

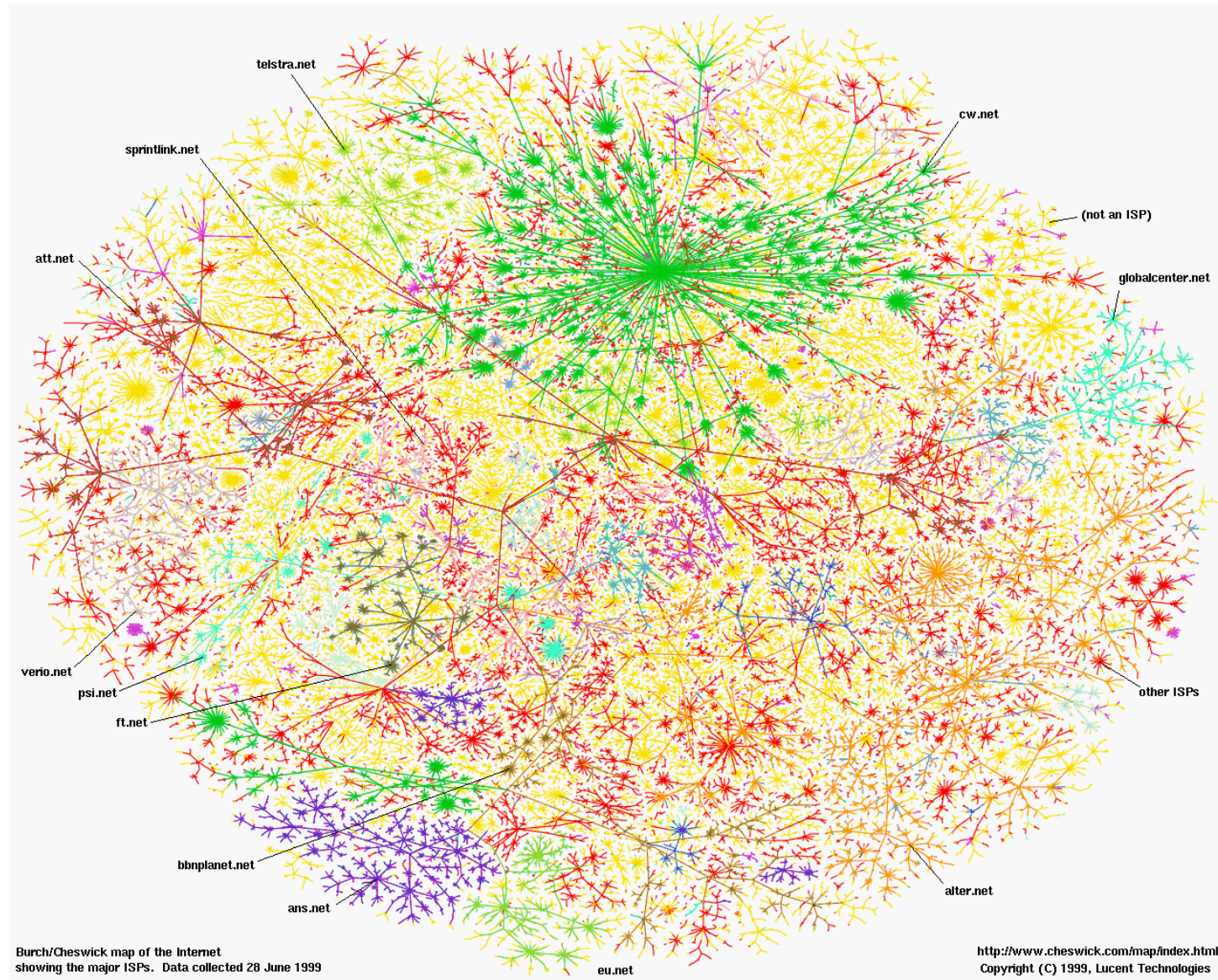
# Chapter 1: Introduction

## Our goal:

- ❑ get “feel” and terminology
- ❑ more depth, detail *later* in course
- ❑ approach:
  - use Internet as example

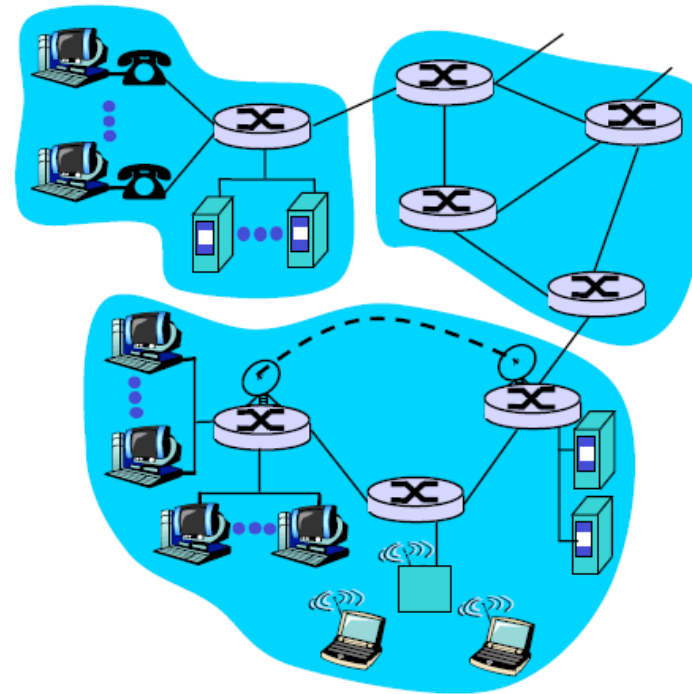
## Overview:

- ❑ what’s the Internet
- ❑ what’s a protocol?
- ❑ network edge
- ❑ access net, physical media
- ❑ network core
- ❑ Internet/ISP structure
- ❑ performance: loss, delay
- ❑ protocol layers, service models
- ❑ network modeling



# A closer look at network structure:

- ❑ network edge:  
applications and hosts
- ❑ network core:
  - routers
  - network of networks
- ❑ access networks,  
physical media:  
communication links



# The network edge:

## □ end systems (hosts):

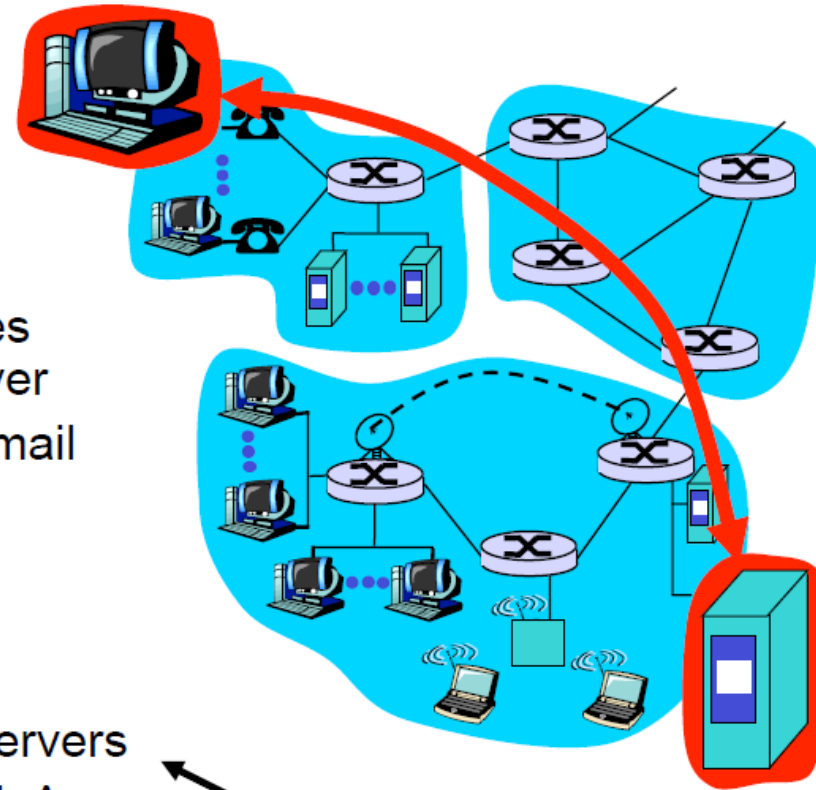
- run application programs
- e.g. Web, email

## □ client/server model

- client host requests, receives service from always-on server
- e.g. Web browser/server; email client/server

## □ peer-peer model:

- minimal use of dedicated servers
- e.g. Skype, BitTorrent, KaZaA



Any idea how?

# Network edge: connection-oriented service

Goal: data transfer between end systems

- **Connection:** prepare for data transfer ahead of time
  - Request / Respond
  - *set up “state”* in two communicating hosts
- TCP - Transmission Control Protocol
  - Internet's connection-oriented service

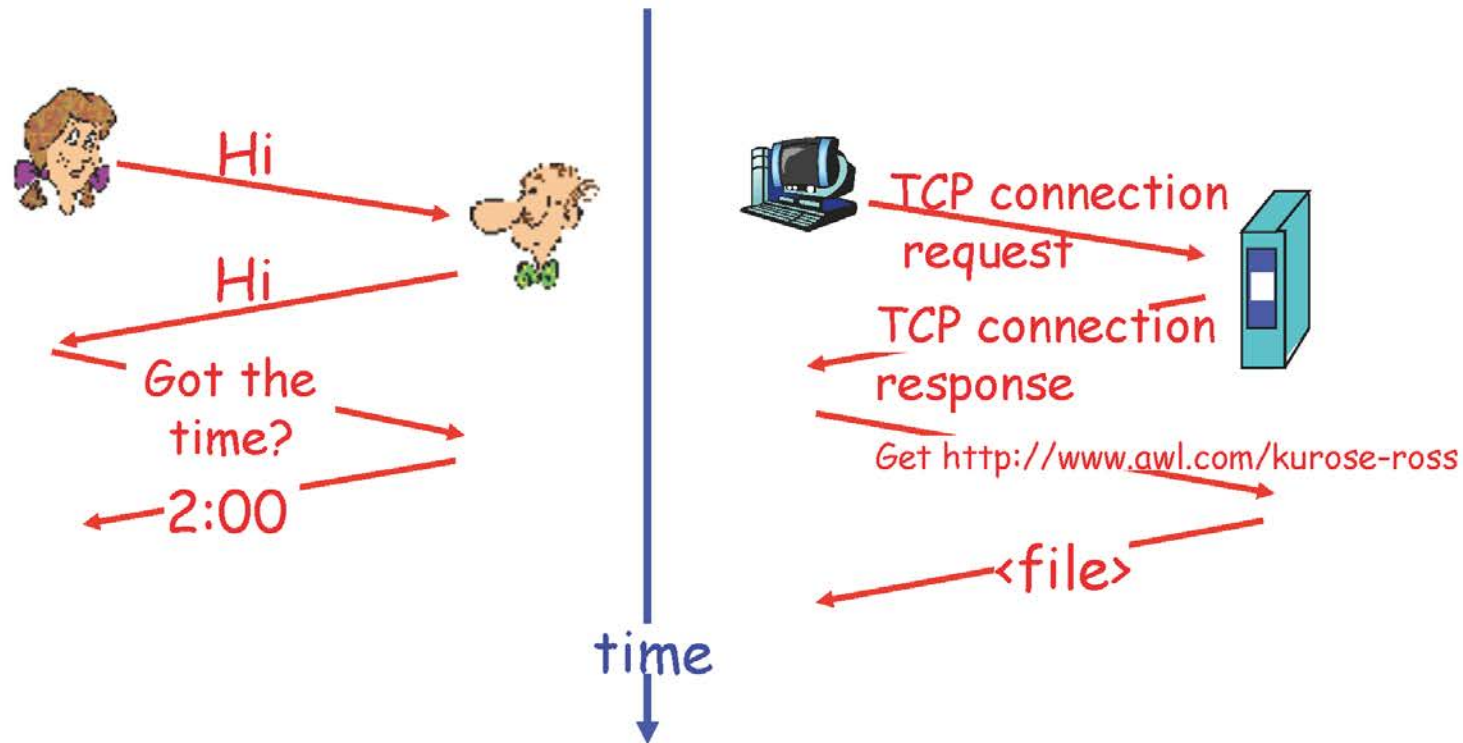
TCP service [RFC 793]

- *reliable, in-order* byte-stream data transfer
  - loss: acknowledgements and retransmissions
- *flow control:*
  - sender won't overwhelm receiver
- *congestion control:*
  - senders “slow down sending rate” when network congested



# What's a protocol?

a human protocol and a computer network protocol:

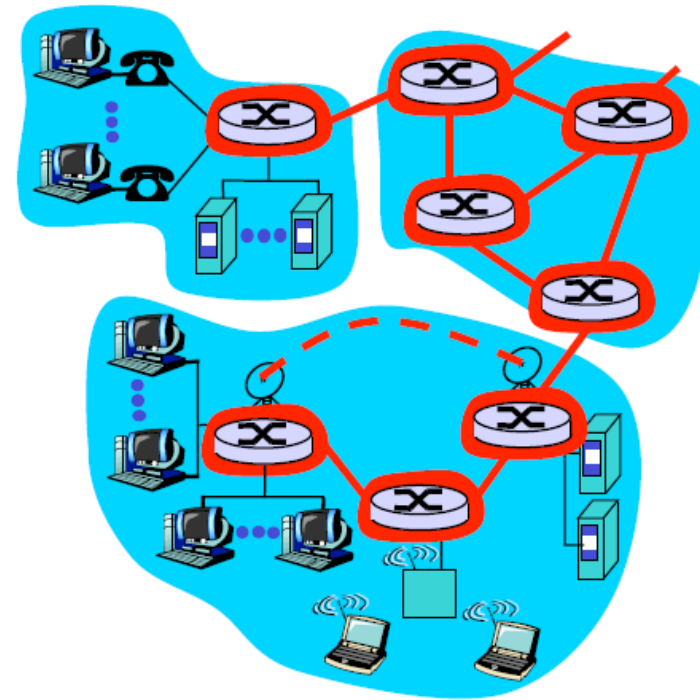


- All communication in Internet **coordinated** by protocols

# The Network Core

- mesh of interconnected routers
- the fundamental question: how is data transferred through net?

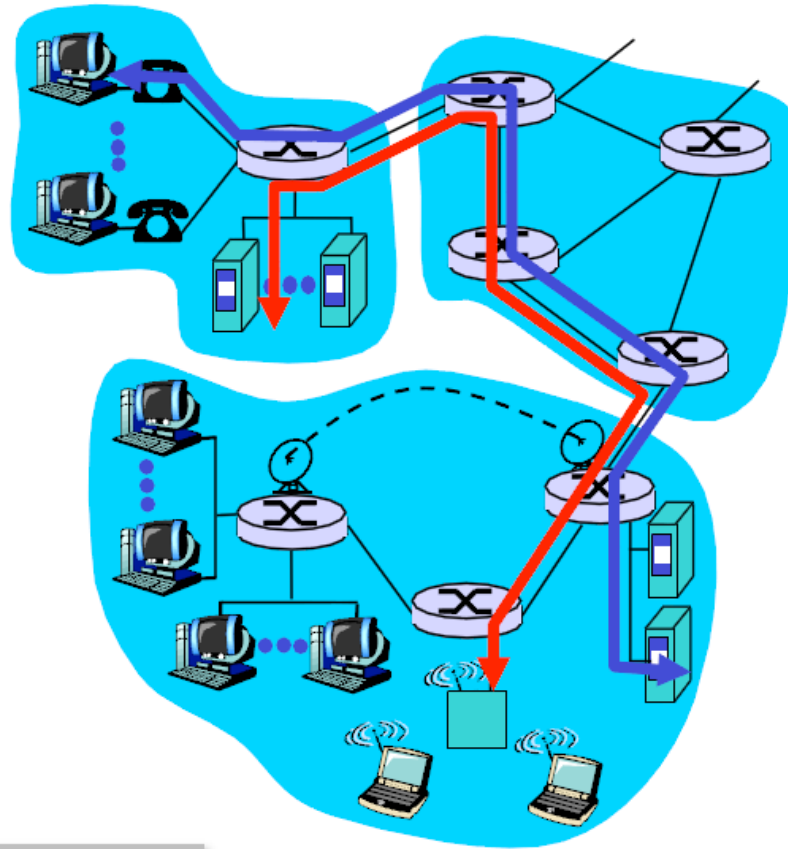
- **circuit switching:** dedicated circuit per call: telephone net
- **packet-switching:** data sent thru net in discrete “chunks”



# Network Core: Circuit Switching

End-end resources reserved for “call”

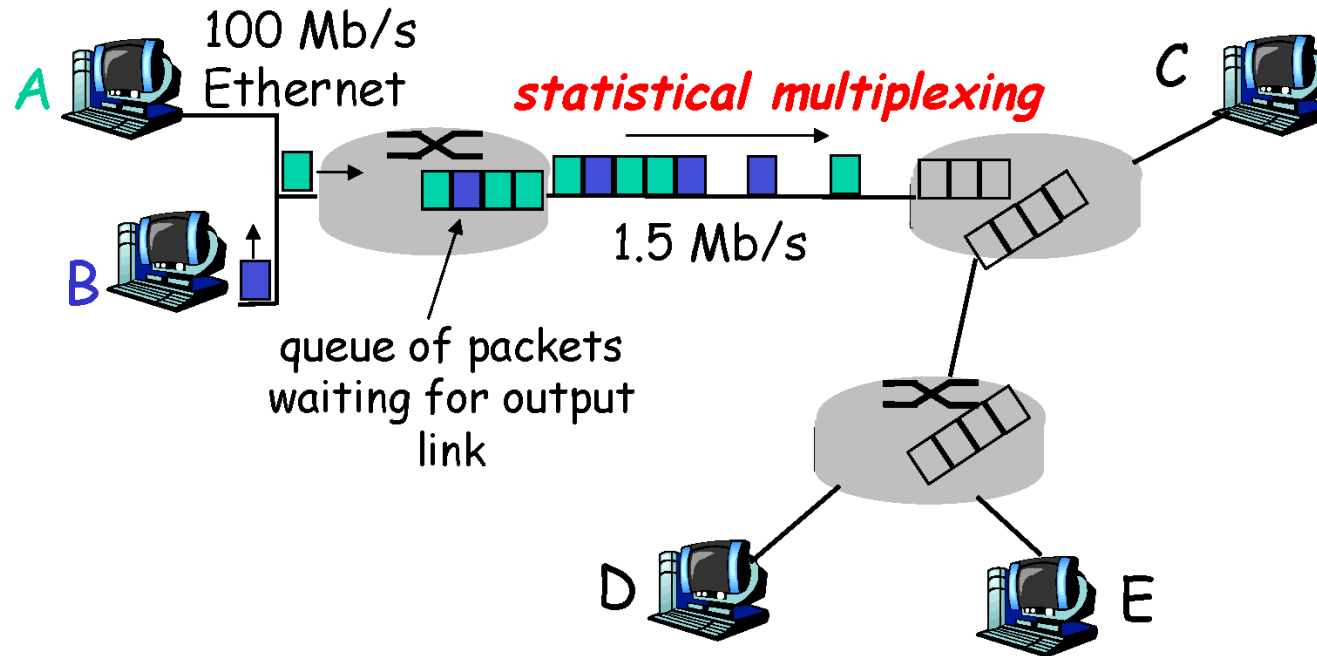
- ❑ link bandwidth, switch capacity
- ❑ dedicated resources: no sharing
- ❑ circuit-like (guaranteed) performance
- ❑ call setup required



Analogy: When president travels, a CS path set up.



# Packet Switching: Statistical Multiplexing



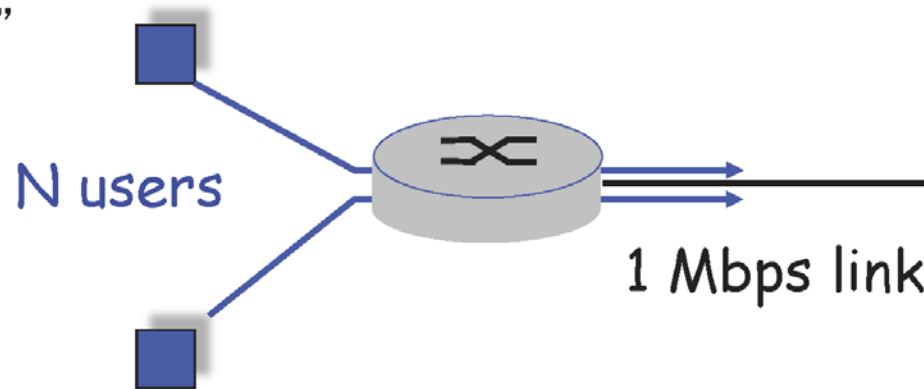
Sequence of A & B packets does not have fixed pattern, shared on demand ➡ **statistical multiplexing**.

TDM: each host gets same slot in revolving TDM frame.

# Packet switching versus circuit switching

Packet switching allows more users to use network!

- 1 Mb/s link
- each user:
  - 100 kb/s when “active”
  - active 10% of time
- circuit-switching:
  - 10 users
- packet switching:
  - with 35 users, probability  $> 10$  active less than .0004



Q: how did we get value 0.0004?

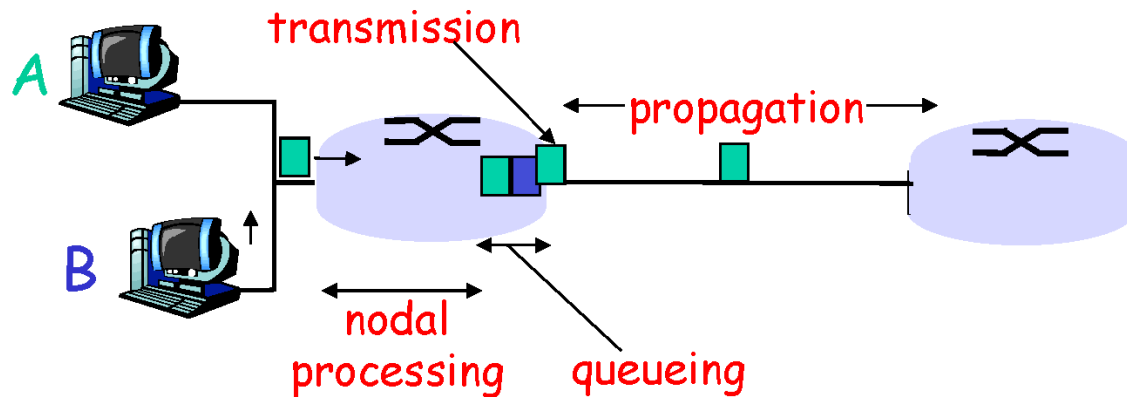
# Four sources of packet delay

## 1. nodal processing:

- check bit errors
- determine output link

## 2. queueing

- time waiting at output link for transmission
- depends on congestion level of router



# Delay in packet-switched networks

## 3. Transmission delay:

- $R$  = link bandwidth (bps)
- $L$  = packet length (bits)
- time to send bits into link =  $L/R$

## 4. Propagation delay:

- $d$  = length of physical link
- $s$  = propagation speed in medium ( $\sim 2 \times 10^8$  m/sec)
- propagation delay =  $d/s$

**Note:**  $s$  and  $R$  are very different quantities!

