

ACM/ICPC Template Manaual

浙江工商大学

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0 头文件

```
// 巨菜的ACMer-Happy233
   #include <bits/stdc++.h>
3
4
5 using namespace std;
6
7 //----
8 typedef double db;
9 typedef long long ll;
10 typedef unsigned int ui;
11 typedef vector<int> vi;
12 typedef pair<int, int> pii;
13 typedef pair<ll, ll> pll;
14 #define fi first
15 #define se second
16 #define pw(x) (111 << (x))
17 #define bt(x, k) (((x) >> k) & 1)
18 #define sz(x) ((int)(x).size())
19 #define all(x) (x).begin(),(x).end()
20 #define rep(i, l, r) for(int i=(l);i<(r);++i)</pre>
21 #define per(i, l, r) for(int i=(r)-1;i>=(l);--i)
22 #define mst(t, v, n) memset(t, v, sizeof(decltype(*(t))) * (n))
23 #define sf(x) scanf("%d", &(x))
24 #ifndef ACM_LOCAL
25 #define endl '\n'
26 #endif
27
28 int main() {
   #ifdef ACM_LOCAL
        freopen("./data/std.in", "r", stdin);
// freopen("./data/std.out", "w", stdout);
30
31
32 #else
        ios_base::sync_with_stdio(false);
33
        cin.tie(0);
34
        cout.tie(0);
35
36 #endif
37
   #ifdef ACM_LOCAL
38
39
        auto start = clock();
40
   #endif
41
        int t = 1;
   //
          cin >> t;
42
        while (t--)
43
44
            solve();
   #ifdef ACM_LOCAL
45
        auto end = clock();
46
        cerr << "Run Time: " << double(end - start) / CLOCKS_PER_SEC << "s" << endl;</pre>
47
   #endif
48
49
        return 0;
50 }
```

1 字串符

1.1 KMP

```
template<class elemType>
    inline void kmp_nxt(elemType &T, vector<int> &nxt) {
3
        nxt[0] = -1;
        for (int i = 1; i < T.size(); i++) {</pre>
4
            int j = nxt[i - 1];
5
            while (j \ge 0 \& T[i - 1] != T[j]) j = nxt[j];
6
            if (j \ge 0 \&\& T[i - 1] == T[j]) nxt[i] = j + 1;
7
8
            else nxt[i] = 0;
9
        }
10
   }
11
   template<class elemType>
12
    inline int kmp_count(elemType &S, elemType &T) {
        vector<int> nxt(T.size());
14
        kmp_nxt(T, nxt);
15
        int index, count = 0;
16
        for (index = 0; index < S.size(); ++index) {</pre>
17
            int pos = 0;
18
19
            int iter = index;
            while (pos < T.size() && iter < S.size()) {</pre>
20
                if (S[iter] == T[pos]) {
21
                     ++iter;
22
23
                     ++pos;
24
                } else {
25
                     if (pos == 0) ++iter;
                     else pos = nxt[pos - 1] + 1;
26
                }
27
28
29
            if (pos == T.size() && (iter - index) == T.size()) ++count;
30
31
        return count;
32
   }
33
   template<class elemType>
   inline void kmp_next(elemType T[], int count, vector<int> &nxt) {
36
        nxt[0] = -1;
37
        for (int i = 1; i < count; i++) {</pre>
38
            int j = nxt[i - 1];
            while (j \ge 0 \&\& T[i - 1] != T[j]) j = nxt[j];
39
            if (j \ge 0 \& T[i - 1] == T[j]) nxt[i] = j + 1;
40
            else nxt[i] = 0;
41
42
        }
   }
43
44
   template<class elemType>
45
   inline int kmp_count(elemType S[], int c1, elemType T[], int c2) {
        vector<int> nxt(c2);
47
48
        kmp_nxt(T, c2, nxt);
        int index, count = 0;
49
50
        for (index = 0; index < c1; ++index) {
51
            int pos = 0;
            int iter = index;
52
            while (pos < c2 && iter < c1) {</pre>
53
                if (S[iter] == T[pos]) {
54
                     ++iter;
55
```

```
56
                         ++pos;
                   }
else {
   if (pos == 0) ++iter;
   else pos = nxt[pos - 1] + 1;
57
58
59
60
61
               }
if (pos == c2 && (iter - index) == c2) ++count;
62
63
          }
64
65
          return count;
66 }
```

2 动态规划

2.1 01Bag

```
void dp(int n, int m) {
1
       // n=物品个数
2
       for (int i = 0; i < n; i++) {
3
          // m=背包最大容量
4
          for (int j = m; j >= wei[i]; j--)
5
6
              // wei=大小 val=价值
7
              f[j] = max(f[j], f[j - wei[i]] + val[i]);
8
       }
9
   }
   2.2 BagProblem
1 #define N 1000
  // val=价值 wei=重量 num=数量
  int val[N], wei[N], num[N], f[N];
4 // n=种类个数 m=背包最大值
5
  // 01背包
6
   void dp1(int n, int m) {
7
8
       for (int i = 0; i < n; i++) {
9
          for (int j = m; j >= wei[i]; j--)
10
              f[j] = max(f[j], f[j - wei[i]] + val[i]);
11
       }
   }
12
13
  // 完全背包
14
   void dp2(int n, int m) {
15
       //初始化看要求
16
       for (int i = 0; i <= m; i++) {
17
          f[i] = INF;
18
19
20
       f[0] = 0;
       //若要求恰好装满背包,那在初始化时除了f[0]=0其它f[1..V]均=-∞
21
22
       //若没要求背包装满,只希望价格大,初始化时应将f[0..V]=0)
23
       for (int i = 0; i < n; i++)
24
          for (int j = wei[i]; j <= m; j++)</pre>
              f[j] = max(f[j], f[j - wei[i]] + val[i]);
25
26 }
27
  // 多重背包
28
   void dp3(int n, int m) {
       for (int i = 0; i < n; i++)
30
31
          for (int k = 0; k < num[i]; k++)
              for (int j = m; j >= wei[i]; j--)
32
                  f[j] = max(f[j], f[j - wei[i]] + val[i]);
33
34 }
   2.3 FullBag
  /*
1
   完全背包问题的特点是,每种物品可以无限制的重复使用,可以选择放或不放。
   完全背包问题描述:
4 有N物品和一个容量为V的背包。第i件物品的重量是wei[i],价值是val[i]。
```

```
*/
5
6
   #include <cstdio>
7
   #define INF 0x3fffffff
8
   #define N 10047
9
  int f[N], val[N], wei[N];
  int min(int a,int b)
11
12
   {
       return x<y?x:y;</pre>
13
   }
14
  int main()
15
16
   {
       int t,i,j,k,E,F,m,n;
17
       scanf("%d",&t);
18
       while(t--)
19
20
          scanf("%d%d",&E,&F);
21
           int c = F-E;
22
           for(i = 0 ; i <= c ; i++)</pre>
23
              f[i]=INF;
24
          scanf("%d",&n);
25
           for(i = 0 ; i < n ; i++)
26
           {
27
28
              scanf("%d%d",&val[i],&wei[i]);//val[i]为面额,wei[i]为重量
29
30
          f[0]=0;//因为此处假设的是小猪储钱罐 恰好装满 的情况
          //注意初始化(要求恰好装满背包,那么在初始化时除了f[0]为0其它f[1..V]均设为-∞,
31
          //这样就可以保证最终得到的f[N]是一种恰好装满背包的最优解。
32
          //如果并没有要求必须把背包装满,而是只希望价格尽量大,初始化时应该将f[0..V]全部设为0)
33
34
          for(i =0; i < n; i++)
35
              for(j = wei[i] ; j <= c ; j++)</pre>
36
37
                  f[j] = min(f[j], f[j-wei[i]]+val[i]); //此处求的是最坏的情况所以用<math>min, 确定最少
38
      的钱,当然最后就用max了, HEHE
39
              }
          }
40
41
           if(f[c] == INF)
              printf("This is impossible.\n");
42
43
              printf("The minimum amount of money in the piggy-bank is %d.\n",f[c]);
44
45
46
       return 0;
47
   //此代码为HDU1114;
48
   2.4 MultiBag
1 //多重背包(MultiplePack): 有N种物品和一个容量为V的背包。
2 //第i种物品最多有n[i]件可用,每件费用是c[i],价值是w[i]。
3 //求解将哪些物品装入背包可使这些物品的费用总和不超过背包容量,
4 //且价值总和最大。
  //HDU 2191
5
6
  #include <cstdio>
7
8 #include <cstring>
9 #define N 247
10 int max(int x,int y){
```

```
11
       return x>y?x:y;
12
   int main() {
13
       int t,n,m,i,j,k;
14
       int w[N],pri[N],num[N],f[N];
15
       while(~scanf("%d",&t)){
16
           while(t--){
17
               memset(f,0,sizeof(f));
18
               scanf("%d%d",&n,&m);//n为总金额, m为大米种类
19
               for(i = 0 ; i < m ; i++){
20
21
                    scanf("%d%d%d",&pri[i],&w[i],&num[i]);//num[i]为每种大米的袋数
22
               for(i = 0 ; i < m ; i++){
23
                    for(k = 0 ; k < num[i] ; k++){
24
                        for(j = n ; j >= pri[i]; j--){
25
                            f[j] = max(f[j], f[j-pri[i]]+w[i]);
26
                        }
27
                    }
28
29
30
               printf("%d\n",f[n]);
           }
31
32
       return 0;
33
34 }
        Maze01
   2.5
1
   struct Maze01 {
       // 原始01矩阵 1-n 1-m
2
3
       bool a[N][N];
4
       // 以(i,j)向上最大高度的最大子矩阵的左右高
5
       int l[N][N], r[N][N], h[N][N];
6
       // 最大连续前缀1的左端位置,如果(i,j)为0则ml=j+1
7
       int ml[N][N];
8
       // 矩阵大小
9
       int n, m;
10
       void prework() {
11
12
            for (int i = 1; i \le m; i++) {
13
               l[0][i] = 0;
               r[0][i] = m;
14
15
16
           for (int i = 1; i <= n; i++) {
                int maxl = 1;
17
                int maxr = m;
18
               for (int j = 1; j <= m; j++) {
19
20
                    if (a[i][j] == 0) {
                        maxl = j + 1;
21
                        h[i][j] = l[i][j] = 0;
22
23
                    } else {
                        h[i][j] = h[i - 1][j] + 1;
24
                        l[i][j] = max(maxl, l[i - 1][j]);
25
26
27
                   ml[i - 1][j] = maxl;
                   ml[i][j] = j + 1;
28
29
               for (int j = m; j >= 1; --j) {
30
31
                    if (a[i][j] == 0) {
```

```
32
                         maxr = j - 1;
                         r[i][j] = m;
33
                    } else {
    r[i][j] = min(maxr, r[i - 1][j]);
34
35
36
37
                }
            }
38
39
        }
40
        // 单调栈
41
42
        pii st[N];
43
        int calc() {
44
            prework();
45
            int ans = 0;
46
            for (int i = 1; i <= n; i++) {
47
                int top = 0;
48
                for (int j = 1; j \le m; j++) {
49
                     if (h[i][j] == 0) continue;
50
                     if (top == 0 || st[top] != make_pair(l[i][j], r[i][j])) {
51
                         st[++top] = make_pair(l[i][j], r[i][j]);
52
53
54
                    while (top && st[top].second == j) {
55
                         int pos = st[top--].first;
56
                         if (pos < ml[i][j]) {</pre>
                             // l[i][j]-r[i][j]为底, h[i][j]为高
57
                             // 可以得到所有的唯一子矩阵,不存在maze1完全属于maze2
58
                             ans++;
59
                         }
60
                    }
61
                }
62
63
64
            return ans;
65
   } dp;
66
```

3 数据结构

3.1 BTree

```
template<class T>
1
2
3
   struct TreeNode {
        T value;
4
        TreeNode *left;
5
        TreeNode *right;
6
   };
7
8
   template<class T>
9
10
   TreeNode<T> *createTree(const T *pre, const T *in, const int len) {
        TreeNode<T> *t = NULL;
11
        if (len > 0) {
12
            t = new TreeNode<T>;
13
            t->value = pre[0];
14
15
            int index;
            for (index = 0; index < len; index++) {</pre>
16
                if (in[index] == pre[0]) {
17
                     break;
18
                }
19
20
            if (index == len) {
21
22
                index = -1;
23
            }
24
            t->left = createTree(pre + 1, in, index);
25
            t->right = createTree(pre + index + 1, in + index + 1, len - index - 1);
26
27
        return t;
28
   }
29
30
   template<class T>
31
   int preOrder(TreeNode<T> *root, queue<T> &out) {
32
        if (root) {
            int count = 1;
33
34
            out.push(root->value);
35
            count += pre0rder(root->left, out);
36
            count += pre0rder(root->right, out);
37
            return count;
        } else {
38
39
            return 0;
40
        }
   }
41
42
   template<class T>
43
   int inOrder(TreeNode<T> *root, queue<T> &out) {
44
        if (root) {
45
            int count = 1;
46
            count += inOrder(root->left, out);
47
48
            out.push(root->value);
            count += inOrder(root->right, out);
49
50
            return count;
51
        } else {
52
            return 0;
53
        }
54
   }
55
```

```
template<class T>
    void postOrder(TreeNode<T> *root, queue<T> &out) {
57
58
         if (root) {
             postOrder(root->left, out);
59
             postOrder(root->right, out);
60
             out.push(root->value);
61
         } else {
62
             return;
63
         }
64
    }
65
66
67
    template<class T>
    T *convertQueueToArray(queue<T> &out, int len) {
68
         T *list = new T[len];
69
         int now = 0;
70
         while (!out.empty() && now < len) {</pre>
71
72
             list[now] = out.front();
             out.pop();
73
74
             now++;
75
         return list;
76
    }
77
78
79
    template<class T>
80
    void destroyTree(TreeNode<T> *root) {
81
         if (root) {
             destroyTree(root->left);
82
             destroyTree(root->right);
83
             delete root;
84
85
         } else return;
    }
86
87
    template<class T>
88
    void insertIntoBSTree(TreeNode<T> *root, const T &value) {
89
         if (!root) {
90
             return;
91
92
93
         if (value < root->value) {
             if (root->left) {
94
                 insertIntoTree(root->left, value);
95
             } else {
96
                 root->left = new TreeNode<T>;
97
                 root->left->value = value;
98
99
                 root->left->left = NULL;
                 root->left->right = NULL;
100
             }
101
         } else if (value > root->value) {
102
             if (root->right) {
103
                 insertIntoTree(root->right, value);
104
105
106
                 root->right = new TreeNode<T>;
107
                 root->right->value = value;
108
                 root->right->left = NULL;
                 root->right->right = NULL;
109
             }
110
111
         }
112
113
114 template<class T>
```

```
TreeNode<T> *createBSTree(T *list, int len) {
        if (len < 1) {
116
            return NULL;
117
118
        TreeNode<T> *root = new TreeNode<char>;
119
        root->value = list[0];
120
        root->left = NULL;
121
        root->right = NULL;
122
        for (int i = 1; i < len; i++) {
123
            insertIntoBSTree(root, list[i]);
124
125
126
        return root;
127 }
    3.2 pbds-bbtree
 1 // RBTree 红黑树
 2 #include <ext/pb_ds/tree_policy.hpp>
 3 #include <ext/pb_ds/assoc_container.hpp>
 4 // 红黑树
   __gnu_pbds::tree<int, null_type, less<int>, rb_tree_tag,
       tree_order_statistics_node_update> t;
 6 // null_type无映射(低版本q++为null_mapped_type)
 7 // 类似multiset
    __gnu_pbds::tree<int, null_type, less_equal<int>, rb_tree_tag,
       tree_order_statistics_node_update> t;
 9 find_by_order(size_t order);
10 // 结点更新
11 tree_order_statistics_node_update
12 insert(p);
13 erase(it);
14 // 求k在树中是第几大:
15 order_of_key(p);
16 // 找到第order小的迭代器
17 find_by_order(order);
18 // 前驱
19 lower_bound(p);
20 // 后驭
21 upper_bound(p);
22 // 合并
23 a.join(b);
24 // 分割 key小于等于v的元素属于a. 其余的属于b
25 a.split(v, b);
27 // 优先队列
28 #include <ext/pb_ds/priority_queue.hpp>
29 #include <ext/pb_ds/assoc_container.hpp>
30 // 操作类似于stl的优先队列
   typedef __qnu_pbds::priority_queue<node, qreater<node>, __qnu_pbds::thin_heap_tag> heap
    heap::point_iterator; // 指向元素的指针
    3.3
         树状数组
 1 const int N = 1000005;
    struct BITree {
        int n;
```

```
ll c[N];
4
5
        void init(int _n) {
6
7
            n = _n + 1;
            memset(c, 0, sizeof(ll) * n);
8
        }
9
10
        void change(int pos, ll v) {
11
            for (int i = pos; i < n; i += i & (-i))
12
13
                c[i] += v;
        }
14
15
        11 query(int x) {
16
            11 \text{ ans} = 0;
17
            for (int i = x; i > 0; i -= i & (-i))
18
19
                ans += c[i];
20
            return ans;
        }
21
22
        void update(int l, int r, ll v) {
23
24
            change(1, v);
25
            change(r + 1, -v);
26
        }
27 } tree;
28
29 // 区间更新区间查询
   struct BITree {
30
31
        int n;
        ll c[N], d[N];
32
33
        void init(int _n) {
34
            n = _n + 1;
35
            memset(c, 0, sizeof(ll) * n);
36
37
            memset(d, 0, sizeof(ll) * n);
        }
38
39
40
        void change(int pos, ll v) {
            for (int i = pos; i < n; i += i & (-i))
41
                c[i] += v, d[i] += v * (pos - 1);
42
43
        }
44
        11 query(int x) {
45
            11 \text{ ans} = 0;
46
            for (int i = x; i > 0; i -= i & (-i))
47
48
                ans += x * c[i] - d[i];
49
            return ans;
50
        }
51
52
        void update(int l, int r, ll v) {
53
            change(1, v);
54
            change(r + 1, -v);
55
        }
56
        ll sum(int l, int r) {
57
            return query(r) - query(l - 1);
58
        }
59
60 } tree;
```

3.4 二维树状数组

```
1 const int N = 2005;
2
3
   inline int lowbit(const int &x) {
4
            return x & -x;
5
6
   struct TdBITree {
7
        int n, m;
8
        11 c[N][N];
9
10
        void init(int n, int m) {
11
12
            this->n = n;
            this->m = m;
13
14
            memset(c, 0, sizeof(c))
        }
15
16
        void init(int n, int m, ll v) {
17
18
            this -> n = n;
            this->m = m;
19
            rep(x, 1, N) {
20
                 rep(y, 1, N) {
21
                     c[x][y] = (x * y + (x - lowbit(x)) * (y - lowbit(y)) - x * (y - lowbit(x))
22
       y)) - (x - lowbit(x)) * y) * v;
23
24
            }
        }
25
26
27
        void change(int x, int y, ll v) {
            for (int i = x; i <= n; i += lowbit(i))</pre>
28
                 for (int j = y; j \leftarrow m; j \leftarrow lowbit(j))
29
30
                     c[i][j] += v;
        }
31
32
        11 query(int x, int y) {
33
            11 ans = 0;
34
            for (int i = x; i \rightarrow 1; i \rightarrow lowbit(i))
35
                 for (int j = y; j >= 1; j -= lowbit(j))
36
37
                     ans += c[i][j];
38
            return ans;
39
        }
40
        ll solve(int x1, int y1, int x2, int y2) {
41
            return query(x2, y2) - query(x1 - 1, y2) - query(x2, y1 - 1) + query(x1 - 1, y1
42
         - 1);
        }
43
44 };
         线段树
   3.5
1 // hdu 6562
2
   struct TreeNode {
3
        int 1, r;
        int lson, rson;
4
        ll sum;
5
        ll len;
6
        ll laz1, laz2, laz3;
```

```
8
        inline void init(int a, int b, int ls, int rs) {
9
            lson = ls;
10
11
            rson = rs;
            l = a, r = b;
12
            sum = 0;
13
            len = 1;
14
            laz1 = laz2 = 0;
15
            laz3 = 1;
16
        }
17
18
19
        inline int mid() {
            return (l + r) >> 1;
20
        }
21
22
        inline int width() {
23
            return r - l + 1;
24
        }
25
26
        inline void add(ll val) {
27
            sum = (sum * 10 + val * len * 10 + val * width()) % MOD;
28
29
            len = (len * 100) % MOD;
            // 右懒惰
30
31
            laz1 = (laz1 * 10 + val) % MOD;
32
            // 左懒惰
            laz2 = (laz2 + val * laz3) % MOD;
33
            // 长度懒惰
34
            laz3 = (laz3 * 10) % MOD;
35
        }
36
   };
37
38
   struct SegTree {
39
40
        int tot;
        TreeNode node[N << 1];</pre>
41
42
        inline void init() {
43
            tot = 0;
44
45
46
        inline void up(int k) {
47
            TreeNode &nd = node[k];
48
            nd.sum = (node[nd.lson].sum + node[nd.rson].sum) % MOD;
49
            nd.len = (node[nd.lson].len + node[nd.rson].len) % MOD;
50
51
        }
52
        inline void push(int k) {
53
            TreeNode &nd = node[k];
54
            if (nd.laz3 == 1) return;
55
            TreeNode &lson = node[nd.lson];
56
57
            TreeNode &rson = node[nd.rson];
58
59
            lson.sum = (nd.laz1 * lson.width() + lson.sum * nd.laz3 + nd.laz2 * lson.len %
       MOD * nd.laz3) % MOD;
            lson.len = (lson.len * nd.laz3 % MOD * nd.laz3) % MOD;
60
            lson.laz1 = (lson.laz1 * nd.laz3 + nd.laz1) % MOD;
61
            lson.laz2 = (nd.laz2 * lson.laz3 + lson.laz2) % MOD;
62
63
            lson.laz3 = (lson.laz3 * nd.laz3) % MOD;
64
```

```
rson.sum = (nd.laz1 * rson.width() + rson.sum * nd.laz3 + nd.laz2 * rson.len %
65
        MOD * nd.laz3) % MOD;
             rson.len = (rson.len * nd.laz3 % MOD * nd.laz3) % MOD;
66
             rson.laz1 = (rson.laz1 * nd.laz3 + nd.laz1) % MOD;
67
             rson.laz2 = (nd.laz2 * rson.laz3 + rson.laz2) % MOD;
68
             rson.laz3 = (rson.laz3 * nd.laz3) % MOD;
69
70
             nd.laz1 = nd.laz2 = 0;
71
             nd.laz3 = 1;
72
        }
73
74
75
        void build(int k, int l, int r) {
             TreeNode &nd = node[k];
76
             nd.init(l, r, tot + 1, tot + 2);
77
             tot += 2;
78
             if (l == r) {
79
80
                 return;
81
             int mid = nd.mid();
82
             build(nd.lson, nd.l, mid);
83
             build(nd.rson, mid + 1, nd.r);
84
85
             up(k);
        }
86
87
88
        void change(int k, int l, int r, ll val) {
             TreeNode &nd = node[k];
89
             if (nd.l == l && nd.r == r) {
90
                 nd.add(val);
91
                 return;
92
93
94
             push(k);
             int mid = nd.mid();
95
             if (r <= mid) {</pre>
96
                 change(nd.lson, l, r, val);
97
             } else if (l > mid) {
98
                 change(nd.rson, l, r, val);
99
100
             } else {
101
                 change(nd.lson, l, mid, val);
                 change(nd.rson, mid + 1, r, val);
102
103
             }
104
             up(k);
        }
105
106
        ll query(int k, int l, int r) {
107
108
             TreeNode &nd = node[k];
109
             if (nd.l == l \& nd.r == r) {
110
                 return nd.sum;
             }
111
             push(k);
112
113
             int mid = nd.mid();
114
             11 ans = 0;
115
             if (r <= mid) {
                 ans += query(nd.lson, l, r);
116
             } else if (l > mid) {
117
                 ans += query(nd.rson, l, r);
118
119
             } else {
120
                 ans += query(nd.lson, l, mid);
121
                 ans += query(nd.rson, mid + 1, r);
             }
122
```

```
123
             return ans % MOD;
124
125 } tree;
    3.6 二维线段树
    const int N = int(2e3 + 48);
 1
 2
 3
    ll p[N][N];
 4
    struct SegTree {
 5
 6
 7
        inline int static get(int l, int r) {
             return (l + r) | (l != r);
 8
        }
 9
 10
        struct node {
11
12
             int 1, r;
13
             inline int mid() { return (l + r) >> 1; }
14
15
             inline bool point() { return l == r; }
16
17
             inline node operator--(int) { return {1, mid()}; }
18
19
20
             inline node operator++(int) { return {mid() + 1, r}; }
21
            bool operator==(node t) const { return l == t.l && r == t.r; }
22
23
            operator int() { return get(l, r); }
24
        };
25
26
        ll s[N * 2][N * 2];
27
28
        inline void upX(node x, node y) {
29
             s[x][y] = max(s[x--][y], s[x++][y]);
30
31
        }
32
33
        inline void upY(node x, node y) {
34
             s[x][y] = max(s[x][y--], s[x][y++]);
        }
35
36
        void buildY(node x, node y) {
37
             if (y.point()) {
38
39
                 s[x][y] = p[x.1][y.1];
                 return;
40
41
             buildY(x, y--);
42
             buildY(x, y++);
43
             upY(x, y);
44
             return;
45
        }
46
47
        void mergeY(node x, node y) {
48
             if (y.point()) {
49
50
                 upX(x, y);
51
                 return;
             }
52
```

```
mergeY(x, y--);
53
            mergeY(x, y++);
54
            upX(x, y);
55
56
            return;
        }
57
58
        void build(node x, node y) {
59
            if (x.point()) {
60
                buildY(x, y);
61
                return;
62
63
64
            build(x--, y);
            build(x++, y);
65
            mergeY(x, y);
66
67
            return;
        }
68
69
        11 queryY(node x, node y, node a, node b) {
70
            if (y == b) {
71
72
                return s[x][y];
73
            int mid = y.mid();
74
            if (b.r <= mid) {</pre>
75
                return queryY(x, y--, a, b);
76
77
            } else if (b.l > mid) {
                return queryY(x, y++, a, b);
78
79
            } else {
                return max(queryY(x, y--, a, {b.l, mid}),
80
                            queryY(x, y++, a, \{mid + 1, b.r\}));
81
82
            }
        }
83
84
        11 query(node x, node y, node a, node b) {
85
            if (x == a) {
86
                return queryY(x, y, a, b);
87
            }
88
89
            int mid = x.mid();
90
            if (a.r <= mid) {
                return query(x--, y, a, b);
91
            } else if (a.l > mid) {
92
                return query(x++, y, a, b);
93
            } else {
94
                return max(query(x--, y, {a.1, mid}, b),
95
96
                            query(x++, y, \{mid + 1, a.r\}, b));
97
            }
98
99
  } tree;
         树状数组求逆序对
1 BITree t;
2 int n:
3 pii a[N];
4
   void solve() {
5
6
        t.init(n);
        for (int i = 1; i <= n; i++) {
7
8
            int x;
```

```
9
            cin >> x;
            a[i] = make_pair(x, i);
10
11
        sort(a + 1, a + n + 1);
12
13
        11 \text{ ans} = 0;
        for (int i = 1; i <= n; i++) {
14
            t.change(a[i].second, 1);
15
            ans += (i - t.query(a[i].second));
16
17
18
        cout << ans << endl;</pre>
19
   }
   3.8 ST
   // 只需要取值
1
   struct ST {
2
3
        int ck[N];
4
        int dp[20][N];
5
6
        void init(int n, int squ[]) {
7
            ++n;
            ck[0] = -1;
8
9
            for (int i = 1; i <= n; i++) {
                ck[i] = ck[i - 1] + ((i & (i - 1)) == 0 ? 1 : 0);
10
11
12
            for (int i = 0; i < n; i++) {
13
                dp[0][i] = squ[i];
14
            }
            for (int k = 1; k <= ck[n]; k++) {
15
                int dk = k - 1;
16
17
                for (int i = 0; i + (1 << k) - 1 < n; i++) {
18
                     dp[k][i] = max(dp[dk][i], dp[dk][i + (1 << dk)]);
19
                }
20
            }
21
        }
22
23
        int query(int 1, int r) {
            if (l > r) swap(l, r);
24
25
            int k = ck[r - l + 1];
26
            return max(dp[k][1], dp[k][r - (1 << k) + 1]);
27
        }
   };
28
29
   // 可得到下标
30
31
   struct ST {
        int ck[N];
32
33
        int rmq[N];
34
        int dp[20][N];
35
        void init(int n, int squ[]) {
36
37
            ck[0] = -1;
38
            for (int i = 1; i <= n; i++) {</pre>
39
                ck[i] = ck[i - 1] + ((i & (i - 1)) == 0 ? 1 : 0);
40
41
            memcpy(rmq, squ, sizeof(int) * n);
42
            for (int i = 0; i < n; i++) {
43
                dp[0][i] = i;
44
```

```
45
           for (int k = 1; k <= ck[n]; k++) {
46
               int dk = k - 1;
47
               for (int i = 0; i + (1 << k) - 1 < n; i++) {
48
                    int a = dp[dk][i];
49
                    int b = dp[dk][i + (1 << dk)];
50
                    dp[k][i] = rmq[a] < rmq[b] ? a : b;
51
               }
52
           }
53
       }
54
55
56
       int query(int 1, int r) {
           if (l > r) swap(l, r);
57
           int k = ck[r - l + 1];
58
           int a = dp[k][1];
59
           int b = dp[k][r - (1 << k) + 1];
60
           return rmq[a] < rmq[b] ? a : b;</pre>
61
       }
62
63
   };
        笛卡尔树
   3.9
   // 笛卡尔树, 静态建树, 区间最值跳转
   struct CartesianTree {
3
       int rt; // 根节点
       pii ch[N]; // 左右儿子
4
       int st[N]; // 单调栈
5
6
       void build(int n, int p□) {
7
8
           rt = 0;
           int t = 0:
9
10
           for (int i = 1; i <= n; i++) {
               ch[i] = \{0, 0\};
11
               // 决定了大于还是小于
12
               while (t && p[st[t]] > p[i]) --t;
13
14
               if (t) {
                    // 上一个点的右儿子作为自己的左儿子
15
                    // 成为上一个点的右儿子
16
                    ch[i].first = ch[st[t]].second;
17
                    ch[st[t]].second = i;
18
               } else { // 自己作为根节点
19
                    ch[i].first = rt;
20
21
                    rt = i;
22
23
               st[++t] = i;
           }
24
25
   } dika;
26
   3.10 DancingLinks
1 // Dancing Links
   struct DLX {
3
       int n, m, size;
       int U[MaxNode], D[MaxNode], L[MaxNode], R[MaxNode], Row[MaxNode], Col[MaxNode];
4
       int H[MaxN], S[MaxM];
5
       int ansd, ans[MaxN];
```

```
7
        void init(int _n, int _m) {
8
9
             n = _n;
10
            m = _m;
             for (int i = 0; i <= m; i++) {
11
12
                 S[i] = 0;
                 U[i] = D[i] = i;
13
                 L[i] = i - 1;
14
                 R[i] = i + 1;
15
16
17
            R[m] = 0;
18
            L[0] = m;
19
            size = m;
            for (int i = 0; i \le n; i++) {
20
                 H[i] = -1;
21
             }
22
        }
23
24
        void Link(int r, int c) {
25
26
             ++S[Col[++size] = c];
             Row[size] = r;
27
             D[size] = D[c];
28
29
            U[D[c]] = size;
30
            U[size] = c;
            D[c] = size;
31
             if (H[r] < 0) {
32
                 H[r] = L[size] = R[size] = size;
33
            } else {
34
                 R[size] = R[H[r]];
35
                 L[R[H[r]]] = size;
36
37
                 L[size] = H[r];
38
                 R[H[r]] = size;
39
            }
        }
40
41
        void remove(int c) {
42
43
             L[R[c]] = L[c];
44
            R[L[c]] = R[c];
             for (int i = D[c]; i != c; i = D[i]) {
45
                 for (int j = R[i]; j != i; j = R[j]) {
46
                     U[D[j]] = U[j];
47
                     D[U[j]] = D[j];
48
                      --S[Col[j]];
49
50
                 }
            }
51
52
        };
53
        void resume(int c) {
54
             for (int i = U[c]; i != c; i = U[i])
55
56
                 for (int j = L[i]; j != i; j = L[j])
57
                     ++S[Col[U[D[j]] = D[U[j]] = j]];
58
             L[R[c]] = R[L[c]] = c;
59
        }
60
        bool Dance(int d) {
61
62
             if (R[0] == 0) {
                 for (int i = 0; i < d; i++) {
    printf("%d%c", ans[i], " \n"[i == d - 1]);
63
64
                 }
65
```

```
66
                return true;
67
            int c = R[0];
68
            for (int i = R[0]; i != 0; i = R[i]) if (S[i] < S[c]) c = i;
69
            remove(c);
70
            for (int i = D[c]; i != c; i = D[i]) {
71
72
                ans[d] = Row[i];
                for (int j = R[i]; j != i; j = R[j])remove(Col[j]);
73
                if (Dance(d + 1))return true;
74
                for (int j = L[i]; j != i; j = L[j])resume(Col[j]);
75
76
            }
77
            resume(c);
            return false;
78
79
       }
   };
80
   3.11
          静态主席树
1 // m=update count, MAXN>=m*log(n)
   const int N = int(2e5 + 10);
3
  const int MAXN = int(1e7 + 10);
4
5
   struct PSegTree {
       const int *a;
6
7
       pii ran; // Tree Range
8
       int c[MAXN];
9
       int tot = 0;
       int lson[MAXN], rson[MAXN];
10
11
       int build(int 1, int r) {
12
13
            int k = ++tot;
14
            c[k] = 0;
15
            if (l == r) {
16
                return k;
17
            int mid = (l + r) \gg 1;
18
19
            lson[k] = build(l, mid);
20
            rson[k] = build(mid + 1, r);
21
            return k;
22
       }
23
       int init(int l, int r, const int num[]) {
24
25
            tot = 0;
26
            a = num;
27
            ran = \{1, r\};
            return build(l, r);
28
29
       }
30
       int update(int rt, int p, int v) {
31
            int k = ++tot, rst = k;
32
            int 1, r;
33
            tie(l, r) = ran;
34
35
            // calc
            c[k] = c[rt] + v;
36
            while (l < r) {
37
                int mid = (l + r) >> 1;
38
                // 下面的逗号表达式顺序不能换
39
                if (p <= mid) {
40
```

```
// go left
41
                     rson[k] = rson[rt], rt = lson[rt], k = lson[k] = ++tot;
42
                     r = mid;
43
                 } else {
44
                     // go right
45
                     lson[k] = lson[rt], rt = rson[rt], k = rson[k] = ++tot;
46
                     l = \overline{mid} + 1;
47
48
                 // calc
49
                 c[k] = c[rt] + v;
50
            }
51
52
            return rst;
        }
53
54
        // r1=right_root, r2=left_root, kth number
55
        int query(int r1, int r2, int k) {
56
57
            int l, r;
            tie(l, r) = ran;
58
            while (l < r) {
59
                 int mid = (l + r) \gg 1;
60
                 int cnt = c[lson[r1]] - c[lson[r2]];
61
                 if (cnt >= k) {
62
                     r1 = lson[r1], r2 = lson[r2];
63
                     r = mid;
64
65
                 } else {
66
                     k -= cnt;
                     r1 = rson[r1], r2 = rson[r2];
67
                     l = mid + 1;
68
69
70
            }
71
            return 1;
72
        }
73
        // r1=right_root, r2=left_root, sum of num<=k</pre>
74
        int sum(int r1, int r2, int k) {
75
            int 1, r;
76
77
            tie(l, r) = ran;
78
            int ans = 0;
            while (l < r) {
79
                 int mid = (l + r) \gg 1;
80
                 int cnt = c[lson[r1]] - c[lson[r2]];
81
                 if (k <= mid) {
82
                     r1 = lson[r1], r2 = lson[r2];
83
84
                     r = mid;
85
                 } else {
86
                     ans += cnt;
87
                     r1 = rson[r1], r2 = rson[r2];
                     l = mid + 1;
88
                 }
89
90
            }
91
            ans += c[r1] - c[r2];
92
            return ans;
93
94 } tree;
```

3.12 动态主席树

1 // m: update count, MAXN>=m*log(n)^2

```
2 const int N = int(2e5 + 10);
  const int MAXN = int(2e7 + 10);
   const int MAXM = int(2e7 + 10);
  const int LN = 40;
6
   struct PSegTree {
7
8
        const int *a;
9
        pii ran;
        int n;
10
        11 c[MAXN];
11
12
        int tot = 0;
13
        int lson[MAXN], rson[MAXN];
        // t: static root, s: dynamic root
14
        int s[N], t[N];
15
16
        // SegTree Range and n points, num can be nullptr
17
        void init(int l, int r, int _n, const int num[]) {
18
19
            tot = 0;
            a = num;
20
            ran = \{l, r\};
21
22
            n = _n;
23
            for (int i = 0; i \leftarrow n; i++) s[i] = t[i] = 0;
            // init node 0
24
25
            lson[0] = rson[0] = 0;
26
            c[0] = 0;
        }
27
28
29
        inline int get(int &k, bool flag) {
            if (k == 0 \mid | flag) {
30
31
                k = ++tot;
32
                lson[k] = rson[k] = 0;
33
34
            return k;
        }
35
36
        // update the root in k
37
38
        void update(int k[], int rt[], int cnt, int p, int v, bool flag = false) {
39
            // calc
            for (int i = 0; i < cnt; i++)
40
                c[k[i]] = c[rt[i]] + v;
41
            int 1, r;
42
            tie(l, r) = ran;
43
            while (l < r) {
44
                int mid = (l + r) >> 1;
45
                // 下面的逗号表达式顺序不能换
46
                if (p <= mid) {
47
                    // go left
48
                    for (int i = 0; i < cnt; i++) {
49
                         rson[k[i]] = rson[rt[i]];
50
51
                         rt[i] = lson[rt[i]], k[i] = get(lson[k[i]], flag);
52
                    }
53
                    r = mid;
54
                } else {
                    // go right
55
                    for (int i = 0; i < cnt; i++) {
56
                         lson[k[i]] = lson[rt[i]];
57
58
                         rt[i] = rson[rt[i]], k[i] = get(rson[k[i]], flag);
59
                    l = mid + 1;
60
```

```
61
                  // calc
62
                 for (int i = 0; i < cnt; i++)
63
64
                      c[k[i]] = c[rt[i]] + v;
65
             }
         }
66
67
         // build static tree
68
         void change(int pos, int p, int v) {
69
             // use int as the int[]
70
71
             // must use variable because I use the pointer
72
             int rt = t[pos - 1];
             int k = t[pos] = ++tot;
73
             update(&k, &rt, 1, p, v, true);
74
         }
75
76
         int use1[LN], use2[LN];
77
78
         // edit dynamic tree
79
         void add(int pos, int p, int v) {
80
             // memory reuse
81
             int *k = use1, *rt = use2;
82
             int cnt = 0;
83
84
             for (int i = pos; i \le n; i += (i \& -i), cnt++) {
85
                 rt[cnt] = s[i], k[cnt] = get(s[i], false);
86
87
             update(k, rt, cnt, p, v);
         }
88
89
         // calc lson value in use
90
         inline ll sum(int use[], int cnt) {
91
             11 ans = 0;
92
             for (int i = 0; i < cnt; i++)
93
                 ans += c[lson[use[i]]];
94
95
             return ans;
         }
96
97
98
         // calc value in use
         inline ll calc(int use[], int cnt) {
99
100
             ll ans = 0;
101
             for (int i = 0; i < cnt; i++)</pre>
                 ans += c[use[i]];
102
103
             return ans;
104
         }
105
106
         // ans=p1-p2
107
         11 querySum(int p1, int p2, int k) {
             int rt1 = t[p1], rt2 = t[p2];
108
             11 \text{ cnt} 1 = 0, \text{ cnt} 2 = 0;
109
             // calc root in need
110
111
             for (int i = p1; i; i -= (i & -i)) use1[cnt1++] = s[i];
112
             for (int i = p2; i; i -= (i \& -i)) use2[cnt2++] = s[i];
             int l, r;
113
             tie(l, r) = ran;
114
             11 \text{ ans} = 0;
115
116
             while (l < r) {
117
                 int mid = (l + r) \gg 1;
118
                 if (k <= mid) {
                      // go left
119
```

```
rt1 = lson[rt1], rt2 = lson[rt2];
120
                      for (int i = 0; i < cnt1; i++) use1[i] = lson[use1[i]];</pre>
121
                      for (int i = 0; i < cnt2; i++) use2[i] = lson[use2[i]];</pre>
122
123
                      r = mid;
                 } else {
124
                      // go right
125
                      ll cnt = sum(use1, cnt1) - sum(use2, cnt2) + c[lson[rt1]] - c[lson[rt2
126
        ]];
127
                      ans += cnt;
                      rt1 = rson[rt1], rt2 = rson[rt2];
128
129
                      for (int i = 0; i < cnt1; i++) use1[i] = rson[use1[i]];</pre>
130
                      for (int i = 0; i < cnt2; i++) use2[i] = rson[use2[i]];</pre>
131
                      l = mid + 1;
                 }
132
133
             ll cnt = calc(use1, cnt1) - calc(use2, cnt2) + c[rt1] - c[rt2];
134
135
             ans += cnt;
136
             return ans;
         }
137
138
         // query k
139
         int queryK(int p1, int p2, int k) {
140
             int r1 = t[p1], r2 = t[p2];
141
142
             int cnt1 = 0, cnt2 = 0;
143
             // calc root in need
             for (int i = p1; i; i -= (i & -i)) use1[cnt1++] = s[i];
144
             for (int i = p2; i; i -= (i & -i)) use2[cnt2++] = s[i];
145
             int 1, r;
146
             tie(l, r) = ran;
147
             while (l < r) {
148
                 int mid = (l + r) \gg 1;
149
                 int cnt = c[lson[r1]] - c[lson[r2]] + sum(use1, cnt1) - sum(use2, cnt2);
150
                 if (cnt >= k) {
151
                      // go left
152
                      r1 = lson[r1], r2 = lson[r2];
153
                      for (int i = 0; i < cnt1; i++) use1[i] = lson[use1[i]];</pre>
154
155
                      for (int i = 0; i < cnt2; i++) use2[i] = lson[use2[i]];</pre>
156
                      r = mid;
                 } else {
157
                      // go right
158
159
                      k -= cnt;
                      r1 = rson[r1], r2 = rson[r2];
160
                      for (int i = 0; i < cnt1; i++) use1[i] = rson[use1[i]];</pre>
161
                      for (int i = 0; i < cnt2; i++) use2[i] = rson[use2[i]];</pre>
162
                      l = mid + 1;
163
164
                 }
165
             }
             return 1;
166
167
   } tree;
168
    3.13
           伸展树 Splay
 1 const int N = int(1e6 + 100);
 2
 3 struct Node;
 4 Node *tail, *null;
 5 // 回收栈
```

```
Node *st[N], **stop;
7
   struct Node {
8
9
        // 内存池
        static Node pool[N];
10
11
12
        Node *ch[2], *fa;
        bool rev, edit;
13
        int size, val;
14
15
        ll sum;
        ll dpl, dpr, dpm; // 左端开始最大, 右端开始最大, 实际最大
16
17
18
        // 初始化内存池
        static void init() {
19
            null = tail = pool;
20
            null->clear();
21
22
            stop = st;
        }
23
24
25
        // 清空当前
        inline void clear(int _val = 0, int _size = 0) {
26
            edit = rev = false;
27
            sum = val = _val;
28
29
            size = _size;
30
            fa = ch[0] = ch[1] = null;
            dpl = dpr = dpm = \_val;
31
32
        }
33
        inline void push_up() {
34
            if (this == null) return;
35
            size = ch[0] -> size + ch[1] -> size + 1;
36
            sum = ch[0] -> sum + ch[1] -> sum + val;
37
            Node *x = ch[0], *y = ch[1];
38
            if (size == 1) {
39
                dpl = dpr = dpm = val;
40
            } else if (y == null) {
41
42
                dpl = max(x->dpl, x->sum + val);
43
                dpr = max(1ll * val, val + x->dpr);
                dpm = max(x->dpm, val + max(x->dpr, 011));
44
            } else if (x == null) {
45
                dpl = max(1ll * val, val + y->dpl);
46
                dpr = max(y->dpr, y->sum + val);
47
                dpm = max(y->dpm, val + max(y->dpl, 0ll));
48
49
            } else {
                dpl = max(x->dpl, max(x->sum + val, x->sum + val + y->dpl));
50
                dpr = max(y->dpr, max(y->sum + val, y->sum + val + x->dpr));
51
52
                dpm = max(max(x->dpm, y->dpm), val + max(x->dpr, 0ll) + max(y->dpl, 0ll));
53
            }
        }
54
55
56
        // 设置儿子
57
        inline void setc(Node *p, int d) { ch[d] = p, p->fa = this; }
58
        // 获取方向
59
        inline bool d() { return fa->ch[1] == this; }
60
61
62
        inline void update_rev() {
63
            if (this == null) return;
            swap(ch[0], ch[1]);
64
```

```
swap(dpl, dpr);
65
             rev ^= 1;
66
67
68
         inline void update_val(int v) {
69
             if (this == null) return;
70
             val = v;
71
             sum = 1ll * size * v;
72
73
             edit = 1;
             dpl = dpr = dpm = max(1ll * val, sum);
74
75
76
         inline void push_down() {
77
             if (this == null) return;
78
             if (rev) {
79
                  ch[0]->update_rev(), ch[1]->update_rev();
80
81
                  rev = 0;
82
             if (edit) {
83
                  ch[0]->update_val(val), ch[1]->update_val(val);
84
                  edit = 0;
85
             }
86
         }
87
88
89
         inline bool isroot() {
              return fa == null || (this != fa->ch[0] && this != fa->ch[1]);
90
         }
91
    };
92
93
    Node Node::pool[N];
94
95
    // 获取第k个
96
    Node *get_kth(Node *x, int k) {
97
         while (x->ch[0]->size + 1 != k) {
98
             x->push_down();
99
             if (x->ch[0]->size >= k) {
100
101
                  x = x -> ch[0];
102
             } else {
                  k -= x -> ch[0] -> size + 1;
103
104
                  x = x - ch[1];
105
             }
106
107
         return x;
108
    }
109
110
    void rotate(Node *x) {
         Node *f = x \rightarrow fa, *ff = f \rightarrow fa;
111
         f->push_down();
112
         x->push_down();
113
114
         int c = x \rightarrow d(), cc = f \rightarrow d();
115
         f->setc(x->ch[!c], c);
116
         x->setc(f, !c);
         if (ff->ch[cc] == f) ff->setc(x, cc);
117
         else x->fa=ff;
118
         f->push_up();
119
120
    }
121
    void splay(Node *&root, Node *x, Node *goal) {
122
         for (Node *f; (f = x \rightarrow fa) != goal; rotate(x)) {
123
```

```
if (f->fa == goal) continue;
124
             f->fa->push_down();
125
             f->push_down();
126
             x->push_down();
127
             rotate(x->d() == f->d() ? f : x);
128
        }
129
130
        x->push_up();
        if (goal == null) root = x;
131
    }
132
133
134
   int a[N];
135
136
    // 获取有效节点
    inline Node *getNode() { return stop == st ? ++tail : *--stop; }
137
138
    Node *build(int l, int r) {
139
        int mid = (l + r) \gg 1;
140
        Node *root = getNode();
141
        root->clear();
142
        root->val = a[mid];
143
        root->size = 1;
144
        if (l < mid) root->setc(build(l, mid - 1), 0);
145
        if (r > mid) root->setc(build(mid + 1, r), 1);
146
147
        root->push_up();
148
        return root;
149 }
150
    // l->root, r->root.ch[1]
151
    inline Node *make(Node *&root, int 1, int r) {
152
        Node *x = get_kth(root, 1);
153
        splay(root, x, null);
154
155
        Node *y = get_kth(root, r);
156
        splay(root, y, root);
        return y;
157
158
    }
159
160
   // 插入p子树
161
    inline void insert(Node *&root, int l, int r, Node *p) {
        Node *x = make(root, 1, r);
162
163
        x->setc(p, 0);
164
        x->push_up();
165
        x \rightarrow fa \rightarrow push_up();
    }
166
167
168
   // 回收内存
169
    inline void clear(Node *&x) {
170
        if (x == null) return;
        auto now = stop;
171
172
        *stop++ = x;
173
        while (now != stop) {
174
             x = *now++;
             if (x->ch[0] != null) *stop++ = x->ch[0];
175
             if (x->ch[1] != null) *stop++ = x->ch[1];
176
177
        }
    }
178
179
180
    // 释放子树
    inline void del(Node *&root, int l, int r) {
181
182
        Node *x = make(root, l, r);
```

```
183
         clear(x->ch[0]);
         x \rightarrow ch[0] = null;
184
         x->push_up();
185
186
         x \rightarrow fa \rightarrow push_up();
187
    }
188
    // 转向
189
    inline void reverse(Node *&root, int l, int r) {
190
         Node *x = make(root, l, r);
191
         x->ch[0]->update_rev();
192
193
         x->push_up();
194
         x \rightarrow fa \rightarrow push_up();
195
    }
196
    // set value
197
    inline void make_same(Node *&root, int l, int r, int val) {
198
         Node *x = make(root, 1, r);
199
200
         x->ch[0]->update_val(val);
         x->push_up();
201
202
         x->fa->push_up();
203 }
204
205 // 求和
206 inline ll sum(Node *&root, int l, int r) {
207
         Node *x = make(root, 1, r);
         return x->ch[0]->sum;
208
209
    }
210
211
    // 求最大子串和
    inline ll maxsum(Node *&root, int l, int r) {
212
213
         Node *x = make(root, l, r);
214
         return x->ch[0]->dpm;
215
    }
216
217
    void debug(Node *root) {
218
         function<void(Node *)> dfs = [&dfs](Node *x) {
219
             if (x == null) return;
220
             x->push_down();
             dfs(x->ch[0]);
221
             cout << x->val << ' ';
222
223
             dfs(x->ch[1]);
         };
224
         dfs(root);
225
226
         cout << endl;</pre>
227 }
            带权并查集
     3.14
 1 const int N = 100005;
    int f[N], deep[N], d[N];
 3
    void init(int n) {
 4
         for (int i = 0; i \le n; i++) f[i] = i, d[i] = 0;
 5
 6
 7
 8
    int fa(int x) {
         if (f[x] == x) return x;
 9
 10
         int fx = fa(f[x]);
```

```
(d[x] += d[f[x]]) \% = 3;
11
        return f[x] = fx;
12
   }
13
14
15
   bool un(int x, int y, int dis) {
        int fx = fa(x), fy = fa(y);
16
        if (fx == fy) {
17
            return ((d[y] - d[x] + 3) \% 3 != dis);
18
19
        int w = (3 + (d[x] - d[y]) + dis) \% 3;
20
21
        // 按轶合并
22
        if (deep[fx] < deep[fy]) {</pre>
23
            swap(fx, fy);
            W = (3 - W) \% 3;
24
25
        f[fy] = fx;
26
        d[fy] = w;
27
28
        // 按轶合并
        deep[fx] += deep[fx] == deep[fy];
29
30
        return true;
31 }
          可持久化并查集
   3.15
   const int N = int(2e5 + 10);
   const int MAXN = N * 150;
3
4
   struct PArray {
        int tot = 0;
5
        struct node {
6
7
            int val, deep;
        } c[MAXN];
8
9
        int lson[MAXN], rson[MAXN];
10
        pii ran;
11
        int build(int 1, int r) {
12
13
            int k = tot++;
14
            if (l == r) {
15
                c[k].val = 1;
16
                return k;
17
            }
            int mid = (l + r) \gg 1;
18
            lson[k] = build(l, mid);
19
            rson[k] = build(mid + 1, r);
20
21
            return k;
22
        }
23
24
        int init(int 1, int r) {
25
            ran = \{1, r\};
            return build(l, r);
26
        }
27
28
29
        int update(int rt, int p, node w) {
30
            int l = ran.first, r = ran.second;
            int k = tot++, rst = k;
31
            while (l_<_r) {
32
                lson[k] = lson[rt], rson[k] = rson[rt];
33
34
                int mid = (l + r) >> 1;
```

```
if (p <= mid) {</pre>
35
                                                             rt = lson[rt];
36
                                                             k = lson[k] = tot++;
37
38
                                                             r = mid;
39
                                                } else {
                                                             rt = rson[rt];
40
                                                             k = rson[k] = tot++;
41
                                                             l = mid + 1;
42
                                                }
43
44
                                    }
                                    c[k].val = w.val;
45
46
                                    c[k].deep = c[rt].deep + w.deep;
47
                                    return rst;
                       }
48
49
                       node query(int rt, int p) {
50
                                    int l = ran.first, r = ran.second;
51
                                    while (l < r) {
52
                                                int mid = (l + r) \gg 1;
53
                                                if (p <= mid) {
54
                                                             rt = lson[rt];
55
                                                             r = mid;
56
                                                } else {
57
58
                                                             rt = rson[rt];
59
                                                             l = mid + 1;
60
                                                }
61
62
                                    return c[rt];
                       }
63
64
          };
65
          struct PDSU {
66
67
                       PArray tree;
                       int root[N];
68
69
                       void init(int n) {
70
71
                                    root[0] = tree.init(1, n);
72
73
74
                       PArray::node fa(int rt, int x) {
75
                                    PArray::node u = tree.query(root[rt], x);
                                   while (u.val != x) {
76
77
                                                x = u.val;
78
                                                u = tree.query(root[rt], x);
79
80
                                    return u;
                       }
81
82
                       inline void change(int rt, int k) {
83
84
                                    root[k] = root[rt];
85
86
87
                       inline void un(int rt, int x, int y) {
                                    auto fx = fa(rt, x), fy = fa(rt, y);
88
                                    if (fx.val == fy.val) return;
89
                                   if (fx.deep > fy.deep) swap(fx, fy);
root[rt] = tree.update(root[rt], fx.val, {fy.val, fx.deep});
if (fx.deep == fy.deep) root[rt] = tree.update(root[rt], fy.val, fy.val
90
91
92
                      deep + 1);
```

```
93 }
94 } dsu;
```

4 图论

4.1 Graph

```
#define forg(i, h, eg) for(int i = (h); ~i; i = (eg[i]).first)
2
3
   struct Edge {
 4
        int to;
5
        11 w;
6
   };
7
   inline bool operator<(Edge a, Edge b) {</pre>
8
        return a.w < b.w;</pre>
9
  }
10
11
12 const ll INF = ll(1e11);
13 const int N = int(1e6 + 10);
14 const int M = int(2e6 + 10);
15
   struct Graph {
16
        pair<int, Edge> eg[M];
int head[N];
17
18
19
        int tot;
20
        void init(int n) {
21
            memset(head, -1, sizeof(int) * ++n);
22
23
            tot = 0;
24
        }
25
        inline void addEdge(int x, Edge g) {
26
            eg[tot] = \{head[x], g\};
27
            head[x] = tot++;
28
        }
29
30
        Edge &operator[](int p) {
31
32
            return eg[p].second;
33
34 } gh;
   4.2 Dijkstra
   int dist[N];
   int path[N];
2
3
   void bfs(int s, int n) {
4
5
        n++;
6
        rep(i, 0, n) dist[i] = INF;
        memset(path, -1, sizeof(int) * n);
7
        dist[s] = 0;
8
9
        path[s] = s;
        // 注意优先队列默认less运算,但选择最大的作为top,注意cmp!!!
10
        priority_queue<Edge, vector<Edge>, greater<Edge>> q;
11
12
        q.push(Edge(s, dist[s]));
13
        while (!q.empty()) {
            Edge f = q.top();
14
15
            q.pop();
            for (int i = gh.head[f.e]; \sim i; i = gh.eg[i].nxt) {
16
17
                Edge \&t = gh.eg[i];
```

```
if (dist[t.e] > f.v + t.v) {
18
                    dist[t.e] = f.v + t.v;
19
                    path[t.e] = f.e;
20
                    q.push(Edge(t.e, dist[t.e]));
21
22
                }
23
           }
24
       }
   }
25
26
27 #include <ext/pb_ds/priority_queue.hpp>
28 #include <ext/pb_ds/assoc_container.hpp>
29 typedef __gnu_pbds::priority_queue<Edge, greater<Edge>> heap;
30 // 使用该模板,需要注意因为使用了greater,所以需要重载大于运算
31 // 默认pairing_heap_tag
32 // push O(1), pop O(logn) modify O(logn) erase O(logn) join O(1)
33 // 可选thin_heap_tag
  // push O(1), pop O(\log n) modify O(1) erase O(\log n) join O(n)
35
36 heap::point_iterator its[N];
   int cnt[N];
37
38
   void bfs(int s, int n) {
39
40
       n++;
       rep(i, 0, n) dist[i] = INF;
41
       memset(cnt, 0, sizeof(int) * n);
42
43
       dist[s] = 0;
       cnt[s] = 1;
44
45
       heap q;
       its[s] = q.push(Edge(s, dist[s]));
46
       while (!q.empty()) {
47
           Edge f = q.top();
48
           q.pop();
49
           for (int i = gh.head[f.e]; ~i; i = gh.eg[i].nxt) {
50
                Edge &t = gh.eg[i];
51
                its[t.e] = 0;
52
                int v = f.v + t.v;
53
                if (dist[t.e] > v) {
54
55
                    dist[t.e] = v;
                    if (its[t.e] != 0) {
56
                        q.modify(its[t.e], Edge(t.e, dist[t.e]));
57
                    } else {
58
                        its[t.e] = q.push(Edge(t.e, dist[t.e]));
59
60
                    cnt[t.e] = cnt[f.e];
61
                } else if (dist[t.e] == v) {
62
63
                    (cnt[t.e] += cnt[f.e]) \% = 100003;
                }
64
           }
65
       }
66
67
  }
   4.3 spfa
1 vector<int> dist;
2
   vector<vector<node>> eg;
3
   vector<int> path;
  bool spfa(int n, int start) {
```

```
dist.assign(n, INF);
6
7
        dist[start] = 0;
        deque<int> q;
8
        q.push_back(start);
9
10
        path.assign(n, -1);
        vector<int> cnt(n, 0);
11
        vector<bool> flag(n, false);
12
        cnt[start] = flag[start] = true;
13
        while (!q.empty()) {
14
            const int now = q.front();
15
16
            q.pop_front();
17
            flag[now] = false;
            for (auto i: eg[now]) {
18
                if (dist[i.x] > dist[now] + i.d) {
19
                     dist[i.x] = dist[now] + i.d;
20
                     path[i.x] = now;
21
                     if (!flag[i.x]) {
22
                         if (n == ++cnt[i.x]) return false;
23
                         //队列非空且优于队首(SLF)
24
25
                         if (!q.empty() && dist[i.x] < dist[q.front()]) {</pre>
                             q.push_front(i.x);
26
27
                         } else {
28
                             q.push_back(i.x);
29
30
                         flag[i.x] = true;
                    }
31
32
                }
33
            }
34
35
        return true;
  }
36
   4.4 Dinic
   struct Graph {
1
        pair<int, Edge> eg[M];
2
3
        int head[N];
4
        int tot;
5
        void init(int n) {
6
            memset(head, -1, sizeof(int) * ++n);
7
8
            tot = 0;
9
        }
10
        inline void addEdge(int x, Edge g) {
11
12
            eg[tot] = \{head[x], g\};
            head[x] = tot++;
13
14
   } gh;
15
16
   struct Dinic {
17
        Graph qh;
18
19
        // 点的范围[0, n)
20
        int n;
        // 弧优化
21
        int cur[N], dis[N];
22
23
24
        Dinic() {};
```

```
25
        // 设置N
26
        void init(int _n) {
27
28
            n = _n + 1;
29
            gh.init(n);
        }
30
31
        // 加流量
32
        void addFlow(int x, int y, ll f) {
33
            printf("%d->%d: %lld\n", x, y, f);
34
35
            gh.addEdge(x, {y, f});
36
            gh.addEdge(y, \{x, 0\});
        }
37
38
        bool bfs(int s, int e) {
39
            memset(dis, -1, sizeof(int) * n);
40
41
            int q[N];
            int l, r;
42
            1 = r = 0;
43
            dis[s] = 0;
44
            q[r++] = s;
45
            while (l < r) {
46
                int f = q[l++];
47
                for (int i = gh.head[f]; ~i; i = gh.eg[i].first) {
48
49
                     auto &eg = gh.eg[i].second;
                     int to = eg.to;
50
                     if (eg.w > 0 \& dis[to] == -1) {
51
                         dis[to] = dis[f] + 1;
52
                         q[r++] = to;
53
                     }
54
55
                }
56
57
            return dis[e] > 0;
        }
58
59
        11 dfs(int s, int e, ll mx) {
60
            if (s == e | | mx == 0) {
61
62
                 return mx;
63
            ll flow = 0;
64
            for (int &k = cur[s]; ~k; k = gh.eg[k].first) {
65
                auto &eg = gh.eg[k].second;
66
67
68
                if (eg.w > 0 \& dis[eg.to] == dis[s] + 1 \& (a = dfs(eg.to, e, min(eg.w, mx))
       )))) {
69
                     eg.w -= a;
70
                     gh.eg[k \land 1].second.w += a;
                     flow += a;
71
                     mx -= a;
72
73
                     if (mx <= 0) break;
74
                }
75
            }
76
            return flow;
77
        }
78
        11 max_flow(int s, int e) {
79
80
            11 \text{ ans} = 0;
            while (bfs(s, e)) {
81
                memcpy(cur, gh.head, sizeof(int) * n);
82
```

```
ans += dfs(s, e, INF);
83
84
85
            return ans;
86
   } dinic;
   4.5 hungry
1 #define N 105
2 #define M 10005
3 int n, m, k;
  pii eg[M * 2];
  int result[N * 2];
6
   int head[N * 2];
   int cnt = 0;
7
8
   void addEdge(int x, int y) {
9
        eg[cnt].first = y;
10
        eg[cnt].second = head[x];
11
12
        head[x] = cnt++;
   }
13
14
   bool vis[M * 2] = {false};
15
16
   int dfs(int x) {
17
18
        for (int i = head[x]; \sim i; i = eg[i].second) {
            int y = eg[i].first;
19
20
            if (!vis[y]) {
21
                vis[y] = true;
                if (result[y] == -1 || dfs(result[y])) {
22
23
                     result[y] = x;
24
                     return 1;
25
                }
26
            }
27
28
        return 0;
29 }
30
31
   int MaxMatch() {
32
        int ans = 0;
33
        memset(result, -1, sizeof(result));
        rep(i, 1, n + 1) {
34
            memset(vis, 0, sizeof(vis));
35
36
            ans += dfs(i);
37
38
        return ans;
   }
39
40
   void solve() {
41
        scanf("%d%d", &m, &k);
42
        memset(head, -1, sizeof(head));
43
        cnt = 0;
44
        rep(i, 0, k) {
45
46
            int x, y;
            scanf("%d%d", &x, &y);
47
48
            addEdge(x, y);
49
50
        int ans = MaxMatch();
```

```
printf("%d\n", ans);
51
52 }
   4.6 MinSpanTree
1
   * Prim 求 MST
2
   * 耗费矩阵 cost[][], 标号从 0 开始, 0~n-1
3
   * 返回最小生成树的权值,返回 -1 表示原图不连通
6 const int INF = 0x3f3f3f3f;
   const int N = 110;
7
8 bool vis[N];
9 int lowc[N]; //点是 0 n-1
  int prim(int cost[][N], int n) {
       int ans = 0;
11
       memset(vis, false, sizeof(vis));
12
       vis[0] = true;
13
       for (int i = 1; i < n; i++)lowc[i] = cost[0][i];</pre>
14
       for (int i = 1; i < n; i++) {
15
            int minc = INF;
16
            int p = -1;
17
            19
18
            for (int j = 0; j < n; j++)
    if (!vis[j] && minc > lowc[j]) {
19
20
21
                    minc = lowc[j];
22
                    p = j;
23
            if (minc == INF)return -1;//原图不连通
24
25
            ans += minc;
            vis[p] = true;
26
27
            for (int j = 0; j < n; j++)
                if (!vis[j] && lowc[j] > cost[p][j])
28
29
                    lowc[j] = cost[p][j];
30
31
       return ans;
32
   }
   4.7 MinCostMaxFlow
   struct Edge {
2
       int e, nxt;
3
       ll flow, cost;
4
       Edge() {};
5
6
       Edge(int a, ll b, ll c, int d = \emptyset) : e(a), flow(b), cost(c), nxt(d) {}
7
   };
8
9
  const ll INF = ll(1e15);
11 const int N = int(2e3 + 10);
  const int M = int(1e6 + 10);
13
14
  //前向星
15
   struct Graph {
16
       Edge eg[M];
17
       int head[N];
```

```
int cnt;
18
19
        void init(int n) {
20
            memset(head, -1, sizeof(int) * ++n);
21
22
            cnt = 0;
        }
23
24
25
        inline void addEdge(int x, int y, ll v, ll c) {
26
            eg[cnt] = Edge(y, v, c, head[x]);
27
            head[x] = cnt++;
28
        }
29
   };
30
   struct MinCostMaxFlow {
31
32
        Graph gh;
        // 点的范围[0, n)
33
34
        int n;
35
        // 设置N
36
37
        void init(int _n) {
38
            n = _n + 1;
            gh.init(n);
39
        }
40
41
42
        // 加流量,反向是负的花费
        void addFlow(int x, int y, ll f, ll c) {
43
            // printf("%d->%d: %lld\t%lld\n", x, y, f, c); fflush(stdout);
44
            gh.addEdge(x, y, f, c);
45
            gh.addEdge(y, x, 0, -c);
46
        }
47
48
49
        // 该pre存的是边
        int pre[N];
50
        11 dis[N];
51
        bool vis[N];
52
53
54
        bool spfa(int s, int e) {
            queue<int> q;
55
            for (int i = 0; i < n; i++) {
56
                dis[i] = INF;
57
                vis[i] = false;
58
                pre[i] = -1;
59
60
61
            dis[s] = 0;
            vis[s] = true;
62
63
            q.push(s);
64
            while (!q.empty()) {
                int u = q.front();
65
                q.pop();
66
67
                vis[u] = false;
                for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
68
69
                     Edge &eg = gh.eg[i];
                     if (eg.flow > 0 && dis[eg.e] > dis[u] + eg.cost) {
70
                         dis[eg.e] = dis[u] + eg.cost;
71
72
                         pre[eg.e] = i;
                         if (!vis[eg.e]) {
73
74
                             vis[eg.e] = true;
75
                             q.push(eg.e);
                         }
76
```

```
}
77
                 }
78
79
80
             return pre[e] != -1;
        }
81
82
        pll cal(int s, int e) {
83
             ll flow = 0, cost = 0;
84
             while (spfa(s, e)) {
85
                 ll f = INF;
86
87
                 for (int i = pre[e]; ~i; i = pre[gh.eg[i ^ 1].e]) {
88
                     f = min(f, gh.eg[i].flow);
89
                 for (int i = pre[e]; ~i; i = pre[gh.eg[i ^ 1].e]) {
90
                     gh.eg[i].flow -= f;
91
                     gh.eg[i \land 1].flow += f;
92
93
                     cost += gh.eg[i].cost;
94
95
                 flow += f;
             }
96
             return make_pair(flow, cost);
97
        }
98
99
100 } network;
101
102 const int N = int(2e3 + 10);
    const ll INF = ll(1e15);
103
104
    struct Edge {
105
106
         int e, nxt;
107
        ll flow, cost;
108
109
        Edge() {};
110
        Edge(int a, ll b, ll c, int d = 0) : e(a), flow(b), cost(c), nxt(d) {}
111
    };
112
113
114
   // vector图存
    struct MinCostMaxFlow {
115
        vector<Edge> g[N];
116
117
        // 点的范围[0, n)
        int n = 0;
118
119
120
        // 设置N
        void init(int _n) {
121
             n = \_n + 1;
122
123
             rep(i, 0, n) {
                 g[i].clear();
124
             }
125
126
        }
127
128
        // 加流量,反向是负的花费
        void addFlow(int x, int y, ll f, ll c) {
129
             g[x].push_back(Edge(y, f, c, g[y].size()));
130
             g[y].push_back(Edge(x, 0, -c, g[x].size() - 1));
131
132
        }
133
134
        // 该pre存的是(点,边)
135
        pii pre[N];
```

```
11 dis[N];
136
         bool vis[N];
137
         ll h[N];
138
139
         bool bfs(int s, int e) {
140
             using pli = pair<ll, int>;
141
             priority_queue<pli, vector<pli>, greater<pli>> q;
142
             fill_n(dis, n, INF);
143
             fill_n(vis, n, false);
144
             fill_n(pre, n, pii(-1, -1));
145
             dis[s] = 0;
146
147
             q.push(pii(0, s));
             while (!q.empty()) {
148
                 pli f = q.top();
149
                 int u = f.second;
150
                 q.pop();
151
                 if (f.first != dis[u]) continue;
152
153
                 for (int i = 0; i < g[u].size(); i++) {</pre>
                      auto &eg = g[u][i];
154
                      if (eg.flow == 0) continue;
155
                      int v = eg.e;
156
                     ll cost = eg.cost + dis[u] + h[u] - h[v];
157
                      if (dis[v] > cost) {
158
                          dis[v] = cost;
159
160
                          pre[v] = pii(u, i);
                          q.push(pii(dis[v], v));
161
                      }
162
                 }
163
164
             for (int i = 0; i < n; i++) {
165
166
                 h[i] += dis[i];
167
             return pre[e].second != -1;
168
         }
169
170
         pll cal(int s, int e, ll limit) {
171
172
             ll flow = 0, cost = 0;
173
             fill_n(h, n, 0);
             while (limit) {
174
                 if (!bfs(s, e)) break;
175
                 ll f = INF;
176
                 for (int i = e; ~pre[i].second; i = pre[i].first) {
177
                      f = min(f, g[pre[i].first][pre[i].second].flow);
178
179
                 for (int i = e; ~pre[i].second; i = pre[i].first) {
180
                      g[pre[i].first][pre[i].second].flow -= f;
181
                      g[i][g[pre[i].first][pre[i].second].nxt].flow += f;
182
183
                 cost += f * h[e];
184
185
                 flow += f;
186
                 limit -= f;
187
             }
188
             return make_pair(flow, cost);
189
    } network;
190
```

4.8 ISAP

```
struct ISAP {
2
        Graph gh;
        // 点的范围[0, n)
3
4
        int n;
        // 弧优化
5
        int cur[N], dis[N];
6
        ISAP() {};
7
        // 设置N
8
9
        void init(int _n) {
10
            n = _n;
            gh.init(n);
11
12
        }
13
       // 加流量
14
        inline void addFlow(int x, int y, ll f) {
15
            gh.addEdge(x, y, f);
16
17
            gh.addEdge(y, x, 0);
        }
18
19
20
        int dep[N]; // 记录距离标号
21
        int gap[N]; // gap常数优化
22
        int q[N]; // 数组模拟队列
23
24
        void bfs(int s, int e) {
25
            memset(dep, -1, sizeof(int) * n);
            memset(gap, 0, sizeof(int) * n);
26
            gap[0] = 1;
27
            dep[e] = 0;
28
            int l = 0, r = 0;
29
            q[r++] = e;
30
            while (l < r) {
31
32
                int u = q[l++];
                for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
33
                     int v = gh.eg[i].e;
34
35
                     if (~dep[v]) continue;
                     q[r++] = v;
36
37
                     dep[v] = dep[u] + 1;
38
                     gap[dep[v]]++;
                }
39
40
            }
        }
41
42
        ll st[N]; // 栈优化
43
44
        11 max_flow(int s, int e) {
45
46
            bfs(s, e);
47
            memcpy(cur, gh.head, sizeof(int) * n);
            int top = 0;
48
            int u = s;
49
50
            11 \text{ ans} = 0;
51
            while (dep[s] < N) {</pre>
52
                if (u == e) {
                     ll mf = INF;
53
                     int sel = 0;
54
                     for (int i = 0; i < top; i++) {
55
                         if (mf > gh.eg[st[i]].v) {
56
                             mf = gh.eg[st[i]].v;
57
                             sel = i;
58
                         }
59
```

```
}
60
61
                     for (int i = 0; i < top; i++) {
62
                         gh.eg[st[i]].v -= mf;
63
64
                         gh.eg[st[i] ^ 1].v += mf;
65
                     ans += mf;
66
                     top = sel;
67
                     u = gh.eg[st[top] ^ 1].e;
68
69
                     continue;
70
                 }
71
                 bool flag = false;
72
                 int v = 0;
                 for (int i = cur[u]; ~i; i = gh.eg[i].nxt) {
73
                     v = gh.eg[i].e;
74
                     if (gh.eg[i].v > 0 \&\& dep[v] + 1 == dep[u]) {
75
                         flag = true;
cur[u] = i;
76
77
                         break;
78
                     }
79
80
                 if (flag) {
81
                     st[top++] = cur[u];
82
83
                     u = v;
84
                     continue;
85
                 int mind = N;
86
                 for (int i = gh.head[u]; ~i; i = gh.eg[i].nxt) {
87
                     if (gh.eg[i].v > 0 \&\& dep[gh.eg[i].e] < mind) {
88
                         mind = dep[gh.eg[i].e];
89
                         cur[u] = i;
90
                     }
91
92
                 }
                 gap[dep[u]]--; // 当前层无法连通, 降层
93
                 if (!gap[dep[u]]) return ans; // 断层结束运算
94
                 dep[u] = mind + 1; // 进入更高层
95
96
                 gap[dep[u]]++;
97
                 if (u != s) u = gh.eg[st[--top] ^ 1].e;
            }
98
99
            return ans;
100
        }
101 } isap;
         树链剖分
    4.9
    struct TreeChain {
 1
        int top[N]; // 链条顶端点ID
 2
 3
        int fa[N]; // 父亲节点
        int son[N]; // 重儿子
 4
        int deep[N]; // 深度
 5
        int num[N]; // 儿子节点数 (包括自己)
 6
 7
 8
        int p[N]; // 在线段树中的ID,
 9
        int fp[N]; // 线段树中ID对应的点
        int tot;
10
11
        void dfs(int u, int pre, int d) {
12
13
            num[u] = 1;
```

```
deep[u] = d;
14
            fa[u] = pre;
15
            son[u] = -1;
16
            for (int i = gh.head[u]; \sim i; i = gh.eg[i].first) {
17
                int v = gh.eg[i].second.to;
18
                if (v == pre) continue;
19
20
                dfs(v, u, d + 1);
                num[u] += num[v];
21
22
                if (son[u] == -1 \mid l \mid num[v] > num[son[u]]) {
                    son[u] = v;
23
24
                }
25
            }
       }
26
27
       void getpos(int u, int sp) {
28
            top[u] = sp;
29
            p[u] = tot++;
30
31
            fp[p[u]] = u;
            if (son[u] == -1) return;
32
            getpos(son[u], sp);
33
            for (int i = gh.head[u]; ~i; i = gh.eg[i].first) {
34
                int v = gh.eg[i].second.to;
35
                if (v == son[u] || v == fa[u]) continue;
36
37
                getpos(v, v);
38
            }
       }
39
40
       // start是线段树中的ID起始数值
41
       void build(int n, int root = 1, int start = 1) {
42
43
            tot = start;
            // deep起始值如果是0,则遇到数据错误超出这棵树,有可能死循环
44
            // dfs(root, 0, 0);
45
46
            dfs(root, 0, 1);
            getpos(root, root);
47
       }
48
49
50
       int lca(int u, int v) {
51
            int f1 = top[u];
            int f2 = top[v];
52
            while (f1 != f2) {
53
                if (deep[f1] < deep[f2]) {</pre>
54
                    swap(f1, f2);
55
56
                    swap(u, v);
57
                }
                u = fa[f1];
58
59
                f1 = top[u];
60
            return deep[u] < deep[v] ? u : v;</pre>
61
62
63 } tch;
64
65
   // 树状数组, 如果无需在线查询可以使用差分树
66
   BITree tree;
67
   // 点权修改
68
   void change1(int u, int v, ll val) {
70
       int f1 = tch.top[u];
71
       int f2 = tch.top[v];
       while (f1 != f2) {
72
```

```
if (tch.deep[f1] < tch.deep[f2]) {</pre>
73
                 swap(f1, f2);
74
75
                 swap(u, v);
76
             tree1.update(tch.p[f1], tch.p[u], val);
77
             u = tch.fa[f1];
78
             f1 = tch.top[u];
79
80
         if (tch.deep[u] > tch.deep[v]) {
81
             swap(u, v);
82
83
84
         tree1.update(tch.p[u], tch.p[v], val);
    }
85
86
87
    // 边权修改
    void change2(int u, int v, ll val) {
88
         int f1 = tch.top[u];
89
90
         int f2 = tch.top[v];
         while (f1 != f2) {
91
             if (tch.deep[f1] < tch.deep[f2]) {</pre>
92
                 swap(f1, f2);
93
                 swap(u, v);
94
95
96
             tree2.update(tch.p[f1], tch.p[u], val);
97
             u = tch.fa[f1];
             f1 = tch.top[u];
98
99
         if (tch.deep[u] > tch.deep[v]) {
100
             swap(u, v);
101
102
         tree2.update(tch.p[tch.son[u]], tch.p[v], val);
103
104 }
    4.10
           倍增 LCA
    const int MAX_DEP = 20;
 2
 3
    // 倍增2^k的父亲
   int fa[N][MAX_DEP];
 5
   int dep[N];
 6
    // 倍增LCA
 7
    int lca(int u, int v) {
   if (dep[u] > dep[v]) {
 8
 9
10
             swap(u, v);
11
         int hu = dep[u], hv = dep[v];
12
13
         int tu = u, tv = v;
         for (int det = hv - hu, i = 0; det; det >>= 1, i++) {
14
             if (det & 1)
15
                 tv = fa[tv][i];
16
17
18
         if (tu == tv) {
19
             return tu;
20
         for (int i = MAX_DEP - 1; i >= 0; i--) {
21
             if (fa[tu][i] == fa[tv][i]) {
22
23
                 continue;
```

```
}
24
            tu = fa[tu][i];
25
26
            tv = fa[tv][i];
27
28
        return fa[tu][0];
29
   }
30
31
   // 动态更新节点的父亲属性
   void lineFa(int u, int v) {
32
33
        fa[u][0] = v;
34
        for (int i = 1; i < MAX_DEP; i++) {</pre>
35
            v = fa[u][i] = fa[v][i - 1];
36
   }
37
   4.11
           Tarjan
   int dfn[N], low[N], st[N], belong[N], num[N];
   bool inst[N];
2
   int idx, top, scc;
3
   void tarjan(int u) {
5
        dfn[u] = low[u] = ++idx;
6
        st[top++] = u;
7
        inst[u] = true;
8
        for (int i = qh.head[u]; i != -1; i = qh.eq[i].nxt) {
9
10
            int v = gh.eg[i].e;
            if (!dfn[v]) {
11
                tarjan(v);
12
                low[u] = min(low[u], low[v]);
13
            } else if (inst[v]) {
14
                low[u] = min(low[u], dfn[v]);
15
16
17
        }
18
        int v;
        if (dfn[u] == low[u]) {
19
20
            SCC++;
21
            do {
                v = st[--top];
22
23
                inst[v] = false;
                belong[v] = scc;
24
25
                num[scc]++;
26
            } while (u != v);
        }
27
   }
28
29
30
   void work() {
        fill_n(dfn, n + 1, 0);
31
        fill_n(num, n + 1, 0);
32
        fill_n(inst, n + 1, false);
33
        idx = top = scc = 0;
34
        for (int i = 1; i <= n; i++) {
35
            if (!dfn[i])
36
37
                tarjan(i);
38
        }
   }
39
```

4.12 支配树

```
const int MAX_DEP = 20;
1
2
3
   // 注意0,1点的边界问题
   struct DominatorTree {
4
       int deg[N]; // 入度
5
6
        int dep[N]; //
7
        int dfn[N];
        int st[N];
8
9
        int tot;
10
        // 拓扑序,要保证root是入度为0
11
12
        void bfs(Graph &gh, int root) {
13
            queue<int> q;
            q.push(root);
14
            tot = 0;
15
            while (!q.empty()) {
16
17
                int u = q.front();
                q.pop();
18
                dfn[u] = ++tot;
19
20
                st[tot] = u;
                forg(i, gh.head[u], gh.eg) {
21
                    int v = gh.eg[i].e;
22
                    if ((--deg[v]) == 0) {
23
24
                        q.push(v);
25
26
                }
27
            }
28
        }
29
30
        // 倍增2^k的父亲
        int fa[N][MAX_DEP];
31
32
33
        // 倍增LCA
34
        int lca(int u, int v) {
            if (dep[u] > dep[v]) {
35
36
                swap(u, v);
37
38
            int hu = dep[u], hv = dep[v];
39
            int tu = u, tv = v;
            for (int det = hv - hu, i = 0; det; det >>= 1, i++) {
40
                if (det & 1)
41
                    tv = fa[tv][i];
42
43
            if (tu == tv) {
44
                return tu;
45
46
            for (int i = MAX_DEP - 1; i >= 0; i--) {
47
                if (fa[tu][i] == fa[tv][i]) {
48
49
                    continue;
                }
50
                tu = fa[tu][i];
51
52
                tv = fa[tv][i];
53
            }
54
            return fa[tu][0];
55
        }
56
        // 动态更新节点的父亲属性
57
```

```
void lineFa(int u, int v) {
58
            fa[u][0] = v;
59
            for (int i = 1; i < MAX_DEP; i++) {</pre>
60
                v = fa[u][i] = fa[v][i - 1];
61
62
       }
63
64
       // 建树, op是gh的反向图, 用来寻找其父亲
65
       void build(Graph &gh, Graph &op, int n, int root) {
66
            memcpy(deg, gh.deg, sizeof(int) * (n + 1));
67
            bfs(gh, root);
68
69
            for (int k = 1; k \le tot; k++) {
                int u = st[k], fath = -1;
70
                dep[u] = 0;
71
                for (int i = op.head[u]; \sim i; i = op.eg[i].nxt) {
72
                    int v = op.eg[i].e;
73
                    if (dfn[v] > dfn[u]) continue;
74
                    fath = (fath == -1 ? v : lca(fath, v));
75
76
                if (fath == -1) fath = u;
77
                lineFa(u, fath);
78
79
                dep[u] = dep[fath] + 1;
           }
80
81
   } dtree;
   4.13 Hopcroft-Karp
1 int dis;
   // linkx: x链接的y, linky: y链接的x
  int linkx[N], linky[N];
   int dx[N], dy[N];
5
   bool vis[N];
6
   bool searchP(int n) {
7
       queue<int> q;
8
9
       dis = INF;
10
       mst(dx, -1, n);
11
       mst(dy, -1, n);
12
       for (int i = 0; i < n; i++) {
            if (linkx[i] == -1) {
13
14
                q.push(i);
15
                dx[i] = 0;
            }
16
17
       while (!q.empty()) {
18
            int u = q.front();
19
            q.pop();
20
            if (dx[u] > dis) break;
21
            for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
22
23
                int v = gh.eg[i].e;
                if (dy[v] == -1) {
24
                    dy[v] = dx[u] + 1;
25
                    if (linky[v] == -1) {
26
                        dis = dy[v];
27
28
                    } else {
                        dx[linky[v]] = dy[v] + 1;
29
30
                        q.push(linky[v]);
```

```
}
31
                 }
32
             }
33
34
35
        return dis != INF;
   }
36
37
   bool dfs(int u) {
38
        for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
39
40
             int v = gh.eg[i].e;
             if (!vis[v] \&\& dy[v] == dx[u] + 1) {
41
42
                 vis[v] = true;
                 if (linky[v] != -1 && dy[v] == dis) continue;
43
                 if (\lim_{v \to \infty} v) = -1 \mid \inf_{v \to \infty} (\lim_{v \to \infty} v) {
44
                      linky[v] = u;
45
                      linkx[u] = v;
46
47
                      return true;
                 }
48
             }
49
50
        return false;
51
   }
52
53
54
   int MaxMatch(int n) {
55
        int rst = 0;
        mst(linkx, -1, n);
56
        mst(linky, -1, n);
57
        while (searchP(n)) {
58
             mst(vis, false, n);
for (int i = 0; i < n; i++) {</pre>
59
60
                 if (linkx[i] == -1 \&\& dfs(i)) {
61
62
                      rst++;
63
                 }
             }
64
65
        return rst;
66
67
   }
    4.14 Kuhn-Munkres
   const int N = 1005;
   const ll INF = ll(2e11);
2
3
   11 cost[N][N];
  // lx, ly 表示当前期望值, slack是得到链接的期望值
ll lx[N], ly[N], slack[N];
5
   // prey->alternating tree, match是y->x的链接
   int prey[N], match[N];
   bool vy[N]; // 是否已经变化的点
9
10
   void augment(int n, int root) {
11
        fill_n(vy + 1, n, false);
12
13
        fill_n(slack + 1, n, INF);
14
        int py = 0;
        match[py] = root;
15
16
        do {
             vy[py] = true;
17
             int x = match[py], yy = 0;
18
```

```
11 delta = INF;
19
             for (int y = 1; y <= n; y++) {
    if (vy[y]) continue;</pre>
20
21
                 ll w = lx[x] + ly[y] - cost[x][y];
22
                 if (w < slack[y])</pre>
23
24
                      slack[y] = w, prey[y] = py;
                 if (slack[y] < delta) delta = slack[y], yy = y;</pre>
25
26
27
             for (int y = 0; y <= n; y++) {
28
                 if (vy[y])
                      lx[match[y]] -= delta, ly[y] += delta;
29
30
                      slack[y] -= delta;
31
            }
32
33
            py = yy;
        } while (match[py] != -1);
34
35
             int pre = prey[py];
36
37
            match[py] = match[pre], py = pre;
38
        } while (py);
39 }
40
   11 KM(int n) {
41
42
        fill_n(match + 1, n, -1);
43
        fill_n(ly + 1, n, 0);
        for (int i = 1; i <= n; i++) {
44
            lx[i] = *max\_element(cost[i] + 1, cost[i] + n + 1);
45
46
        il ans = 0;
47
        for (int root = 1; root <= n; root++) augment(n, root);</pre>
48
        for (int i = 1; i <= n; i++) ans += lx[i], ans += ly[i];
49
50
        return ans;
51 }
```

5 博弈

5.1 GameProblem

```
// 巴什博奕,是否先手必胜
   inline bool bash_game(int n, int m) {
       //一堆东西, n个物品,最多选m个
3
       return n % (m + 1);
4
   }
5
6
7
  // 威佐夫博弈,是否先手必胜
  // 有两堆各若干的物品,两人轮流从其中一堆取至少一件物品,至多不限,或从两堆中同时取相同件物品,规定最后
      取完者胜利。
   inline bool wythoff_game(int n, int m) {
9
       if (n > m) {
10
          swap(n, m);
11
12
       int temp = floor((n2 - n1) * (1 + sqrt(5.0)) / 2.0);
13
       return temp != n1;
14
   }
15
16 // SG函数
17 #define N 1001
18 //f[]:可以取走的石子个数
19 //sg[]:0~n的SG函数值
20 int f[N], sg[N], mex[N];
21
  void getSG(int n) {
22
23
       int i, j;
       memset(sq, 0, sizeof(sq));
24
25
       for (i = 1; i <= n; i++) {
          memset(mex, 0, sizeof(mex));
26
          for (j = 1; f[j] \le i; j++)
27
              mex[sg[i - f[j]]] = 1;
28
          for (j = 0; j <= n; j++) { //求mes{}中未出现的最小的非负整数 if (mex[j] == 0) {
29
30
31
                  sg[i] = j;
32
                  break;
33
              }
34
          }
       }
35
36 }
37
38 // Auti-nim 反尼姆游戏
39 // 当先拿完所有石子时候输
40 // 当如下条件时, 先手必胜
41 // [: 所有堆的石子数均=1, 且有偶数堆。
42 // □:至少有一个堆的石子数>1,且石子堆的异或和≠0。
```

6 分治

6.1 IntegerFastPower

```
1 ll fpow(ll x, ll k) {
2          ll base = x, r = 1;
3          for (; k; k >>= 1) {
4              if (k & 1) r = r * base;
5              base = base * base;
6          }
7          return r;
8 }
```

6.2 MatrixFastPower

```
#define MAX_N 10
2
   #define mod_num 9973
3
   struct Mat {
4
5
        long long mat[MAX_N][MAX_N];
        long long n;
6
7
        Mat() {
8
            memset(mat, 0, sizeof(mat));
9
            n = 0;
10
11
        Mat(long long n) {
12
            memset(mat, 0, sizeof(mat));
13
            this->n = n;
14
        void init() {
15
            for (int i = 0; i < n; ++i) {</pre>
16
17
                mat[i][i] = 1;
18
19
        Mat(const long long ** list, long long n) {
20
21
            this->n = n;
22
            for (int i = 0; i < n; ++i) {
                for (int j = 0; j < n; ++j) {
23
24
                     mat[i][j] = list[i][j];
25
26
            }
        }
27
28
   };
29
   Mat operator * (Mat a, Mat b) {
30
        long long n = a.n;
31
        Mat c(n);
32
        memset(c.mat, 0, sizeof(c.mat));
33
        for (int i = 0; i < n; ++i) {
34
            for (int j = 0; j < n; ++j) {
35
                for (int k = 0; k < n; ++k) {
36
                     c.mat[i][j] += (a.mat[i][k] * b.mat[k][j]) % mod_num;
37
38
                     c.mat[i][j] %= mod_num;
39
                }
            }
40
41
42
        return c;
43
   }
```

```
44
   Mat operator ^ (Mat a, int k) {
45
       long long n = a.n;
46
47
       Mat c(n);
48
       c.init();
       for (; k; k >>= 1) {
49
            if (k & 1)
50
                        c = c * a;
            a = a * a;
51
       }
52
53
       return c;
54 }
   6.3 三维 CDQ
   struct node {
2
       // time: 时间 | id: 0,1 是否修改 | f: 正负 | x是下标 | y是权值
3
       int time, id, f, x, y;
   };
4
5
6
   bool cmp(const node &a, const node &b) {
       return a.x < b.x;</pre>
7
   }
8
9
10 int ans[N];
   node p[N \ll 2], et[N \ll 2];
12
13
   void cdq(int 1, int r) {
14
       if (l + 1 == r) return;
       int mid = (l + r) \gg 1;
15
       cdq(l, mid), cdq(mid, r);
16
17
       int t = 1;
       for (int i = mid; i < r; i++) {
18
19
            // if edit continue
20
            if (p[i].id)continue;
            for (; t < mid && p[t].x <= p[i].x; t++) {</pre>
21
22
                if (p[t].id) tree.change(p[t].y, p[t].f);;
23
            int f = p[i].f;
24
            int cnt = tree.query(p[i].y);
25
            ans[p[i].time] += f * cnt;
26
27
       // 逆操作p[t].y
28
       while (--t >= 1) {
29
            if (p[t].id) tree.change(p[t].y, -p[t].f);
30
31
       // 归并排序
32
       int t1 = 1, t2 = mid, k = 0;
33
       while (t1 < mid && t2 < r) \{
34
            et[k++] = p[t1].x < p[t2].x ? p[t1++] : p[t2++];
35
       }
36
       copy(p + t1, p + mid, et + k);
37
38
       copy(p + t2, p + r, et + k);
39
       copy(et, et + (r - 1), p + 1);
40
   }
```

6.4 树分治-点分治

```
1 // 题意: n个节点的树, 存在边权, 范围1e18
2 // 求任意两点之间点集的子集中两点之间路径异或和为0的个数
  // u<v,u'<v',(u',v') [ path(u,v),求path(u', v')异或和==0
5
   struct Edge {
       int to, nxt;
6
       11 w;
7
   };
8
9 const int N = int(1e5 + 10);
10 const int M = N \ll 1;
12 struct Grahp {
       int head[N];
13
       Edge eg[M];
14
       int tot;
15
16
       void init(int n) {
17
           memset(head, -1, sizeof(int) * ++n);
18
19
       }
20
       inline void addEdge(int u, int v, ll w) {
21
           eg[tot] = \{v, head[u], w\};
22
23
           head[u] = tot++;
25 } gh;
26
   bool vis[N];
27
   // q队列,fa祖先,sz是子树大小,smx是子树最大int q[N],fa[N],sz[N],smx[N];
30
   int froot(int s) {
31
32
       int l, r, mn = N, rt = 0;
33
       q[l = r = 1] = s;
       while (l <= r) {</pre>
34
           int u = q[l++];
35
           sz[u] = 1;
36
37
           smx[u] = 0;
38
           for (int i = gh.head[u]; ~i; i = gh.eg[i].nxt) {
               int v = gh.eg[i].to;
39
               if (v == fa[u] || vis[v]) continue;
40
41
               fa[v] = u;
               q[++r] = v;
42
           }
43
44
       // 反向遍历所有点算size
45
46
       while (--1) {
           int u = q[l];
47
           int mx = max(smx[u], r - sz[u]);
48
           if (mx < mn) mn = mx, rt = u;
49
           if (l == 1) break; // 根节点没有fa
51
           sz[fa[u]] += sz[u];
52
           smx[fa[u]] = max(smx[fa[u]], sz[u]);
53
       return rt;
54
   }
55
56
  // sons子树方向节点个数,val根到该节点异或和,gc边后继方向的节点个数
  int sons[N], gc[M];
59 ll val[N];
```

```
60 11 ans = 0;
    int n;
61
62
    const int MOD = int(1e9 + 7);
63
64
    11 nums[N];
65
    int cnt[N];
66
67
    void go(int s, int rt) {
68
         fa[s] = rt;
69
70
         val[s] = 0;
71
         int 1, r;
         // 不计算s
72
         q[l = r = 0] = s;
73
         int m = 0;
74
         while (l \ll r) {
75
             int u = q[l++];
76
77
             nums[m++] = val[u];
             for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
78
79
                 int v = gh.eg[i].to;
                 if (v == fa[u] || vis[v]) continue;
80
                 fa[v] = u;
81
82
                 q[++r] = v;
83
                 val[v] = val[u] \wedge gh.eg[i].w;
84
                 // 这个点方向后面有多少点
                 sons[v] = gc[i];
85
             }
86
87
         sort(nums, nums + m);
88
         m = unique(nums, nums + m) - nums;
89
90
         mst(cnt, 0, m);
91
         // 遍历分支
         for (int j = gh.head[s]; \sim j; j = gh.eg[j].nxt) {
92
             // 分支的根
93
             int du = gh.eg[j].to;
94
             if (vis[du]) continue;
95
96
             a[l = r = 1] = du;
97
             while (l <= r) {</pre>
                 int u = q[l++];
98
                 int k = lower_bound(nums, nums + m, val[u]) - nums;
99
                 (ans += 111 * sons[u] * cnt[k] % MOD) %= MOD;
100
                 if (val[u] == 0) {
101
                      (ans += 111 * sons[u] * (n - gc[j]) % MOD) %= MOD;
102
103
                 for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
104
105
                      int v = gh.eg[i].to;
106
                      if (v == fa[u] || vis[v]) continue;
107
                      q[++r] = v;
                 }
108
109
             }
110
             // 增加这个方向的值
111
             while (--1) {
                 int u = q[1];
112
                 int k = lower_bound(nums, nums + m, val[u]) - nums;
113
114
                 (cnt[k] += sons[u]) \% = MOD;
115
             }
116
         }
117
    }
118
```

```
119 void work(int u) {
        // 换根
120
        u = froot(u);
121
122
        vis[u] = true;
123
        go(u, 0);
        for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
124
125
            int v = gh.eg[i].to;
            if (vis[v]) continue;
126
            work(v);
127
        }
128
129
    }
130
131
   // 预处理边后继节点个数
    int pdfs(int u, int f) {
132
        int fg_id = -1;
133
        int s = 1;
134
        for (int i = gh.head[u]; \sim i; i = gh.eg[i].nxt) {
135
136
            int v = gh.eg[i].to;
            if (v == f) { // 记录父边ID
137
                fg_id = i;
138
                continue;
139
140
            int c = pdfs(v, u);
141
142
            gc[i] = c;
143
            S += C;
        }
144
        .
// 存在父边
145
        if (\sim fg_id) gc[fg_id] = n - s;
146
147
        return s;
    }
148
149
    void solve() {
150
        while (cin >> n) {
151
152
            gh.init(n);
            for (int i = 2; i \le n; i++) {
153
                int u, v;
154
155
                11 w;
                u = i;
156
                cin >> v >> w;
157
                gh.addEdge(u, v, w);
158
                gh.addEdge(v, u, w);
159
160
            mst(vis, false, n + 1);
161
162
            pdfs(1, 0);
            ans = 0;
163
164
            work(1);
165
            cout << ans << endl;</pre>
166
        }
   }
167
        DsuOnTree1
 1 /*
 2 DSU On Tree 树上启发式合并
 3 进行轻重链剖分,每次先递归轻链,然后消除轻儿子影响,然后计算重儿子,并且保留重儿子结果
   然后以常数级别复杂度将重儿子状态修改成当前节点状态!!! (<-----重点)
 4
    然后递归计算当前子树结果
    复杂度0(nlogn)*0(计算/更新)
```

```
7
   [CF741D] 一棵有根树, 边上有字母a~v, 求每个子树中最长的边, 满足这个边上的所有字母重排后可以构成回文。
8
   很容易想到点分治,但由于是有根树,所以点分治是做不了的。
9
   先对没给点记录xn[x]表示x到根的路径的各个字母出现奇偶性,然后b[S]记录满足xn[x]=S的所有点x的最大深度。
   类似点分治,一棵子树一棵子树地处理,每次先对子树内每个点查找之前子树中是否有能与它拼成回文串的点与它的最
       大深度,再更新b[]。
12
   注意查找和更新不能同时进行,要整棵子树都查询完毕之后再进行更新操作。注意查找与更新当前子树的根。
13
   我试图去写了一个一次进出完整查询和更新的代码,逻辑是先查询后进人最后更新,但没有合适的写法,难度太大
14
15
  const int N = 500005;
16
17
  const int M = N \ll 1;
   struct Edge {
18
       int to, next;
19
20
21
   struct Gragh {
22
23
       int head[N];
       Edge eg[M];
24
25
       int tot;
26
27
       void init(int n) {
          tot = 0;
28
29
          memset(head, -1, sizeof(int) * ++n);
30
       }
31
32
       inline void addEdge(int u, int v) {
          eg[tot] = \{v, head[u]\};
33
          head[u] = tot++;
34
35
36
   } gh;
37
   int tsz[N], son[N], cson;
38
   int dep[N] = \{0\};
39
40
   void pdfs(int u, int f) {
41
42
       dep[u] = dep[f] + 1;
43
       tsz[u] = 1;
       son[u] = -1;
44
       for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
45
          int v = gh.eg[i].to;
46
          if (v == f) continue;
47
          pdfs(v, u);
48
49
          tsz[u] += tsz[v];
          if (son[u] == -1 \mid | tsz[v] > tsz[son[u]]) son[u] = v;
50
       }
51
52
   }
53
54 bool ans[N];
   char s[N];
56
   int p[N];
57
58
   void calc(int u, int f, int d) {
       p[dep[u]] ^= (1 << (s[u] - 'a'));
59
       for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
60
61
           int v = gh.eg[i].to;
           if (v == f || v == cson) continue;
62
63
          calc(v, u, d);
       }
64
```

```
65 }
66
    void dfs(int u, int f, int d) {
67
         // 遍历轻儿子, 结果删除
68
         for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
69
             int v = gh.eg[i].to;
70
             if (v == f || v == son[u]) continue;
71
72
             dfs(v, u, 0);
73
         if (~son[u]) { // 计算重儿子, 结果保留
74
75
             dfs(son[u], u, 1);
             cson = son[u];
76
         }
77
         calc(u, f, 1);
78
         cson = -1;
79
         for (int i = qry.head[u]; \sim i; i = qry.eg[i].next) {
80
             int h = qry.eg[i].to;
81
82
             ans[i] = __builtin_popcount(p[h]) <= 1;</pre>
83
         if (d == 0) { // 删除当前子树结果
84
             calc(u, f, -1);
85
86
         }
87 }
88
89
    void solve() {
90
         int n, m;
         while (cin >> n >> m) {
91
             gh.init(n);
92
             qry.init(n);
93
             for (int i = 2; i <= n; i++) {
94
95
                 int u;
96
                 cin >> u;
                 gh.addEdge(u, i);
97
             }
98
             s[0] = ' ';
99
             cin >> (s + 1);
100
101
             rep(i, 0, m) {
102
                 int v, h;
                 cin >> v >> h;
103
                 qry.addEdge(v, h);
104
105
             }
             pdfs(1, 0);
106
             dfs(1, 0, 0);
rep(i, 0, m) {
107
108
109
                 cout << (ans[i] ? "Yes" : "No") << endl;</pre>
110
             }
111
         }
112 }
113
114 const int N = 500005;
115 const int M = N << 1;
116 struct Edge {
117
         int to, next;
118
    };
119
    struct Gragh {
120
121
         int head[N];
122
         Edge eg[M];
         int tot;
123
```

```
124
         void init(int n) {
125
             tot = 0;
126
             memset(head, -1, sizeof(int) * ++n);
127
128
129
         inline void addEdge(int u, int v) {
130
             eg[tot] = \{v, head[u]\};
131
             head[u] = tot++;
132
133
         }
134
    } gh, qry;
135
int tsz[N], son[N], dep[N] = \{0\}, fa[N], cson;
    int xn[N];
137
    int a[N];
138
139
    void pdfs(int u, int f) {
140
141
         fa[u] = f;
         dep[u] = dep[f] + 1;
142
143
         tsz[u] = 1;
         son[u] = -1;
144
145
         xn[u] = xn[f] \wedge a[u];
146
147
         for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
148
             int v = gh.eg[i].to;
             if (v == f) continue;
149
             pdfs(v, u);
150
             tsz[u] += tsz[v];
151
             if (son[u] == -1 \mid | tsz[v] > tsz[son[u]]) son[u] = v;
152
         }
153
    }
154
155
156
    int b[1 << 22];
157
    int rst, ans[N], mx, root;
158
   int num[N], tot = 0;
159
160
161
    void check(int u) {
         int bs = xn[u];
162
163
         if (b[bs]) {
             rst = max(rst, dep[u] + b[bs] - dep[root] * 2);
164
165
         for (int i = 0; i < 22; i++) {
166
167
             int bt = bs ^(1 \ll i);
168
             if (b[bt]) {
169
                 rst = max(rst, dep[u] + b[bt] - dep[root] * 2);
170
             }
171
172
         bs ^= xn[root];
173
         if (bs == (bs \& (-bs))) {
174
             rst = max(rst, dep[u] - dep[root]);
175
         }
176
    }
177
    void calc(int u) {
178
179
         check(u);
180
         for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
181
             int v = gh.eg[i].to;
             if (v == fa[u]) continue;
182
```

```
183
             calc(v);
         }
184
    }
185
186
    void add(int u) {
187
         if (b[xn[u]]) {
188
189
             b[xn[u]] = max(b[xn[u]], dep[u]);
         } else {
190
             b[xn[u]] = dep[u];
191
             num[tot++] = xn[u];
192
193
         }
         if (u == cson) return;
194
         for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
195
             int v = gh.eg[i].to;
196
             if (v == fa[u]) continue;
197
             add(v);
198
         }
199
200
    }
201
    void go(int u) {
202
203
         rst = 0;
204
         if (~son[u]) {
205
             rst = 1;
206
             root = son[u];
207
             check(son[u]);
             add(son[u]);
208
         }
209
         root = u;
210
         check(u);
211
         for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
212
213
             int v = gh.eg[i].to;
214
             if (v == fa[u] || v == cson) continue;
             calc(v);
215
216
             add(v);
217
         }
218
    }
219
220
    void del() {
         for (int i = 0; i < tot; b[num[i++]] = 0);
221
222
         tot = 0;
223
    }
224
225
    void dfs(int u, int d) {
226
         // 遍历轻儿子, 结果删除
227
         for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
             int v = gh.eg[i].to;
228
229
             if (v == fa[u] || v == son[u]) continue;
             dfs(v, 0);
230
231
         if (~son[u]) { // 计算重儿子, 结果保留
232
233
             dfs(son[u], 1);
234
             cson = son[u];
235
236
         go(u); // 计算当前
237
         cson = -1;
         ans[u] = rst;
238
         if (d == 0) {
239
             del();
240
         }
241
```

```
242 }
243
    void dfs2(int u) {
244
        for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
245
246
             int v = gh.eg[i].to;
247
            dfs2(v);
            ans[u] = max(ans[u], ans[v]);
248
        }
249
    }
250
251
252
    void input(int n) {
253
        a[1] = 0;
        for (int i = 2; i \le n; i++) {
254
255
            int f;
256
            char c;
            cin >> f >> c;
257
            gh.addEdge(f, i);
258
            a[i] = 1 << (c - 'a');
259
        }
260
261 }
262
263
    void solve() {
264
        int n;
265
        while (cin >> n) {
            gh.init(n);
266
            input(n);
267
            memset(b, 0, sizeof(b));
268
            pdfs(1, 0);
269
            dfs(1, 0);
270
            dfs2(1);
271
            for (int i = 1; i <= n; i++) {
272
                printf("%d%c", ans[i], " \n"[i == n]);
273
274
            }
275
        }
276 }
    6.6 DsuOnTree2
 1
 2
    [CF570D]一棵树,询问某棵子树指定深度(对于树根的深度)的点能否构成回文。
 3
    用二进制记录每个深度的每个字母出现次数的奇偶性即可。
 4
 5
    const int N = 500005;
 7
    const int M = N \ll 1;
    struct Edge {
 8
 9
        int to, next;
10
   };
11
    struct Gragh {
12
        int head[N];
13
        Edge eg[M];
14
15
        int tot;
16
        void init(int n) {
17
18
            tot = 0;
            memset(head, -1, sizeof(int) * ++n);
19
20
        }
```

```
21
22
        inline void addEdge(int u, int v) {
            eg[tot] = \{v, head[u]\};
23
24
            head[u] = tot++;
25
26
   } gh, qry;
27
  int tsz[N], son[N], cson;
28
   int dep[N] = \{0\};
29
30
31
   void pdfs(int u, int f) {
32
        dep[u] = dep[f] + 1;
        tsz[u] = 1;
33
        son[u] = -1;
34
        for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
35
            int v = gh.eg[i].to;
36
            if (v == f) continue;
37
            pdfs(v, u);
38
            tsz[u] += tsz[v];
39
            if (son[u] == -1 \mid | tsz[v] > tsz[son[u]]) son[u] = v;
40
        }
41
   }
42
43
44 bool ans[N];
45
   char s[N];
   int p[N];
46
47
   void calc(int u, int f, int d) {
   p[dep[u]] ^= (1 << (s[u] - 'a'));</pre>
48
49
        for (int i = gh.head[u]; ~i; i = gh.eg[i].next) {
50
51
            int v = gh.eg[i].to;
52
            if (v == f || v == cson) continue;
53
            calc(v, u, d);
        }
54
   }
55
56
57
   void dfs(int u, int f, int d) {
58
        // 遍历轻儿子, 结果删除
        for (int i = gh.head[u]; \sim i; i = gh.eg[i].next) {
59
60
            int v = gh.eg[i].to;
            if (v == f || v == son[u]) continue;
61
            dfs(v, u, 0);
62
63
        if (~son[u]) { // 计算重儿子, 结果保留
64
            dfs(son[u], u, 1);
65
66
            cson = son[u];
67
        calc(u, f, 1);
68
        cson = -1;
69
70
        for (int i = qry.head[u]; ~i; i = qry.eg[i].next) {
71
            int h = qry.eg[i].to;
72
            ans[i] = __builtin_popcount(p[h]) <= 1;</pre>
73
        if (d == 0) { // 删除当前子树结果
74
            calc(u, f, -1);
75
76
        }
77
   }
78
79 void solve() {
```

```
80
         int n, m;
         while (cin >> n >> m) {
81
              gh.init(n);
82
83
              qry.init(n);
              for (int i = 2; i <= n; i++) {
84
85
                   int u;
86
                   cin >> u;
87
                   gh.addEdge(u, i);
88
              }
              s[0] = ' ';
89
              cin >> (s + 1);
90
              rep(i, 0, m) {
91
92
                   int v, h;
                   cin >> v >> h;
93
                   qry.addEdge(v, h);
94
              }
95
              pdfs(1, 0);
dfs(1, 0, 0);
rep(i, 0, m) {
96
97
98
99
                  cout << (ans[i] ? "Yes" : "No") << endl;</pre>
100
              }
101
         }
102 }
```

7 数论

7.1 线性基

```
#define rep(i, l, r) for(int i=(l);i<(r);++i)
#define per(i, l, r) for(int i=(r)-1;i>=(l);--i)
   #define pw(x) (111 << (x))
4 #define bt(x, i) ((x >> i) & 1)
5 const int LN = 61;
    struct LB {
6
7
        ll d[LN] = \{0\}, p[LN] = \{0\};
8
        int g[LN] = \{0\};
9
        int cnt = 0;
10
11
        LB() = default;
12
        11 & operator[](int pos) {
13
             return d[pos];
14
        }
15
16
17
        const ll &operator[](int pos) const {
             return d[pos];
18
        }
19
20
        void insert(ll val, int pos) {
21
22
             per(i, 0, LN) {
23
                  if (val & pw(i)) {
                      if (!d[i]) {
24
                           d[i] = val;
25
                           g[i] = pos;
26
                           return;
27
                      }
28
                      // 贪心保留最右
29
30
                      if (pos > g[i]) {
31
                           swap(pos, g[i]);
32
                           swap(val, d[i]);
33
                      val ^= d[i];
34
35
                  }
36
             }
37
        }
38
        11 query_max(int 1) {
39
40
             11 \text{ ret} = 0;
             per(i, 0, LN) {
41
                  if (g[i] >= l)
42
                      ret = max(ret, ret ^ d[i]);
43
             }
44
45
             return ret;
        }
46
47
        11 query_max() {
48
             ll ret = 0;
49
50
             per(i, 0, LN) {
                  ret = max(ret, ret ^ d[i]);
51
52
53
             return ret;
        }
54
55
```

```
11 query_min() {
56
              rep(i, 0, LN) {
    if (d[i]) return d[i];
57
58
59
60
              return 0;
         }
61
62
         bool test(ll x) const {
63
              per(i, 0, LN) {
64
                  if (bt(x, i)) {
65
66
                       if (!d[i]) return false;
67
                       x \wedge = d[i];
                  }
68
              }
69
70
              return true;
         }
71
72
         void rebuild() {
73
              per(i, 0, LN) {
74
75
                  per(j, 0, i) {
                       if (d[i] & (1LL << j)) d[i] ^= d[j];</pre>
76
77
              }
78
79
              rep(i, 0, LN) {
80
                  if (d[i]) p[cnt++] = d[i];
81
         }
82
83
         ll kth_query(ll k) {
84
              int ret = 0;
85
              if (k >= pw(cnt)) return -1;
86
87
              per(i, 0, LN) {
                  if (bt(k, 1)) ret ^= p[i];
88
89
90
              return ret;
         }
91
92
    };
93
    // 求并集
    LB operator+(const LB &n1, const LB &n2) {
94
         LB ret = n1;
95
         per(i, 0, LN)
96
              if (n2.d[i])
97
                  ret.insert(n1.d[i], n1.g[i]);
98
99
         return ret;
100 }
101 // 求交集
102 LB operator^(const LB &n1, const LB &n2) {
103
         LB ans = \{\}, c = n2, d = n2;
104
         rep(i, 0, LN) {
105
              ll x = n1[i];
106
              if (!x) continue;
107
              int p = i;
              11 T = 0;
108
              per(j, 0, p + 1) {
109
                  if (bt(x, j)) {
    if (c[j]) {
110
111
                           x ^= c[j];
112
                           T \stackrel{\wedge}{=} d[j];
113
114
                       } else {
```

```
115
                         p = j;
                         break;
116
                     }
117
                 }
118
119
            if (!x) {
120
                 ans[i] = T;
121
122
            } else {
                 c[p] = x;
123
124
                 d[p] = T;
125
            }
126
        }
127
        return ans;
128 }
    7.2 FWT
    void FWT(int p[], int n) { // 如果要取模运算记得-x+MOD,*inv
 2
        for (int i = 1; i < n; i <<= 1) {
            for (int j = 0; j < n; j += (i << 1)) {
 3
                 for (int k = 0; k < i; k++) {
 4
                     // or{
 5
                     p[i + j + k] += +p[j + k];
 6
 7
                     // and
                     p[j + k] += p[i + j + k];
 8
 9
                     // xor
                     tie(p[j + k], p[i + j + k]) = make_pair(p[j + k] + p[i + j + k], p[j + k]
10
        k] - p[i + j + k]);
11
12
            }
        }
13
    }
14
15
    void IFWT(int p[], int n) { // 如果要取模运算记得-x+MOD,*inv
16
        for (int i = 1; i < n; i <<= 1) {
17
            for (int j = 0; j < n; j + = (i << 1)) {
18
                 for (int k = 0; k < i; k++) {
19
20
                     p[i + j + k] -= p[j + k];
21
                     // and
22
                     p[j + k] -= p[i + j + k];
23
24
                     // xor
                     tie(p[j + k], p[i + j + k]) = make_pair(p[j + k] + p[i + j + k], p[j + k])
25
        k] - p[i + j + k]);
                     p[j + k] = p[j + k] / 2, p[i + j + k] = p[i + j + k] / 2;
26
27
28
            }
        }
29
30 }
    7.3 FFT-NTT
 1 const int N = int(1e6 + 48);
 2 const int MOD = 998244353;
 3 const int pr = 3;
 4 const int phi = MOD - 1;
 5 const int sq2 = 116195171;
```

```
int fpow(ll a, ll k) {
7
       11 c = 1;
8
       for (; k; k >>= 1) {
9
           if (k & 1) (c *= a) %= MOD;
10
           (a *= a) \%= MOD;
11
12
13
       return c;
   }
14
15
   int calc(int n) {
16
17
       int k = 0;
       while ((1 << k) < n) k++;
18
       return k;
19
   }
20
21
  int rk[N], rn = 0;
22
23
24 template<class T>
25
   void change(T p[], int n) {
       int k = calc(n);
26
27
       if (rn != n) {
28
           rn = n;
29
           for (int i = 0; i < n; i++) rk[i] = (rk[i >> 1] >> 1) | ((i & 1) << (k - 1));
30
       for (int i = 0; i < n; i++) if (i < rk[i]) swap(p[i], p[rk[i]]);
31
32 }
33
34 // FFT
   const double pi = acos(-1.0);
   using Complex = complex<double>;
37
   void FFT(Complex p[], int n, int opt) {
38
       change(p, n);
39
       for (int i = 1; i < n; i <<= 1) { //待合并区间的长度的一半
40
           Complex wn(cos(pi / i), opt * sin(pi / i)); //单位根
41
42
            for (int R = i << 1, j = 0; j < n; j += R) { //R是区间的长度, j表示前已经到哪个位置了
43
                Complex w(1, 0); //幂
                for (int k = 0; k < i; ++k, w = w * wn) { //枚举左半部分
44
                    Complex x = p[j + k], y = w * p[j + i + k]; //蝴蝶效应
45
                    p[j + k] = x + y;
46
                    p[j+i+k] = x - y;
47
                }
48
49
           }
       }
50
51
   }
52
   void NTT(ll p[], int n, int opt) {
53
       change(p, n);
54
55
       for (int i = 1; i < n; i <<= 1) {
56
           ll wn = fpow(pr, phi / (i << 1));
57
            for (int R = i \ll 1, j = 0; j \ll n; j += R) {
58
                11 w = 1;
                for (int k = 0; k < i; ++k, w = w * wn % MOD) {
59
                    ll x = p[j + k], y = w * p[i + j + k] % MOD;
60
                    p[j + k] = (x + y) \% MOD;
61
62
                    p[i + j + k] = (x - y + MOD) \% MOD;
63
                }
64
           }
```

```
65
       if (opt == -1)reverse(p + 1, p + n);
66
  }
67
   7.4 圆卜整点
1 // x^2+y^2=r^2 --> y^2=(r-x)*(r+x)
2 // d = gcd(r-x,r+x), n=(r-x)/d, m=(r+x)/d, y^2=d^2*m*n
3 // 因为y^2和d^2为完全平方数, 所以 n*m 为完全平方数
4 // 又因为n,m互质, 所以 n=u^2, m =v^2
5 // 所以 r-y=d*u^2, r+y=d*v^2
6 // 可得 x=d*(v^2-u^2), y=d*u*v
7
  ll gcd(ll a, ll b) { return !b ? a : gcd(b, a % b); }
8
9
10 inline int work(vector<pll> &p, ll d, ll t) \{ // d * t == 2 * r \}
       int sum = 0;
11
       for (ll u = 1; 2 * u * u < t; u++) {
12
           ll v = ll(sqrt(t - u * u));
13
           if (v == u \mid | v * v + u * u != t \mid | gcd(u, v) != 1) continue;
14
           11 x = d * (v * v - u * u) / 2;
15
           11 y = d * u * v;
16
           p.push_back({x, y});
17
           sum++;
18
19
20
       return sum;
   }
21
22
   int calc(vector<pll> &p, ll r) {
23
24
       int sum = 0;
25
       r <<= 1;
26
       for (ll i = 1; i * i <= r; i++) { // sqrt 枚举因子
27
           if (r % i) continue;
           sum += work(p, i, r / i);
28
           if (i * i < r) sum += work(p, r / i, i);
29
30
       for (int i = 0, cnt = p.size(); i < cnt; i++) {</pre>
31
           int x = p[i].fi, y = p[i].se;
32
33
           p.push_back(\{x, -y\}), p.push_back(\{-x, y\}), p.push_back(\{-x, -y\});
       }
34
       sum <<= 2;
35
36
       r >>= 1;
37
       sum += 4;
       p.push_back(\{r, 0\}), p.push_back(\{-r, 0\}), p.push_back(\{0, r\}), p.push_back(\{0, -r\})
38
       });
39
       sort(p.begin(), p.end());
       return sum;
40
41 }
   7.5 线性逆元
1 // 逆元, 阶乘, 阶乘逆元
2 ll inv[N] = \{1, 1\}, fac[N] = \{1, 1\}, ifac[N] = \{1, 1\};
  // 线性求逆元
3
4
   void init() {
       // 如果用的Int记得下面inv[i] = 1ll*...
5
       for (int i = 2; i < N; i++) {
```

8 其他

8.1 BigInteger

```
// base and base_digits must be consistent
   constexpr int base = 1000000000;
   constexpr int base_digits = 9;
3
4
   struct bigint {
5
        // value == 0 is represented by empty z
6
7
        vector<int> z; // digits
8
        // sign == 1 <==> value >= 0
9
10
        // sign == -1 <==> value < 0
        int sign;
11
12
        bigint() : sign(1) {}
13
14
        bigint(ll v) { *this = v; }
15
16
        bigint &operator=(ll v) {
17
            sign = v < 0 ? -1 : 1;
18
            v *= sign;
19
            z.clear();
20
            for (; v > 0; v = v / base) z.push_back((int) (v \% base));
21
            return *this;
22
23
        }
24
25
        bigint(const string &s) { read(s); }
26
        bigint &operator+=(const bigint &other) {
27
            if (sign == other.sign) {
28
29
                for (int i = 0, carry = 0; i < other.z.size() || carry; ++i) {
30
                    if (i == z.size())
31
                        z.push_back(0);
                    z[i] += carry + (i < other.z.size() ? other.z[i] : 0);
32
                    carry = z[i] >= base;
33
34
                    if (carry)
35
                        z[i] -= base;
36
            } else if (other != 0 /* prevent infinite loop */) {
37
38
                *this -= -other;
39
40
            return *this;
41
42
        friend bigint operator+(bigint a, const bigint &b) { return a += b; }
43
44
        bigint &operator-=(const bigint &other) {
45
            if (sign == other.sign) {
46
                if (sign == 1 && *this >= other || sign == -1 && *this <= other) {
47
                    for (int i = 0, carry = 0; i < other.z.size() | | carry; ++i) {
48
                        z[i] = carry + (i < other.z.size() ? other.z[i] : 0);
49
                        carry = z[i] < 0;
50
51
                        if (carry)
52
                             z[i] += base;
53
54
                    trim();
                } else {
55
```

```
*this = other - *this;
56
                      this->sign = -this->sign;
57
58
59
             } else {
                 *this += -other;
60
61
             return *this;
62
         }
63
64
         friend bigint operator-(bigint a, const bigint &b) {
65
             return a -= b;
66
67
         }
68
         bigint &operator*=(int v) {
69
             if (v < 0) sign = -sign, v = -v;
70
             for (int i = 0, carry = 0; i < z.size() || carry; ++i) {
71
                 if (i == z.size()) z.push_back(0);
 72
                 ll cur = (ll) z[i] * v + carry;
73
                 carry = (int) (cur / base);
74
                 z[i] = (int) (cur % base);
75
             }
76
             trim();
77
             return *this;
78
79
         }
80
         bigint operator*(int v) const { return bigint(*this) *= v; }
81
82
         friend pair<bigint, bigint> divmod(const bigint &a1, const bigint &b1) {
83
             int norm = base / (b1.z.back() + 1);
84
             bigint a = a1.abs() * norm;
85
             bigint b = b1.abs() * norm;
86
87
             bigint q, r;
             q.z.resize(a.z.size());
88
89
             for (int i = (int) \ a.z.size() - 1; \ i >= 0; \ i--) {
90
                 r *= base;
91
                 r += a.z[i];
92
93
                 int s1 = b.z.size() < r.z.size() ? r.z[b.z.size()] : 0;</pre>
                 int s2 = b.z.size() - 1 < r.z.size() ? r.z[b.z.size() - 1] : 0;</pre>
94
                 int d = (int) (((ll) s1 * base + s2) / b.z.back());
95
                 r -= b * d;
96
                 while (r < 0) r += b, --d;
97
98
                 q.z[i] = d;
             }
99
100
             q.sign = a1.sign * b1.sign;
101
             r.sign = a1.sign;
102
             q.trim();
103
104
             r.trim();
             return {q, r / norm};
105
106
107
108
         friend bigint sqrt(const bigint &a1) {
             bigint a = a1;
109
             while (a.z.empty() \parallel a.z.size() \% 2 == 1) a.z.push_back(0);
110
111
112
             int n = a.z.size();
113
             int firstDigit = (int) ::sqrt((double) a.z[n - 1] * base + a.z[n - 2]);
114
```

```
int norm = base / (firstDigit + 1);
115
116
             a *= norm;
             a *= norm;
117
             while (a.z.empty() \mid | a.z.size() \% 2 == 1) a.z.push_back(0);
118
119
             bigint r = (ll) a.z[n - 1] * base + a.z[n - 2];
120
             firstDigit = (int) ::sqrt((double) a.z[n - 1] * base + a.z[n - 2]);
121
             int q = firstDigit;
122
             bigint res;
123
124
125
             for (int j = n / 2 - 1; j >= 0; j--) {
126
                 for (;; --q) {
                      bigint r1 = (r - (res * 2 * base + q) * q) * base * base +
127
                                  (j > 0? (ll) a.z[2 * j - 1] * base + a.z[2 * j - 2] : 0);
128
                      if (r1 >= 0) {
129
                          r = r1;
130
131
                          break;
                     }
132
133
                 (res *= base) += q;
134
135
                 if (j > 0) {
136
                      int d1 = res.z.size() + 2 < r.z.size() ? r.z[res.z.size() + 2] : 0;</pre>
137
                      int d2 = res.z.size() + 1 < r.z.size() ? r.z[res.z.size() + 1] : 0;
138
139
                      int d3 = res.z.size() < r.z.size() ? r.z[res.z.size()] : 0;</pre>
                      q = (int) (((ll) d1 * base * base + (ll) d2 * base + d3) / (firstDigit
140
        * 2));
141
                 }
             }
142
143
144
             res.trim();
             return res / norm;
145
146
        }
147
        bigint operator/(const bigint &v) const {
148
             return divmod(*this, v).first;
149
150
        }
151
        bigint operator%(const bigint &v) const {
152
             return divmod(*this, v).second;
153
154
155
        bigint &operator/=(int v) {
156
             if (v < 0) sign = -sign, v = -v;
157
             for (int i = (int) z.size() - 1, rem = 0; i >= 0; --i) {
158
                 ll cur = z[i] + rem * (ll) base;
159
                 z[i] = (int) (cur / v);
160
                 rem = (int) (cur % v);
161
162
             trim();
163
164
             return *this;
165
        }
166
        bigint operator/(int v) const {
167
             return bigint(*this) /= v;
168
169
        }
170
171
        int operator%(int v) const {
172
             if (v < 0) v = -v;
```

```
int m = 0;
173
             for (int i = (int) z.size() - 1; i >= 0; --i)
174
                 m = (int) ((z[i] + m * (ll) base) % v);
175
176
             return m * sign;
         }
177
178
179
         bigint &operator*=(const bigint &v) {
             return *this = *this * v;;
180
181
182
183
         bigint &operator/=(const bigint &v) {
184
             return *this = *this / v;
185
186
         bool operator<(const bigint &v) const {</pre>
187
             if (sign != v.sign)
188
                  return sign < v.sign;</pre>
189
190
             if (z.size() != v.z.size())
                 return z.size() * sign < v.z.size() * v.sign;</pre>
191
             for (int i = (int) z.size() - 1; i >= 0; i--)
192
                  if (z[i] != v.z[i])
193
                      return z[i] * sign < v.z[i] * sign;</pre>
194
             return false;
195
196
         }
197
         bool operator>(const bigint &v) const { return v < *this; }</pre>
198
199
         bool operator<=(const bigint &v) const { return !(v < *this); }</pre>
200
201
         bool operator>=(const bigint &v) const { return !(*this < v); }</pre>
202
203
204
         bool operator==(const bigint &v) const { return !(*this < v) && !(v < *this); }</pre>
205
206
         bool operator!=(const bigint &v) const { return *this < v | | v < *this; }</pre>
207
208
         void trim() {
209
             while (!z.empty() && z.back() == 0) z.pop_back();
210
             if (z.empty()) sign = 1;
         }
211
212
213
         bool isZero() const {
             return z.empty();
214
         }
215
216
217
         friend bigint operator-(bigint v) {
218
             if (!v.z.empty()) v.sign = -v.sign;
219
             return v;
220
         }
221
222
         bigint abs() const {
223
             return sign == 1 ? *this : -*this;
224
         }
225
         ll longValue() const {
226
227
             11 \text{ res} = 0;
             for (int i = (int) z.size() - 1; i >= 0; i--)
228
229
                  res = res * base + z[i];
             return res * sign;
230
         }
231
```

```
232
233
         friend bigint gcd(const bigint &a, const bigint &b) {
234
             return b.isZero() ? a : gcd(b, a % b);
235
236
         friend bigint lcm(const bigint &a, const bigint &b) {
237
238
             return a / gcd(a, b) * b;
         }
239
240
         void read(const string &s) {
241
242
             sign = 1;
             z.clear();
243
244
             int pos = 0;
             while (pos < s.size() && (s[pos] == '-' || s[pos] == '+')) {</pre>
245
                 if (s[pos] == '-') sign = -sign;
246
247
                 ++pos;
248
             for (int i = (int) s.size() - 1; i \ge pos; i = base\_digits) {
249
                 int x = 0;
250
                 for (int j = max(pos, i - base_digits + 1); j <= i; j++)
251
                      x = x * 10 + s[j] - '0';
252
                 z.push_back(x);
253
254
255
             trim();
256
         }
257
258
         friend istream &operator>>(istream &stream, bigint &v) {
259
             string s;
             stream >> s;
260
             v.read(s);
261
262
             return stream;
263
         }
264
265
         friend ostream &operator<<(ostream &stream, const bigint &v) {</pre>
266
             if (v.sign == -1)
                 stream << '-';
267
268
             stream << (v.z.empty() ? 0 : v.z.back());</pre>
269
             for (int i = (int) v.z.size() - 2; i >= 0; --i)
                 stream << setw(base_digits) << setfill('0') << v.z[i];</pre>
270
271
             return stream;
272
         }
273
         static vector<int> convert_base(const vector<int> &a, int old_digits, int
274
        new_digits) {
             vector<ll> p(max(old_digits, new_digits) + 1);
275
276
             p[0] = 1;
277
             for (int i = 1; i < p.size(); i++)</pre>
                 p[i] = p[i - 1] * 10;
278
             vector<int> res;
279
280
             11 \text{ cur} = 0;
281
             int cur_digits = 0;
282
             for (int v : a) {
                 cur += v * p[cur_digits];
283
                 cur_digits += old_digits;
284
                 while (cur_digits >= new_digits) {
285
                      res.push_back(int(cur % p[new_digits]));
286
287
                      cur /= p[new_digits];
288
                      cur_digits -= new_digits;
                 }
289
```

```
}
290
             res.push_back((int) cur);
291
             while (!res.empty() && res.back() == 0) res.pop_back();
292
293
             return res;
294
295
296
        typedef vector<ll> vll;
297
        static vll karatsubaMultiply(const vll &a, const vll &b) {
298
299
             int n = a.size();
300
             vll res(n + n);
301
             if (n <= 32) {
                 for (int i = 0; i < n; i++)
302
303
                     for (int j = 0; j < n; j++)
                          res[i + j] += a[i] * b[j];
304
305
                 return res;
             }
306
307
308
             int k = n \gg 1;
             vll a1(a.begin(), a.begin() + k);
309
             vll a2(a.begin() + k, a.end());
310
             vll b1(b.begin(), b.begin() + k);
311
             vll b2(b.begin() + k, b.end());
312
313
             vll a1b1 = karatsubaMultiply(a1, b1);
314
             vll a2b2 = karatsubaMultiply(a2, b2);
315
316
             for (int i = 0; i < k; i++) a2[i] += a1[i];
317
             for (int i = 0; i < k; i++) b2[i] += b1[i];
318
319
             vll r = karatsubaMultiply(a2, b2);
320
             for (int i = 0; i < a1b1.size(); i++) r[i] -= a1b1[i];
321
322
             for (int i = 0; i < a2b2.size(); i++) r[i] -= a2b2[i];
323
             for (int i = 0; i < r.size(); i++) res[i + k] += r[i];
324
             for (int i = 0; i < a1b1.size(); i++) res[i] += a1b1[i];</pre>
325
326
             for (int i = 0; i < a2b2.size(); i++) res[i + n] += a2b2[i];
327
             return res;
        }
328
329
330
        bigint operator*(const bigint &v) const {
             vector<int> a6 = convert_base(this->z, base_digits, 6);
331
             vector<int> b6 = convert_base(v.z, base_digits, 6);
332
333
             vll a(a6.begin(), a6.end());
             vll b(b6.begin(), b6.end());
334
335
             while (a.size() < b.size()) a.push_back(0);</pre>
             while (b.size() < a.size()) b.push_back(0);</pre>
336
             while (a.size() & (a.size() - 1)) a.push_back(0), b.push_back(0);
337
             vll c = karatsubaMultiply(a, b);
338
             bigint res;
339
340
             res.sign = sign * v.sign;
341
             for (int i = 0, carry = 0; i < c.size(); i++) {
342
                 ll cur = c[i] + carry;
                 res.z.push_back((int) (cur % 1000000));
343
                 carry = (int) (cur / 1000000);
344
345
346
             res.z = convert_base(res.z, 6, base_digits);
347
             res.trim();
             return res;
348
```

```
349
        }
350 };
    8.2 FastIO
 1
     * FastIO
 2
 3
     * 代码模板 !
 4
       如有雷同 !
     * 纯属巧合!
 5
     */
 6
    namespace FastI0 {
 7
    #define BUF_SIZE 10000000
    #define OUT_SIZE 10000000
    #define ll long long
10
11
         //fread->read
        bool IOerror = 0;
12
13
        inline char nc() {
14
             static char buf[BUF_SIZE], *p1 = buf + BUF_SIZE, *pend = buf + BUF_SIZE;
15
             if (p1 == pend) {
16
                 p1 = buf;
17
                 pend = buf + fread(buf, 1, BUF_SIZE, stdin);
18
19
                 if (pend == p1) {
                     I0error = 1;
20
21
                     return -1;
22
23
                 //{printf("IO error!\n");system("pause");for (;;);exit(0);}
24
25
             return *p1++;
26
        }
27
28
        inline bool blank(char ch) { return ch == ' ' | | ch == '\n' | | ch == '\r' | | ch ==
        '\t'; }
29
        inline void read(int &x) {
30
31
             bool sign = 0;
32
             char ch = nc();
33
             x = 0;
34
             for (; blank(ch); ch = nc());
35
             if (I0error)return;
             if (ch == '-')sign = 1, ch = nc();
36
             for (; ch >= '0' \&\& ch <= '9'; ch = nc())x = x * 10 + ch - '0';
37
38
             if (sign)x = -x;
39
        }
40
        inline void read(ll &x) {
41
42
             bool sign = 0;
             char ch = nc();
43
44
             x = 0;
             for (; blank(ch); ch = nc());
45
             if (I0error)return;
46
             if (ch == '-')sign = 1, ch = nc();
47
             for (; ch \ge '0' \&\& ch \le '9'; ch = nc())x = x * 10 + ch - '0';
48
49
             if (sign)x = -x;
        }
50
51
        inline void read(double &x) {
52
```

```
bool sign = 0;
53
             char ch = nc();
54
             x = 0;
55
             for (; blank(ch); ch = nc());
56
             if (I0error)return;
57
             if (ch == '-')sign = 1, ch = nc();
58
             for (; ch \ge 0' \& ch \le 9'; ch = nc()x = x * 10 + ch - 0';
59
             if (ch == '.') {
60
                 double tmp = 1;
61
62
                 ch = nc();
                 for (; ch >= '0' && ch <= '9'; ch = nc())tmp /= 10.0, x += tmp * (ch - '0')
63
64
             if (sign)x = -x;
65
66
67
        inline void read(char *s) {
68
             char ch = nc();
69
             for (; blank(ch); ch = nc());
70
             if (I0error)return;
71
             for (; !blank(ch) && !I0error; ch = nc()*s++ = ch;
72
             *s = 0;
73
        }
74
75
76
        inline void read(char &c) {
             for (c = nc(); blank(c); c = nc());
77
             if (I0error) {
78
                 c = -1;
79
                 return;
80
81
             }
        }
82
83
        //fwrite->write
84
        struct Ostream_fwrite {
85
             char *buf, *p1, *pend;
86
             Ostream_fwrite() {
87
88
                 buf = new char[OUT_SIZE];
89
                 p1 = buf;
                 pend = buf + OUT_SIZE;
90
91
             void out(char ch) {
92
                 if (p1 == pend) {
93
                     fwrite(buf, 1, OUT_SIZE, stdout);
94
95
                     p1 = buf;
96
                 *p1++ = ch;
97
98
             void print(int x) {
99
                 static char s[15], *s1;
100
101
                 s1 = s;
102
                 if (!x)*s1++ = '0';
                 if (x < 0)out('-'), x = -x;
103
                 while (x)*s1++ = x % 10 + '0', x /= 10;
104
                 while (s1-- != s)out(*s1);
105
106
107
             void println(int x) {
108
                 static char s[15], *s1;
                 s1 = s;
109
                 if (!x)*s1++ = '0';
110
```

```
if (x < 0)out('-'), x = -x;
111
                while (x)*s1++ = x \% 10 + '0', x /= 10;
112
                while (s1-- != s)out(*s1);
113
                out('\n');
114
115
            void print(ll x) {
116
                static char s[25], *s1;
117
                s1 = s;
118
                if (!x)*s1++ = '0';
119
                if (x < 0)out('-'), x = -x;
120
                while (x)*s1++ = x % 10 + '0', x /= 10;
121
122
                while (s1-- != s)out(*s1);
123
            }
            void println(ll x) {
124
125
                static char s[25], *s1;
                s1 = s;
126
                if (!x)*s1++ = '0';
127
                if (x < 0)out('-'), x = -x;
128
                while (x)*s1++ = x % 10 + '0', x /= 10;
129
                while (s1-- != s)out(*s1);
130
                out('\n');
131
132
            void print(double x, int y) {
133
                static ll mul[] = {1, 10, 100, 1000, 10000, 100000, 1000000, 10000000,
134
        100000000,
                                   135
       LL, 10000000000000LL,
                                   136
        LL, 1000000000000000000LL};
                if (x < -1e-12)out('-'), x = -x;
137
                x *= mul[y];
138
                ll x1 = (ll) floor(x);
139
                if (x - floor(x) >= 0.5) ++ x1;
140
                ll x2 = x1 / mul[y], x3 = x1 - x2 * mul[y];
141
                print(x2);
142
                if (y > 0) {
143
                    out('.');
144
145
                    for (size_t i = 1; i < y && x3 * mul[i] < mul[y]; out('0'), ++i);</pre>
146
                    print(x3);
                }
147
            }
148
            void println(double x, int y) {
149
                print(x, y);
150
                out('\n');
151
152
            void print(char *s) { while (*s)out(*s++); }
153
            void println(char *s) {
154
                while (*s)out(*s++);
155
                out('\n');
156
            }
157
158
            void flush() {
159
                if (p1 != buf) {
160
                    fwrite(buf, 1, p1 - buf, stdout);
                    p1 = buf;
161
162
163
164
            ~Ostream_fwrite() { flush(); }
165
        } Ostream;
        inline void print(int x) { Ostream.print(x); }
166
```

```
inline void println(int x) { Ostream.println(x); }
167
        inline void print(char x) { Ostream.out(x); }
168
        inline void println(char x) {
169
            Ostream.out(x);
170
171
            Ostream.out('\n');
172
        inline void print(ll x) { Ostream.print(x); }
173
        inline void println(ll x) { Ostream.println(x); }
174
        inline void print(double x, int y) { Ostream.print(x, y); }
175
        inline void println(double x, int y) { Ostream.println(x, y); }
176
        inline void print(char *s) { Ostream.print(s); }
177
        inline void println(char *s) { Ostream.println(s); }
178
        inline void println() { Ostream.out('\n'); }
179
        inline void flush() { Ostream.flush(); }
180
181
    };
182
    using namespace FastIO;
    8.3 InputOutputSpeedUp
 1
 2
    template <class T>
 3
    inline bool read(T &x) {
 4
 5
        x = 0;
        char c = getchar();
 6
        if(c == EOF) return false;
 7
        bool f = false;
 8
        for (; !isdigit(c); c = getchar()) f ^= (c == '-');
 9
        for (; isdigit(c); c = getchar()) x = x * 10 + (c - '0');
10
        x = f ? -x : x;
11
        return true;
12
    }
13
14
    template <class T>
15
    inline void write(T x) {
16
        if (x < 0) {
17
            putchar('-');
18
19
            X = -X;
20
        T y = 1;
21
        int len = 1;
22
23
        for (; y \le x / 10; y *= 10) ++len;
        for (; len; --len, x = y, y = 10) putchar(x / y + '0');
24
25 }
    8.4 gcd
    ll gcd(ll x, ll y) { // 循环版
 2
        ll t;
        while (y){
 3
            t = x \% y;
 4
            x = y;
 5
 6
            y = t;
 7
 8
        return x;
    }
 9
10
```

```
11 ll gcd(ll a, ll b) { // 递归版
       return b == 0? a : gcd(b, a % b);
12
13 }
14
   // 扩展欧几里得
15
   ll exgcd(ll a, ll b, ll &x, ll &y) {
16
       if (b == 0) {
17
            x = 1, y = 0;
18
19
            return a;
20
21
       ll q = exgcd(b, a \% b, y, x);
22
       y -= a / b * x;
23
       return q;
24 }
   8.5 myItoa
   char * myItoa(int value, char* result, int base = 10);
3
   char * myItoa(int value, char* result, int base) {
       // check that the base if valid
4
5
       if (base < 2 || base > 16) { *result = 0; return result; }
6
       char* out = result;
7
       int quotient = abs(value);
8
9
       do {
            const int tmp = quotient / base;
10
            *out = "0123456789abcdef"[quotient - (tmp*base)];
11
12
           ++out;
            quotient = tmp;
13
       } while (quotient);
14
15
       // Apply negative sign
       if (value < 0) *out++ = '-';</pre>
16
       std::reverse(result, out);
17
       *out = 0;
18
       return result;
19
20 }
   8.6 Permutation
1 // 错排问题
  // D(n) = n! [(-1)^2/2! + \cdots + (-1)^n(n-1)/(n-1)! + (-1)^n/n!].
3 \quad long \quad long \quad table[1000] = \{0, 0, 1\};
   void init() {
4
       for (int i = 3; i \le 20; i++) {
5
            table[i] = (i - 1) * (table[i - 1] + table[i - 2]);
6
7
8
   }
   8.7 prime
1 // 普通素数筛
   const int PMAX = 1000000;
2
3
4 int prime_count = 0;
5 bool prime_list[PMAX] = { false };//元素值为0代表是素数
```

```
int prime_table[PMAX] = { 0 };
7
   void initPrime() {
8
       for (int i = 2; i < PMAX; i++) {
9
            if (!prime_list[i])
10
                prime_table[prime_count++] = i;
11
12
            for (int j = 0, e = (PMAX - 1) / i;
                 j < prime_count && prime_table[j] <= e; j++) {</pre>
13
                prime_list[i * prime_table[j]] = true;
14
                if (i % prime_table[j] == 0) break;
15
16
            }
17
       }
   }
18
19
20
   // 可以得到其中一个质因子的素数筛
21
   const int PMAX = 1000005;
22
23
24 int prime_count = 0;
25 int prime_list[PMAX] = {0}; //元素值为0代表是素数
26 int prime_table[PMAX] = {0};
27
28
   void initPrime() {
29
       for (int i = 2; i < PMAX; i++) {
30
            if (!prime_list[i])
                prime_list[i] = prime_table[prime_count++] = i;
31
            for (int j = 0, e = (PMAX - 1) / i, now;
32
                 j < prime_count && prime_table[j] <= e; j++) {</pre>
33
                prime_list[i * now] = now;
34
                if (i % now == 0) break;
35
36
            }
37
       }
38 }
   8.8 Hash
1
   struct Hash {
2
       int num[N];
       int tot;
3
 4
       void init() { tot = 0; }
5
6
7
       void insert(int x) { num[tot++] = x; }
8
       void build() {
9
            sort(num, num + tot);
10
            tot = unique(num, num + tot) - num;
11
12
13
14
       inline int operator[](int x) { return lower_bound(num, num + tot, x) - num; }
15 } hs;
```