

Model-theoretic Semantics: Names and verbs

Models

A model M for predicate logic is an ordered pair $\langle U, \llbracket \cdot \rrbracket \rangle$, in which:

- U is a nonempty set of individuals;
- $\llbracket \cdot \rrbracket$ is an interpretation function.

Ling-exp $\rightarrow \llbracket \cdot \rrbracket \rightarrow$ some object built from U

A sample model

$U = \{a, b, c\}$

a: a girl named *Ada*
b: a girl named *Becky*
c: a girl named *Cindy*

Facts

- a, b and c are girls.
- a and b danced, but c didn't.
- a saw b, but b didn't see a.
- c saw a, and a saw c, too.
- a and c are Americans, but b is a British.

Names

$\llbracket \text{Ada} \rrbracket = a$ $\llbracket \text{Becky} \rrbracket = b$ $\llbracket \text{Cindy} \rrbracket = c$

Type e

Intransitive verbs

$\llbracket \text{dance} \rrbracket = \{x \mid x \text{ dances}\} = \{a, b\} \subseteq U$

Characteristic functions

$\llbracket \text{dance} \rrbracket = \begin{bmatrix} a \rightarrow 1 \\ b \rightarrow 1 \\ c \rightarrow 0 \end{bmatrix}$

Type $e \rightarrow t$

Transitive verbs

$\llbracket \text{see} \rrbracket = \{\langle x, y \rangle \mid x \text{ saw } y\}$
 $= \{\langle a, b \rangle, \langle c, a \rangle, \langle a, c \rangle\} \subseteq U \times U$

Characteristic functions

$\llbracket \text{see} \rrbracket = \begin{bmatrix} \langle a, b \rangle \rightarrow 1 \\ \langle a, c \rangle \rightarrow 1 \\ \langle c, a \rangle \rightarrow 1 \\ \langle a, a \rangle \rightarrow 0 \\ \langle c, b \rangle \rightarrow 0 \\ \langle c, c \rangle \rightarrow 0 \\ \langle b, a \rangle \rightarrow 0 \\ \langle b, c \rangle \rightarrow 0 \\ \dots \end{bmatrix} = \begin{bmatrix} a \rightarrow \begin{bmatrix} a \rightarrow 1 \\ b \rightarrow 0 \\ c \rightarrow 0 \end{bmatrix} \\ b \rightarrow \begin{bmatrix} a \rightarrow 1 \\ b \rightarrow 0 \\ c \rightarrow 0 \end{bmatrix} \\ c \rightarrow \begin{bmatrix} a \rightarrow 1 \\ b \rightarrow 0 \\ c \rightarrow 0 \end{bmatrix} \end{bmatrix}$

Type $e \rightarrow e \rightarrow t$

Combination

$\begin{array}{c} t \\ \swarrow \quad \searrow \\ e \quad e \rightarrow t \\ \text{Ada} \quad \text{danced} \end{array}$

$\begin{array}{c} t \\ \swarrow \quad \searrow \\ e \quad e \rightarrow t \\ \text{Ada} \quad \begin{array}{c} \swarrow \quad \searrow \\ e \rightarrow e \rightarrow t \quad e \\ \text{saw} \quad \text{Becky} \end{array} \end{array}$

Functional Application

$\frac{f :: a \rightarrow b \quad x :: a}{f(x) :: b}$

1. $\llbracket \text{Ada danced} \rrbracket = \llbracket \text{danced} \rrbracket(\llbracket \text{Ada} \rrbracket)$
2. $\llbracket \text{Ada saw Becky} \rrbracket = \llbracket \text{saw} \rrbracket(\llbracket \text{Becky} \rrbracket)(\llbracket \text{Ada} \rrbracket)$

Formal concepts

Set “a collection of distinct objects, considered as an object in its own right”

Defining sets

- Extensionally: $\{a, b, c\}$
- Intensionally: $\{x \mid x \text{ is a girl}\}$

Empty set: \emptyset or $\{\}$

Relations between sets

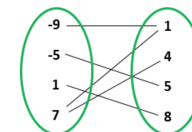
- Proper subset: $A \subset B$ iff $A \subseteq B$ and $B \not\subseteq A$
- Equivalence: $A = B$ iff $A \subseteq B$ and $B \subseteq A$
- Superset: $B \supseteq A$ iff $A \subseteq B$

Power sets $\wp(A) = \{B \mid B \subseteq A\}$

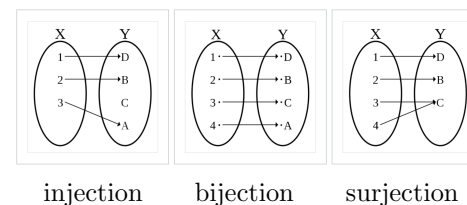
Pairs $\langle a, b \rangle \neq \langle b, a \rangle$ **n-tuples** $\langle a, b, c, \dots \rangle$

Relations as sets of pairs (Cartesian products)

$A \times B := \{\langle x, y \rangle \mid x \in A \text{ and } y \in B\}$



Function any relation where each input is paired with at most one output



injection

bijection

surjection

Type a ‘name tag’ of a set (D_e, D_t)

Basic types: e (individuals); t (truth values)

Function types: $a \rightarrow b$