

## CS 344: OPERATING SYSTEMS LABS

### GROUP 12

### ASSIGNMENT\_0A-1

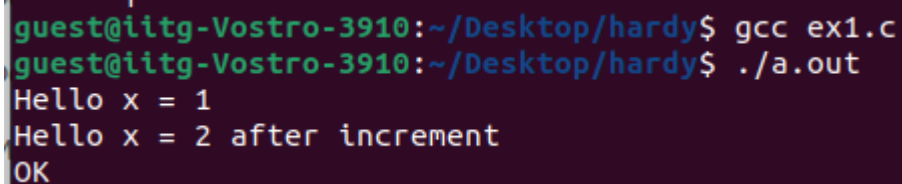
220123029 KAVURI VEDA VARSHA  
220123030 KODIBOYINA LEELA VARSHINI  
220123031 KOJJA VAMSI KRISHNA  
220123052 RAMINENI HARDHIKA

#### EXERCISE 1:

Modified code for ex1.c is as follows:

```
1 #include <stdio.h>
2 int main(int argc, char **argv)
3 {
4     int x = 1;
5     printf("Hello x = %d\n", x);
6     asm ("movl %1, %0;":"=r"(x):"r"(x+1):);
7     printf("Hello x = %d after increment\n", x);
8     if(x == 2) printf("OK\n");
9     else printf("ERROR\n");
10 }
```

Output:



```
guest@iitg-Vostro-3910:~/Desktop/hardy$ gcc ex1.c
guest@iitg-Vostro-3910:~/Desktop/hardy$ ./a.out
Hello x = 1
Hello x = 2 after increment
OK
```

Explanation:

The format of GNU assembler assembly code in C is:

asm [volatile]( AssemblerTemplate

: Output Operands

[: Input Operands

[: Clobbers ] ] )

In asm: %0 -> first variable(here x), %1 -> second variable(here x+1)

movl src, dst moves the contents of source to destination.

The “r” in the operand constraint string indicates that the operand must be located in a register. Each output operand must have “=” in its constraint.

#### EXERCISE 2:

When executing the ‘si’ command multiple times, we observed that:

Upon a processor reset, the processor switches to real mode, setting the Code Segment (CS) to 0xf000 and the IP to 0xffff0. This configuration causes execution to start at the address specified by these segment registers (CS:IP). Then the BIOS jumps backward to an earlier location in the BIOS.

'lidtl' command sets up an interrupt descriptor table and 'lgdtl' sets up a global descriptor table. Along with this, it initializes various devices such as the VGA display.

```
(gdb) source .gdbinit
+ target remote localhost:26000
warning: No executable has been specified and target does not support
determining executable automatically. Try using the "file" command.
The target architecture is assumed to be i8086
[f000:ffff] 0xffff0: ljmp $0x3630,$0xf000e05b
0x0000ffff in ?? ()
+ symbol-file kernel
warning: A handler for the OS ABI "GNU/Linux" is not built into this
configuration
of GDB. Attempting to continue with the default i8086 settings.

(gdb) si
[f000:e05b] 0xfe05b: cmpw $0xffc8,%cs:(%esi)
0x0000e05b in ?? ()
(gdb) si
[f000:e062] 0xfe062: jne 0xd241d416
0x0000e062 in ?? ()
(gdb) si
[f000:e066] 0xfe066: xor %edx,%edx
0x0000e066 in ?? ()
(gdb) si
[f000:e068] 0xfe068: mov %edx,%ss
0x0000e068 in ?? ()
(gdb) si
[f000:e06a] 0xfe06a: mov $0x7000,%sp
0x0000e06a in ?? ()
(gdb) si
[f000:e070] 0xfe070: mov $0x2d4e,%dx
0x0000e070 in ?? ()
(gdb) si
[f000:e076] 0xfe076: jmp 0x5575ff02
0x0000e076 in ?? ()
```

```
(gdb) si
[f000:ff00]    0xffff00: cli
0x0000ff00 in ?? ()
(gdb) si
[f000:ff01]    0xffff01: cld
0x0000ff01 in ?? ()

(gdb) si
[f000:ff02]    0xffff02: mov    %ax,%cx
0x0000ff02 in ?? ()
(gdb) si
[f000:ff05]    0xffff05: mov    $0x8f,%ax
0x0000ff05 in ?? ()
(gdb) si
[f000:ff0b]    0xffff0b: out    %al,$0x70
0x0000ff0b in ?? ()
(gdb) si
[f000:ff0d]    0xffff0d: in     $0x71,%al
0x0000ff0d in ?? ()
(gdb) si
[f000:ff0f]    0xffff0f: in     $0x92,%al
0x0000ff0f in ?? ()
(gdb) si
[f000:ff11]    0xffff11: or     $0x2,%al
0x0000ff11 in ?? ()
(gdb) si
[f000:ff13]    0xffff13: out    %al,$0x92
0x0000ff13 in ?? ()
(gdb) si
[f000:ff15]    0xffff15: mov    %cx,%ax
0x0000ff15 in ?? ()
(gdb) si
[f000:ff18]    0xffff18: lidt1  %cs:(%esi)
0x0000ff18 in ?? ()
```

```

(gdb) si
[f000:ff1e] 0xffff1e: lgdtl lgdtl %cs:(%esi)
0x0000ff1e in ?? ()
(gdb) si
[f000:ff24] 0xffff24: mov %cr0,%ecx
0x0000ff24 in ?? ()
(gdb) si
[f000:ff27] 0xffff27: and $0xffff,%cx
0x0000ff27 in ?? ()
(gdb) si
[f000:ff2e] 0xffff2e: or $0x1,%cx
0x0000ff2e in ?? ()
(gdb) si
[f000:ff32] 0xffff32: mov %ecx,%cr0
0x0000ff32 in ?? ()
(gdb) si
[f000:ff35] 0xffff35: jmpw $0xf,$0xff3d
0x0000ff35 in ?? ()
(gdb) si
The target architecture is assumed to be i386
=> 0xffff3d: mov $0x10,%ecx
0x0000ffff3d in ?? ()
(gdb) si
=> 0xffff42: mov %ecx,%ds

```

### EXERCISE 3:

Observations:

In tracing the `readsect()` and `bootmain()` functions in `bootmain.c` and the `bootblock.asm` assembly code:

1. The loop starts at address `0x7d8f`. The entry condition is checked at `0x7d83`; if true, the `jmp` at `0x7d8d` directs execution to the loop at `0x7d96`. Subsequent iterations of the loop run from `0x7d8f` to `0x7db4`.
2. The loop ends at address `0x7db4`. Just before exiting, the `jbe` instruction at `0x7d94` is executed if the loop continuation condition at `0x7d92` is not met.
3. After the loop, execution moves to address `0x7d87`, where a call to the `entry()` function is made. Setting a breakpoint at `0x7d87`, and continuing execution, shows that the last instruction executed is `call *0x10018`. Using `si` after this instruction leads to entering the kernel at address `0x0010000c`, as shown in the output image.

```

for(; ph < eph; ph++){
7d83: 39 f3          cmp    %esi,%ebx
7d85: 72 0f          jb     7d96 <bootmain+0x5b>
entry();
7d87: ff 15 18 00 01 00 call   *0x10018
7d8d: eb d5          jmp    7d64 <bootmain+0x29>
for(; ph < eph; ph++){
7d8f: 83 c3 20       add    $0x20,%ebx
7d92: 39 de          cmp    %ebx,%esi
7d94: 76 f1          jbe    7d87 <bootmain+0x4c>
pa = (uchar*)ph->paddr;
7d96: 8b 7b 0c       mov    0xc(%ebx),%edi
readseg(pa, ph->filesz, ph->off);
7d99: ff 73 04       pushl  0x4(%ebx)
7d9c: ff 73 10       pushl  0x10(%ebx)
7d9f: 57             push   %edi
7da0: e8 53 ff ff ff call    7cf8 <readseg>
if(ph->memsz > ph->filesz)
7da5: 8b 4b 14       mov    0x14(%ebx),%ecx
7da8: 8b 43 10       mov    0x10(%ebx),%eax
7dab: 83 c4 0c       add    $0xc,%esp
7dae: 39 c1          cmp    %eax,%ecx
7db0: 76 dd          jbe    7d8f <bootmain+0x54>
    stosb(pa + ph->filesz, 0, ph->memsz - ph->filesz);
7db2: 01 c7          add    %eax,%edi
7db4: 29 c1          sub    %eax,%ecx

```

1. The processor begins executing 32-bit code at address 0x7c31, where it encounters the instruction `mov $0x10, %ax`. The CR0 register, a 32-bit control register in the 386 processor, contains the PE flag, which determines whether the processor is in Protected mode (PE = 1) or Real mode (PE = 0). To transition from 16-bit to 32-bit mode, several steps are performed: first, a global descriptor table (GDT) is set up and loaded using the `lgdt` command. Next, the PE flag in the CR0 register is set to 1 to enable Protected mode. Finally, a long jump (`ljmp`) is executed to address `start32`, ensuring that the processor correctly reloads the code segment (`%cs`) and instruction pointer (`%eip`) in 32-bit mode. This setup uses identity mapping for the segment descriptors, which means no address translation is applied.

```

# Switch from real to protected mode. Use a bootstrap GDT that makes
# virtual addresses map directly to physical addresses so that the
# an effective memory map doesn't change during the transition.
lgdt    gdt_desc
movl    %cr0, %eax
orl     $CR0_PE, %eax
movl    %eax, %cr0

//PAGEBREAK!
# Complete the transition to 32-bit protected mode by using a long jmp
# to reload %cs and %eip. The segment descriptors are set up with no
# translation, so that the mapping is still the identity mapping.
ljmp    $(SEG_KCODE<<3), $start32

.code32 # Tell assembler to generate 32-bit code now.
start32:
# Set up the protected-mode data segment registers
movw    $(SEG_KDATA<<3), %ax    # Our data segment selector

```

2. The last instruction of the bootloader that was executed is `call *0x10018` which calls the entry function and hence the bootloader transfers control to the kernel which is at address `0x0010000c`. The first instruction of the kernel it just loaded is `mov %cr4, %eax` at address `0x0010000c`. The below image shows the outcomes.

```

(gdb) source .gdbinit
+ target remote localhost:26000
warning: No executable has been specified and target does not support
determining executable automatically. Try using the "file" command.
The target architecture is assumed to be i8086
[f000:ffff] 0xffff0: jmp    $0x3630,$0xf000e05b
0x0000ffff in ?? ()
+ symbol-file kernel
warning: A handler for the OS ABI "GNU/Linux" is not built into this configuration
of GDB. Attempting to continue with the default i8086 settings.

(gdb) b *0x7c00
Breakpoint 1 at 0x7c00
(gdb) c
Continuing.
[ 0:7c00] => 0x7c00: cli

Thread 1 hit Breakpoint 1, 0x00007c00 in ?? ()
(gdb) si 18
[ 0:7c29] => 0x7c29: mov    %eax,%cr0
0x00007c29 in ?? ()
(gdb) si
[ 0:7c2c] => 0x7c2c: jmp    $0xb866,$0x87c31
0x00007c2c in ?? ()
(gdb) si
The target architecture is assumed to be i386
=> 0x7c31: mov    $0x10,%ax
0x00007c31 in ?? ()
(gdb) b *0x7d87
Breakpoint 2 at 0x7d87
(gdb) c
Continuing.
=> 0x7d87: call   *0x10018

Thread 1 hit Breakpoint 2, 0x00007d87 in ?? ()
(gdb) si
=> 0x10000c: mov    %cr4,%eax
0x0010000c in ?? ()
(gdb) x/1x 0x10018
0x10018: 0x0010000c
(gdb)

```

3. The disk sector size is 512 bytes, as specified in `bootmain.c` on line 13. When the bootloader loads the kernel, it first reads the 4096-byte ELF header, which spans 8 sectors (implemented in `bootmain.c` on line 28). After loading the ELF header, the



bootloader proceeds to load each kernel segment. Running `readelf -a` kernel revealed 16 sections, each aligned to 512 bytes. To determine the total number of sectors read by the bootloader, we calculate the sum of the ceiling values of each section's size divided by the sector size (512 bytes). This calculation showed that 436 sectors are needed to fetch the entire kernel. The terminal screenshot below demonstrates the command output used to compute the number of sectors read.

```
guest@itg-Vostro-3910:~/Downloads/xv6/xv6-public-master$ readelf -a -e kernel
ELF Header:
  Magic:   7f 45 4c 46 01 01 01 00 00 00 00 00 00 00 00 00
  Class:                               ELF32
  Data:                                   2's complement, little endian
  Version:                               1 (current)
  OS/ABI:                                UNIX - System V
  ABI Version:                           0
  Type:                                  EXEC (Executable file)
  Machine:                                Intel 80386
  Version:                                0x1
  Entry point address:                    0x10000c
  Start of program headers:                52 (bytes into file)
  Start of section headers:                201244 (bytes into file)
  Flags:                                   0x0
  Size of this header:                     52 (bytes)
  Size of program headers:                 32 (bytes)
  Number of program headers:                3
  Size of section headers:                 40 (bytes)
  Number of section headers:               17
  Section header string table index:       16

Section Headers:
 [Nr] Name                Type              Addr             Off             Size             ES Flg  Lk  Inf  Al
 [ 0]                      NULL              00000000          000000          000000          00  0   0   0   0
 [ 1] .text                  PROGBITS          80100000          001000          0071e8          00  AX   0   0  16
 [ 2] .rodata                PROGBITS          80107200          008200          0009cf          00  A    0   0  32
 [ 3] .data                  PROGBITS          80108000          009000          002576          00  WA   0   0 4096
 [ 4] .bss                   NOBITS            8010a580          00b576          00afb0          00  WA   0   0  32
 [ 5] .debug_line             PROGBITS          00000000          00b576          006aee          00  0    0   0   1
 [ 6] .debug_info             PROGBITS          00000000          012064          010ee3          00  0    0   0   1
 [ 7] .debug_abbrev           PROGBITS          00000000          022f47          0044aa          00  0    0   0   1
 [ 8] .debug_aranges          PROGBITS          00000000          0273f8          0003b0          00  0    0   0   8
 [ 9] .debug_str              PROGBITS          00000000          0277a8          000e16          01  MS   0   0   1
[10] .debug_loclists          PROGBITS          00000000          0285be          0050b1          00  0    0   0   1
[11] .debug_rnglists          PROGBITS          00000000          02d66f          000845          00  0    0   0   1
[12] .debug_line_str          PROGBITS          00000000          02deb4          000146          01  MS   0   0   1
[13] .comment                 PROGBITS          00000000          02dffa          00002b          01  MS   0   0   1
[14] .symtab                  SYMTAB            00000000          02e028          001fa0          10  15  66   4
[15] .strtab                  STRTAB            00000000          02ffc8          0011a4          00  0    0   0   1
[16] .shstrtab                STRTAB            00000000          03116c          0000ad          00  0    0   0   1
```

## EXERCISE 4:

Output of pointers.c:

```
guest@iitg-Vostro-3910:~/Downloads$ gcc pointers.c
guest@iitg-Vostro-3910:~/Downloads$ ./a.out
1: a = 0x7ffd6a2564a0, b = 0x55bd593412a0, c = 0x7ffd6a2568d9
2: a[0] = 200, a[1] = 101, a[2] = 102, a[3] = 103
3: a[0] = 200, a[1] = 300, a[2] = 301, a[3] = 302
4: a[0] = 200, a[1] = 400, a[2] = 301, a[3] = 302
5: a[0] = 200, a[1] = 128144, a[2] = 256, a[3] = 302
6: a = 0x7ffd6a2564a0, b = 0x7ffd6a2564a4, c = 0x7ffd6a2564a1
```

First an array `a` is declared as `int a[4]` so `a` would store the address of the first element of the array `a`. `int *b = malloc(16)`; would allocate a 16 byte block of memory and store the address of this newly allocated memory. `int *c`; declares the pointer variable `c` is assigned garbage value as it is not assigned a value.

Before line 2, `c` is set to `a` and it will store the address of the first element of array `a`. Values of array `a` are set to 100, 101, 102, 103.

Before line 3, `a[1] = 301` and set to `c[1]`. `*(c+2)` is equivalent to `c[2] = 302` changes the value of `a[3] = 302` then the value of array is 200, 300, 301, 302.

Before line 4, `c = c+1` changes `c` to the second element of `a`. So `*c = 400` sets `a[1] = 400`. Array `a` is 200, 400, 301, 302.

Just after printing line 4, `c` points to the memory location `0x7ffd6a2564a4`, the integer `a[1]` is stored from `0x7ffd6a2564a4 - 0x7ffd6a2564a7`. The statement, `c = (int *) ((char *) c + 1)`; changes `c` to `0x7ffd6a2564a5` as `sizeof(char) = 1`, and typecasting `c` into a `char*` and incrementing it would increase the memory location by 1 instead of 4 locations. Now `*c = 500` would change contents of the memory locations from `0x7ffd6a2564a5 - 0x7ffd6a2564a8`, which would change the contents of `a[1]` and `a[2]` both, as the memory locations `0x7ffd6a2564a5 - 0x7ffd6a2564a7` is a part of `a[1]` and the location `0x7ffd6a2564a8` is a part of `a[2]`. So the values stored in `a[1]` and `a[2]` are not garbage although they seem to be corrupted.

As `b` is of type `int*`, and `a` points to the location `0x7ffd6a2564a0`, `b` would point to `0x7ffd6a2564a4` as `sizeof(int) = 4`. `(char*)a` points to the memory location `0x7ffd6a2564a0` and `(char*)a + 1` would be `0x7ffd6a2564a1`, so finally `b` points to `0x7ffd6a2564a1`.



objdump commands:

```
guest@iitg-Vostro-3910:~/Downloads/xv6/xv6-public-master$ objdump -h kernel

kernel:      file format elf32-i386

Sections:
Idx Name          Size      VMA           LMA           File off  Algn
  0 .text          000071e8  80100000  00100000  00001000  2**4
    CONTENTS, ALLOC, LOAD, READONLY, CODE
  1 .rodata         000009cf  80107200  00107200  00008200  2**5
    CONTENTS, ALLOC, LOAD, READONLY, DATA
  2 .data           00002576  80108000  00108000  00009000  2**12
    CONTENTS, ALLOC, LOAD, DATA
  3 .bss            0000afb0  8010a580  0010a580  0000b576  2**5
    ALLOC
  4 .debug_line     000006ae  00000000  00000000  0000b576  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  5 .debug_info     00010ee3  00000000  00000000  00012064  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  6 .debug_abbrev   000044aa  00000000  00000000  00022f47  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  7 .debug_aranges  000003b0  00000000  00000000  000273f8  2**3
    CONTENTS, READONLY, DEBUGGING, OCTETS
  8 .debug_str      00000e16  00000000  00000000  000277a8  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  9 .debug_loclists 000050b1  00000000  00000000  000285be  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
10 .debug_rnglists 00000845  00000000  00000000  0002d66f  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
11 .debug_line_str  00000146  00000000  00000000  0002deb4  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
12 .comment        0000002b  00000000  00000000  0002dffa  2**0
    CONTENTS, READONLY
```

objdump -h kernel: As we can see in the above screenshot, VMA and LMA of .text section is different indicating that it loads and executes from different addresses.

```
guest@iitg-Vostro-3910:~/Downloads/xv6/xv6-public-master$ objdump -h bootblock.o

bootblock.o:  file format elf32-i386

Sections:
Idx Name          Size      VMA           LMA           File off  Algn
  0 .text          000001c3  00007c00  00007c00  00000074  2**2
    CONTENTS, ALLOC, LOAD, CODE
  1 .eh_frame       000000b0  00007dc4  00007dc4  00000238  2**2
    CONTENTS, ALLOC, LOAD, READONLY, DATA
  2 .comment        0000002b  00000000  00000000  000002e8  2**0
    CONTENTS, READONLY
  3 .debug_aranges  00000040  00000000  00000000  00000318  2**3
    CONTENTS, READONLY, DEBUGGING, OCTETS
  4 .debug_info     00000585  00000000  00000000  00000358  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  5 .debug_abbrev   0000023c  00000000  00000000  000008dd  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  6 .debug_line     00000283  00000000  00000000  00000b19  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  7 .debug_str      0000021a  00000000  00000000  00000d9c  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  8 .debug_line_str  00000055  00000000  00000000  00000fb6  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
  9 .debug_loclists 0000018d  00000000  00000000  0000100b  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
10 .debug_rnglists 00000033  00000000  00000000  00001198  2**0
    CONTENTS, READONLY, DEBUGGING, OCTETS
```

objdump -h bootblock.o: As we can see in the above screenshot, VMA and LMA of .text section is same indicating that it loads and executes from the same address.

### EXERCISE 5:

In the Makefile, we updated the link address to 0x7c04 from the original 0x7c00. This change caused an error: the bootloader, which initially switched from 16-bit real mode to 32-bit protected mode, started executing from the BIOS memory again instead of entering 32-bit mode. This issue arises because the 'ljmp' instruction requires a correct absolute address, and if the link and load addresses are mismatched, the jump target address becomes incorrect.

```
(gdb) source .gdbinit
+ target remote localhost:26000
warning: No executable has been specified and target does not support
determining executable automatically. Try using the "file" command.
The target architecture is assumed to be i8086
[f000:fff0] 0xffff0: ljmp $0x3630,$0xf000e05b
0x0000fff0 in ?? ()
+ symbol-file kernel
warning: A handler for the OS ABI "GNU/Linux" is not built into this configuration
of GDB. Attempting to continue with the default i8086 settings.

(gdb) si
[f000:e05b] 0xfe05b: cmpw $0xffc8,%cs:(%esi)
0x0000e05b in ?? ()
(gdb) si
[f000:e062] 0xfe062: jne 0xd241d416
0x0000e062 in ?? ()
(gdb) b *0x7c00
Breakpoint 1 at 0x7c00
(gdb) c
Continuing.
[ 0:7c00] => 0x7c00: cli

Thread 1 hit Breakpoint 1, 0x00007c00 in ?? ()
(gdb) si 18
[ 0:7c29] => 0x7c29: mov %eax,%cr0
0x00007c29 in ?? ()
(gdb) si
[ 0:7c2c] => 0x7c2c: ljmp $0xb866,$0x87c35
0x00007c2c in ?? ()
(gdb) si
[f000:e05b] 0xfe05b: cmpw $0xffc8,%cs:(%esi)
0x0000e05b in ?? ()
(gdb) si
[f000:e062] 0xfe062: jne 0xd241d416
0x0000e062 in ?? ()
(gdb) □
```

Output of objdump -f kernel:

The start address is the starting point (0x0010000c) at which the bootloader enters the kernel.

```

guest@iitg-Vostro-3910:~/Downloads/xv6/xv6-public-master$ objdump -f kernel

kernel:      file format elf32-i386
architecture: i386, flags 0x00000112:
EXEC_P, HAS_SYMS, D_PAGED
start address 0x0010000c

```

## EXERCISE 6:

After restarting the machine and examining 8 words of memory at address 0x00100000:

1. When the BIOS enters the bootloader, the memory at this location is all zeros.
2. When the bootloader enters the kernel, the memory at this location contains non-zero values.

This is because, before the kernel is loaded, the bootloader has not yet populated this memory area, resulting in zero values. Once the bootloader has loaded the kernel, this memory area contains data from the kernel, so the values become non-zero.

```

(gdb) source .gdbinit
+ target remote localhost:26000
warning: No executable has been specified and target does not support
determining executable automatically. Try using the "file" command.
The target architecture is assumed to be i8086
[f000:fff0] 0xffff0: jmp $0x3630,$0xf000e05b
0x0000fff0 in ?? ()
+ symbol-file kernel
warning: A handler for the OS ABI "GNU/Linux" is not built into this configuration
of GDB. Attempting to continue with the default i8086 settings.

(gdb) b *0x7c00
Breakpoint 1 at 0x7c00
(gdb) c
Continuing.
[ 0:7c00] => 0x7c00: cli

Thread 1 hit Breakpoint 1, 0x00007c00 in ?? ()
(gdb) x/8x 0x00100000
0x100000: 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
0x100010: 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
(gdb) b *0x7d87
Breakpoint 2 at 0x7d87
(gdb) c
Continuing.
The target architecture is assumed to be i386
=> 0x7d87: call *0x10018

Thread 1 hit Breakpoint 2, 0x00007d87 in ?? ()
(gdb) si
=> 0x10000c: mov %cr4,%eax
0x0010000c in ?? ()
(gdb) x/8x 0x00100000
0x100000: 0x1badb002 0x00000000 0xe4524ffe 0x83e0200f
0x100010: 0x220f10c8 0x9000b8e0 0x220f0010 0xc0200fd8
(gdb) 

```