

Rareskills hack solutions

Ethernaut

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The exploit lies in the fact that the `block.number` is not random

If you write a contract which has the exact same logic to generate the side and call the main contract flip with this generated side it will work

This is because the whole thing runs in one transaction and hence the `block.number` will be the same

The contract doesn't allow 10 flip guesses within the same block though

So you do need to call the exploit contract 10 times each submitted after the last block is done

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The exploit lies in the fact that the contract allows anyone to change the owner as long as `tx.origin != msg.sender`

If A calls contract B which calls contract C, `tx.origin` is A while `msg.sender` is B

So if we deploy a contract which can call the ethernaut contract's `changeOwner` method, `tx.origin` and `msg.sender` will differ and we get to gain ownership

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Ethernaut deploys the contract with an `_initialSupply` of 20

So us as the `msg.sender` now has a balance of 20

This contract uses solidity `^0.6.0` which means there is no arithmetic overflow protection

So see what happens if you call the transfer function with `_to` as some random

address and `_value` being 21

`balances[msg.sender] - _value` becomes `20 - 21` which causes an underflow and results in $2^{256}-1$ as the value

This means that the require check in line 14 passes as $2^{256}-1$ is greater than or equal to 0

And it sets our balance to $2^{256}-1$ in line 15 which is the largest possible uint256 value in solidity. Exploited :(

Also the `_to` address in the transfer call shouldn't be our address as that would cause the underflow caused in line 15 to overflow again in line 16 back to value 20

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The exploit here is reentrancy

The contract does a `msg.sender.call` before updating the `balances` object

This means we can write an exploiter contract with a `exploit` method which calls the `withdraw` method and has a `receive` callback which again triggers the `exploit` method

This cycle will keep repeating till the entire balance of the contract is drained and then the `exploit` method terminates

Note: this approach would also have failed if the balance reduction logic used `SafeMath` or had overflow protection

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The solution lies in the fact that we get to:

1. implement the `isLastFloor()` method for the building
2. the elevator has a `goTo()` method which calls `isLastFloor()` twice:
 - a. once to check if the floor is the top floor of the building
 - b. setting value of top

See how the elevator sets the floor state variable between the above 2 events

So all we need to do is implement the `isLastFloor()` view method such that it returns two different values depending on the value of floor in the elevator

Return true if the floor state variable is equal to the top floor for our building (we just take that as 101 here) and false otherwise

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``transfer`` is guarded by the `lockTokens` modifier but ``transferFrom`` isn't!

All we need to do is approve ourselves as a spender for the entire supply and then use the `transferFrom` method to move the tokens

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There are two aspects to this problem:

1. Figuring out the token address as it's supposed to be lost
2. How to drain the ether stored in the token

Solving 1 is not hard. The address of a new contract is deterministic and only depends on the deployer address (aka the factory) and the nonce

Solving 2 is easy too. Notice the token has a unprotected `destroy` function which anyone can call. This function calls `selfdestruct` and hence anyone can drain the token of its ether provided they calculate the address from the step above

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This is a very interesting exploit

We need to make sure that the owner can't call the `withdraw` method

But the only thing we can control is who the partner is

The `withdraw` function does send ETH to the partner

But it uses a `call` method so us writing a contract to act like a partner and reverting on ETH received won't work. The contract doesn't check the success bool of this `call` method

The ethernaut problem saying: whilst the contract still has funds, and the transaction is of 1M gas or less gives us two hints:

Reentrancy is not the answer here though we can use it to drain the funds in the contract

But gas limit of upto 1M shows us the solution

When a `call` method happens, the person making the transaction, i.e. the owner has to spend gas to execute the logic executed by the partner contract on receiving ETH

So if we write a contract which on receiving ETH executes logic which spends more than 1M

gas and make this contract the partner, the owners withdraw call will fail as it is called with a gaslimit < 1M

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The solution lies in the fact that we get to:

1. implement the price() method for the buyer
2. the shop has a buy() method which calls price() twice:
 - a. once to check if the price is acceptable for the shop
 - b. setting the actual price paid

See how the shop sets the isSold state variable between the above 2 events

So all we need to do is implement the price() view method such that it returns two different values depending on the value of isSold in the store

Return a large amount for the min-price check and then return 0 for the final call

Capture the Ether

Guess the new number

All you need is a contract to call the guess function with the guess being the same logic as what's in the above contract

The block.number and now will be the same in the same transaction so the answer will be the same

Guess the number

The answer is 42

It's hardcoded in the contract :P

Guess the random number

Notice how the answer is calculated and stored in the constructor

It just uses blockhash and block.timestamp both of which can be easily queried

i.e. after this contract is deployed, you can easily get all the data required to calculate this answer by just looking at the deployed and prev block

So calculate it and call the guess function with that value

Guess the secret number

I thought this was impossible till I saw that the guess was supposed to be a uint8

So the possible guesses are just 0 to 255

Just calculate keccak256 for all of them and see which one matches the hash in the contract

Predict the block hash

The trick here lies in the fact that blockhash will return 0 for block numbers which are more than 256 blocks older than the current block

So lock in a guess of 0 as the guess, wait for 256 blocks to be mined and then call the settle method

The blockhash will return 0 as the settlementBlockNumber is now more than 256 blocks old and your guess of 0 will match

Predict the future

Here you are supposed to guess a future value

But that guess can only be one of 0-9

This future value is represented by a fixed logic for each block

We write an exploiter contract and make it guess a value 0 initially

Then for every block we make it calculate the correct guess for that block

If it matches the initially guessed value, we make the exploiter contract call the settle function

We repeat this logic for every block till a match eventually happens

Token Bank

The bank contract has a reentrancy bug

Notice how in the withdraw method, the token transfer method is called before the balanceOf object is updated

If we can find a way to call the withdraw method again through the transfer method before the balanceOf object is updated we can effectively drain the bank of its token backing

But the token is a ERC223 Token

Which means all we need to do is to make the player a contract which implements a tokenFallback method which calls the withdraw method on the bank again till its fully drained

Token Sale

The max value of a uint256 is $2^{256}-1$

2^{256} overflows and becomes 0 because this contract is still on pragma ^0.4.21

The buy logic in line 16 does `numTokens * PRICE_PER_TOKEN`

Price per token is 1 ether i.e. 10^{18} wei

So if the user is able to make `numTokens * PRICE_PER_TOKEN` equal 0 when calling the buy function with 0 wei, the require check will pass

So call the function with numTokens as 2^{238}

Multiplying with 1 ether the require statement will calculate 2^{256} which becomes 0

The require check passes and the user is allotted 2^{238} tokens

Now withdrawing the 1 ether present in the contract is as simple as calling sell with the numTokens as 1

Token Whale

The max value of a uint256 is $2^{256}-1$

2^{256} overflows and becomes 0 because this contract is still on pragma ^0.4.21

Similarly 0 underflows and becomes $2^{256}-1$ too

The exploit is with how transferFrom calls the _transfer method

Notice how _transfer reduces the balance from msg.sender not the from address which should be the case when called by transferFrom

So the exploit works as follows:

1. Assume there are 2 users: user1 and user2. user1 starts with 1000 tokens and is the player
2. First user1 calls transfer(user2, 1000) and transfers all their tokens to user2
3. Then user2 calls approve(user1, 1) and gives user1 approval to spend 1 token they hold
4. So user1 has 0 tokens, user2 has 1000 tokens and user1 has approval to spend 1 token of user2
5. Now all user1 has to call is transferFrom(user2, address(0), 1)
6. All the require checks pass
 1. balance[user2] == 1000 which is > 1
 2. balanceOf[to] + value \geq balanceOf[to] works fine because we used address(0)
Any address here works for value 1 as long as the to address doesn't have a balance of $2^{256}-1$
 3. allowance[from][msg.sender] == 1 which is ≥ 1
 4. line: 27 now does balanceOf[msg.sender] -= value
but balanceOf[msg.sender] is 0 as msg.sender is user1
value is 1
So $0-1$ happens which underflows to $2^{256}-1$ and balance[user1] gets set to this

Damn Vulnerable Defi

Free rider

There are two bugs in the marketplace that makes this exploit possible:

1. The buyMany method calls the _buyOne method which checks msg.value, but this means you can buy multiple nft's and it will only check if the sent ETH is enough to buy 1 nft
i.e. if you try to buy 10 nft's worth 10 ETH each
you can call the buyMany method to buy all of them

- and just need to send 10 ETH instead of the required 100 ETH
2. The pay seller part of the `_buyOne` method is actually paying the buyer the ETH not the seller
This is because it assumes `token.ownerOf(tokenId)` is the seller which is infact the buyer because the previous line already transfered the nft from the seller to the buyer

Using the above 2 bugs, we can get a balance of the following in our attacker address:

0.5 ETH(initial balance) +
90 ETH(total ETH in marketplace) +
45 ETH(buyer reward) -
15 ETH(used to buy all the nft's from the marketplace) -
uniswapFees
to get a total of ~120.45 ETH

The steps to do this is as follows:

1. Write a exploit contract which implements `IUniswapV2Callee` and `IERC721Receiver` and has a `receive()` method to accept ETH
2. Use uniswap to get a flashloan of 15 WETH
3. In the uniswap callback method, withdraw the 15 WETH for 15 ETH
4. Use the 15 ETH to buy all 6 nft's exploiting bug #1
Now you have ownership of all 6 nft's
Bug #2 also means the marketplace sent you 90 ETH
5. Transfer all 6 nft's to the buyer hence getting 45 ETH from the buyer as the reward
6. Calculate the amount to pay back to uniswap to settle the flash loan
This will be 15 WETH + 0.3% fees
7. We only posses ETH right now
Deposit ETH equal to the amount calculated in step #6 to the WETH contract to gain the same amount of WETH
8. Transfer the above WETH to the uniswap pair to settle the flash loan
9. Send all remaining ETH to the attacker addres
For the above tests this amount should be ~120.45 ETH

Naive receiver

The exploit lies in the fact that the `flashLoan` method in the pool has `borrower` as an argument

This means that you can borrow loans on behalf of anyone

And every loan causes a 1 ETH fee to be paid back to the pool

So all we need is to call the `flashLoan` method with the `borrower` as the receiver address to drain ether from the receiver

As the receiver has 10 ETH and the fixed fee for the pool is 1 ETH we need to call the `flashLoan` method with the receiver as the borrower

10 times

This can be done in 1 transaction if you do all 10 calls using one method in an exploit contract

Side entrance

Another fun exploit

The flashLoan method only checks if the total ETH in the pool is same at the end

It doesn't check if the balances value of the borrower is the same

The pool also allows ETH you deposited into the pool to be withdrawn

So the exploit works like this:

1. write an exploit contract which implements IFlashLoanEtherReceiver and can receive ETH
 2. in the execute method called by the pool, deposit all ETH borrowed back into the pool
 3. call the flashLoan method borrowing all the ETH in the pool followed by calling the withdraw method followed by sending the withdrawn ETH to the attacker address we want the final ETH to reach
 4. this causes the exploit contract to first borrow all the ETH and then deposit it back into the pool because the flashLoan triggers the execute method
- The execute method deposits the borrowed ETH and updates the balances of the exploit contract
- Now all the ETH in the pool is withdrawable by the exploit contract
- Hence the following method calls like withdraw work as expected draining the pool

Truster

This is a fun exploit. The exploit lies in the fact that the pool does a target.functionCall(data) where both target and data is something you can pass in

But the pool does a balance check after that

So what we want to do is do some sort of call which grants you the ability to drain the tokens from the pool after calling flashLoan

Wait the token is a ERC20 contract and hence has an approve call

So the exploit works like this:

1. write an exploit contract
we want to do all this in 1 transaction
2. call the flashLoan method with
borrowAmount=0,
borrower=attacker(any address is fine in this case),
target=erc20TokenAddress and
data=approve(address(exploitContract),fundsToDrain)

Essentially we make the pool call the approve method on the token to grant our exploit contract approval to move all the funds in the pool later via transferMany

As this is only an approval and not an actual transfer, the balance checks towards the end of the flashLoan method passes just fine

3. call the transferFrom method on the token to transfer all the pools funds to the attacker

Unstoppable

The exploit lies in the fact that the flashLoan method in the lender contains a check saying `assert(poolBalance == balanceBefore)`

It's assuming that the lender will only receive tokens via its `depositTokens` method which updates `poolBalance` too

But the token is a normal ERC20 contract

Anyone can send tokens to the lenders balance and increase it thus causing `poolBalance` to not match the token balance of the lender

This causes `assert` to revert the flashLoan transaction and the lender breaks!