

# Delay Analysis and Buffer Sizing for Priority-Aware Networks-on-Chip

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## Abstract

Priority-aware wormhole-switched Networks-on-Chip (NoC) is promising to meet both the worst-case and average-case performance requirements of on-chip communication. To deploy real-time applications on this platform, the worst-case delay and buffer requirement of each flow must be analyzed and guaranteed. In this paper, we first build a Real-Time Calculus (RTC) based performance model for the priority-aware wormhole-switched NoC. Based on this model, we then propose a delay analysis algorithm and a buffer sizing algorithm. The delay analysis algorithm can give tighter delay bound than the deterministic network calculus-based method, because our performance model takes the maximum service capability of routers and minimum arrival rate of each flow into consideration. The buffer sizing algorithm tries to reduce the buffer space required for each flow without violating the deadline constraint, which improves the result obtained by the link-level buffer-space analysis method. Both algorithms are topology-independent. Taking the architecture parameters and flow specifications as input, they give the end-to-end delay bound and buffer requirement for each application. The tightness and correctness of our algorithms are verified by comparing with simulation results.

*Keywords:* Networks-on-Chip (NoC), delay bound, buffer sizing

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## 1. Introduction

Networks-on-Chip (NoC) is proposed to meet the strict and complex communication requirements of modern large-scale Chip-MultiProcessor (CMP) and System-on-Chip (SoC). Although various NoC proposals have emerged, most of  
5 them are designed to improve the average-case performance of on-chip interconnects. However, there are a variety of on-chip applications, which are sensitive

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to the worst-case communication performance of NoC. Designing the on-chip real-time communication infrastructure for these applications and analyzing its feasibility is a major challenge for researchers.

10 To meet the rigorous on-chip real-time communication requirement, some special hardware implementations, e.g. Time-Division Multiplexing-Access (T-DMA) [1], circuit-switch [2] and time-triggered switch [3], have been proposed. However, the average performance and resource utilization of these proposals are very poor when compared with the conventional wormhole-switched NoC. Thus,  
 15 providing real-time communication support based on the wormhole-switched NoC becomes the most promising solution to meet both average-case and worst-case communication requirements.

A key step before the wormhole-switched NoC is adopted as the platform of real-time communication is to check whether the deadline constraints of all the  
 20 real-time flows are met. In addition, an effective buffer analysis approach is also needed to optimize the buffer allocation under real-time constraint, since the on-chip buffer usually contributes to a significant portion of the entire router's power and area [4, 5]. The tightness of worst-case delay and buffer requirement analysis is crucial for the application of wormhole switched NoCs in real-time  
 25 communication, because an overoptimistic analysis will lead to the violation of deadline, while an overly pessimistic analysis will make the utilization of on-chip resource very low.

The conventional simulation-based method is not appropriate for the worst-case analysis, because the worst-case scenarios are difficult to be captured by  
 30 simulation. In contrast, analytical methods can directly establish the relationship between performance metrics and design parameters, and give the worst-case performance bounds. Thus, most of the existing research focusing on the worst-case delay bound and buffer requirement of priority-aware wormhole-switched NoC are based on the analytical method [6, 7, 8, 9, 10, 11, 12]. Among  
 35 all these analytical methods, the Link-Level Analysis (LLA) method [7, 12] and the Deterministic Network Calculus (DNC)-based method [8] outperform the others when the tightness of delay bound and buffer requirement are considered. The LLA method assumes that the buffer size is large enough to eliminate the influence of flow control on the delay bound. The DNC-based method [8]  
 40 overcomes this limitation by utilizing the advanced operators and properties of the DNC theory. However, the delay bound obtained in [8] can be further improved if the maximum service capability of routers and the minimum arrival rate of input traffic are taken into consideration.

In this paper, we try to improve the delay bound and buffer requirement  
 45 derived by the DNC [8] and LLA [7, 12] method. The theoretical background of our performance model is the Real-Time Calculus (RTC) [13], which is originally used for the schedulability analysis of real-time system. To the best of our knowledge, it is the first time that this theory is used to model and analyze the performance of NoCs. The reason we choose RTC lies in the fact that it  
 50 is the extension of DNC theory, which considers the maximum service capability and minimum arrival rate while deriving the performance bounds. The main contribution of this paper is three-fold: (1) We construct a novel perfor-

mance model for the priority-aware wormhole-switched NoC with credit-based flow control. (2) We propose an end-to-end delay analysis algorithm for the priority-aware wormhole-switched NoC based on our performance model. The delay bound obtained by our algorithm is much tighter than the DNC-based method [8]. (3) We propose a buffer sizing algorithm based on our performance model. This algorithm considers the impact of flow control on the delay bound, and allocates just enough buffer space at each router for the flows to meet their deadline. When applied to guide the buffer allocation of priority-aware wormhole-switched NoCs, our algorithm can save more buffer space than the Link-Level Buffer-space Analysis (LLBA) method [12]. Our delay analysis and buffer sizing algorithms can be used for the design space exploration, IP core mapping, task mapping, routing selection, etc.

The rest of this paper is organized as follows: we present the existing performance analysis methods for the wormhole-switched NoC in Section 2. In Section 3, a brief introduction to the priority-aware wormhole-switched NoC and RTC theory are presented. The detailed modeling process is presented in Section 4, where we also propose the end-to-end delay analysis algorithm and buffer sizing algorithm. We present the comparison results with simulation and other analytical methods in Section 5. Finally, we summarize our paper in Section 6.

## 2. Related Work

Network-on-Chip is designed to be either best-effort or guaranteed-service to meet different on-chip communication requirements. Best-effort NoC can make better use of the on-chip shared resource, but it does not necessarily provide any performance guarantee for the applications. To provide the guaranteed services for different applications, a simple and effective solution is to classify these applications into several service classes, each with different priorities, and the network provides services according to the priority of each class. Representative priority-aware implementations include QNoC [14], fixed-priority NoC [6] and Aetheral [1] etc. The performance evaluation methods for the priority-aware NoC include the average-case analysis and worst-case analysis. For the average-case analysis, simulation and analytical methods hold the dominant position. However, for the worst-case analysis, simulation is not sufficient due to the difficulty in covering all the corner cases.

The worst-case performance bound of the priority-aware wormhole-switched networks has been extensively studied in literature. In [9], the contention tree model was proposed to analyze the feasibility of real-time traffic delivered by priority-aware wormhole-switched NoC. It improves the previous results, e.g. lumped link model [10] and dependency graph model [11], by allowing the concurrent link usage. The Flow-Level Analysis (FLA) proposed in [6] improves the results obtained by contention tree model [9], lumped link model [10] and dependency graph model [11]. A comprehensive comparison between FLA and the other method can be found in [15], in which several defects of the previous method are illustrated and the advantages of FLA are highlighted. The LLA [7] improves the FLA by analyzing each link segment separately. Two buffer

sizing methods based on FLA and LLA, i.e. Flow-Level Buffer-space Analysis (FLBA) and Link-Level Buffer-space Analysis (LLBA), are proposed in [12] to estimate the buffer size of priority-aware wormhole-switched NoC. Both FLA and LLA assume that the traffic arrives periodically or sporadically, and that routers have sufficiently large buffer size, which is a significant simplification to the realistic traffic pattern and router implementation. In addition, the FLBA and LLBA can only compute the minimum buffer size at each router to avoid the back-pressure caused by flow control. We can further reduce the buffer size until the deadline constraint is violated.

On the other hand, although the DNC based performance model for best-effort NoC proposed in [16] can also be applied to the analysis of priority-aware wormhole-switched NoC, the obtained performance bounds are very conservative, especially for the high-priority flows. This is because the model does not take the priority into consideration. To overcome this limitation, a revised DNC performance model was proposed to analyze the worst-case delay of priority-aware wormhole-switched NoC in [8]. We found that the DNC method in [8] can be further improved if we take the maximum service capability of each router and minimum arrival rate of each flow into consideration. Motivated by this observation, we adopt the RTC theory [13] to build the worst-case performance model for the priority-aware wormhole-switched NoC. Real-time calculus extends the DNC theory [17] by integrating the minimum arrival curve and maximum service curve to characterize more detailed information about the traffic and service processes. Due to its ability of analyzing complex system, RTC has been widely used in the modeling and analysis of Controller Area Networks [18], FlexRay [19], etc. To ease the application of RTC, a toolbox has been implemented in [20] to support the numerical calculation.

### 3. Preliminaries

#### 3.1. Basic Assumptions

In this paper, the priority-aware NoC is represented as a directional network topology graph  $G : V \times E$ , where  $V$  and  $E$  represent the set of routers and links respectively. Each link  $e_{i,j} \in E$  corresponds to a physical channel connecting the two routers  $R_i$  and  $R_j$ . A flow is a sequence of packets with the same transmission path, source address and destination address. Packet of different flows generated by a Intellectual Property (IP) core are buffered at different queues within the corresponding Network Interface (NI). Each packet is comprised of one head flit, one tail flit and several body flits. The path of a flow  $f_i$  traversed is defined as a router chain starting from the injection router (denoted as  $start_i$ ) and ending at the ejection router (denoted as  $end_i$ ). The set of all the flows in the network is denoted as  $\mathcal{F}$ , and each flow  $f_i \in \mathcal{F}$  has a fixed-priority  $P_i$  and deadline  $D_i$ . The set of routers along the path of  $f_i$  is denoted as  $\mathcal{R}_i$ , and the set of links a flow  $f_i$  traversed is denoted as  $\Gamma_i$ . There exists contention between flow  $f_i$  and  $f_j$ , if and only if  $\Gamma_i \cap \Gamma_j \neq \emptyset$ . For router  $R_j$  along the path of flow  $f_i$ , denote the set of contending flows at  $R_j$  sharing the same priority with  $f_i$  as

140  $\Theta_{R_j, f_i}$ , the set of contending flows at  $R_j$  with lower priorities than  $f_i$  as  $\Omega_{R_j, f_i}$ , and the buffer size reserved at  $R_j$  for  $f_i$  as  $B_{R_j, f_i}$ .

The router we considered is the priority-aware wormhole-switched router proposed in [6] and further discussed in [21][10][7]. Each router has the same number of input and output ports, and each input port have sufficient Virtual  
145 Channel (VC) to accommodate all the incoming packets of different priority levels. The allocation of VC is determined by the VC allocator. The buffer depth of each VC is finite, and the credit-based flow control [22] is adopted between adjacent routers to prevent buffer overflow. To ensure the predicable transmission delay, a deterministic routine computation module is used to de-  
150 termine the output port of each packet. The crossbar is utilized to switch traffic from input ports to the output ports, and the switch operation is determined by the switch allocator. The switch allocator is priority-aware. If multiple flits from different input ports or different VCs of the same input port contend for the same output port, it will only grant the flit with the highest priority. Flits  
155 from a lower priority can transmit a flit, if and only if there are no contending flits from higher priority or the higher-priority flits are self-blocked due to the insufficiency of VC buffer at the downstream router.

The micro-architecture of the priority-aware router considered in this paper has standard pipeline stages, i.e. Buffer-Write (BW), Route Computation (RC),  
160 VC Allocation (VA), Switch Allocation (SA), Switch Traversal (ST) and Link Traversal (LT). For the detailed description, please refer to [23]. Each head flit should go through all these stages to determine the path and reserve a VC for the following non-head flits. Non-head flits skip the RC and VA stages since the routine and VC have been determined by the head flit. Router resource  
165 and control information reserved for a packet will be released only after the tail flit of the packet has departed from the router. An additional priority field in the head flit is required for the routers to schedule multiple contending flows according to their priority. Although we focus on standard router, our method can be easily modified to support other router micro-architecture, e.g. single-  
170 cycle router [21][6][10][7] and speculation-based router [23]. We will demonstrate the adoption of our model in a single-cycle router in subsection 5.1. To simplify our analysis, we also assume that the entire chip is synchronous, with clock frequency  $f$  and period  $T$ . Our method can also be applied to analyze Global Asynchronous Local Synchronous (GALS) NoC with little modification, because  
175 the routers located in different voltage-frequency islands can be synchronized with a half cycle synchronizer [24], which can be abstracted as a fixed-latency element in DNC theory [17].

Our performance model is topology-independent, but to demonstrate the basic idea, we take the mesh topology shown in Fig. 1 as an example through-  
180 out this paper. Routers in the mesh topology have five input/output ports, corresponding to the four cardinal directions (West, East, North and South) and the Local IP core. Although there are only four flows, i.e.  $f_1$ ,  $f_2$ ,  $f_3$  and  $f_4$ , in the network, it is sufficient to demonstrate the idea of our method, and our method can handle more traffic flows efficiently. Our method extends the  
185 methods in [7][8] to allow multiple flows to share the same priority. Flits of

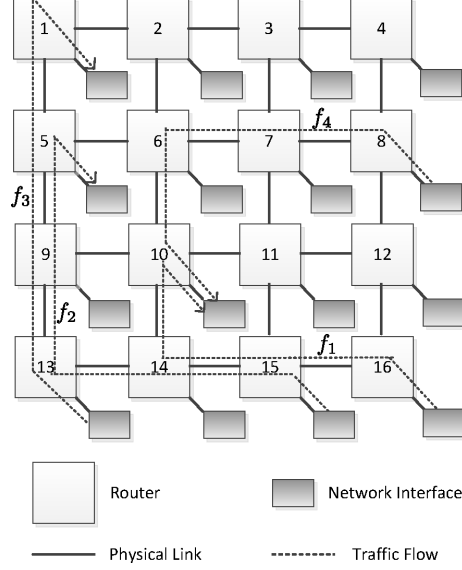


Figure 1: Mesh topology with four real-time traffic flows.

different flows sharing the same priority are served in round-robin order when they compete the same output port. Since the minimum transmission unit in the priority-aware wormhole-switched NoC is flit and a high-priority flit can preempt the transmission of a low-priority flit, the NoC architecture considered in this paper is flit-level preemptive [25].

### 3.2. Introduction to Real-Time Calculus

Real-time calculus [13] is the extension of the DNC theory [17], by adding the upper service curve and lower arrival curve to describe the maximum service capability of a system and the minimum arrival rate of a event stream. Due to the space limitation, we only present the definitions of the RTC arrival curve and service curve in this subsection. For more details, please refer to [13].

**Definition 1 (Real-Time Arrival Curve [13]).** Denote by  $R[s, t]$  the number of events within the time interval  $[s, t]$ . The lower and upper bounds on  $R[s, t]$  are called the lower arrival curve  $\alpha^l$  and upper arrival curve  $\alpha^u$ , which satisfy

$$\alpha^l(t - s) \leq R[s, t] \leq \alpha^u(t - s), \forall s < t$$

and  $\alpha^l(0) = \alpha^u(0) = 0$ . The RTC arrival curve for a event stream is denoted as  $\langle \alpha^l, \alpha^u \rangle$  for short.

**Definition 2 (Real-Time Service Curve [13]).** Denote by  $S[s, t]$  the number of events that can be processed by the system in time interval  $[s, t]$ . The

lower and upper bounds on  $S[s, t)$  are called the lower service curve  $\beta^l$  and upper service curve  $\beta^u$ , which satisfy

$$\beta^l(t - s) \leq S[s, t) \leq \beta^u(t - s), \forall s < t$$

and  $\beta^l(0) = \beta^u(0) = 0$ . The RTC service curve for a system is denoted as  $\langle \beta^l, \beta^u \rangle$  for short.

From these two definitions, we find that the upper arrival curve and lower service curve correspond to the arrival curve and service curve of DNC theory [17]. Similarly, the upper service curve corresponds to the maximum service curve of DNC theory. Thus, the two concatenation theorems for service curve (see Theorem 1.46 in [17]) and maximum service curve (see Theorem 1.6.1 in [17]) together form the concatenation theorem for the RTC service curve. Assume an event stream traverses two systems  $S_1$  and  $S_2$  in sequence, and  $S_i$  offers an RTC service curve  $\langle \beta_i^l, \beta_i^u \rangle$  ( $i = 1, 2$ ) to this event stream. The concatenation theorem can give the equivalent RTC service curve offered by these two systems to the event stream, which is  $\langle \beta_1^l \otimes \beta_2^l, \beta_1^u \otimes \beta_2^u \rangle$ .

In this paper, we will utilize the discrete time RTC arrival curve and service curve to characterize the arrived traffic and service capability of the wormhole-switched NoC, since the minimum time unit of this system is the clock period  $T$ . Events in the definitions of arrival curve and service curve refer to the arrival and service of flits, respectively. If we obtain the arrival curve  $\langle \alpha^l, \alpha^u \rangle$  of a specific flow at specific router and the service curve  $\langle \beta^l, \beta^u \rangle$  provided by this router, we can get the output arrival curve  $\langle \alpha^{l'}, \alpha^{u'} \rangle$  of this flow and leftover service curve  $\langle \beta^{l'}, \beta^{u'} \rangle$  of this router with the following equations [13]:

$$\alpha^{l'} = \min\{(\alpha^l \otimes \beta^u) \otimes \beta^l, \beta^l\} \quad (1)$$

$$\alpha^{u'} = \min\{(\alpha^u \otimes \beta^u) \otimes \beta^l, \beta^u\} \quad (2)$$

$$\beta^{l'} = (\beta^l - \alpha^u) \bar{\otimes} 0 \quad (3)$$

$$\beta^{u'} = \max\{(\beta^u - \alpha^l) \bar{\otimes} 0, 0\} \quad (4)$$

where  $\otimes$ ,  $\ominus$ ,  $\bar{\otimes}$ ,  $\bar{\ominus}$  correspond to the min-plus convolution, min-plus de-convolution, max-plus convolution and max-plus de-convolution [17].

After obtaining the arrival curve  $\langle \alpha_f^l, \alpha_f^u \rangle$  of flow  $f$  and the equivalent service curve  $\langle \beta_f^l, \beta_f^u \rangle$  offered by the system to flow  $f$ , we can get the delay bound by the following equation [17]

$$Delay(f) = H(\alpha_f^u, \beta_f^l) \quad (5)$$

where operator  $H(\cdot, \cdot)$  computes the maximal horizontal deviation between its two operands.

## 4. Delay Analysis and Buffer Sizing

230 In this section, we first build an RTC-based performance model for the priority-aware wormhole-switched NoC with credit-based flow control. Based on this performance model, we then propose an end-to-end delay analysis algorithm and a buffer sizing algorithm.

235 The performance model comprises two parts, i.e. traffic model and service model. The traffic model utilizes the RTC arrival curve to describe the arrival process of each flow. We will introduce two methods to obtain the arrival curve in subsection 4.1. The service model characterizes the services offered by the priority-aware NoC to each flow. The construction of the service model is much more complicated than the traffic model, and the following three issues should be considered: (1) Only the head flit needs to traverse the RC and VA stage 240 of a router, because the non-head flits of a packet follow the data-path built by the head flit. To simplify our service model, we need a special mechanism to characterize the service offered to head and non-head flits in a unified way. (2) Our performance model supports priority sharing among flows. Thus, the leftover service curve provided to the lower-priority flows can be derived only 245 when all the service curves of high-priority flows have been computed. (3) The cyclic-dependence between the adjacent routers caused by flow control prevents us from deriving the end-to-end delay bound with Eq.(5), since the RTC theory is generally applicable to the feed-forward system. Thus, we should first break 250 this cyclic-dependence before analyzing the performance bound. We will discuss the first two issues in subsection 4.2, and the last issue is discussed in subsection 4.3. Finally, we present the delay analysis algorithm and buffer sizing algorithm in subsection 4.4 and subsection 4.5, respectively.

### 4.1. Traffic Model

255 The communication in priority-aware wormhole-switched NoC is realized by transmitting packets, and each packet is divided into flits, which is the minimum transmission unit. Denote by  $\langle \alpha^l(\Delta), \alpha^u(\Delta) \rangle$  the flit arrival curve of a flow, namely, the minimum and maximum number of flits can be seen within any time window of length  $\Delta$ . We can extract the flit arrival curve from the 260 synthetic traffic or communication trace with the sliding window method [13]. For each window length  $\Delta$ , this method tries to find the maximal and minimal number of arrived flits (corresponding to  $\alpha^l(\Delta)$  and  $\alpha^u(\Delta)$ ) by analyzing the time series of flits. However, the obtained flit arrival curve can only be applied to compute the worst-case performance bound at the flit level, because 265 the service model we constructed in the following subsection characterizes the services provided by the system at the flit level. To obtain the packet level delay bound, this arrival curve must be  $L$ -packetized [17]. Denote by  $L(n)$  the cumulative packet length (in flits) of the first  $n$  packets in a flow,  $R(t)$  the cumulative arrived flits by time  $t$ . Then, the  $L$ -packetizer operator  $\mathcal{P}^l(\cdot)$  is defined



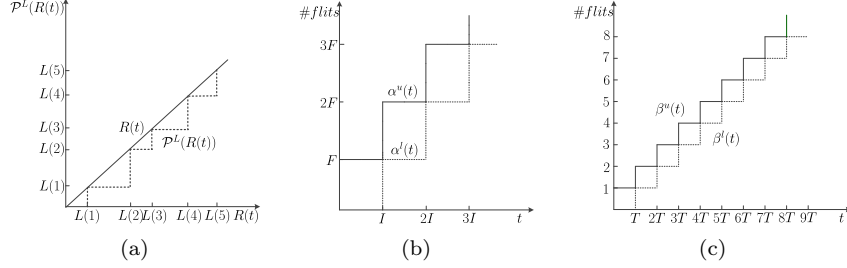


Figure 2: Traffic model and service model. (a) Definition of  $\mathcal{P}^L(R(t))$ . Cumulative arrival function  $R(t)$  and the  $L$ -packetized cumulative arrival function  $\mathcal{P}^L(R(t))$  are represented by the solid line and dotted line, respectively. (b) Real-time calculus arrival curve for periodically arrived traffic with period  $I$  and packet length  $F$ . The solid line and dotted line represent the upper arrival curve and lower arrival curve. (c) Service model for each pipeline stage. The solid lines and dotted lines represent the upper service curves and lower service curves, respectively.

270 as  $\mathcal{P}^L(R(t)) = \sup_{n \in \mathcal{N}} \{L(n)1_{L(n) \leq R(t)}\}^1$ . Intuitively,  $\mathcal{P}^L(\cdot)$  can be interpreted as the largest cumulative packet length contained in  $R(t)$ , as shown in Fig. 2a. For any flit arrival curve  $\langle \alpha^l(\Delta), \alpha^u(\Delta) \rangle$ , the  $L$ -packetized arrival curve can be obtained by applying the following theorem.

**Theorem 1 ( $L$ -packetized arrival curve).** Suppose a flow has a flit arrival curve  $\langle \alpha^l(\Delta), \alpha^u(\Delta) \rangle$ , and the maximum packet length (in flits) of this flow is  $l_{max}$ . Then, it has an  $L$ -packetized arrival curve  $\langle \alpha^l(\Delta) - l_{max}1_{\{\Delta > 0\}}, \alpha^u(\Delta) + l_{max}1_{\{\Delta > 0\}} \rangle$ .

PROOF. For  $\forall t \geq 0, \Delta \geq 0$ , according to the basic properties of  $L$ -packetizer [17], we have

$$R(t) - l_{max} < \mathcal{P}^L(R(t)) \leq R(t)$$

and

$$R(t + \Delta) - l_{max} < \mathcal{P}^L(R(t + \Delta)) \leq R(t + \Delta).$$

The inequalities above indicate

$$R(t + \Delta) - R(t) - l_{max} < \mathcal{P}^L(R(t + \Delta)) - \mathcal{P}^L(R(t))$$

and

$$\mathcal{P}^L(R(t + \Delta)) - \mathcal{P}^L(R(t)) < R(t + \Delta) - R(t) - l_{max}.$$

Based on the definition 1, the  $L$ -packetized flow has a packet arrival curve  $\langle \alpha^l(\Delta) - l_{max}1_{\{\Delta > 0\}}, \alpha^u(\Delta) + l_{max}1_{\{\Delta > 0\}} \rangle$ , which ends the proof. ■

<sup>1</sup> $\mathcal{N}$  is the set of natural numbers and  $1_{\{val\}}$  is the indicator function,  $1_{\{val\}} = 1$  if and only if  $val$  is true.

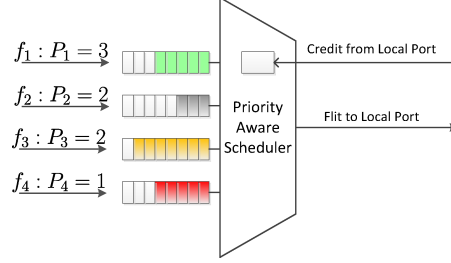


Figure 3: Structure of priority-aware network interface. Each flow has its dedicated buffer, and the scheduler schedules the flit with the highest-priority to go through the output link at each cycle.

280 For some special cases, we can also obtain the  $L$ -packetized arrival curve directly instead of transformation from flit arrival curve. For example, suppose all the packets in a flow have the same length  $F$  and arrived periodically with period  $I$ . By applying the sliding window method [13], we can obtain the flit arrival curve of this flow, which is a pair of staircase functions<sup>2</sup>  $< F \cdot u_{I,0}, F \cdot u_{I,I} >$ . As shown in Fig. 2b, the obtained flit arrival curve which is equal to the  $L$ -packetized arrival curve  $\mathcal{P}^L(\alpha)$ , since  $\mathcal{P}^L(R(t)) = R(t)$ ,  $\mathcal{P}^L(\alpha^l(t)) = \alpha^l(t)$  and  $\mathcal{P}^L(\alpha^u(t)) = \alpha^u(t)$ .  
285

#### 4.2. Basic Feed-forward Service Model

290 The service model characterizes the service obtained by each flow along its entire transmission path. The transmission path of a flow comprises of NI, routers and physical links. The service curve offered by a physical link is a pair of staircase function  $< u_{T,0}, u_{T,T} >$  since it serves one flit per cycle. The service curves provided by the NI and routers will be discussed as follows.

##### 4.2.1. Service Curve Provided by Source NI

295 Figure 3 illustrates the internal structure of this priority-aware NI, where the IP core generates four flows simultaneously. Messages of different flows are encapsulated and stored in the dedicated buffer for that flow. The priority-aware scheduler selects one flit with the highest-priority at a time for transmission. Unscheduled flits will be imposed an additional latency  $T$ . The RTC service curve provided by the source NI (denoted as  $< \beta_{NI}^l, \beta_{NI}^u >$ ) to all these flows  
300 can be obtained by applying the sliding window method [13], which is a pair of staircase functions  $< u_{T,0}, u_{T,T} >$ , as shown in Fig. 2c.

Given the set of flow specifications, the service curve obtained by each flow at the source NI can be derived by applying Algorithm 1. The flow specification

<sup>2</sup>A staircase function  $u_{T,\tau} = \lceil \frac{t+\tau}{T} \rceil$  for  $t > 0$  and 0 otherwise, where  $0 \leq \tau \leq T$  and  $\lceil \cdot \rceil$  is the ceiling operator.

305 of  $f_i$  is a quadruple  $(\langle \alpha^l, \alpha^u \rangle, \mathcal{R}_i, D_i, P_i)$  that specifies the arrival curve, routine, deadline and priority of a flow.

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**Algorithm 1** Compute the service curve at source NI

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**Require:**  $(\langle \alpha^l, \alpha^u \rangle, \mathcal{R}_i, D_i, P_i)$  // the set of flow specifications

**Ensure:**  $\langle \beta_{NI, f_j}^l, \beta_{NI, f_j}^u \rangle$  // the service curve obtained by each flow

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1: Group the flows with priority  $P_i$  into subset  $\mathcal{F}_i$ .
2:  $\beta_{NI}^l = \lfloor \frac{t}{T} \rfloor$ ;  $\beta_{NI}^u = \lceil \frac{t}{T} \rceil$ .
3: for each  $\mathcal{F}_i$  from highest priority to lowest priority do
4:   for each flow  $f_j \in \mathcal{F}_i$  do
5:      $\beta_{NI, f_j}^l = \lfloor \frac{\beta_{NI}^l}{|\mathcal{F}_i|} \rfloor$ ;  $|\mathcal{F}_i|$  denotes the cardinality of  $\mathcal{F}_i$ 
6:      $\beta_{NI, f_j}^u = \lceil \frac{\beta_{NI}^u}{|\mathcal{F}_i|} \rceil$ .
7:   end for
8:    $\alpha^l = \sum_{f_j \in \mathcal{F}_i} \alpha_{f_j}^l$ ;  $\alpha^u = \sum_{f_j \in \mathcal{F}_i} \alpha_{f_j}^u$ ;
9:    $\beta_{NI}^l = (\beta_{NI}^l - \alpha^u) \bar{\otimes} 0$ ;  $\beta_{NI}^u = \max\{(\beta_{NI}^u - \alpha^l) \bar{\otimes} 0, 0\}$ .
10: end for
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#### 4.2.2. Service Curve Provided by Routers

While modeling the service capability of a router with RTC, we can analyze the data-path of a flow in the router stage-by-stage. Once we obtain the service curves offered by each stage, the service curve provided by the router to a flow can be derived by concatenating all these service curves. This is significantly different from the existing DNC based model [16, 8], where they treat the entire router as a whole and designate a Latency-Rate (LR) service curve [17] to simplify the performance bounds derivation. Whereas, our model uses the staircase functions to characterize the detailed behavior of this discrete time system. The advantage of our method is that, it can be easily modified to characterize the non-standard router micro-architectures, by simply letting the service curve of non-existed stages to be a burst delay function  $\delta_0(t)$ <sup>3</sup>. Next, we try to derive the service curves of all these stages:

320 (1) BW stage, SA stage and LT stage: all the flits within a traffic flow will go through these three stages, and experience a fixed delay  $T$  at each stage. The service curves provided by these stages, i.e.  $\langle \beta_{BW}^l, \beta_{BW}^u \rangle$ ,  $\langle \beta_{SA}^l, \beta_{SA}^u \rangle$  and  $\langle \beta_{LT}^l, \beta_{LT}^u \rangle$ , can be derived by applying the sliding window method [13], which are the same as the service curve of source NI, as shown in Fig. 2c.

325 (2) RC stage and VA stage: the latency of head flit experienced at these two stages is  $T$ . Although the non-head flits do not go through these two stages, they have to wait for at most two cycles before entering the SA stage, e.g. the first three body flits of a packet shown in Fig. 4. Thus, a sophisticated solution to construct a unified lower service curve for head flit and non-head flits at these

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<sup>3</sup> $\delta_{val}(t) = +\infty$  if  $t > val$ , and 0 otherwise.

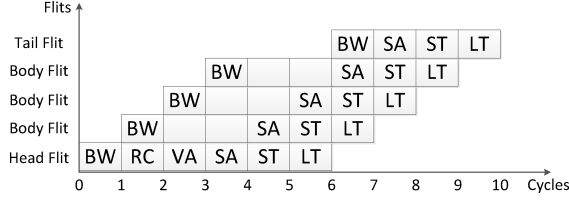


Figure 4: Time-line graph of a packet going through the standard router pipeline. The delayed tail flit enters the ST stage immediately after it is written into the dedicated buffer.

two stages comes from viewing each of these two stages impose an additional delay  $T$  for all the flits. Thus, the equivalent lower service curve of these two stages, i.e.  $\beta_{RC}^l$  and  $\beta_{VA}^l$ , can be obtained by applying the sliding window method [13], which are equal to the staircase function  $u_{T,0}(t)$ . To derive the upper service curve of these two stages, let us consider the most ‘lucky’ flits of a flow, e.g. the tail flit shown in Fig. 4. This flit can enter the SA stage immediately after it was written into the dedicated VC buffer. For this case, the RC and VA stages impose a zero latency to it. Thus, we can utilize the burst delay function  $\delta_0(t)$  to represent the upper service curve of these two stages.

(3) ST stage: each output port of the wormhole-switched NoC has a switch allocator to schedule the switch traversal among all the contending flows at each clock cycle. The notation  $\langle \beta_{ST,R_i^p}^l, \beta_{ST,R_i^p}^u \rangle$  is used to identify the service curve obtained by all the contending flows injected into switch port  $p$ . For the mesh topology, let the port indicator  $p$  be  $W$  (West port),  $E$  (East port),  $S$  (South port),  $N$  (North port) or  $L$  (Local port). Thus, Following the same procedure as BW stage, we can get the service curves provided to all the contending flows, which is  $\langle u_{0,T}(t), u_{T,T}(t) \rangle$ , as shown in Fig. 2c.

It can be found that the contention of different flows within a router only might occur at ST stage. For the fixed-priority scheduling policy, switch allocators schedule the flow with the highest priority first, flows with the same priority will be served with round-robin order. All the unscheduled flows will impose an additional latency  $T$  due to the failure of switch arbitration.

Denote by  $\langle \beta_{ST,R_i^p}^l, \beta_{ST,R_i^p}^u \rangle$  the total service curve provided by the ST stage,  $\langle \beta_{ST,R_i,f_j}^l, \beta_{ST,R_i,f_j}^u \rangle$  the service curve provided to flow  $f_j$  by ST stage of router  $R_i$ , and  $\langle \beta_{ST,R_i^p}^{l'}, \beta_{ST,R_i^p}^{u'} \rangle$  the leftover service curve after serving the flows with higher priority than  $f_j$ . In order to obtain the service curve  $\langle \beta_{ST,R_i,f_j}^l, \beta_{ST,R_i,f_j}^u \rangle$ , we should consider the following two cases:

- (a) All the flows contending with  $f_j$  at  $R_i$  have lower priorities than  $f_i$ . For the synchronized router architecture, flow  $f_j$  gets the total service curve  $\langle \beta_{ST,R_i^p}^{l'}, \beta_{ST,R_i^p}^{u'} \rangle$ .
- (b) There exist some contention flows with the same priority as  $f_j$ . Since all the flows in  $\Theta_{R_i,f_j}$  got serviced in round-robin order, the service curve provided

to  $f_j$  is  $< \lfloor \beta_{ST,R_i^p}^{l'} / (|\Theta_{R_i,f_j}| + 1) \rfloor, \lceil \beta_{ST,R_i^p}^{u'} / (|\Theta_{R_i,f_j}| + 1) \rceil >$ . After serving all the flows in  $\Theta_{R_i,f_j}$ , the leftover service curve for low-priority flows can be derived by applying Eq.(3) and Eq.(4).

After we obtained the service curve provided by ST stage to flow  $f_j$ , we can get the service curve of the router directly. The equivalent feed-forward service curve of router  $R_i$  provided to  $f_j$ , i.e.  $< \beta_{R_i,f_j}^l, \beta_{R_i,f_j}^u >$ , can be obtained by concatenating the service curves of all these stages together:

$$\begin{aligned}\beta_{R_i,f_j}^l &= \beta_{BW}^l \otimes \beta_{RC}^l \otimes \beta_{VA}^l \otimes \beta_{SA}^l \otimes \beta_{ST,R_i,f_j}^l, \\ \beta_{R_i,f_j}^u &= \beta_{BW}^u \otimes \beta_{RC}^u \otimes \beta_{VA}^u \otimes \beta_{SA}^u \otimes \beta_{ST,R_i,f_j}^u.\end{aligned}$$

#### 365 4.3. Feedback Service Model

To this end, we have constructed the traffic model and the basic feed-forward service model. However, these two models can only be applied to analyze the performance of wormhole-switched NoC without credit-based flow control. Flow control introduces cyclic-dependence between the adjacent routers, and blocks  
370 a flow when there is no buffer space available at the downstream router. The cyclic-dependence between the adjacent routers prevents us from deriving the performance bound directly even after we have obtained the service curve reserved at each router for the target flow. In existing literature, this problem is addressed by fixed-point iteration [26] or transformation from marked dataflow graph [27]. In this subsection, we will try to tackle the same problem with  
375 another solution.

Our solution is inspired by [16], where the authors abstract the flow control as a network element (called flow controller) providing a service curve (corresponding to the lower service curve of RTC theory), and this service curve can  
380 be obtained by applying some basic properties of DNC theory [17]. To make the discussion concrete, we take flow  $f_2$  in Fig. 1 as an example. We also utilize the scheduling network model [13] in RTC theory to visualize the credit-based flow control and complex relationship among  $f_2$  and the other flows, as shown in Fig. 5. We ignore flow  $f_4$  and the flow control of the other flows for brevity and  
385 clarity. We also assume that, all the destination IP cores can consume the arrived flits immediately, thus there is no flow control between the ejection router and destination NI. However, to prevent the buffer overflow, the flow control between source NI and injection router is necessary.

A flow in the wormhole-switched router with credit-based flow control will  
390 be blocked if there is no credit available. We can treat the blocking of a flow caused by credit-insufficiency as traversing a virtual pipeline stage, called Flow Control (FC) stage, as shown in Fig. 5. The equivalent service curve for this virtual stage can be obtained by the following theorem, which enables us to break the cyclic-dependence caused by flow control and build a comprehensive  
395 performance model for the wormhole-switched router.

**Theorem 2.** Suppose that a router provides a feed-forward service curve  $< \beta^l, \beta^u >$ , and the physical link provides service curve  $< \beta_{LT}^l, \beta_{LT}^u >$ , the buffer

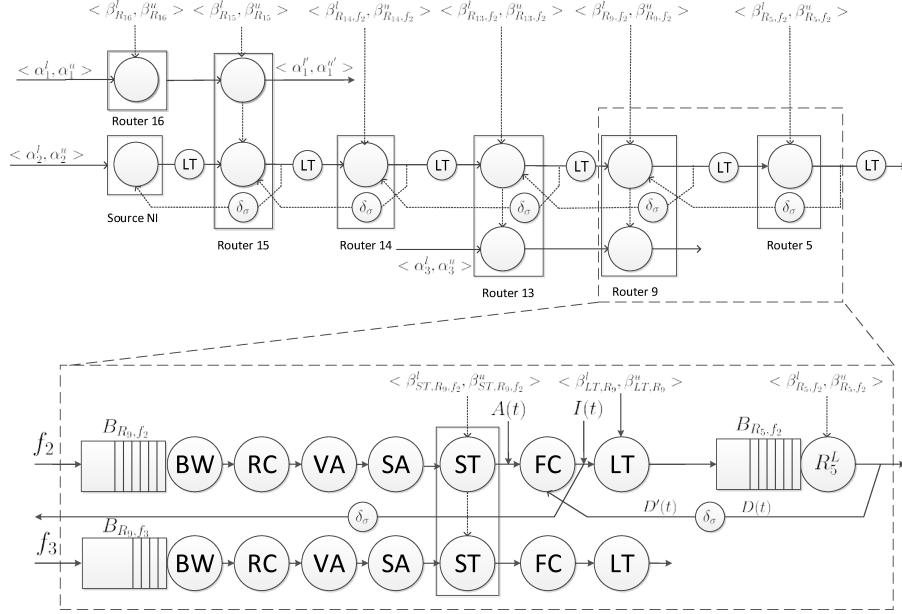


Figure 5: Scheduling network model for flow  $f_2$

size and credit feedback delay are denoted as  $B$  and  $\sigma$ , respectively. Then, the corresponding flow controller provides an service curve  $\langle \beta^l \otimes \beta_{LT}^l \otimes \delta_\sigma + B, \beta^u \otimes \beta_{LT}^u \otimes \delta_\sigma + B \rangle$ , where  $f$  is the sub-additive closure<sup>4</sup> of  $f$ .

PROOF. We will take the flow controller between  $R_9$  and  $R_5$  in Fig. 5 as an example to derive the service curve of flow controller. Denote the amount of injected and departed flits at  $R_5$  by time  $t$  as  $I(t)$  and  $D(t)$ , and the amount of flits served by  $R_9$  by time  $t$  as  $A(t)$ . The feedback link can be represented as a network element providing upper service curve  $\delta_\sigma(t)$ . The DNC service curve (i.e. lower service curve of RTC) has been derived in [16], which is  $\beta^l \otimes \beta_{LT}^l \otimes \delta_\sigma + B_{R_5, f_2}$ .

In the rest of this proof, we will derive the upper service curve for the flow controller. For the flow control between router  $R_9$  and  $R_5$ , we have  $I(t) \leq A(t)$  for causality and  $I(t) \leq D'(t) + B_{R_5, f_2}$  due to the effect of flow control, where  $D'(t) \leq D \otimes \delta_\sigma(t)$ . Thus,

$$I(t) \leq \min\{A(t), D'(t) + B_{R_5, f_2}\}.$$

Since the upper service curve corresponds to the maximum service curve in DNC theory [17], the following inequality holds by the definition of maximum

<sup>4</sup>  $\bar{f} = \inf_{n \geq 0} \{f^{(n)}\}$ , where  $f^{(n)}$  represents the  $n$ th-fold convolution of  $f$  and  $f^{(0)}(t) = \delta_0(t)$ .

service curve (see definition 1.6.1 in [17])

$$D(t) \leq I \otimes \beta_{LT}^u(t) \otimes \beta_{R_5, f_2}^u.$$

Bring  $I(t)$  and  $D'(t)$  into the above inequality, we get

$$\begin{aligned} D(t) &\leq I \otimes \beta_{R_5, f_2}^u \otimes \beta_{LT}^u(t) \\ &\leq \min\{A \otimes \beta_{R_5, f_2}^u \otimes \beta_{LT}^u(t), D \otimes \delta_\sigma \otimes \beta_{R_5, f_2}^u \otimes \beta_{LT}^u(t) + B\}. \end{aligned}$$

By applying Theorem 4.31 in [17], we have

$$D \leq A \otimes \beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \overline{\beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma} + B_{R_5, f_2}.$$

Thus,

$$\begin{aligned} I &\leq \min\{A, D' + B_{R_5, f_2}\} \\ &\leq \min\{A, D \otimes \delta_\sigma + B_{R_5, f_2}\} \\ &\leq \min\{A, A \otimes \beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \overline{\beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma} + B_{R_5, f_2} \otimes \delta_\sigma + B_{R_5, f_2}\} \\ &= \min\{A \otimes \delta_\sigma, A \otimes \overline{\delta_\sigma \otimes \beta_{R_5, f_2}^u \otimes \beta_{LT}^u} + B_{R_5, f_2}\} \\ &= A \otimes \min\{\delta_\sigma, \overline{\beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma} + B_{R_5, f_2}\} \\ &= A \otimes \overline{\beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma} + B_{R_5, f_2} \end{aligned}$$

410 where the steps from the third line to the fifth line hold due to the general properties of  $\otimes$  operator (see Rule 6 and Rule 7 of Theorem 3.1.5 in [17]), and the last step follows from the definition of sub-additive closure.

The inequality  $I \leq A \otimes \overline{\beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma} + B_{R_5, f_2}$  implies that the flow controller at  $R_9$  provides an upper service curve  $\overline{\beta_{R_5, f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma} + B_{R_5, f_2}$ . Thus, 415 we can conclude that for any router providing an upper service curve  $\beta^u$ , the corresponding flow controller offers an upper service curve  $\overline{\beta^u \otimes \beta_{LT}^u \otimes \delta_\sigma(t)} + B$ , which ends the proof. ■

On obtaining the equivalent service curve of flow controller for  $f_i$  at router  $R_j$  (denoted as  $\langle \beta_{FC, R_j, f_i}^l, \beta_{FC, R_j, f_i}^u \rangle$ ), we get the equivalent service curve of router  $R_j$  after breaking the cyclic-dependence loop:

$$\begin{aligned} \beta_{R_j, f_i}^l &= \beta_{BW}^l \otimes \beta_{RC}^l \otimes \beta_{VA}^l \otimes \beta_{SA}^l \otimes \beta_{ST, R_j, f_i}^l \otimes \beta_{FC, R_j, f_i}^l, \\ \beta_{R_j, f_i}^u &= \beta_{BW}^u \otimes \beta_{RC}^u \otimes \beta_{VA}^u \otimes \beta_{SA}^u \otimes \beta_{ST, R_j, f_i}^u \otimes \beta_{FC, R_j, f_i}^u. \end{aligned}$$

Theorem 2 can be utilized to derive the RTC service curve of a single flow controller, and we can get the service curves of all the flow controllers along 420 the router chain of any flow by applying Theorem 2 iteratively. As shown in Fig. 5, the service curve of a flow controller is determined by the service curves of the downstream flow controllers and routers. Hence, for each flow, we should compute the service curves of flow controllers from the ejection router to the source NI. Take flow  $f_2$  as an example, we have  $\beta_{FC, R_5, f_2}^l(t) = \delta_0(t)$  and

425  $\beta_{FC,R_5,f_2}^u(t) = \delta_0(t)$  since there is no flow control between  $R_5$  and destination NI. Then, we can compute the service curve of flow controller at  $R_9$ , which is  $< \beta_{R_5,f_2}^l \otimes \beta_{LT}^l \otimes \delta_\sigma + B_{R_5,f_2}, \overline{\beta_{R_5,f_2}^u \otimes \beta_{LT}^u \otimes \delta_\sigma + B_{R_5,f_2}} >$ . By applying the concatenation theorem, we can obtain the equivalent service curve provided to  $f_2$  by router  $R_9$ , which can be utilized to derive  $< \beta_{FC,R_{13},f_2}^l, \beta_{FC,R_{13},f_2}^u >$  further.

430 Following the same procedure, the service curve  $< \beta_{FC,R_{14},f_2}^l, \beta_{FC,R_{14},f_2}^u >$ ,  $< \beta_{FC,R_{15},f_2}^l, \beta_{FC,R_{15},f_2}^u >$  and  $< \beta_{FC,NI,f_2}^l, \beta_{FC,NI,f_2}^u >$ <sup>5</sup> can be derived.

#### 4.4. End-to-End Delay Analysis

In this subsection, we present the delay analysis algorithm, as shown in Algorithm 2. This algorithm takes the architecture parameters and flow specifications as input, and gives the worst-case end-to-end delay for all the flows.

435 The architecture parameters specify the network topology graph, buffer size of each VC and the service curve of each pipeline stage. The flow specifications describes the arrival curve, routine, deadline and priority of each flow. The arrival curve of flow  $f_i$  at the source NI and ST stage of router  $R_j$  are denoted as  $< \alpha_{f_i}^l, \alpha_{f_i}^u >$  and  $< \alpha_{R_j,f_i}^l, \alpha_{R_j,f_i}^u >$ , respectively. The leftover service curve of ST stage at output port  $p$  is represented as  $< \beta_{ST,R_j^p}^l, \beta_{ST,R_j^p}^u >$  (Initially, let  $\beta_{ST,R_j^p}^l = \beta_{ST,R_j^p}^l$  and  $\beta_{ST,R_j^p}^u = \beta_{ST,R_j^p}^u$ ).

In the fixed-priority flit-level preemptive NoC, only the leftover service curve can be used by the low-priority flows. Thus, our algorithm compute the leftover service curve and delay bound from high-priority flows to low-priority flows.

445 For each iteration, it performs the following four steps in sequence: (1) Calculate the service curves provided by the routers (lines 6-7) and flow controllers (lines 8-9) along the path. (2) Compute the worst-case end-to-end delay of the flow (lines 11-15), where the service curve provided by the source NI to  $f_i$ , i.e.  $< \beta_{NI,f_i}^l, \beta_{NI,f_i}^u >$  has been computed with Algorithm 1. (3) The highlight of performance model when compared with the LLA method [7] and DNC method [8] is that our algorithm supports the priority-sharing. Thus, the leftover service curve at each router for low-priority flows can only be calculated when all the flows sharing the same priority have been calculated. To calculate the leftover service curve at ST stage, we have to first derive the equivalent service curve from source NI to  $R_j$  (lines 21-22) and the arrival curve at  $R_j$  (lines 24-25). Then, derive the leftover service curve with the aggregate arrival curve of the same priority-level (lines 26-31). The overall algorithm has two-level embedded loops, and the computation complexity for this algorithm is  $O(HN)$ , where  $N$  and  $H$  is the number of flows and the hop count of each flow. This algorithm can be easily implemented in the RTC toolbox [20] to compute the end-to-end delay bound automatically. Since our algorithm takes the upper service curve and lower arrival curve into consideration, the delay bound obtained by our

<sup>5</sup> $< \beta_{FC,NI,f_2}^l, \beta_{FC,NI,f_2}^u >$  denotes the service curve of flow controller between source NI and injection router.



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**Algorithm 2** End-to-end delay analysis algorithm
 

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**Require:** Architecture parameters and flow specifications

**Ensure:** Worst-case end-to-end delay for all the flows

```

1: for each flow  $f_i \in \mathcal{F}$  with priority order do
2:   // Compute the service curve at each router for  $f_i$ 
3:    $\beta_\tau^l = \delta_0(t)$ ;
4:    $\beta_\tau^u = \delta_0(t)$ ;
5:   for each router  $R_j \in \mathcal{R}_i$  from  $end_i$  to  $start_i$  do
6:      $\beta_{R_j, f_i}^l = \beta_{BW}^l \otimes \beta_{RC}^l \otimes \beta_{VA}^l \otimes \beta_{SA}^l \otimes \lfloor \frac{\beta_{ST, R_j^p}^{l'}}{|\Theta_{R_j, f_i}|+1} \rfloor \otimes \beta_\tau^l$ ;
7:      $\beta_{R_j, f_i}^u = \beta_{BW}^u \otimes \beta_{RC}^u \otimes \beta_{VA}^u \otimes \beta_{SA}^u \otimes \lceil \frac{\beta_{ST, R_j^p}^{u'}}{|\Theta_{R_j, f_i}|+1} \rceil \otimes \beta_\tau^u$ ;
8:      $\beta_\tau^l = \frac{\beta_{R_j, f_i}^l \otimes \beta_{LT}^l \otimes \delta_\sigma(t) + B_{R_j, f_i}}{\beta_{R_j, f_i}^l \otimes \beta_{LT}^l \otimes \delta_\sigma(t) + B_{R_j, f_i}}$ ;
9:      $\beta_\tau^u = \frac{\beta_{R_j, f_i}^u \otimes \beta_{LT}^u \otimes \delta_\sigma(t) + B_{R_j, f_i}}{\beta_{R_j, f_i}^u \otimes \beta_{LT}^u \otimes \delta_\sigma(t) + B_{R_j, f_i}}$ ;
10:  end for
11:   $\beta_{FC, NI, f_i}^l = \beta_\tau^l$ ;
12:   $\beta_{FC, NI, f_i}^u = \beta_\tau^u$ ;
13:  // Compute the end-to-end delay for  $f_i$ 
14:   $\beta_{f_i} = \beta_{NI, f_i}^l \otimes \beta_{FC, NI, f_i}^l \otimes (\bigotimes_{R_k \in \mathcal{R}_i} (\beta_{R_k, f_i}^l \otimes \beta_{LT}^l))$ ;
15:   $Delay(f_i) = H(\alpha_{f_i}^u, \beta_{f_i})$ ;
16:  // Compute the leftover service curve for low-priority flows
17:   $\beta_{f_i}^l = \beta_{NI, f_i}^l \otimes \beta_{FC, NI, f_i}^l$ ;
18:   $\beta_{f_i}^u = \beta_{NI, f_i}^u \otimes \beta_{FC, NI, f_i}^u$ ;
19:  for  $\forall R_j \in \mathcal{R}_i$  from  $start_i$  to  $end_i$  do
20:    if  $\Omega_{R_j, f_i} \neq \emptyset$  then
21:       $\beta^l = \beta_{f_i}^l \otimes \beta_{BW}^l \otimes \beta_{RC}^l \otimes \beta_{VA}^l \otimes \beta_{SA}^l$ ;
22:       $\beta^u = \beta_{f_i}^u \otimes \beta_{BW}^u \otimes \beta_{RC}^u \otimes \beta_{VA}^u \otimes \beta_{SA}^u$ ;
23:      // Compute the arrival curve of  $f_i$  at  $R_j$ 
24:       $\alpha_{R_j, f_i}^l = \min\{(\alpha_{f_i}^l \otimes \beta^u) \otimes \beta^l, \beta^l\}$ ;
25:       $\alpha_{R_j, f_i}^u = \min\{(\alpha_{f_i}^u \otimes \beta^u) \otimes \beta^l, \beta^u\}$ ;
26:      if  $Delay(f_k)$  has been calculated for  $\forall f_k \in \Theta_{R_j, f_i}$  then
27:         $\alpha_{R_j, f_i}^l = \alpha_{R_j, f_i}^l + \sum_{f_k \in \Theta_{R_j, f_i}} \alpha_{R_j, f_k}^l$ ;
28:         $\alpha_{R_j, f_i}^u = \alpha_{R_j, f_i}^u + \sum_{f_k \in \Theta_{R_j, f_i}} \alpha_{R_j, f_k}^u$ ;
29:         $\beta_{ST, R_j^p}^{l'} = (\beta_{ST, R_j^p}^{l'} - \alpha_{R_j, f_i}^u) \bar{\otimes} 0$ ;
30:         $\beta_{ST, R_j^p}^{u'} = \max\{(\beta_{ST, R_j^p}^{u'} - \alpha_{R_j, f_i}^l) \bar{\otimes} 0, 0\}$ ;
31:      end if
32:    end if
33:     $\beta_{f_i}^l = \beta_{f_i}^l \otimes \beta_{R_j, f_i}^l \otimes \beta_{LT}^l$ ;
34:     $\beta_{f_i}^u = \beta_{f_i}^u \otimes \beta_{R_j, f_i}^u \otimes \beta_{LT}^u$ ;
35:  end for
36: end for

```

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algorithm is much tighter than that of DNC-based delay analysis algorithm  
 465 proposed in [8].

#### 4.5. Buffer Sizing

The priority-aware wormhole-switched NoC [6, 11, 21] requires the same  
 amount of VCs as the priorities to prevent priority inversion, which refers to  
 the blocking of high-priority flows when the low priority flows occupy all the  
 470 VCs [10]. To reduce the buffer area and power consumption of priority-aware  
 wormhole-switched NoC, priority sharing [28] and buffer optimization [12] tech-  
 niques have been proposed. However, the backlog bound of each router derived  
 in [12] is the minimum buffer size that does not trigger the flow control. Reduc-  
 ing this buffer size further will cause the back-pressure between adjacent routers  
 475 and lead to a larger end-to-end delay. However, it is allowed to do so as long  
 as the deadline constraint of each flow is not violated. Thus, our buffer sizing  
 algorithm tries to reduce the initial buffer size of each router iteratively until  
 the end-to-end delay of a flow violates its deadline or the buffer size is equal to  
 one. The initial buffer size to avoid flow control can be obtained by applying  
 480 the following theorem.

**Theorem 3.** Denote by  $\beta_{R_j, f_i}^l$  the feed-forward lower service curve obtained at  
 $R_j \in \mathcal{R}_i$  by flow  $f_i$ ,  $\beta_{LT}^l$  the lower service curve provided by the physical link  
 between two adjacent routers,  $B_{R_j, f_i}$  the VC buffer size reserved for  $f_i$  at  $R_j$   
 and  $\sigma$  the credit feedback delay. Then, for  $\forall R_j \in \mathcal{R}_i$ , the buffer size

$$B_{R_j, f_i} = \lceil \inf \{ B | \beta_{R_j, f_i}^l \otimes \beta_{LT}^l \otimes \overline{\beta_{R_j, f_i}^l \otimes \beta_{LT}^l \otimes \delta_\sigma(t) + B} \geq \beta_{R_j, f_i}^l \otimes \beta_{LT}^l \} \rceil$$

is large enough to avoid flow control.

PROOF. We take the flow  $f_2$  in Fig. 5 as an example to verify this theorem. As  
 stated by Theorem 2, we have

$$\beta_{FC, R_9, f_2}^l = \overline{\beta_{LT}^l \otimes \beta_{R_5, f_2}^l \otimes \delta_\sigma + B_{R_5, f_2}}$$

and

$$\begin{aligned} \beta_{FC, R_{13}, f_2}^l &= \overline{\overline{\beta_{LT}^l \otimes \beta_{R_9, f_2}^l \otimes \overline{\beta_{LT}^l \otimes \beta_{R_5, f_2}^l \otimes \delta_\sigma + B_{R_5, f_2}} \otimes \delta_\sigma + B_{R_9, f_2}} \otimes \delta_\sigma + B_{R_9, f_2}} \\ &= \overline{(\beta_{LT}^l \otimes \beta_{R_9, f_2}^l \otimes \delta_\sigma + B_{R_9, f_2}) \otimes \overline{\beta_{LT}^l \otimes \beta_{R_5, f_2}^l \otimes \delta_\sigma + B_{R_5, f_2}}} \\ &= \overline{\beta_{LT}^l \otimes \beta_{R_9, f_2}^l \otimes \delta_\sigma + B_{R_9, f_2} \otimes \overline{\beta_{LT}^l \otimes \beta_{R_5, f_2}^l \otimes \delta_\sigma + B_{R_5, f_2}}} \end{aligned}$$

where the step from the first line to the second line holds due to the basic  
 property of min-plus convolution (see Rule 7 of Theorem 3.1.5 in [17]). By  
 485 applying Theorem 3.1.11 in [17], the last line holds.

Similarly, we can prove that

$$\beta_{FC, R_{14}, f_2}^l = \overline{\beta_{LT}^l \otimes \beta_{R_{13}, f_2}^l \otimes \delta_\sigma + B_{R_{13}, f_2} \otimes \beta_{FC, R_{13}, f_2}^l},$$

$$\begin{aligned}\beta_{FC,R_{15},f_2}^l &= \overline{\beta_{LT}^l \otimes \beta_{R_{14},f_2}^l \otimes \delta_\sigma + B_{R_{13},f_2}} \otimes \beta_{FC,R_{14},f_2}^l, \\ \beta_{FC,N_{I},f_2}^l &= \overline{\beta_{LT}^l \otimes \beta_{R_{15},f_2}^l \otimes \delta_\sigma + B_{R_{15},f_2}} \otimes \beta_{FC,R_{15},f_2}^l.\end{aligned}$$

Then, the equivalent feedback service curve obtained by  $f_2$  is

$$\begin{aligned}\beta_{f_2}^l &= \beta_{N_{I},f_2}^l \otimes \beta_{FC,N_{I},f_2}^l \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} (\beta_{R_j,f_2}^l \otimes \beta_{FC,R_j,f_2}^l \otimes \beta_{LT}^l) \right) \\ &= \beta_{N_{I},f_2}^l \otimes (\beta_{LT}^l \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} \beta_{R_j,f_2}^l \right)) \otimes (\beta_{FC,N_{I},f_2}^l \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} \beta_{FC,R_j,f_2}^l \right)) \\ &= \beta_{N_{I},f_2}^l \otimes (\beta_{LT}^l \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} \beta_{R_j,f_2}^l \right)) \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} \overline{\beta_{R_j,f_2}^l \otimes \beta_{LT}^l \otimes \delta_\sigma + B_{R_j,f_2}} \right) \\ &= \beta_{N_{I},f_2}^l \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} (\beta_{R_j,f_2}^l \otimes \beta_{LT}^l \otimes \overline{\beta_{LT}^l \otimes \beta_{R_j,f_2}^l \otimes \delta_\sigma + B_{R_j,f_2}}) \right)\end{aligned}$$

where the second line and the fourth line hold due to the commutativity of min-plus convolution, and the step from the second line to the third line holds due to the basic property of sub-additive closure (see Corollary 3.1.1 in [17]).

490 According to the isotonicity of min-plus convolution (see Theorem 3.1.7 in [17]), we know that

$$\beta_{R_j,f_2}^l \otimes \beta_{LT}^l \otimes \overline{\beta_{R_j,f_2}^l \otimes \beta_{LT}^l \otimes \delta_\sigma(t) + B_{R_j,f_2}} \geq \beta_{R_j,f_2}^l \otimes \beta_{LT}^l, \forall R_j \in \mathcal{R}_2 \quad (6)$$

is a sufficient condition for

$$\beta_{f_2}^l \geq \beta_{N_{I},f_2}^l \otimes \left( \bigotimes_{R_j \in \mathcal{R}_2} (\beta_{R_j,f_2}^l \otimes \beta_{LT}^l) \right)$$

where the right side is the equivalent end-to-end feed-forward service curve provided to  $f_2$ . Thus, to avoid flow control, the buffer size at each router reserved for flow  $f_2$  can be assigned to the minimum integer value to make  
495 Eq.(6) hold, which verifies the correctness of this theorem. ■

Suppose that the applications have been mapped onto the NoC, and each flow  $f_i$  has been assigned to their corresponding priority  $P_i$  and deadline  $D_i$ . Following the same notation as Algorithm 2, we propose the buffer sizing algorithm to allocate just enough buffer for each flow to meet their deadline constraint, as  
500 shown in Algorithm 3. We assume that the service curve provided by the source NI of  $f_i$ , i.e.  $\langle \beta_{N_{I},f_i}^l, \beta_{N_{I},f_i}^u \rangle$ , has been computed with Algorithm 1.

Our algorithm tries to reduce the buffer size for each flow from high-priority to low-priority iteratively. For each iteration, it performs the following four steps: (1) Calculate the equivalent feed-forward service curves provided by the  
505 routers (lines 3-5). (2) Calculate the initial buffer size to avoid flow control for each router (line 6). (3) Reduce the initial buffer size gradually as long as the constraint of deadline is not being violated (line 9). This step is abstracted as a procedure, i.e. `Buffer_Reduction( $f_i$ )`. This procedure tries to reduce the buffer

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**Algorithm 3** Buffer sizing algorithm
 

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**Require:** Architecture parameters and flow specifications

**Ensure:** Buffer required for each flow to meet their deadline

```

1: for each flow  $f_i \in \mathcal{F}$  with priority order do
2:   for each router  $R_j \in \mathcal{R}_i$  do
3:      $\beta_{FC,R_j,f_i}^l = \delta_0(t)$ ;  $\beta_{FC,R_j,f_i}^u = \delta_0(t)$ ;
4:      $\beta_{R_j,f_i}^l = \beta_{BW}^l \otimes \beta_{RC}^l \otimes \beta_{VA}^l \otimes \beta_{SA}^l \otimes \lfloor \frac{\beta_{ST,R_j}^{l'}}{|\Theta_{R_j,f_i}|+1} \rfloor \otimes \beta_{FC,R_j,f_i}^l$ ;
5:      $\beta_{R_j,f_i}^u = \beta_{BW}^u \otimes \beta_{RC}^u \otimes \beta_{VA}^u \otimes \beta_{SA}^u \otimes \lceil \frac{\beta_{ST,R_j}^{u'}}{|\Theta_{R_j,f_i}|+1} \rceil \otimes \beta_{FC,R_j,f_i}^u$ ;
6:     compute the initial buffer size  $B_{R_j,f_i}$  according to Theorem 3;
7:   end for
8:    $\beta_{FC,NI,f_i}^l = \delta_0(t)$ ;  $\beta_{FC,NI,f_i}^u = \delta_0(t)$ ;
9:   Buffer_Reduction( $\alpha_{f_i}^u, \mathcal{R}_i, D_i$ );  $\triangleright$  Call the Buffer_Reduction procedure
10:   $\beta_{f_i}^l = \beta_{NI,f_i}^l \otimes \beta_{FC,NI,f_i}^l$ ;  $\beta_{f_i}^u = \beta_{NI,f_i}^u \otimes \beta_{FC,NI,f_i}^u$ ;
11:  for  $\forall R_j \in \mathcal{R}_i$  from  $start_i$  to  $end_i$  do
12:    if  $\Omega_{R_j,f_i} \neq \emptyset$  then
13:       $\beta^l = \beta_{f_i}^l \otimes \beta_{BW}^l \otimes \beta_{RC}^l \otimes \beta_{VA}^l \otimes \beta_{SA}^l$ ;
14:       $\beta^u = \beta_{f_i}^u \otimes \beta_{BW}^u \otimes \beta_{RC}^u \otimes \beta_{VA}^u \otimes \beta_{SA}^u$ ;
15:       $\alpha_{R_j,f_i}^l = \min\{(\alpha_{f_i}^l \otimes \beta^u) \otimes \beta^l, \beta^l\}$ ;
16:       $\alpha_{R_j,f_i}^u = \min\{(\alpha_{f_i}^u \otimes \beta^u) \otimes \beta^l, \beta^u\}$ ;
17:      if  $\forall f_k \in \Theta_{R_j,f_i}$ ,  $Delay(f_k)$  has been calculated then
18:         $\alpha_{R_j,f_i}^l = \alpha_{R_j,f_i}^l + \sum_{f_k \in \Theta_{R_j,f_i}} \alpha_{R_j,f_k}^l$ ;
19:         $\alpha_{R_j,f_i}^u = \alpha_{R_j,f_i}^u + \sum_{f_k \in \Theta_{R_j,f_i}} \alpha_{R_j,f_k}^u$ ;
20:         $\beta_{ST,R_j}^{l'} = (\beta_{ST,R_j}^{l'} - \alpha_{R_j,f_i}^u) \bar{\otimes} 0$ ;
21:         $\beta_{ST,R_j}^{u'} = \max\{(\beta_{ST,R_j}^{u'} - \alpha_{R_j,f_i}^l) \bar{\otimes} 0, 0\}$ ;
22:      end if
23:    end if
24:     $\beta_{f_i}^l = \beta_{f_i}^l \otimes \beta_{R_j,f_i}^l \otimes \beta_{LT}^l$ ;  $\beta_{f_i}^u = \beta_{f_i}^u \otimes \beta_{R_j,f_i}^u \otimes \beta_{LT}^u$ ;
25:  end for
26: end for

```

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size of each router reserved for the target flow  $f_i$  from its ejection router to its injection router. For each router on the path of  $f_i$ , it decreases the buffer size iteratively, until the buffer size is reduced to one or the deadline constraint is violated. If the iteration stops due to the violation of deadline, the buffer size and related service curves are restored to their previous value. To ease the description, we define the operator  $Pre(R_j)$  to identify the predecessor of router  $R_j$ , and let  $Pre(start_i) = NI$ . (4) Calculate the leftover service curve at each router for low-priority flows (lines 10-25). The computation complexity of the Buffer\_Reduction( $f_i$ ) procedure and the entire buffer sizing algorithm are  $O(H^2)$  and  $O(NH^2)$ , where  $N$  and  $H$  denote the number of flows analyzed and

the maximum hop count of these flows. This algorithm can be implemented in RTC toolbox [20] to optimize the buffer size automatically.

---

```

1: procedure BUFFER_REDUCTION( $\alpha_{f_i}^u, \mathcal{R}_i, D_i$ )
2:   for each router  $R_j \in \mathcal{R}_i$  from  $end_i$  to  $start_i$  do
3:      $Delay(f_i) = H(\alpha_{f_i}^u, \beta_{NI, f_i}^l \otimes \beta_{FC, NI, f_i}^l \otimes (\bigotimes_{R_k \in \mathcal{R}_i} (\beta_{R_k, f_i}^l \otimes \beta_{LT}^l)))$ ;
4:     while  $Delay(f_i) \leq D_i$  and  $B_{R_j, f_i} > 1$  do
5:        $B_{R_j, f_i} = B_{R_j, f_i} - 1$ ;
6:       for all  $R_k \in \mathcal{R}_i$  from  $Pre(R_j)$  to  $start_i$  do
7:         compute  $\langle \beta_{FC, R_k, f_i}^l, \beta_{FC, R_k, f_i}^u \rangle$  according to Theorem 2;
8:         compute the service curve  $\langle \beta_{R_k, f_i}^l, \beta_{R_k, f_i}^u \rangle$ ;
9:       end for
10:      compute the service curve  $\langle \beta_{FC, NI, f_i}^l, \beta_{FC, NI, f_i}^u \rangle$ ;
11:       $Delay(f_i) = H(\alpha_{f_i}^u, \beta_{NI, f_i}^l \otimes \beta_{FC, NI, f_i}^l \otimes (\bigotimes_{R_k \in \mathcal{R}_i} (\beta_{R_k, f_i}^l \otimes \beta_{LT}^l)))$ ;
12:    end while
13:    if  $Delay(f_i) > D_i$  then
14:       $B_{R_j, f_i} = B_{R_j, f_i} + 1$ ;
15:      recompute  $\langle \beta_{FC, Pre(R_j), f_i}^l, \beta_{FC, Pre(R_j), f_i}^u \rangle$ ;
16:      recompute  $\langle \beta_{Pre(R_j), f_i}^l, \beta_{Pre(R_j), f_i}^u \rangle$  if  $R_j \neq start_i$ ;
17:    end if
18:  end for
19: end procedure

```

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The number of cycles that a packet can be delayed in the network without violating its deadline is referred to ‘slack’ in [29]. Thus, the Buffer\_Reduction procedure of our buffer sizing algorithm is the process of slack minimization. However, while used for power reduction, our buffer sizing algorithm is significantly different from the DNC-based slack optimization algorithm proposed in [29]. In [29], the energy optimization is achieved by adjusting the voltage, frequency and link bandwidth of on-chip routers for the fixed configuration and deadline. In contrast, our method tries to reduce the buffer size under the deadline constraint, and the buffer reduction directly leads to the area and power saving. In addition, our algorithm can be also used in conjunction with the priority sharing techniques [28] to minimize the hardware cost of priority-aware wormhole-switched NoC.

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## 5. Experiments

In this section, we validate the correctness and tightness of our performance model by comparing with simulation results and other analytical methods. Several analytical methods exist for the delay analysis of priority-aware NoC, examples include contention tree model [9], lumped link model [10], dependency graph model [11], FLA [6], LLA [7] and DNC [8], etc. There is also extensive

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research on the buffer sizing problem of the priority-aware NoC, representative  
 540 methods include shaping delay analysis [30] and LLBA [12]. Among all these  
 analytical methods, DNC [8] and LLBA [12] based models outperform the oth-  
 ers when the tightness of delay bound and buffer requirement is considered.  
 Thus, we will only perform the comparison with LLBA and DNC, as presented  
 in subsection 5.1 and subsection 5.2 respectively. We also present the simulation  
 545 results to validate the correctness and tightness of our method in subsection 5.3.

### 5.1. Comparison with Link Level Buffer-space Analysis

The traffic pattern discussed in this subsection is shown in Fig. 1. There  
 are four flows (i.e.  $f_1, f_2, f_3$  and  $f_4$ ) in the network, with different priorities  
 $P_4 > P_1 > P_2 > P_3$ . We perform the comparison on a set of periodical traffic  
 550 due to the restriction of LLBA method [12]. The packet length (in flits) and  
 injection period (in cycles) of flow  $f_i$  ( $i = 1, 2, 3, 4$ ) are denoted as  $F_i$  and  $I_i$ ,  
 respectively. To ease the analysis of buffer requirement, the LLBA method  
 assumes the number of bits in a flit is the same as the physical channel width,  
 and the latency of a router is one cycle [12]. Thus, our performance model for  
 555 the standard wormhole-switched router should be specialized, which is achieved  
 by letting the service curve of BW, RC, VA, SA and LT stage be a pair of burst  
 delay function  $\langle \delta_0(t), \delta_0(t) \rangle$ . Under this condition, the service curve of the  
 entire router is equal to the service curve provided by the ST stage, which is  
 $\langle \beta_{ST, R_i^p}^l, \beta_{ST, R_i^p}^u \rangle$ . The traffic jitter for all the flows are assumed to be zero  
 560 for brevity and clarity. In addition, we set the credit feedback delay  $\sigma = 0$  cycle  
 in our model since the LLBA method does not consider the influence of credit  
 feedback delay on the buffer requirement.

Let us suppose that all the flows have the same packet injection period  
 $I_i = 50$  cycles and packet length  $F_i = 8$  flits ( $i = 1, 2, 3, 4$ ). The delay bounds  
 565 computed with the LLA method [7] for these four flows are 21, 22, 21 and 13  
 cycles, respectively. We can also obtain the total buffer size reserved for the  
 four flows (i.e.  $f_1, f_2, f_3$  and  $f_4$ ) with LLBA method [12], which are 11, 12, 18  
 and 4 flits, respectively. The total buffer size required by these four flows can  
 be obtained by summing up the buffer size reserved for each flow along its path,  
 570 which is 45 flits. This value is the minimum buffer space required for all these  
 four flows to avoid the back-pressure caused by flow control. By allowing the  
 flow control to be triggered, our buffer sizing algorithm can be utilized to reduce  
 this buffer size further as long as the deadline constraint is not violated. Figure 6  
 demonstrates the normalized buffer requirement calculated with LLBA and our  
 575 method under different deadline constraint. It can be found that our method  
 can give much tighter buffer estimation than LLBA, and the improvement of  
 our method over LLBA becomes more and more significant as the deadline  
 constraint is relaxed.

### 5.2. Comparison with Network Calculus based Delay Analysis

580 In this subsection, we present the numerical results to demonstrate the im-  
 provement of our method over the DNC method proposed in [8]. The traffic

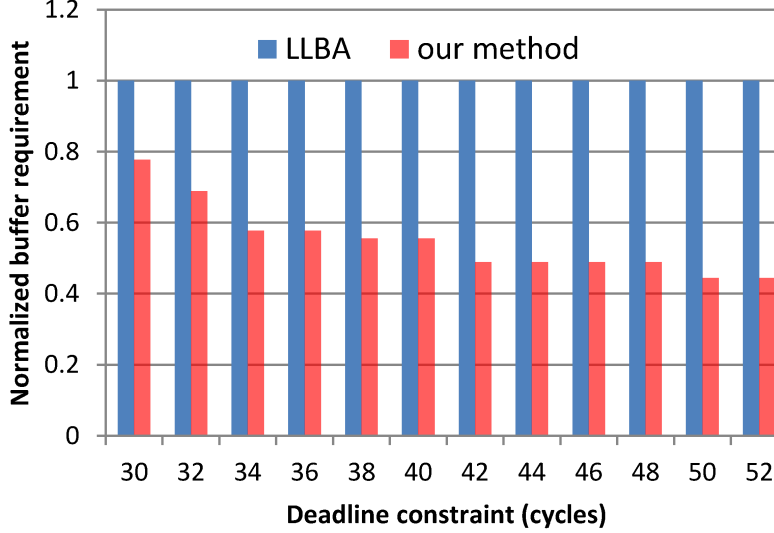


Figure 6: Buffer requirement computed with our method for different deadline

pattern discussed in this subsection is shown in Fig. 1. The priorities of these four flows (i.e.  $f_1, f_2, f_3$  and  $f_4$ ) in the network satisfy  $P_4 > P_1 > P_2 > P_3$ . We perform the comparison on a set of periodical traffic. The packet length (in flits) and injection period (in cycles) of flow  $f_i$  ( $i = 1, 2, 3, 4$ ) are denoted as  $F_i$  and  $I_i$ , respectively. The router architecture we considered in this subsection adopts the lookahead pipeline [23], which removes the RC stage from the critical path of the router pipeline. Thus, our service model is customized by letting the service curve of RC stage to be  $\langle \delta_0(t), \delta_0(t) \rangle$ . The RTC arrival curve  $\langle \alpha_{f_i}^l, \alpha_{f_i}^u \rangle$  of flow  $f_i$  can be obtained according to the method introduced in subsection 4.1. The DNC arrival curve for  $f_i$  is  $\alpha_{f_i} = V_i t + F_i$ , where  $V_i = F_i/I_i$  represents the average arrival rate.

We assume the VC buffer size of each router is 32 flits, and the credit feedback delay  $\sigma = 0$  cycle. We change the injection rate  $V_i$  ( $i = 1, 2, 3, 4$ ) from  $1/3$  to  $1/6$  (flits/cycle) and the packet length  $F_i$  from 1 to 8 flits. The end-to-end delay of flow  $f_3$  calculated with the DNC method and our method are presented in Fig. 7. By comparison, we find that our method can give a much tighter delay bound than the DNC-based method proposed in [8]. The root cause for this improvement lies in the fact that our method utilizes the upper service curve to limit the output upper arrival curve, which further leads to a tighter leftover service curve for the low-priority flows.

### 5.3. Comparison with Simulation

Our performance model is also verified by simulation. We modified the Booksim 2.0 [31], a cycle-accurate NoC simulator, to support the specified traffic

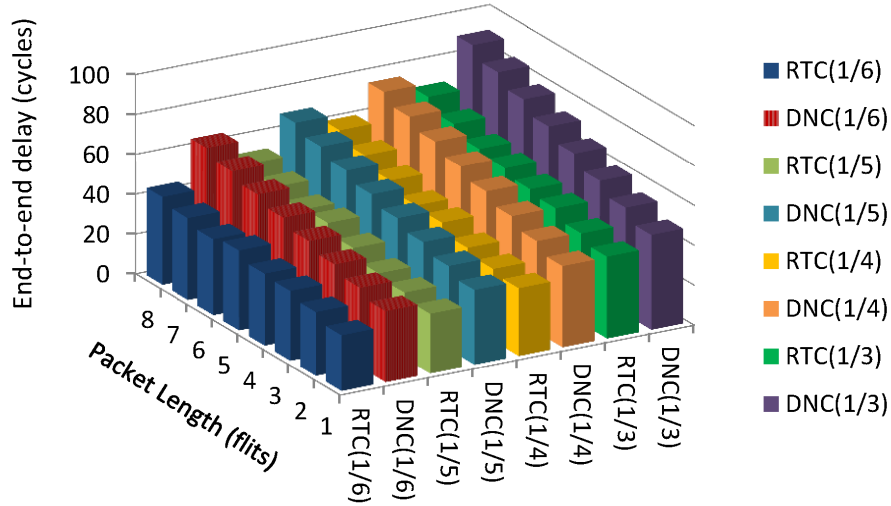


Figure 7: Delay bound comparison with network calculus under different injection rate and packet length

pattern and injection process. The traffic patterns investigated in this subsection include the example shown in Fig. 1 and an real application provided by Ericsson Radio Systems [32][33]. We adopt the optimized lookahead router [23] to construct the mesh network. To fit this optimization, our service model is customized by letting the service curve of RC stage be  $\langle \delta_0(t), \delta_0(t) \rangle$ . Other architecture and simulation parameters used in the simulation are listed in the Table 1.

Table 1: Architecture parameters used in the simulation

network topology	$4 \times 4$ mesh	routing algorithm	X-Y routing
credit delay	0 cycle	channel width	128 bits
buffer size	32 flits	switch allocator	priority-based
link latency	1 cycle	sampling period	$1 \times 10^5$ cycles
clock cycle	1 ns	warmup period	$3 \times 10^5$ cycles

We first investigate the traffic pattern shown in Fig. 1. Suppose that the injection process of each flow is the strictly periodical traffic without any jitter, and the priorities of these four flows satisfy  $P_4 > P_1 > P_2 > P_3$ . However, even under this assumption, determining the worst-case delay by simulation for general scenarios is still a non-trivial task, because the worst-case delay of each flow depends on the traffic pattern and offset (i.e. the injection time of the first packet in a flow  $f_i$ , denoted as  $O_i$ ) of all the flows in the network.

We set the injection rate  $V_i$  ( $i = 1, 2, 3, 4$ ) to  $1/6$  (flits/cycle), and change the packet length  $F_i$  from 2 flits to 9 flits. Under these conditions, we find



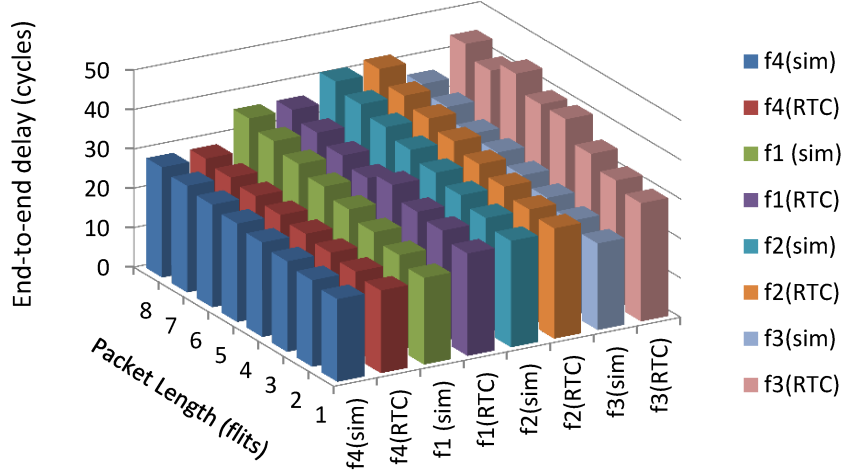
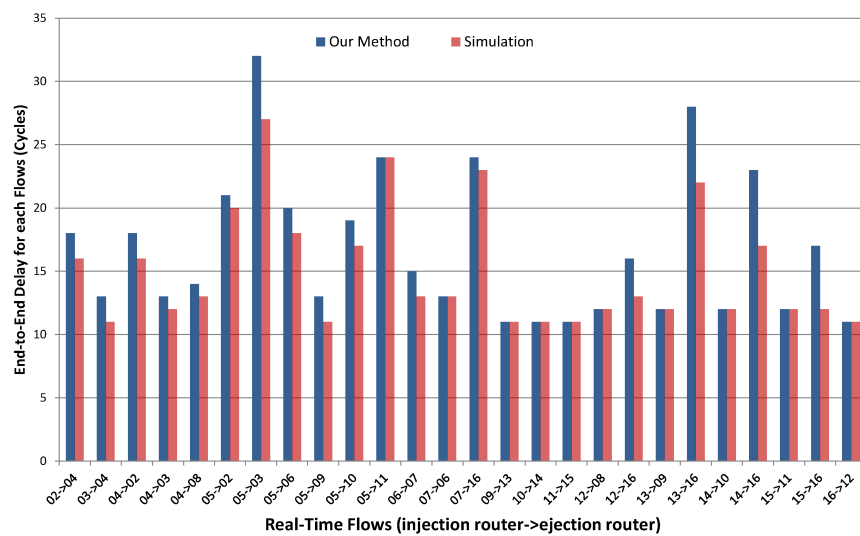
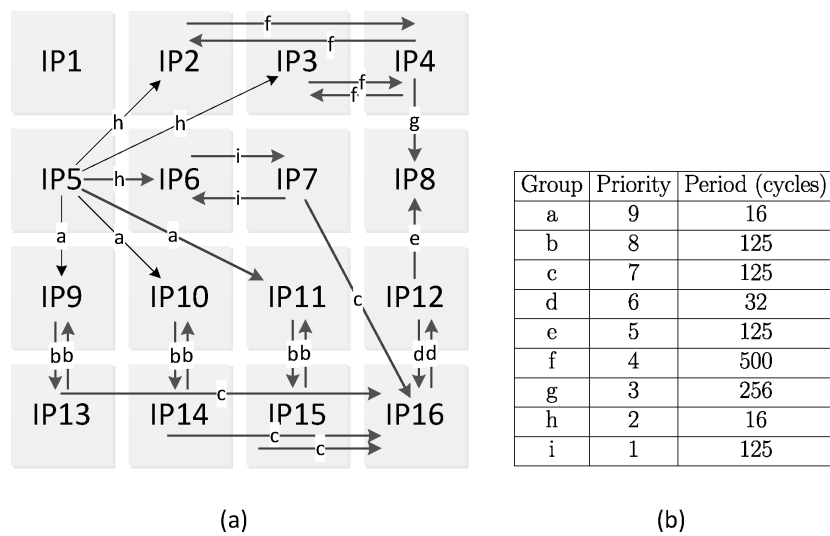


Figure 8: Delay comparison with simulation under different packet length

that  $f_1$  will experience the maximum delay in the network if its offset  $O_1$  is equal to that of  $f_4$ . Similarly, the worst-case interference of  $f_2$  occurs when its offset  $O_2 = O_1 + (4 + 1)$ , where  $O_1 + (4 + 1)$  is the time instance that the first packet of  $f_1$  leaves router  $R_{16}$ , since the basic latency of a lookahead router and physical channel are 4 cycles and 1 cycles, respectively. Flow  $f_3$  experiences the worst-case delay when  $O_3 = O_2 + F_1 + 2 \times (4 + 1)$ . Thus, we set  $O_1 = O_4 = 0$  cycle,  $O_2 = 5$  cycles,  $O_3 = 15 + F_1$  cycles, and run the simulation. The collected maximum end-to-end delay of these four flows under the given offset combination are compared with our RTC method, as shown in Fig. 8. As indicated in this figure, for the given configurations, the delay bounds calculated with our method are indeed the upper bounds of simulation results, which verifies the correctness of our method.

We also take a real application discussed in [32][33] as an example to demonstrate the tightness and ability of our method to analyze complex scenarios. This application is comprised of 16 IP cores. The 26 communication flows among these 16 IPs are classified into nine groups, and each group has their bandwidth requirement. When mapped to a  $4 \times 4$  mesh network, the traffic pattern of this application is demonstrated in Fig. 9(a). We assume that all the flows in a group send packets periodically with the same injection period and priority, as listed in Table. 9(b). We set the packet size to 128 bits, and collect the maximum end-to-end delay of each flow obtained by simulation. The comparison results between one simulation run and the delay bound calculated with our method are shown in Fig. 10. The offset of each flow in this simulation run is zero cycle. We can see that the calculated delay bounds constrain the simulation results well, which verifies the correctness of our method. This comparison also demonstrates the ability of our method to analyze the real system with large number of flows.



## 6. Conclusion

The priority-aware wormhole-switched NoC is a promising platform for the on-chip real-time communication if the worst-case performance can be accurately analyzed and guaranteed. Simulation is not well suited for this purpose because it is difficult to cover all the corner cases. In this paper, we proposed an RTC-based performance model to achieve this goal. We first built the traffic model and service model for the priority-aware NoC, and then proposed a novel method to derive the upper service curve of credit-based flow control. Based on the proposed RTC model, we then proposed an end-to-end delay analysis algorithm and a buffer sizing algorithm. The delay analysis algorithm can be implemented to compute the end-to-end delay for each flow automatically, and verify whether all these flows meet their deadline under the given configuration. Compared with the DNC-based performance model, our model can give tighter delay bound since it takes the upper service curve and lower arrival curve into consideration. The proposed buffer sizing algorithm can optimize the buffer size from high-priority flows to low-priority flows. It can also be implemented to perform the buffer reduction automatically under the constraint of deadline. Compared with the LLBA method, our method can give much tighter buffer bound. Experimental results also illustrate that our method indeed outperforms the other analytical methods, e.g. LLBA and DNC, when the tightness of delay bound and buffer requirement are considered. Our method can be applied to the task mapping, routing and power reduction of priority-aware NoC.

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