# Scientific Computing (M3SC) Project 1

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### 1 MAIN SOLUTION

The code in *solution.py* contains the main program which carries out the process outlined by the question, i.e modelling the process of the cars moving across the city of Rome using the rules described. I have added scripts *to be added* and *to be added* to help answer the related questions at the end of the project.

Below are the imports used by the main program.

```
# Imports
import numpy as np
import csv
import sys
import math as ma

# This import is needed for the last question
from solution_accident_occurs import max_index_tracker_no30
```

the variable  $max\_index\_tracker\_no30$  is imported from another python script,  $solution\_accident\_occurs.py$  in order to answer one of the questions. This will be discussed in detail later.

Below are the functions required by the program. The docstring's explain their use.

```
1 # -
                            FUNCTIONS USED
3
4
5
   def calcWei(RX, RY, RA, RB, RV):
6
7
       This function is taken from Tutorials. It calculates the weight matrix
8
       given information about each node in the system.
9
       :param RX: The x coordinates of each node in the system
10
       :param RY: The y coordinates of each node in the system
11
       :param RA: the connectivity of each node in the system
12
       :param RB: the connectivity of each node in the system
13
       :param RV: the speed limits across each edge in the system
14
       :return: usable weight matrix
15
16
17
       n = len(RX)
18
       wei = np.zeros((n, n), dtype=float)
19
       m = len(RA)
20
       for i in range(m):
           xa = RX[RA[i] - 1]
21
22
            ya = RY[RA[i] - 1]
23
            xb = RX[RB[i] - 1]
24
            yb = RY[RB[i] - 1]
25
            dd = ma.sqrt((xb - xa) ** 2 + (yb - ya) ** 2)
26
            tt = dd / RV[i]
27
            wei[RA[i] - 1, RB[i] - 1] = tt
28
       return wei
29
30
  def Dijkst(ist, isp, wei):
31
32
       This Dijkstra's algorithm implementation is taken from tutorials.
33
34
       :param ist: the index of the starting node
35
       :param isp: the index of the node to reach
       :param wei: the assosciated weight matrix
36
37
       :return:
       .....
38
39
40
       # exception handling (start = stop)
41
       if ist == isp:
42
            shpath = [ist]
43
```

return shpath

```
44
45
       # initialization
46
       N = len(wei)
47
       Inf = sys.maxint
48
       UnVisited = np.ones(N, int)
49
       cost = np.ones(N) * 1.e6
50
       par = -np.ones(N, int) * Inf
51
52
       # set the source point and get its (unvisited) neighbors
53
       jj = ist
       cost[jj] = 0
54
55
       UnVisited[jj] = 0
56
       tmp = UnVisited * wei[jj, :]
57
       ineigh = np.array(tmp.nonzero()).flatten()
       L = np.array(UnVisited.nonzero()).flatten().size
58
59
60
       # start Dijkstra algorithm
61
       while (L != 0):
62
           # step 1: update cost of unvisited neighbors,
63
                      compare and (maybe) update
64
           for k in ineigh:
65
                newcost = cost[jj] + wei[jj, k]
66
                if (newcost < cost[k]):</pre>
67
                    cost[k] = newcost
68
                    par[k] = jj
69
70
           # step 2: determine minimum-cost point among UnVisited
71
                      vertices and make this point the new point
72
           icnsdr = np.array(UnVisited.nonzero()).flatten()
           cmin, icmin = cost[icnsdr].min(0), cost[icnsdr].argmin(0)
73
74
           jj = icnsdr[icmin]
75
76
           # step 3: update "visited"-status and determine neighbors of new point
77
           UnVisited[jj] = 0
78
           tmp = UnVisited * wei[jj, :]
79
           ineigh = np.array(tmp.nonzero()).flatten()
80
           L = np.array(UnVisited.nonzero()).flatten().size
81
82
       # determine the shortest path
83
       shpath = [isp]
84
       while par[isp] != ist:
85
           shpath.append(par[isp])
86
           isp = par[isp]
87
       shpath.append(ist)
```

```
88
89
        return shpath[::-1]
90
91
    def next_node(path):
        """ Returns the next index (after the node itself) in the path.
92
93
            If the path contains only one node, returns the node itself.
94
95
        if len(path) == 1:
            return path[0]
96
97
        else:
98
            return path[1]
99
100
101 def update_weight_matrix(epsilon, c, original_weight_matrix, noNodes=58):
102
103
        This function updates the weight matrix according to step 5 of the
104
        Project. Note the added fix - the weight matrix is not changed if
105
        the original entry was 0.
106
107
108
109
        :param epsilon: given in question
110
        :param c: the vector containing number of cars at each node
111
        :param original_weight_matrix: the weight matrix given by RomeEdges
112
        :param noNodes: number of nodes in the system
113
        :return: the updated weight matrix
114
115
        new_weight_matrix = np.zeros((noNodes, noNodes))
        for i in range(noNodes):
116
117
            for j in range(noNodes):
118
                 if original_weight_matrix[i, j] != float(0):
119
                     new_weight_matrix[i, j] = original_weight_matrix[i, j] + \
120
                                                (epsilon * (float(c[i]) +
121
                                                             float(c[j]))) / float(2)
122
        return new_weight_matrix
123
124
125
    def extract_data():
126
127
        This function opens the RomeVertices and RomeEdges files, and creates
128
        global variables RomeX, RomeY, RomeA, RomeB and RomeV. These are variables
129
        used to create the original weight matrix.
130
        .....
131
```

```
132
        global RomeX, RomeY, RomeA, RomeB, RomeV
133
        RomeX = np.empty(0, dtype=float)
134
        RomeY = np.empty(0, dtype=float)
        with open('./data/RomeVertices', 'r') as file:
135
            AAA = csv.reader(file)
136
137
             for row in AAA:
                 RomeX = np.concatenate((RomeX, [float(row[1])]))
138
139
                 RomeY = np.concatenate((RomeY, [float(row[2])]))
140
        file.close()
        RomeA = np.empty(0, dtype=int)
141
        RomeB = np.empty(0, dtype=int)
142
143
        RomeV = np.empty(0, dtype=float)
144
        with open('./data/RomeEdges2', 'r') as file:
            AAA = csv.reader(file)
145
            for row in AAA:
146
                 RomeA = np.concatenate((RomeA, [int(row[0])]))
147
148
                 RomeB = np.concatenate((RomeB, [int(row[1])]))
                 RomeV = np.concatenate((RomeV, [float(row[2])]))
149
150
        file.close()
```

Now using these functions, we can execute the main program.

```
1
2
                               Main program
3
4
5
6
  if __name__ == '__main__':
7
8
       # Import the rome edges file
9
       extract_data()
10
       # Use the calcWei function from tutorials, along with the data set given
11
12
       # to calculate the weight matrix. Also create a copy which is the
13
       # temporary weight matrix.
14
       weight_matrix = misc.calcWei(RomeX, RomeY, RomeA, RomeB, RomeV)
15
       temp_wei = weight_matrix.copy()
16
17
       # Initialise minutes and number of nodes
18
       minutes = 200
19
       total_nodes = weight_matrix.shape[0]
20
21
       # Need a vector carNumbers which stores the number of cars at each vertex
22
       # in the graph.
23
       cars_at_node = np.zeros(total_nodes, dtype=int)
```

```
24
       cars_at_node_updated = cars_at_node.copy() # cars_at_node updated is simil
25
       max_cars_at_node = cars_at_node.copy() # max_cars_at_node is similar
26
27
       # To find the edges utilised, we need a 58x58 matrix of
28
       # False's. We will set each element to True if we move
29
       # cars from node i to node j.
       edge_utilised = np.zeros((total_nodes, total_nodes), dtype=bool)
30
31
32
       # Iterate through the 200 minutes
33
       for i in range(minutes):
34
35
           # Apply Dijkstra's algorithm to find the fastest path to node 52 in
36
           # the system. Then use next_node to find the next node in the given
37
           # path. (step 1)
38
           next_nodes = [next_node(Dijkst(node, 51, temp_wei))
39
                          for node in range(total_nodes)]
40
41
           # Move all cars as in steps 2,3. Iterate through every node in the
42
           # system to do this.
43
           for j_node in range(total_nodes):
44
45
               if j_node == 51:
                    # We remove 40% of cars from node 52.
46
47
                    cars_at_node_updated[51] += int(round(cars_at_node[51] * 0.6))
48
                else:
49
50
                    # Initialise the number of cars at node j_node.
51
                   number_of_cars = cars_at_node[j_node]
52
53
                    # Initialise the next node to move to.
54
                   node_to_move_to = next_nodes[j_node]
55
56
                    # 70% of cars will move. to keep the total conserved,
57
                    # the amount staying is just
58
                    # number_of_cars - amount_moving
59
                    amount_moving = int(round(0.7 * number_of_cars))
                    amount_staying = number_of_cars - amount_moving
60
61
62
                    # We now update cars_at_node.
63
                    cars_at_node_updated[j_node] += amount_staying
64
                    cars_at_node_updated[node_to_move_to] += amount_moving
65
66
                    if amount_moving > 0:
67
                        # Update edges_utilised matrix
```

```
68
                        edge_utilised[j_node, node_to_move_to] = True
69
           # Now all cars have moved where they need to, we set cars_at_node
70
71
           # to this updated vector, and empty the updated vector for the next
72
           # iteration.
73
           cars_at_node = cars_at_node_updated.copy()
74
           cars_at_node_updated = np.zeros(total_nodes, dtype=int)
75
76
           # For the first 180 minutes, 20 cars are injected into node 13.
77
           if i <= 179:
78
                cars_at_node[12] += 20
79
80
           # The temporary weight matrix is updated.
81
           temp_wei = update_weight_matrix(0.01, cars_at_node, weight_matrix)
82
83
           # We have finished an iteration.
84
85
          # Now we calculate the maximum number of cars at each node in the system
86
           max_cars_at_node = [max(cars_at_node[node], max_cars_at_node[node])
87
                                for node in range(total_nodes)]
```

## 2 QUESTIONS

Determine for each node the maximum load (maximum number of cars) over the 200 iterations.

As we have already incorporated the calculation of maximums in the loop, this is easy:

Thus the i'th element of this array gives [node i, maximum number of cars of node i over the 200 iterations]. The code to print is below along with the output.

```
print('max_index_tracker is')
print(max_index_tracker[0:10])
print(max_index_tracker[10:20])
print(max_index_tracker[20:30])
print(max_index_tracker[30:40])
print(max_index_tracker[40:50])
print(max_index_tracker[50:(len(max_index_tracker)+1)])
```

```
max_index_tracker is
[[1, 1], [2, 2], [3, 0], [4, 7], [5, 0], [6, 8], [7, 8], [8, 0], [9, 16], [10, 14]]
[[11, 0], [12, 8], [13, 28], [14, 0], [15, 28], [16, 22], [17, 7], [18, 28], [19, 11], [20, 28]]
[[21, 38], [22, 11], [23, 7], [24, 24], [25, 40], [26, 23], [27, 6], [28, 10], [29, 13], [30, 32]]
[[31, 12], [32, 19], [33, 15], [34, 15], [35, 23], [36, 9], [37, 3], [38, 14], [39, 19], [40, 30]]
[[41, 30], [42, 11], [43, 31], [44, 28], [45, 4], [46, 0], [47, 0], [48, 11], [49, 0], [50, 23]]
[[51, 18], [52, 63], [53, 17], [54, 15], [55, 12], [56, 14], [57, 11], [58, 11]]
```

#### Which are the five most congested nodes?

Given this array *max\_index\_tracker*, we can simply sort by the second argument to find the top five most congested nodes

```
the five most congested nodes are [[52, 63], [25, 40], [21, 38], [30, 32], [43, 31]]
```

In other words, the top five most congested nodes (from highest to lowest) is node 52 with 63 cars, node 25 with 40 cars, node 21 with 38 cars, node 30 with 32 cars and node 43 with 31 cars.

#### Which edges are not utilized at all? Why?

In the main program we defined a 58x58 boolean matrix of False's, where throughout the process if cars moved from index i to index j, we made the [i,j] element of the matrix True.

Thus, we have a matrix where indices (corresponding to edges) are *True* if they are traversed by some amount of cars (> 0) in the 200 minute process.

All that is left to do is count the number of False's, making sure that we don't count a False if the edge couldn't be traversed to begin with (in the original weight matrix). This corresponds to an element of the original weight matrix being 0.

So an element [l, m] belonging to this vector  $non\_utilised\_edges$  implies that the edge  $l \rightarrow m$  is not utilised.

The code to print this and the output is below.

```
print('the non utilised edges are')
print(non_utilised_edges[0:10])
print(non_utilised_edges[10:20])
print(non_utilised_edges[20:30])
print(non_utilised_edges[30:40])
print(non_utilised_edges[40:50])
print(non_utilised_edges[50:60])
print(non_utilised_edges[60:(len(non_utilised_edges)+1)])
```

```
[1, 7], [2, 3], [3, 2], [3, 5], [4, 1], [4, 7], [5, 8], [7, 1], [7, 6]]
[[8, 9], [8, 11], [9, 8], [10, 9], [11, 14], [11, 16], [12, 4], [14, 15], [14, 18], [15, 13]]
[[16, 11], [17, 19], [18, 15], [19, 10], [21, 18], [23, 12], [27, 23], [28, 17], [29, 34], [30, 21], [31, 27], [31, 36], [32, 22], [32, 26], [33, 22], [33, 32].
                        [39, 38],
            [38, 35],
                                    [41, 29], [41, 39], [41, 43], [42, 31],
                                                                                   [42, 44],
                                                                                              [43, 30],
                                                                                                           [43, 48]]
            [44, 36],
                        [44, 41],
                                    [44, 42],
                                                [45, 30], [45, 46], [46, 37],
                                                                                   [46, 45],
                                                                                               [47, 48],
                                                                                                           [48, 45]]
                                                           [52, 58],
            [49, 43], [49, 47],
                                    [52, 42], [52, 44],
                                                                                               [56, 54],
                                                                       [54,
                                                                            491.
                                                                                   [55, 56],
                                                                                                           [57, 55], [58, 57]]
```

Now I answer *why* these edges aren't utilised.

Firstly we note that from the first question where we found the maximum load for all cars over the 200 iterations, it is clear that nodes with maximum load 0 (i.e. node 3, 5, 8 etc) will always appear on this list since they were never visited. However, the real question is *why* they were never visited.

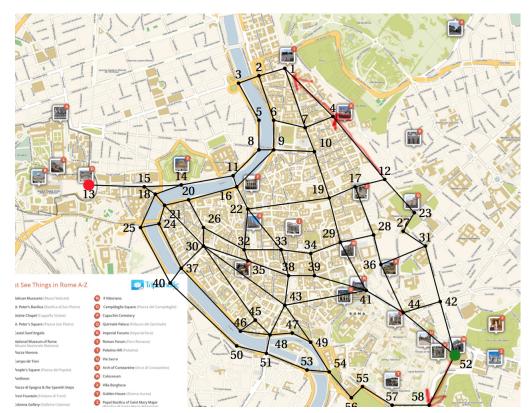
If we perform Dijkstra's algorithm on the original weight matrix and find the fastest route form node 13 to 52, we find this:

```
1 print(Dijkst(12,51,weight_matrix))
[12, 14, 17, 24, 39, 49, 50, 52, 53, 55, 54, 56, 57, 51]
```

note that these are the python indices and not the actual node numbers.

If we look at the diagram, this means without taking the effect of congestion, the fastest path throughout the city is almost completely at the bottom of the map, by the river. Such as  $40 \rightarrow 50 \rightarrow 51$ . This means that it is unlikely for a car to need to travel through the upper west side of the city, such as through nodes 1,3,5 and 2. The only way the shortest path for a car could include such nodes is if the rest of the city is *so* congested that the shortest path will indeed require these upper left nodes. We can see that edges like [3,2], [8,9] are not optimal for this reason, and are thus unused.

However for the vast majority of unused edges, it is intuitively clear that the reason they are included is because they are pointing completely away from the direction towards node 52. For example, I have highlighted edges  $4 \rightarrow 1$ ,  $12 \rightarrow 4$  and  $52 \rightarrow 58$  in the diagram below.



For a lot of these edges [a,b], they are moving completely in the opposite direction that the straight line drawn from node a to node 52 would go.