Symbol Tables Binary Search Trees

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Your Lab Class

数据结构与算法分析B Data Structures and Algorithm Analysis

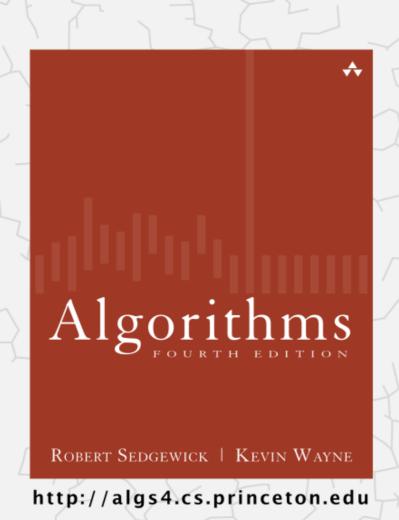
Lecture 10

- > Symbol Tables (3.1 of Text A)
- Binary Search Trees (3.2 of Text A)

To be discussed in Lecture 11:

- ➤ Balanced Search Trees (3.3 of Text A)
- > Hash Tables (3.4 of Text A)

Algorithms



3.1 SYMBOL TABLES

- API
- elementary implementations
- ordered operations

Algorithms

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Symbol tables

Key-value pair abstraction.

- Insert a value with specified key.
- Given a key, search for the corresponding value.

Ex. DNS lookup.

Insert domain name with specified IP address.

key

Given domain name, find corresponding IP address.

domain name	IP address
www.cs.princeton.edu	128.112.136.11
www.princeton.edu	128.112.128.15
www.yale.edu	130.132.143.21
www.harvard.edu	128.103.060.55
www.simpsons.com	209.052.165.60
↑	↑

Symbol table applications

application	purpose of search	key	value			
dictionary	find definition	word	definition			
book index	find relevant pages	term	list of page numbers			
file share	find song to download	name of song	computer ID			
financial account	process transactions	account number	transaction details			
web search	find relevant web pages	keyword	list of page names			
compiler	find properties of variables	variable name	type and value			
routing table	route Internet packets	destination	best route			
DNS	find IP address	domain name	IP address			
reverse DNS	find domain name	IP address	domain name			
genomics	find markers	DNA string	known positions			
file system	find file on disk	filename	location on disk			

Symbol tables: context

Also known as: maps, dictionaries, associative arrays.

Generalizes arrays. Keys need not be between 0 and N-1.

Language support.

- External libraries: C, VisualBasic, Standard ML, bash, ...
- Built-in libraries: Java, C#, C++, Scala, ...
- Built-in to language: Awk, Perl, PHP, Tcl, JavaScript, Python, Ruby, Lua.

every array is an every object is an table is the only associative array associative array primitive data structure

hasNiceSyntaxForAssociativeArrays["Python"] = true
hasNiceSyntaxForAssociativeArrays["Java"] = false

legal Python code

Basic symbol table API

Associative array abstraction. Associate one value with each key.

public class ST <key, value=""></key,>								
	ST()	create an empty symbol table						
void	put(Key key, Value val)	put key-value pair into the table ←	_ a[key] = val;					
Value	get(Key key)	value paired with key ←	_ a[key]					
boolean	contains(Key key)	is there a value paired with key?						
void	delete(Key key)	remove key (and its value) from table						
boolean	isEmpty()	is the table empty?						
int	size()	number of key-value pairs in the table						
Iterable <key></key>	keys()	all the keys in the table						

Conventions

- Values are not null. ← Java allows null value
- Method get() returns null if key not present.
- Method put() overwrites old value with new value.

Intended consequences.

Easy to implement contains().

```
public boolean contains(Key key)
{ return get(key) != null; }
```

Can implement lazy version of delete().

```
public void delete(Key key)
{  put(key, null); }
```

Keys and values

Value type. Any generic type.

specify Comparable in API.

Key type: several natural assumptions.

- Assume keys are Comparable, use compareTo().
- Assume keys are any generic type, use equals() to test equality.
- Assume keys are any generic type, use equals() to test equality;
 use hashCode() to scramble key.

built-in to Java (stay tuned)

Best practices. Use immutable types for symbol table keys.

- Immutable in Java: Integer, Double, String, java.io.File, ...
- Mutable in Java: StringBuilder, java.net.URL, arrays, ...

Equality test

All Java classes inherit a method equals().

Java requirements. For any references x, y and z:

- Reflexive: x.equals(x) is true.
- Symmetric: x.equals(y) iff y.equals(x).
- Transitive: if x.equals(y) and y.equals(z), then x.equals(z).
- Non-null: x.equals(null) is false.

do x and y refer to the same object?

Default implementation. (x == y)

Customized implementations. Integer, Double, String, java.io.File, ...

User-defined implementations. Some care needed.

equivalence relation

Implementing equals for user-defined types

Seems easy.

```
public
            class Date implements Comparable<Date> {
   private final int month;
   private final int day;
   private final int year;
   public boolean equals (Date that) {
                                                           check that all significant
      if (this.day != that.day ) return false;
                                                           fields are the same
      if (this.month != that.month) return false;
      if (this.year != that.year ) return false;
      return true;
```

Implementing equals for user-defined types

typically unsafe to use equals() with inheritance Seems easy, but requires some care. (would violate symmetry) public final class Date implements Comparable<Date> { private final int month; must be Object. private final int day; Why? Experts still debate. private final int year; public boolean equals (Object y) { optimize for true object equality if (y == this) return true; check for null if (y == null) return false; objects must be in the same class if (y.getClass() != this.getClass()) (religion: getClass()vs. instanceof) return false; Date that = (Date) y; cast is guaranteed to succeed if (this.day != that.day) return false; check that all significant if (this.month != that.month) return false; fields are the same if (this.year != that.year) return false; return true;

Equals design

"Standard" recipe for user-defined types.

- Optimization for reference equality.
- Check against null.
- Check that two objects are of the same type and cast.
- Compare each significant field:
 - if field is a primitive type, use ==
 if field is an object, use equals()
 apply rule recursively
 - if field is an array, apply to each entry ← can use Arrays.deepEquals(a, b)
 but not a.equals(b)

Best practices.

- No need to use calculated fields that depend on other fields.
- Compare fields mostly likely to differ first.
- Make compareTo() consistent with equals().

```
x.equals(y) if and only if (x.compareTo(y) == 0)
```

e.g., cached Manhattan distance

ST test client for traces

Build ST by associating value i with ith string from standard input.

```
public static void main(String[] args) {
   ST<String, Integer> st = new ST<String, Integer>();

for (int i = 0; !StdIn.isEmpty(); i++) {
   String key = StdIn.readString();
   st.put(key, i);
   }
   for (String s : st.keys())
      StdOut.println( s + " " + st.get(s));
}
```

```
keys S E A R C H E X A M P L E values 0 1 2 3 4 5 6 7 8 9 10 11 12
```

output

```
A 8
C 4
E 12
H 5
L 11
M 9
P 10
R 3
S 0
X 7
```

ST test client for analysis

Frequency counter. Read a sequence of strings from standard input and print out one that occurs with highest frequency.

```
% more tinyTale.txt
it was the best of times
it was the worst of times
it was the age of wisdom
it was the age of foolishness
it was the epoch of belief
it was the epoch of incredulity
it was the season of light
it was the season of darkness
it was the spring of hope
it was the winter of despair
                                                        tiny example
% java FrequencyCounter 1 < tinyTale.txt</pre>
                                                        (60 words, 20 distinct)
it 10
                                                        real example
% java FrequencyCounter 8 < tale.txt</pre>
                                                        (135,635 words, 10,769 distinct)
business 122
                                                        real example
% java FrequencyCounter 10 < leipzig1M.txt
                                                        (21,191,455 words, 534,580 distinct)
government 24763
```

Frequency counter implementation

```
public class FrequencyCounter {
   public static void main (String[] args) {
      int minlen = Integer.parseInt( args[0]);
                                                                               create ST
      ST<String, Integer> st = new ST<String, Integer>();
      while (!StdIn.isEmpty()) {
         String word = StdIn.readString();
                                                      ignore short strings
         if (word.length() < minlen) continue;</pre>
                                                                               read string and
         if (!st.contains(word)) st.put( word, 1);
                                                                               update frequency
                                   st.put( word, st.get(word) + 1);
         else
      String max = "";
      st.put(max, 0);
                                                                               print a string
      for (String word : st.keys())
                                                                               with max freq
         if (st.get(word) > st.get(max))
            max = word;
      StdOut.println( max + " " + st.get(max));
```

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APH

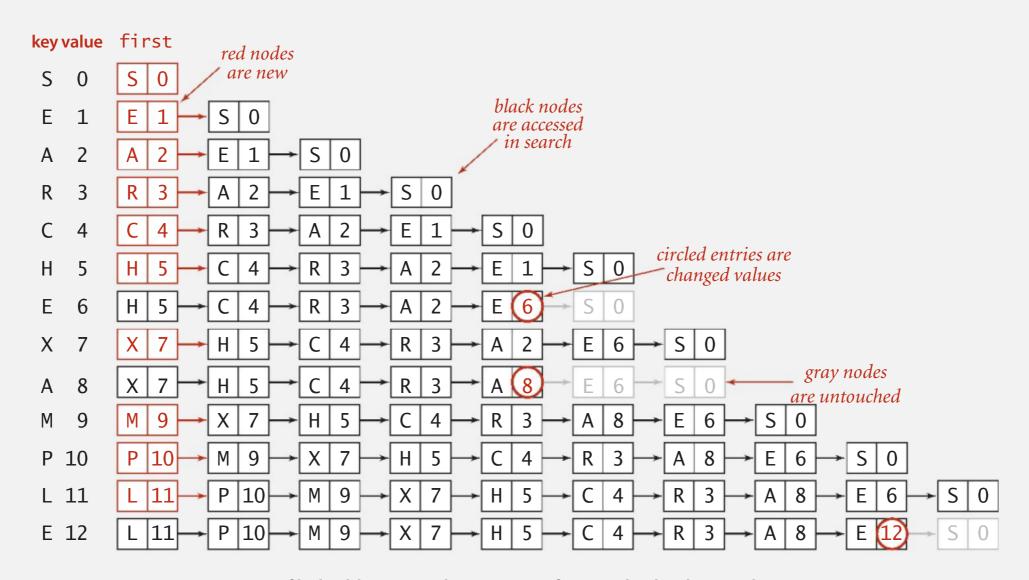
- elementary implementations
- ordered operations

Sequential search in a linked list

Data structure. Maintain an (unordered) linked list of key-value pairs.

Search. Scan through all keys until find a match.

Insert. Scan through all keys until find a match; if no match add to front.



Trace of linked-list ST implementation for standard indexing client

Elementary ST implementations: summary

ST implementation	guara	ntee	avera	key	
31 implementation	search	insert	search hit	insert	interface
sequential search (unordered list)	N	N	N / 2	N	equals()

Challenge. Efficient implementations of both search and insert.

Binary search in an ordered array

Data structure. Maintain an ordered array of key-value pairs.

Rank helper function. How many keys < k?

```
keys[]
                                         4 5 6 7 8
                           ACEHLMPRSX
successful search for P
              lo hi m
                                                                     entries in black
                                                          X
                                                                     are a [lo..hi]
                                                              entry in red is a [m]
                                                      loop exits with keys[m] = P: return 6
       unsuccessful search for Q
unsuccessful search for Q
              lo hi m
                       loop exits with lo > hi: return 7
```

Binary search: Java implementation

```
public Value get (Key key) {
   if (isEmpty()) return null;
   int i = rank(key);
   if (i < N && keys[i].compareTo(key) == 0) return vals[i];
   else return null;
}
private int rank (Key key) {
                                            number of keys < key
   int lo = 0, hi = N-1;
  while (lo <= hi) {</pre>
       int mid = 10 + (hi - 10) / 2;
       int cmp = key.compareTo( keys[mid]);
       if (cmp < 0) hi = mid - 1;
       else if (cmp > 0) lo = mid + 1;
       else if (cmp == 0) return mid;
  return lo;
```

Binary search: trace of standard indexing client

Problem. To insert, need to shift all greater keys over.

						key	's []										va	ls[]				
key	value	0	1	2	3	4	5	6	7	8	9	N	0	1	2	3	4	5	6	7	8	9
S	0	S										1	0									
Ε	1	Ε	S			0	ntrie	c in 1	ed.			2	1	0					tries ved to			<u>.</u>
Α	2	Α	Ε	S			vere i					3	2	1	0		/	1110	veu ii) lile	rigiii	
R	3	Α	Е	R	S							4	2	1	3	0						
C	4	Α	C	Ε	R	S			en	tries	in gra	_{iv} 5	2	4	1	3	0					
Н	5	Α	C	Е	Н	R	S				ot mov		2	4	1	5	3	0		tled e lange		s are
Ε	6	Α	C	Е	Н	R	S					6	2	4	6	5	3	0	CII	unge	u vu	iucs
X	7	Α	C	Е	Н	R	S	X				7	2	4	6	5	3	0	7			
Α	8	Α	C	Е	Н	R	S	X				7	(8)	4	6	5	3	0	7			
M	9	Α	C	Е	Н	M	R	S	X			8	8	4	6	5	9	3	0	7		
Р	10	Α	C	Е	Н	M	P	R	S	X		9	8	4	6	5	9	10	3	0	7	
L	11	Α	C	Е	Н	L	M	Р	R	S	X	10	8	4	6	5	11	9	10	3	0	7
Ε	12	Α	C	Е	Н	L	M	P	R	S	X	10	8	4 (12)	5	11	9	10	3	0	7
		Α	C	Ε	Н	L	M	Р	R	S	Χ		8	4	12	5	11	9	10	3	0	7

Elementary ST implementations: summary

ST implementation	guara	ıntee	avera	key	
31 implementation	search	insert	search hit	insert	interface
sequential search (unordered list)	N	N	N/2	N	equals()
binary search (ordered array)	$\log N$	N	log N	N/2	compareTo()

Challenge. Efficient implementations of both search and insert.

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Examples of ordered symbol table API

```
values
                                  keys
                     min() \longrightarrow 09:00:00
                                           Chicago
                                         Phoenix
                              09:00:03
                              09:00:13 → Houston
            get(09:00:13) 09:00:59
                                          Chicago
                              09:01:10
                                           Houston
          floor(09:05:00) \longrightarrow 09:03:13
                                           Chicago
                                           Seattle
                              09:10:11
                 select(7) \longrightarrow 09:10:25
                                          Seattle
                              09:14:25 Phoenix
                              09:19:32
                                           Chicago
                              09:19:46
                                           Chicago
keys(09:15:00, 09:25:00) \longrightarrow 09:21:05
                                           Chicago
                                           Seattle
                              09:22:43
                              09:22:54
                                          Seattle
                              09:25:52
                                          Chicago
        ceiling(09:30:00) \longrightarrow 09:35:21
                                           Chicago
                                           Seattle
                              09:36:14
                     \max() \longrightarrow 09:37:44
                                          Phoenix
size(09:15:00, 09:25:00) is 5
     rank(09:10:25) is 7
```

Ordered symbol table API

public class	ST <key comparable<key="" extends="">, Value></key>							
Key	min()	smallest key						
Key	max()	largest key						
Key	floor(Key key)	largest key less than or equal to key						
Key	ceiling(Key key)	smallest key greater than or equal to key						
int	rank(Key key)	number of keys less than key						
Key	<pre>select(int k)</pre>	key of rank k						
void	<pre>deleteMin()</pre>	delete smallest key						
void	deleteMax()	delete largest key						
int	size(Key lo, Key hi)	number of keys between lo and hi						
Iterable <key></key>	keys()	all keys, in sorted order						
Iterable <key></key>	keys(Key lo, Key hi)	keys between lo and hi, in sorted order						

Binary search: ordered symbol table operations summary

	sequential search	binary search
search	N	$\log N$
insert / delete	N	N
min / max	N	1
floor / ceiling	N	$\log N$
rank	N	$\log N$
select	N	1
ordered iteration	$N \log N$	N

order of growth of the running time for ordered symbol table operations

Summary

- > Symbol Tables (3.1 of Text A)
- Binary Search Trees (3.2 of Text A)

To be discussed in Lecture 11:

- ➤ Balanced Search Trees (3.3 of Text A)
- > Hash Tables (3.4 of Text A)