The categorial approach

Core assumption

The meaning of questions is defined based on possible **short** answers.

- (1) [which girl smiled]]^{w₀} = $\left\{ \begin{array}{l} \langle \mathbf{a}, w_0 \rangle, & \langle \mathbf{b}, w_1 \rangle, \\ \langle \mathbf{c}, w_1 \rangle, & \langle \mathbf{c}, w_2 \rangle \end{array} \right\}$
- (2) [which girl smiled]] $^{w_0} = \lambda x \lambda w : girl_{w_0}(x).smiled_w(x)$

Answerhood

- (3) A: Which girl smiled?
 - B: (i) Anna.
- (ii) Anna smiled.

Deriving true short answers:

$$\mathbf{Ans}_{w_0}^{\mathcal{S}}(P) = \iota x.x \in \mathsf{Dom}(P) \wedge P(x)(w_0)$$

i.e., THE entity that makes true the function *P* denoted by a question (see Xiang 2021 for a sophisticated version)

Deriving true full answers:

- (4) $\mathsf{Ans}(w_0)(P) = P(\mathsf{Ans}^S(w_0)(P))$
- ✓ Full answers are derived from short answers.

Debate: Semantics vs. Syntax

Short answers are derived from full answers via ellipsis.

Merchant (2004)

- (5) A: [Which book]₁ did Jill read t_1 ?
 - B: $Emma_1$ [Jill read t_1]

Jacobson (2016): Short answers are not semantically equivalent to full answers.

- (6) Which math professor left the party at midnight?
 - a. Jill.

→ Jill is a math professor

- b. Jill did.
- c. Jill left the party at midnight.

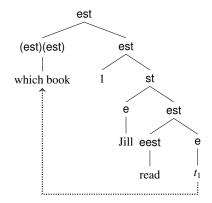
Weir (2014): clausal ellipsis must inherit a presuppositional restriction.

✓ The debate is still there.

Composition

Wh-expressions denote restrictions of functions.

(7) [which book]]^{w₀} = $\lambda P \lambda x \lambda w$: book_{w₀}(x).P(x) (est)(est)



Free relatives

Semantic similarities between *wh*-questions and free relatives (Chierchia & Caponigro 2013)

Mention-all

- (8) What did Jill cook? Pasta.→ All food that Jill cooked is pasta.
- (9) Peter ate what Jill cooked.→ Peter ate all food that Jill cooked.

Mention-some

- (10) Where can I buy coffee? Starbucks.→ Starbucks is just one coffee shop.
- (11) Jill went to where she could buy coffee.→ Jill went to a place where she could buy coffee.

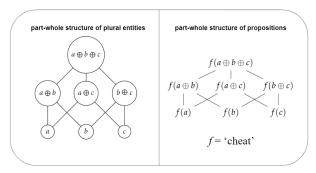
✓ Unification: The meaning of a free relative equals to the true short answer to the corresponding *wh*-question (see also Jacobson 1995; Caponigro 2003).

Question: why are some *wh*-expressions not used to form free relatives? (Caponigro 2003)

- (12) *Peter ate {what / which} food Jill cooked.
 - (13) *I did it why you did it.

Quantificational Variability

- (14) Sarah knows, for the most part, who cheated. → For most people who cheated, Sarah knows they cheated.
- (15) The school paper recorded, **in part**, what made the dean's list.
 - → For some people who made the dean's list, the school paper recorded they made the dean's list.
- H1: Quantification over short answers (Berman 1991)
 most x. x ≤_{atom} Ans^S_{w₀} [who cheated]: know(cheat(x))(s)
- H2: Quantification over full answers (Lahiri 2000)
 most p. p ≤_{atom} Ans_{wo} [who cheated]: know(p)(s)



Williams' (2000) puzzle

- (16) For the most part, Sarah knows who formed the committee.
 - → For most committee members, Sarah knows they are **part of** the committee.

The full answer to the embedded *wh*-question does not have atomic parts. **H2** is challenged.

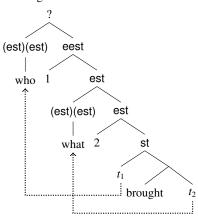


Under **H1**, the meaning of (16) can be formulated as:

Question: How is (17) compositionally derived? Where does the underlined part come from? (see Cremers (2018) for a solution)

Multiple-wh questions

(18) Who brought what?



Xiang's (2021) solution

Karttunen (1977): wh-expressions are existential quantifiers

(19)
$$[\![\mathbf{w}ho]\!]^{w_0} = \lambda P \lambda w \exists x. \mathsf{human}_{w_0}(x) \land P(x)(w) \qquad (\mathsf{est})\mathsf{st}$$

Partee (1986): Shifting existential quantifiers to properties

(20)
$$\mathbf{Be}(\mathcal{P}) = \lambda x. \mathcal{P}(\lambda y \lambda w. y = x)$$

A new type shifter **BeDom**:

BeDom(\mathcal{P}) = $\lambda f \iota g . g$ is a function such that:

- 1. $\mathsf{Dom}(g) = \mathsf{Dom}(f) \cap \mathsf{BE}(\mathcal{P})$ and
- 2. $\forall x \in \mathsf{Dom}(g).g(x) = f(x)$

i.e., for any function f, **BeDom** restricts the domain of f with $\mathbf{BeDom}(\mathcal{P})$.

(21)
$$\mathbf{BeDom}[\![what]\!]^{w_0} = \lambda f_{\mathsf{e}a} \underbrace{\lambda y : y \in \mathsf{thing}_{w_0}.f(y)}_{\mathsf{e}a}$$

(22)
$$\mathbf{BeDom}[\![who]\!]^{w_0} = \lambda f_{\mathsf{e}a} \underbrace{\lambda y : y \in \mathsf{human}_{w_0}.f(y)}_{\mathsf{e}a}$$

 \checkmark A shifted *wh*-expression can combine with a function of any type (Note: a is a variable of type).

Comparison: Chung & Ladusaw (2003) propose a new compositional mode **Predicate Restriction**.

(23) **Restrict** $(\lambda x \lambda y. f(x)(y), \lambda x. g(x)) = \lambda x \lambda y. f(x)(y) \wedge g(x)$

Wh-coordination

- (24) Jenny knows who came and who brought what.
- [who came]: est
- [who brought what]: eest

However, two elements of different types cannot be coordinated.

- (25) a. *Peter brought and came a cake.
 - b. *Susan and student came.

Xiang's (2021) solution

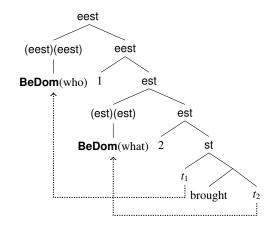
Questions take scope (see also Krifka 2001).

- (26) $\underbrace{ [\text{who came and who brought what}]}_{\lambda f. f(Q_1) \land f(Q_2)} \lambda Q \text{ Jenny knows } Q$
- \odot Both $f(Q_1)$ and $f(Q_2)$ are of type t, so they are conjoinable.

However, ...

© f doesn't have a fixed type! (see Xiang's response on pg 640, ft 34)

Note: No matter which technique is used to resolve the problem, it shouldn't be applicable to the examples in (25).



(27) The denotation of (18): $\lambda x \lambda y \lambda w : x \in \text{thing}_{w_0} \land y \in \text{human}_{w_0}.\text{brought}_w(x)(y)$

The type of (18) is eest.

 \checkmark A single-wh question and a multiple-wh question don't have **different** types.

Other types of questions

- (28) [Did Mary meet Jane or Peter] $= \lambda x \lambda w : x \in \{j, p\}.met_w(x)(m)$
- (29) [Did Mary leave] $= \lambda f : f \in \{\lambda p. p, \lambda p. \overline{p}\}. f(\lambda w. \mathsf{left}_w(\mathsf{m}))$

Other variants

Structured meaning approach

von Stechow (1991); Krifka (2001)

Questions denote ordered pairs $\langle D, f \rangle$

- (30) [who left]]^{w₀} = $\langle \text{human}_{w_0}, \lambda x \lambda w. \text{left}_w(x) \rangle$
- (31) $[\![John_F left]\!] = \langle j, \lambda x \lambda w. left_w(x) \rangle$

Question-answer congruence: A question-answer pair is congruent only if the focus component of the answer is in D and the background component of the answer is equal to f.

- (32) A: Who left?
 - B: (i) John_F left.
 - (ii) #John [left]_F.

Variable-free semantics

Jacobson (1999); Xiang (2019)

Variable-free Semantics can 'pass up' the information about an unfilled syntactic argument.

Geach Rule:
$$\mathbf{g}(f) = \lambda g \lambda x. f(g(x))$$

- i.e., **g** shifts a function of type ab to a function of type (ca)(cb).
- 1. [and]: $ttt \xrightarrow{g} (et)(ett) \xrightarrow{g} (eet)(eett)$
- 2. $gg([and]) \circ [who brought what]: eett$
- 4. [who left]: et $\xrightarrow{\text{lift}}$ ((et)(et))(et) \xrightarrow{g} (e(et)(et))(eet) \xrightarrow{g} (ee(et)(et))(eeet)

A potential problem: It is predicted that a *wh*-expression can be bound as a pronoun.