B.K. DHARANEETSH

CSE-B T

230701072

DAA ASSIGNMENT-2 roblem conference session management using

greedy approach

Session	stat	time		end	time
1	1		15. 40.	_3	
3	2		4	4	
4	6			7	
5	5	a7 -		8	

i) Sort by end time

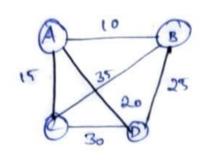
- · Iterate through each session cheek
  - · if start time operates leaned to

end time

- · If it is then select the session
  - i) Add to relected serion
  - ii) Update and time to relected session and time
- i) det end time = 0

Stat 170 End fime og Start 2 23 2 3 Start 3 = 3 End fine = 3 4 1,3 end time = 3 Start 6 >5 5 Start 543 1,3,5 The relected session are 1,3,5. Thus the maximum or non over lapping Session is 3. Items Weight Value 2) 60 10 20 100 2 3 30 120 weight in desending Doder Sort the NOW, the weight becomes 30 20 10 60 120 100

Take the Volue 1 by 1 First 30 : check 30 £ 50 if true then sum: sum + Value decrement current i from total Now, 50 - 30 =20 Total Weight 20 4 20 truc NOW, sum = 220 Next, 30-20 =0 10 2 10 . Foles -> exits So, the total Value can be Stored is 220 With items 2, 2. A B 3) City 0 A 10 10 B 35 15 C 25 20 b



- start a city A
- Visit the nearest unisited oily 2)

from current city.

Repeat step 2 until ell cities have been 3)

Visited.

After visiting all cities, return to 41

the starting city.

Start with A

From A:

A -> B : 10

A -> C : 15

-B is short

Character Frequency 4) 45 Step-1: create leaf node for each Character and their corresponding frequency A (5), B(a), c(12), D(13), E(16). F ( 45). Datep - 2: Build by taking any two lowest nodes i.e (A=5 and B=9) merge them create AB as 14. Next, new nodes (C = 12) (D= 15) (E= 16) (F = 45) (AB = 14)

From B B -> A = to (Already visited) B -> D : 25 i.D is shortest From .. C is shortest C -> A: 15 · A -> B -> D -> · Total distance = 80

A — B

( ) / )

( ) / )

Ans Sort all edges in increasing order of their weights

Pick the smallest edge: It if downto form a cycle add it to the MST.

Repeat step & until you have added (V-11 edges (where V is the number of Vertices).

Steps:  $A \rightarrow B : I \quad MST \quad A \rightarrow B$   $B \rightarrow D : A \quad MST \quad A \rightarrow B$   $B \rightarrow D : A \quad MST \quad A \rightarrow B$   $B \rightarrow D : A \quad MST \quad B \rightarrow B$ 

A -> D:3 forms ayole skip

 $A \rightarrow c : 4$   $\begin{cases} A - B , B - D , A - c \end{cases}$ 

777 Total weight of the MST: 1+2+4=7

· From above the lowest frequency are (C=12 and D=13.) merge them With frequency · Create a node New nodes (F=16) (F=45) (AB =14) (CDF 25) Take & and & and Combine Now nodes (AQ=14) (DR= 25) (YEF. (AB and E) ---- (F = 45) (CD = 25) (ABF = 36) (CD and ABE) merge than New nodes (F=45) (ABECD = 55) Morge last two The root node is loo >>> Resulting: Huffman codes: A 1100 D 10 11 01 6 01 CIII FO

6)

i) T (n) = 5T (n-1) + ws

Forward Substitution

T (n) = 5T (n-1) +13

T (n-1) = 57 (n-2) + (n-1)3

T (n-2) = 5T (n-3) + n(-2)3

After substituting back into the

original recumence.

T(n) = 5 [5T (n-2) + (n-13] + n3

Substituting again:

T (n)= 52+ (n-3) +52 (n-2)3 +...

Backward Substitution:

For Backensel substitution we start

by expressing the recurrence beckward do

lock for a general form . The key is de

oppress each subsequent tem using

the prior team.

T (n) = 4T ( = ) + x = ii) Forward Substitution  $T (n) = 4T \left(\frac{1}{2}\right) + n^{2}$   $T \left(\frac{n}{2}\right) = 2T \left(\frac{n}{4}\right) + \left(\frac{n}{2}\right)^{2}$ Sub T ( N/2) into T (n) T(n) = 4 (41 (1/4) + (2) + + 12 2 427 i) a=3 b=2 K=1 P=1 3 > 2 T (n) = 0 (log 23) T(n) = 4T (n/2) + n2 (15 b=2 de2 a=4 bd Tenie ne logn / 4