

# Theory of Computer Science

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## Exercise Sheet 2

**Due: Wednesday, March 9, 2016**

*Note:* Submissions that are exclusively created with L<sup>A</sup>T<sub>E</sub>X will receive a bonus mark. Please submit only the resulting PDF file (or a printout of this file).

### Exercise 2.1 (Syntax; 0.5+0.5+0.5+0.5 Points)

Formalize the following statements as propositional formulas. In order to do so, also define appropriate atomic propositions. Take care to fully parenthesize all formulas.

- (a) If it does not rain, it is warm.
- (b) If Bob is going for a swim, then he always eats ice cream and it is not raining.
- (c) Bob is going for a swim exactly if he eats ice cream and it is warm or does not rain.
- (d) Either the sun is shining or it is raining (but not both).

### Exercise 2.2 (Truth tables; 1+1+1+1 Points)

Let  $A = \{X, Y\}$  be a set of propositional variables. Specify a propositional formula  $\varphi$  over  $A$  for each of the following properties and then use a truth table to prove that the formula has the property.

- (a)  $\varphi$  is satisfiable and falsifiable.
- (b)  $\varphi$  has exactly three models.
- (c)  $\varphi$  is valid and uses both variables.
- (d)  $\varphi$  is unsatisfiable.

### Exercise 2.3 (Semantics; 0.5+0.25+1+1+0.5 Points)

Consider the propositional formula  $\varphi$  over  $\{A, B, C, D, E, F\}$ :

$$\varphi = ((F \vee ((\neg B \leftrightarrow ((C \wedge A) \rightarrow \neg B)) \vee (D \rightarrow E))) \rightarrow (A \rightarrow \neg F))$$

- (a) How many lines would be needed for a truth table for  $\varphi$ ?
- (b) Formula  $\varphi$  is an implication. Specify the truth table for the general implication formula  $\varphi \rightarrow \psi$  (see chapter B1, slide 30). Attention: You should **not** specify the truth table of  $\varphi$ .
- (c) Specify a model  $\mathcal{I}$  for  $\varphi$  and prove without truth table that  $\mathcal{I} \models \varphi$ .
- (d) Specify an assignment  $\mathcal{I}$  with  $\mathcal{I} \not\models \varphi$  and prove that  $\mathcal{I}$  has the desired property without a truth table.
- (e) Which of the properties *satisfiable*, *unsatisfiable*, *valid*, and *falsifiable* are true for  $\varphi$ ? Justify your answer for each of the four properties.

*Hint:* The proofs for this exercises are fairly short (4 and 6 steps, respectively). If you need a considerably larger amount of steps, rethink your solution and try to find an easier proof. The solution of part (b) may help you identify the requirements for  $\mathcal{I}$ .