \right) - \sum_{M_{1}, M_{2} \subset \mathcal{M}} \mu \left( V_{M_{1}}(x) \cap V_{M_{2}}(x) \right) + \cdots$$

$$= \sum_{M \subset \mathcal{M}} \mu \left( V_{M}(x) \right) - \sum_{M_{1}, M_{2} \subset \mathcal{M}} \mu \left( V_{M_{1} \cup M_{2}}(x) \right) + \cdots$$

To be proofed

**Lemma 1.** The area of the intersection of circles is convex w.s.t. the first candidate's position x in  $\mathbb{R}^2$ 

*Proof.* To prove the lemma, we need several steps. Step 1: the area of is convex if the movement of x does not change the shape of the region, i.e., it's a region bounded by k arcs, and k does not change w.s.t. to the movement of x. Step 2: changes of k happen only if x is on the line of two voter's positions (hence, each domain keeping k constant is a convex set). Step 3: to complete the lemma (working).

Proof of Step 1. The area of the "cap" shape region from a circle with radius r and the "cap" with chord length L is,

$$f(r,L) = \sin^{-1}\frac{L}{2r} - \frac{L}{2}\sqrt{r^2 - L^2/4},$$

which is convex and easily checked in  $\mathbb{R}^2$ . Such a vertex in the intersection region is symmetric to the first candidate's position x w.s.t. the line of two voters' ideal. As a result, it could be represented by a linear transform Ax + b, the length of the chord  $x_1x_2$  is

$$||A_1x+b+1-(A_2x+b_2)||_2$$

which is convex. (to be done).

#### 2.1 Weaker Version

If all voters stand at an extreme point of the set of voters' ideal position, the conclusion is true in  $\mathbb{R}^2$ . (could be extended to  $\mathbb{R}^m$ )

## 3 Probablistic Model

## 3.1 Model given by [Sch07]

The utility function of voter i with ideal position  $v_i$  and a candidate position  $z_j$  is

$$u_{ij}(v_i, z_j) = u_{ij}^*(v_i, z_j) + \varepsilon_{ij},$$

where

$$u_{ij}^*(v_i, z_j) = \lambda_i - \beta ||v_i - z||^2.$$

Here,  $u_{ij}^*$  is the observable utility for voter i associated with party/candidate j.  $\lambda_j$  is the valence of agent/candidate/party j, and  $\beta$  is a positive constant. The terms  $\{\varepsilon_{ij}\}$  are the stochastic errors. We assume that all  $\varepsilon_{ij}$ 's are iid drawed from distribution  $\Psi$ , which is the Type I extreme value distribution and takes the closed form

$$\Psi(h) = \exp(-(\exp -h)).$$

The probability for a voter i to choose party j is that

$$p_{ij} = \Pr[u_{ij}(v_i, z_j) > u_{il}(x_i, z_l)], l \neq j.$$

The expected vote share for agent j is  $V_j(\mathbf{z}) = \sum_{i=1}^n p_{ij}$ , where the **z** is the vector of all candidates' positions.

In the model given above, the expected voter share for candidate j is,

$$V_j(\mathbf{z}) = \sum_{i=1}^n \frac{\exp(u_{ij}^*(v_i, z_j))}{\sum_j \exp(u_{ik}^*(v_i, z_k))}.$$

**Remark 3.1.** (The norm is convex (by definition)) and (the function  $x^2$  is non-decreasing (in positive half-axis) and convex) jointly imply that the square of a norm is convex, which implies that  $u_{ij}^*$  is concave. The function  $e^x$  is convex and non-decreasing We denote  $G: \mathbb{R}^+ \to \mathbb{R}, x \mapsto \frac{x}{x+k}$  for some positive k. It's easy to verify that G is strictly concave and increasing. We recall that  $u_i$  is strictly concave. Hence  $G \circ f_i$  is strictly concave, which implies  $\mathbb{E}_z(x)$  is concave.

$$P_{ij} = \frac{e^{f_i(x_j)}}{e^{f_i(x_1)} + e^{f_i(x_2)}}, \quad j = 1, 2$$

For a position  $z \in \mathbb{R}^m$  that  $P_1$  chooses,  $P_2$  can choose a position to maximize his expectation:

$$\mathbb{E}_z : \mathbb{R}^m \to \mathbb{R}^+, \quad \mathbb{E}_z(x) = \sum_{i=1}^n \frac{e^{f_i(x)}}{e^{f_i(x)} + e^{f_i(z)}}$$

However, the objective function may not be concave.

#### 3.2 Modificated Model

All utility functions  $f_i$  are strictly concave, positive, and second-differentiable.

If the two candidate choose position  $x_1$  and  $x_2$ , denote the probability for voter i to voter candidate  $P_j$  is  $P_{ij}$ ,  $1 \le i \le n, j = 1, 2$ ,

$$P_{ij} = \frac{f_i(x_j)}{f_i(x_1) + f_i(x_2)}, \quad j = 1, 2$$

For a position  $z \in \mathbb{R}^m$  that  $P_1$  chooses,  $P_2$  can choose a position to maximize his expectation:

$$\mathbb{E}_z : \mathbb{R}^m \to \mathbb{R}^+, \quad \mathbb{E}_z(x) = \sum_{i=1}^n \frac{f_i(x)}{f_i(x) + f_i(z)}$$

**Theorem 3.1.** Given a fixed position z, the function  $\mathbb{E}_z(x)$  is (strictly) concave.

*Proof.* We denote  $G: \mathbb{R}^+ \to \mathbb{R}, x \mapsto \frac{x}{x+k}$  for some positive k. It's easy to verify that G is strictly concave and increasing. We recall that  $f_i$  is strictly concave. Hence  $G \circ f_i$  is strictly concave, which implies  $\mathbb{E}_z(x)$  is concave.

**Theorem 3.2.** (Pure Strategy Nash Equilibrium) There exist a unique  $x^* \in \mathbb{R}^m$  such that  $(x^*, x^*)$  is the pure strategy Nash equilibrium, i.e.

$$\arg \max_{x \in \mathbb{R}^m} \mathbb{E}_{x^*}(x) = \arg \max_{x \in \mathbb{R}^m} \sum_{i=1}^n \frac{f_i(x)}{f_i(x) + f_i(x^*)} = x^*$$

which implies  $\mathbb{E}_{x^*}(x^*) = n/2$ 

*Proof.* Step 1: Let  $V \subset \mathbb{R}^m$  be the convex hull of the ideal positions of voters. Denote  $h: V \to \mathbb{R}^m$ ,  $h(z) = \arg\max_{x \in \mathbb{R}^m} \mathbb{E}_z(x)$ . h is well-defined due to the strictly convexity of  $\mathbb{E}_z(x)$ . Claim  $h(z) = \arg\max_{x \in V} \mathbb{E}_z(x)$  such that  $h: V \to V$ .

Step 2: Prove h is continous w.r.t. z. Done by implicit function theorem with the condition  $\mathbb{E}_z(x)$  is second differentiable w.r.t. x.

Step 3: h is a continuous function from a convex compact set to itself; by Brouwer's fixed-point theorem, there exists a fixed point  $x^*$  such that  $h(x^*) = x^*$ . Uniqueness is trivial.

Brouwer fixed-point theorem: Every continuous function from a nonempty convex compact subset K of a Euclidean space to K itself has a fixed point.

**Remark 3.2.** The theorem could be proved by the contraction mapping theorem.

Remark 3.3. Example:

$$\mathbb{E}_x(y) = y^2 - x^2 + \frac{n}{2}, -1 < x, y < 1$$

where  $x, y \in \mathbb{R}^1$  represents the two candidates' position.(0,0) is a PNE here.  $\mathbb{E}_x(y) + \mathbb{E}_y(x) = n, E_x(x) = \frac{n}{2}$ .

**Proposition 3.2.1.** the gradient of the objective function  $\nabla \mathbb{E}_z$  is lipschiz continous.

**Remark 3.4.** The gradient descent algorithm solves the maximization point of  $\mathbb{E}_z$ .

# 4 Deterministic Dynamic Game

#### 4.1 Discrete-Time Game + Probablistic Model

Cor 6.2 states a sufficient and necessary condition for the existence of an optimal strategy. The strategies  $\gamma_k^{j*}, k \in K, j = 1, 2$ , where K is the total time phase, provide a saddle-point solution if and only if, there exists function  $V(k, \cdot) : \mathbb{R}^{2m} \to \mathbb{R}, k \in K$  s.t.

$$V(k,x) = \min_{u_k^1} \max_{u_k^2} \{V(k+1,x+u_k^1+u_k^2)\}, \ V(K,x) = q(x), \ V(K+1,x) = 0$$

where  $x_{k+1} = f_k(x_k, u_k^1, u_k^2) = x_k + u_k^1 + u_k^2$ . The unique saddle point value of the game is  $V(1, x_1)$ , where  $x_1$  is the initial state.

#### 4.2 Continuous-Time Game + Probablistic Model

We fit the probabilistic model in the dynamic game. The set of n voters is  $N = \{1, 2, \dots, n\}$  and the policy space is  $\mathbb{R}^m$ . Predisribed fixed time T to end the voting game. The state space is  $\mathbb{R}^{2m}$ , which describes the two candidates' position positions (first m-entries for  $P_1$  and the latter for  $P_2$ ). The strategy space,

$$\Gamma^1 = \{ \gamma : [0, T] \times \mathbb{R}^{2m} \to \mathbb{R}^m \oplus \mathbf{0} \} | \gamma \text{ is continuously differentiable, Image}(\gamma) \subset B_{\mathbf{0}}(1) \}$$

$$\Gamma^2 = \{ \gamma : [0, T] \times \mathbb{R}^{2m} \to \mathbf{0} \oplus \mathbb{R}^m | \gamma \text{ is continuously differentiable, Image}(\gamma) \subset B_{\mathbf{0}}(1) \}$$

For any pure strategy the candidates choose, it's continuously differentiable, and the speed of action is less than 1. For any given pure strategy  $\gamma^1, \gamma^2$ , the evolution function is defined as,

$$\frac{dx}{dt} := f(t, x, u_1, u_2) = u_1 + u_2 = \gamma^1(t, x) + \gamma^2(t, x)$$

the f is the general notation used in dynamic games. We denote the objective function,

$$L = \int_0^T g(t, x, u^1, u^2) dt + q(T, x(T))$$

at here,  $g \equiv 0$  and T is fixed. To make it simple, we just write L = q(x(T)). We consider the probabilistic voting model and the objective function L is  $P_2$ 's expectation. For example, let  $f_i$  be the utility functions for each voter and,  $p_1, p_2 \in \mathbb{R}^m$  are the final position of the two candidates,

$$L = q(x(T)) = q(p_1, p_2) = \sum_{i=1}^{n} \frac{f_i(p_2)}{f_i(p_1) + f_i(p_2)}$$

then  $P_1$  is the minimizer, and  $P_2$  is the maximizer, since it's a zero-sum game.

**Theorem 4.1.** Cor6.6 (Equation 6.75, for feedback pattern) in Tamer Basar: The sufficient condition for  $\gamma^{1*}, \gamma^{2*}$  to be the optimal (the saddle strategy) is, there exists continuously differentiable function  $V: [0,T] \times \mathbb{R}^{2m} \to \mathbb{R}$  satisfies the Bellman's equation:

$$\begin{split} -\frac{\partial V}{\partial t} &= \min_{u^1} \max_{u^2} \langle \frac{\partial V}{\partial x}, u^1 + u^2 \rangle \\ &= \max_{u^2} \min_{u^1} \langle \frac{\partial V}{\partial x}, u^1 + u^2 \rangle \\ &= \langle \frac{\partial V}{\partial x}, \gamma^{1*}(t, x) + \gamma^{2*}(t, x) \rangle \end{split}$$

where the  $\langle \cdot \rangle$  denotes the general inner product in the Euclidean space.

**Remark 4.1.** The interchangeability is named Issac's condition. Denote the first m-entries variables is  $x_1$ , the latter m-entries is  $x_2$ . The Issac's condition is automatically satisfied,

$$\min_{u^1}\max_{u^2}\langle\frac{\partial V}{\partial x},u^1+u^2\rangle=\min_{u^1}\max_{u^2}\left[\langle\frac{\partial V}{\partial x_1},u^{1'}\rangle+\langle\frac{\partial V}{\partial x_2},u^{2'}\rangle\right]=\min_{u^1}\langle\frac{\partial V}{\partial x_1},u^{1'}\rangle+\max_{u^2}\langle\frac{\partial V}{\partial x_2},u^{2'}\rangle$$

where  $u^{1'}$  is the vector in  $\mathbb{R}^m$  consisted of first m-entries of  $u^1$ . A similar definition is applied on  $u^{2'}$ .

**Remark 4.2.** We consider the constraint that  $||u^1||_2, ||u^2||_2 \le 1$ , we can simplifize

$$\begin{split} \min_{u^1} \langle \frac{\partial V}{\partial x_1}, u^{1'} \rangle &\geq -||u^1||_2 \cdot ||\frac{\partial V}{\partial x_1}||_2 \geq -||\frac{\partial V}{\partial x_1}||_2 \\ &\max_{u^2} \langle \frac{\partial V}{\partial x_2}, u^{2'} \rangle \leq ||\frac{\partial V}{\partial x_2}||_2 \end{split}$$

which are both from Cauchy-Schwarz inequality and take equality when  $u^{1'}$  lies on the oppsite direction of  $\frac{\partial V}{\partial x_1}$  and  $u^{2'}$  lis on the same direction  $\frac{\partial V}{\partial x_2}$ . The PDE is simplified as,

$$\begin{split} -||\frac{\partial V}{\partial x_1}||_2 + ||\frac{\partial V}{\partial x_2}||_2 &= -\frac{\partial V}{\partial t}\\ initial\ condition: V(T,x) = q(x), \forall x\\ boundary\ condition: V(x_1,x_2,t) = 1\ for\ large\ |x_1|,\\ V(x_1,x_2,t) = 0\ for\ large\ |x_2|, \end{split}$$

the boundary condition is due to  $x_1$  being the minimizer, and  $x_2$  being the maximizer. The ideal position should be in the convex hull of voters' ideal positions.

#### Remark 4.3.

$$V(x,T) = q(x) = q(x_1, x_2) = \sum_{i=1}^{n} \frac{g_i(x_2)}{g_i(x_1) + g_i(x_2)}, q_{x_1} = \sum_{i=1}^{n} \frac{-g_i(x_2)}{(g_i(x_1) + g_i(x_2))^2} \times g'(x_1)$$
$$q_{x_2} = \sum_{i=1}^{n} \frac{g'_i(x_2)g_i(x_1)}{(g_i(x_1) + g_i(x_2))^2}$$

in high dimensions, it's  $||q_{x_1}||_2 = \sum_{i=1}^n g_i(x_2)||\nabla g_i(x_1)||_2$ . In the case that the initial state  $(x_0^1, x_0^2)$  satisfies  $||q_{x_1}||_2 = ||q_{x_2}||_2$ , then V is time independent, i.e.,  $\frac{\partial V}{\partial t} \equiv 0$ . This means the value of the game is given by  $q(x_0^1, x_0^2)$  and it holds for any election period T.

#### Remark 4.4.

$$V(x_1, x_2, t)$$

the optimal value start from time t and position  $x_1, x_2$ , given fixed duration T.

## 4.3 The Step-function model

Suppose each voter gets fixed utility only if the candidate close enough to her position. Formally, we define the distribution of voters by a density function p(v) for the probability density of voters with ideal position v in the policy space. We define the threshold r(v) for voters with ideal position v. The shifted expeted value (to shift the constant-sum game into a zero-sum game) from voters with ideal positions v for candidate 2 is:

$$w(v) = \begin{cases} p(v)/2, & \text{if } ||P_2 - v||_2 \le r(v) \text{ and } |P_1 - v||_2 > r(v), \\ -p(v)/2, & \text{if } ||P_2 - v||_2 > r(v) \text{ and } |P_1 - v||_2 \le r(v), \\ 0, else \end{cases}$$

**Remark 4.5.** No open-loop or closed-loop pattern here. All state evolsion are deterministic, instead of including a stochastic error. Theoretically, all candidates know the outcome at the beginning with given utility model, voters' distribution, initial position and voting duration.

**Theorem 4.2.** One candidate's initial position make no difference to the other's strategy.

*Proof.* Recall that it's a zero-sum dynamic game and  $P_1$  is minimizer with final position  $x_1$ ,  $P_2$  is maximizer with final position  $x_2$ , the objective function is

$$\int w(v)p(v)dv = \frac{1}{2}\mathbf{P}(\{v: ||x_2 - v||_2 \le r(v), ||x_1 - v||_2 > r(v)\}) - \frac{1}{2}\mathbf{P}(\{v: ||x_1 - v||_2 \le r(v), ||x_2 - v||_2 > r(v)\})$$

$$= \frac{1}{2}\left(\mathbf{P}(\{v: ||x_2 - v||_2 \le r(v)\}) - \mathbf{P}(v: ||x_1 - v||_2 \le r(v), ||x_2 - v||_2 \le r(v))\right) - \frac{1}{2}\left(\mathbf{P}(\{v: ||x_1 - v||_2 \le r(v)\}) - \mathbf{P}(v: ||x_1 - v||_2 \le r(v), ||x_2 - v||_2 \le r(v))\right)$$

$$= \frac{1}{2}\mathbf{P}(\{v: ||x_2 - v||_2 \le r(v)\}) - \frac{1}{2}\mathbf{P}(\{v: ||x_1 - v||_2 \le r(v)\}).$$

Hence, the value of the game is

$$\min_{P_1} \max_{P_2} \frac{1}{2} \mathbf{P}(\{v : ||x_2 - v||_2 \le r(v)\}) - \frac{1}{2} \mathbf{P}(\{v : ||x_1 - v||_2 \le r(v)\}) 
= \max_{P_2} \frac{1}{2} \mathbf{P}(\{v : ||x_2 - v||_2 \le r(v)\}) - \max_{P_1} \frac{1}{2} \mathbf{P}(\{v : ||x_1 - v||_2 \le r(v)\}).$$

The two candidates will try their best to an admissible position with maximized influence separately. As a result, the candidate  $P_1$  and  $P_2$  with limited moving speed  $\varepsilon$  and prescribed voting duration T and initial position  $x_{10}, x_{20}$  will move to the position

$$\arg \max_{y \in B_{x_10}(T\varepsilon)} \frac{1}{2} \mathbf{P}(\{v : ||y - v||_2 \le r(v)\}),$$
  
$$\arg \max_{y \in B_{x_20}(T\varepsilon)} \frac{1}{2} \mathbf{P}(\{v : ||y - v||_2 \le r(v)\}).$$

The initial and final positions determine the candidates' behaviors.

**Remark 4.6.** Given the centers and radius, the problem is finding the point/region to be covered by the most circles. Questions: in discrete cases, the problem is NP/NP-hard. Approximation/randomized algorithm.

**Remark 4.7.** In this case, why does the candidate not initially stand at the optimal point instead of doing dynamic programming?

**Remark 4.8.** What if the candidate gains benefit from each step? (instead of just considering the final value, in the sense that the current supporting rate impacts the candidates' sponsor/funding).

**Remark 4.9.** What if the threshold is dynamic? For example, the threshold with initial value w is  $\max\{0, w + B_{t\varepsilon}\}\$  for  $t \in [0, T]$ , where T is the voting duration and  $\varepsilon$  is a small number, and B refers to the standard Brownian motion.

Remark 4.10. In case all voters share the same threshold of 1, the two final positions change into

$$\arg \max_{y \in B_{x_10}(T\varepsilon)} \frac{1}{2} \int_{B_{y}(1)} p(v) \, dv, \quad \arg \max_{y \in B_{x_20}(T\varepsilon)} \frac{1}{2} \int_{B_{y}(1)} p(v) \, dv.$$

**Remark 4.11.** The step-function model makes sense to the more general model  $u \sim \exp(-d^2)$ , where d is the distance from the voters' ideal to the candidate's.

# 5 Implementation

We consider the discrete-time version and based on the result in 4.1, for K time-step,  $k = 0, 1, \dots K$ ,  $x_1, x_2 \in \mathbb{R}^2$  represents the candidates' position separately.  $V : \{0, 1, \dots, K\} \times \mathbb{R}^4 \to \mathbb{R}^+$  means the saddle value from input time-step k and the input candidates' positions  $x_1, x_2$  at this stage.

$$V(k,x_1,x_2) = \min_{u_k^1} \max_{u_k^2} \{V(k+1,x_1+u_k^1,x_2+u_k^2)\}, \ V(K,x_1,x_2) = q(x_1,x_2),$$

$$V(0, x_1, x_2) = \underbrace{\min_{\substack{u_1^1 \quad u_2^1 \quad u_1^2 \quad u_1^2 \quad u_2^2 \\ K \text{ min max}}}_{K \text{ min max}} q \left( x_1 + \sum_{i=1}^K u_i^1, x_2 + \sum_{i=1}^K u_i^2 \right)$$

We assume the speed of each time-step is small  $\epsilon$  and  $u_k^j, 1 \leq k \leq K, j = 1, 2$ . For numerical approximation, finite random choices for each  $u_k^j$  and brute-force search for the optimal strategy.

## 5.1 Example

The policy space is  $[-1/2, 1/2] \times [-1/2, 1/2] \subset \mathbb{R}^2$ ,

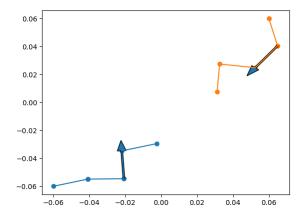
$$q_1(x_1, x_2) = \cos^2(||x_1||_2\pi) + \sin^2(||x_2||_2\pi)$$

$$q_2(x_1, x_2) = \frac{\exp(x_{21}^2 - x_{22}^2 - 1)}{\exp(x_{21}^2 - x_{22}^2 - 1) + \exp(x_{11}^2 - x_{12}^2 - 1)}$$

$$q_3(x_1, x_2) = \frac{\exp(\frac{1}{100||x_2||_2})}{\exp(\frac{1}{100||x_1||_2}) + \exp(\frac{1}{100||x_2||_2})}$$

$$q = \frac{1}{4} \cdot (q_1 + q_2 + q_3) \in [0, 1]$$

where  $x_1 = (x_{11}, x_{12})^T$ ,  $x_2 = (x_{21}, x_{22})^T$  and q represents the payoff function of candidate  $P_2$  (corresponding to  $x_2$ ).  $P_1$  (corresponding to  $x_1$ ) is the minimizer and  $P_2$  is the maximizer. We also have  $q(x_1, x_2) + q(x_2, x_1) = 1$ . In the case K = 4,  $\epsilon = 0.02$  (moving distance in each step), the following 2 figures show the different behaviour of the two candidates.



0.2 - 0.1 - 0.0 - - 0.1 - 0.2 - 0.1 0.0 0.1 0.2

Figure 1: Initial position:  $x_1 = (-0.06, -0.06)$  and  $x_2 = (0.06, 0.06)$ . Two candidates simultaneously move to the center (closer and closer)

Figure 2: Initial position:  $x_1 = (-0.2, -0.2)$  and  $x_2 = (0.2, 0.2)$ . Two candidates simultaneously move further and further

# References

[Sch07] Norman Schofield. The mean voter theorem: Necessary and sufficient conditions for convergent equilibrium. The Review of Economic Studies, 74(3):965–980, 2007.