

# **MATHEMATICS**

## **UPSC**

## **CIVIL SERVICES MAINS**

18

Linear Programming (LPP)

*(n)*

$$\left\{ \begin{array}{l} x_1 + x_2 + x_3 + \dots + x_m \leq 2 \\ 2x_1 - 3x_2 + 4x_3 + \dots \leq 3 \end{array} \right.$$

*m-variables.*

- Optimal solution = Most Arithmetical solution  
it can occur more than once.
- Basic solution : Remove inequality & put <sup>(m-n)</sup> equal to zero & solve.

Terminology

- |                           |                   |
|---------------------------|-------------------|
| → Basic Solution          | → Slack variable. |
| → Basic Variable          | → Surplus "       |
| → Non- " "                |                   |
| → Basic feasible solution |                   |
| → Degenerate solution     |                   |
| → Non- " "                |                   |

Example: Consider the following LPP:

$$\text{Max. } Z = 5x + 8y$$

$$x + y \leq 4 \quad \text{--- (1)}$$

$$2x + y \leq 6 \quad \text{--- (2)}$$

$$x, y \geq 0$$

Sol.

$$\text{No. of eq.} = 2$$

$$\text{Add slack variables i.e. } x + y + S_1 = 4$$

$$2x + y + S_2 = 6$$

$$\text{Now, No. of eq.} = 2$$

$$4^{\text{th}} \text{ variable} = 4$$

$$(4-2) = 2 \text{ variable}$$

Put them equal to zero

← which are non-basic variables

No. of Basic solution	Non-Basic variables (each = 0)	Basic variable	Basic solution
1	$s_1, s_2$	$x, y$	$x = 2, y = 2 \checkmark$
2	$x, s_1$	$y, s_2$	$y = 4, s_2 = 2 \checkmark$
3	$x, s_2$	$y, s_1$	$y = 6, s_1 = -2$
4	$y, s_1$	$x, s_2$	$x = 4, s_2 = -2$
5	$y, s_2$	$x, s_1$	$x = 3, s_1 = 1 \checkmark$
6	$x, y$	$s_1, s_2$	$s_1 = 4, s_2 = 6 \checkmark$

→ Check everyt = Basic feasible solution.

$$(2,2), (0,4), (3,0), (0,0)$$

→ Degenerate solution : if one or more zero is present in B.F.S. i.e.

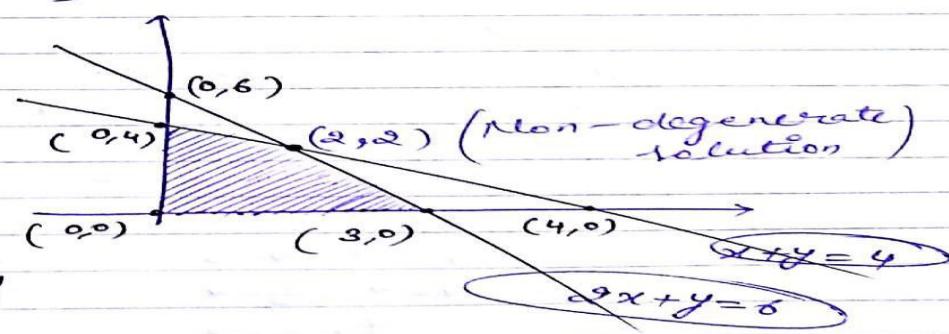
$$\boxed{(0,4), (3,0), (0,0)}$$

→ Non-degenerate solution :-

i.e.  $(2,2)$  is

if all +ve variables are present. i.e. B.F.S.

All solution lies on  $x-y$ -axis  $\Rightarrow$  Degen.



→ Graphical Method

→ Simplex Method

→ Dual-Simplex Method

• Duality

→ Transportation.

(Mains 2015)

Que:- Max.  $Z = x_1 + 2x_2 - 3x_3 + 4x_4$

(20 Marks)

$$x_1 + x_2 + 2x_3 + 3x_4 = 12$$

$$x_2 + 2x_3 + x_4 = 8$$

- i) Using definition, find its all basic solutions. which of these are degenerate basic feasible solution and which are non-degenerate basic feasible solution.

Ans

No of solution	Non-Basic Variable	Basic Variable	Basic solution
1	$x_1, x_2$	$x_3, x_4$	$\checkmark x_3 = 3, x_4 = 2$
2	$x_1, x_3$	$x_2, x_4$	$\checkmark x_2 = 6, x_4 = 2$
3	$x_1, x_4$	$x_3, x_2$	No solution
4	$x_2, x_3$	$x_1, x_4$	$x_1 = -12, x_4 = 8$
5	$x_2, x_4$	$x_1, x_3$	$\checkmark x_1 = 4, x_3 = 4$
6	$x_3, x_4$	$x_2, x_1$	$\checkmark x_1 = 4, x_2 = 8$

→ Checked sign solution = feasible solution and are non-degenerate solution.

→ There is No degenerate " "

optimal but not feasible.

(ii) Max. ( $Z$ ) at  $x_3 = 3, x_4 = 2 = -9 + 8 = -1$

"  $x_2 = 6, x_4 = 2 = 12 + 8 = 20$

"  $x_4 = -12, x_4 = 8 = -12 + 32 = 20$

"  $x_1 = 4, x_3 = 4 = 4 - 12 = -8$

"  $x_1 = 4, x_2 = 8 = 4 + 16 = 20$

} optimal solution, but feasible

Ans

# # Graphical Method :-

Que: (Mains 2014) Max.  $Z = 6x_1 + 5x_2$

$$2x_1 + x_2 \leq 16 \quad \text{--- (1)}$$

$$x_1 + x_2 \leq 11 \quad \text{--- (2)}$$

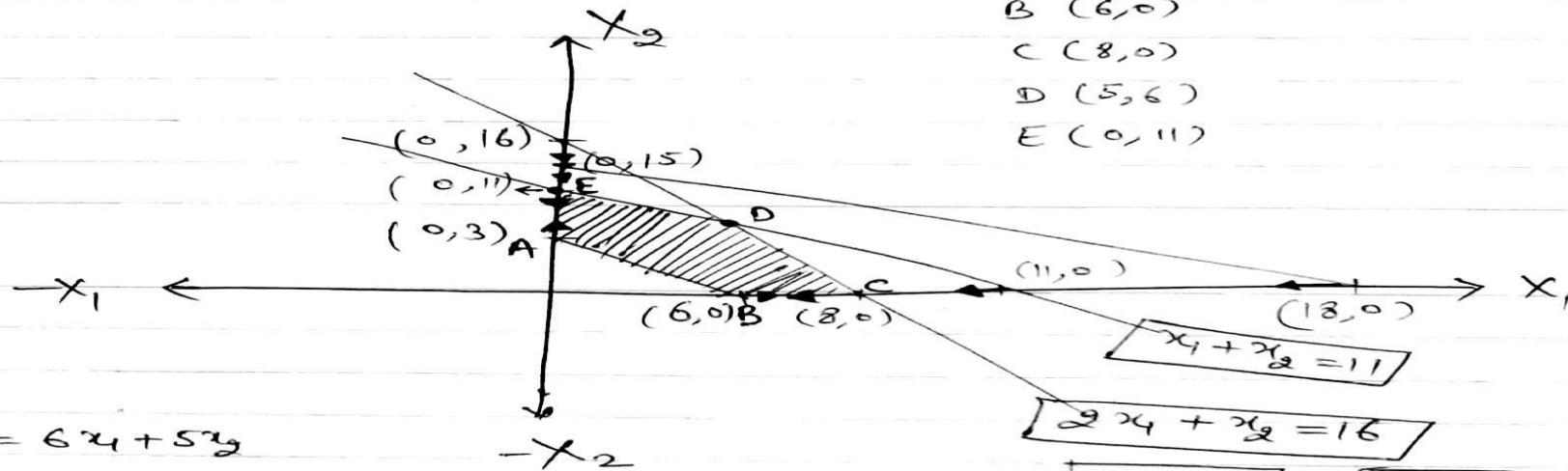
$$x_1 + 2x_2 \geq 6 \quad \text{--- (3)}$$

$$5x_1 + 6x_2 \leq 90 \quad \text{--- (4)}$$

$$x_1, x_2 \geq 0$$

10 Marks

Sol.



$$\text{Max. } Z = 6x_1 + 5x_2$$

$$Z|_{A(0,3)} = 15$$

$$Z|_{B(6,0)} = 36$$

$$Z|_{C(8,0)} = 48$$

$$Z|_{D(5,6)} = 30 + 30 = 60 = \text{Max } Z.$$

$$Z|_{E(0,11)} = 55$$

$$A(0,3)$$

$$B(6,0)$$

$$C(8,0)$$

$$D(5,6)$$

$$E(0,11)$$

so  $\boxed{\text{Max } Z = 60 \text{ at } x_1 = 5, x_2 = 6}$

Ans.

# # Simplex Method :-

MAINS - 2015 ; 20 Marks

Ques:- Solve LPP by using simplex method :-

$$\text{Max. } Z = 2x_1 - 4x_2 + 5x_3$$

$$x_1 + 4x_2 - 2x_3 \leq 2 \quad \text{--- (1)}$$

$$-x_1 + 2x_2 + 3x_3 \leq 1 \quad \text{--- (2)}$$

$$x_1, x_2, x_3 \geq 0$$

Sol.

Add two slack variable in eq. (1) & eq. (2) as follows:

$$\begin{array}{l} x_1 + 4x_2 - 2x_3 + s_1 = 2 \\ -x_1 + 2x_2 + 3x_3 + s_2 = 1 \end{array} \quad \text{Max. } Z = 2x_1 - 4x_2 + 5x_3 + 0s_1 + 0s_2$$

$C_B$	Basic	$C_j^0$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	Ratio
0	$s_1$	2	1	4	-2	1	0	$\frac{2}{2} = -1$
0	$s_2$	1	-1	2	3	0	1	$\frac{1}{3} = \frac{1}{3}$ Min ratio in <u>+</u> ve
		$Z_j$	0	0	0	0	0	
		$C_j = C_j^0 - Z_j$	2	-4	5	↑ 0	0	

Highlight Most +ve No.

$$\text{Now (new) } R_2 \leftrightarrow \frac{1}{3} R_2(\text{old})$$

$$R_1 \leftrightarrow R_1 + 2R_2(\text{new})$$

$\frac{2}{3}$	$\frac{1}{3}$	$\frac{16}{3}$	0	1	$\frac{2}{3}$	
$\frac{1}{3}$	$\frac{-1}{3}$	$\frac{2}{3}$	1	0	$\frac{1}{3}$	$\leftarrow R_2(\text{new})$
$\frac{2}{3}$	$\frac{-2}{3}$	$\frac{4}{3}$	2	0	$\frac{2}{3}$	
+2	1	4	-2	1	0	

$C_B$	Basic	$C_j^0$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	Ratio
0	$s_1$	$\frac{8}{3}$	$\boxed{\frac{1}{3}}$	$\frac{16}{3}$	0	1	$\frac{2}{3}$	(8) $\rightarrow$
5	$x_3$	$\frac{1}{3}$	$\frac{-1}{3}$	$\frac{2}{3}$	1	0	$\frac{1}{3}$	-1
		$Z_j$	$-\frac{5}{3}$	$\frac{10}{3}$	5	0	$\frac{5}{3}$	
		$C_j$	$\uparrow \frac{1}{3}$	$-2\frac{2}{3}$	0	0	$-\frac{5}{3}$	

$$\begin{array}{c}
 R_1 \longleftrightarrow 3R_1 (019) \\
 \xrightarrow{\quad} \begin{matrix} & 8 & 1 & 16 & 0 & 3 & 2 \end{matrix} \\
 R_2 \longleftrightarrow R_1 (019) + R_2 (01a) \\
 \xrightarrow{\quad} \begin{matrix} & 3 & 0 & 6 & 1 & 1 & 1 \end{matrix}
 \end{array}$$

$C_B$	Basis	$\text{val.}$	$z_j$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$
2	$x_1$	8	$z_j$	1	16	0	3	2
5	$x_3$	3	$z_j$	0	6	1	1	1
		$C_j = z_j - z_i$	$z_j$	2	62	5	11	9
		$C_j = z_j - z_i$	0	-66	0	-11	-9	

Since all the elements in the  $C_j = z_j - z_i$  are -ve.  
Hence optimal solution is obtained.

$$\Rightarrow x_1 = 8, x_3 = 3, x_2 = 0, s_1 = 0, s_2 = 0$$

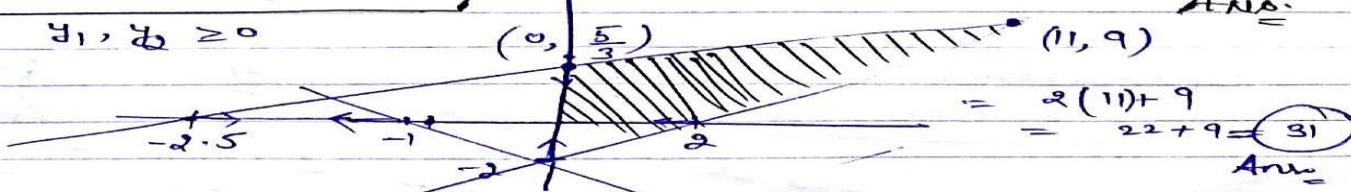
$$\begin{aligned}
 \Rightarrow \text{Max } Z &= 2(8) - 4(0) + 5(3) \\
 &= 16 + 15 = \boxed{31} \text{ Ans.}
 \end{aligned}$$

How to write the dual of the given question:

$$\left[ \begin{matrix} 1 & 4 & -2 \\ -1 & 2 & 3 \end{matrix} \right] \leftrightarrow \left[ \begin{matrix} 1 & -1 \\ 4 & 2 \\ -2 & 3 \end{matrix} \right]$$

$$\left\{
 \begin{array}{l}
 y_1 - y_2 \leq 2 \\
 4y_1 + 2y_2 \leq -4 \\
 -2y_1 + 3y_2 \leq 5
 \end{array}
 \right. \quad \text{and} \quad \boxed{\text{Max. } Z = 2y_1 + y_2}$$

Ans.



Que:- Mains (2014)

20 Marks

Using Simplex method solve :-

$$\text{Max } Z = 30x_1 + 24x_2$$

$$5x_1 + 4x_2 \leq 200$$

$$x_1 \leq 32 \quad x_1, x_2 \geq 0$$

$$x_2 \leq 40$$

$$5x_1 + 4x_2 + S_1 = 200$$

$$x_1 + S_2 = 32$$

$$x_2 + S_3 = 40$$

Sol.

CB	Basis	$C_j$	30	24	0	0	0	Min. Ratio
		Sol <sup>n</sup>	$x_1$	$x_2$	$S_1$	$S_2$	$S_3$	
0	$S_1$	200	5	1	1	0	0	40
0	$S_2$	32	1	0	0	1	0	32 →
0	$S_3$	40	0	1	0	0	1	$\infty$
		$Z_j$	0	0	0	0	0	
		$C_j$	30↑	24	0	0	0	

Operation:  $R_1 \leftrightarrow R_1 - 5R_2$

CB	Basis	$C_j$	30	24	0	0	0	Min. Ratio
		Sol <sup>n</sup>	$x_1$	$x_2$	$S_1$	$S_2$	$S_3$	
0	$S_1$	40	0	1	1	-5	0	40
30	$x_1$	32	1	0	0	1	0	$\infty$ } Tie
0	$S_3$	40	0	1	0	0	1	40
		$Z_j$	30	0	0	30	0	
		$C_j$	0	24↑	0	-30	0	

Operation:

$$R_1 \leftarrow R_1 - P_3$$

And it is found in  
i.e. 2nd Row

$$\begin{bmatrix} S_1 & \frac{1}{1} = 1 & S_3 \\ S_3 & \frac{0}{1} = 0 & \frac{1}{1} = 1 \end{bmatrix}$$

$\xrightarrow{\text{Column wise}}$

$C_B$	Basis	$c_j^*$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
0	$s_1$	0	0	0	1	-5	-1
30	$x_1$	32	1	0	0	1	0
24	$x_2$	40	0	1	0	0	1
	$Z_j$	30	24	0	30	24	
	$C_j = c_j^* - Z_j$	0	0	0	-30	-24	

Hence all the elements of ( $C_j = c_j^* - Z_j$ ) are -ve,  
we get optimal solution.

And the optimal solution is :

$$x_1 = 32 \text{ and } x_2 = 40 \text{ and } [s_1 = s_2 = s_3 = 0]$$

$$\Rightarrow \text{Max. } Z = 30(32) + 24(40) = 960 + 960 = [1920] \text{ Ans}$$

## # Dual-Simplex Method :-

$$\text{Que: - Max. } Z = -2x_1 - x_2 ; \quad 3x_1 + 2x_2 \geq 3$$

$$4x_1 + 3x_2 \geq 6 ; \quad x_1, x_2 \geq 0$$

Not. In this case we will make this "greater than" inequality into "less than" equality.  
And after that we will add "surplus variable"

$$-3x_1 - x_2 + s_1 = -3$$

$$-4x_1 - 3x_2 + s_2 = -6$$

$$-x_1 - 2x_2 + s_3 = -3$$

$C_B$	Basis	$c_i$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
0	$s_1$	-3	-3	-1	1	0	0
0	$s_2$	-6	-4	(-3)	0	1	0
0	$s_3$	-3	-1	-2	0	0	1
	$Z_j$	0	0	0	0	0	0
	$C_j$	-9	-1	0	0	0	0

$$\left[ \begin{array}{l} \frac{-2}{-4} = \frac{1}{2} = 0.5 \\ \frac{-1}{-3} = \frac{1}{3} = 0.3 \end{array} \right] \checkmark$$

$R_2 \leftrightarrow \frac{1}{3} R_2 (\text{old})$  : Operation

Q26.

$$R_1 \leftrightarrow R_1 + R_2, R_3 \leftrightarrow R_3 + 2R_2$$

$C_B$	Basis	$C_j$	-2	-1	0	0	0
			$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
0	$s_1$	-1		$\boxed{-5/3}$	0	1	$-y_3$
-1	$x_2$	2	$4/3$		1	0	$-y_3$
0	$s_3$	1	$5/3$		0	0	$-2/3$
		$\bar{x}_j$	$-4/3$		-1	0	$y_3$
		$C_j$	$-2/3 \uparrow$		0	0	$-y_3$

$$\frac{-2}{\frac{3}{5}} = \left( \frac{2}{5} \right) \text{ MIN.}$$

$$\frac{-5}{\frac{3}{5}} = \left( -\frac{5}{3} \right)$$

$$\frac{1}{3}/-y_3 = (1)$$

$$R_1 \leftrightarrow R_1 - \frac{3}{5}, R_2 \leftrightarrow R_2 - \frac{4}{3}R_1, R_3 \leftrightarrow R_3 - \frac{5}{3}R_1$$

$C_B$	Basis	$G^o$	-2	-1	0	0	0
		$\text{sol.}^n$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
-2	$x_1$	$3/5$		1	0	$-3/5$	$1/5$
-1	$x_2$	$6/5$		0	0	$4/5$	$-3/5$
0	$s_3$	0		0	0	-1	-1

Since all the elements in  $\text{sol.}^n$  column are +ve,  
hence we get optimal solution and no further  
iteration is required.

$$\Rightarrow x_1 = \frac{3}{5} \text{ and } x_2 = \frac{6}{5} \text{ and } s_1 = s_2 = s_3 = 0$$

$$\Rightarrow \text{Max. } Z = -2x_1 - x_2$$

$$x_1 = \frac{3}{5}, x_2 = \frac{6}{5}$$

$$= -2\left(\frac{3}{5}\right) - \frac{6}{5} = \frac{-6}{5} - \frac{6}{5} = \left(-\frac{12}{5}\right) \text{ Ans.}$$

THANKS