

> Repeating Subproblems means it needs ^{DATE} Dynamic Programming
+ Overlapping problems.

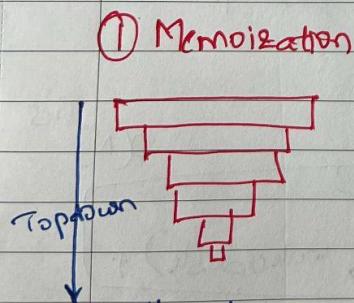
+ problems have optimal substructure property, - Overall soln can be constructed from optimal solutions of its subproblems.

Dynamic Programming solve every subproblem just once and saves the answer in a lookup table, thereby avoiding recalculating the answer everytime the subproblem is encountered.

2 ways / patterns of DP: 1) memoization 2) tabulation.

How to recognize DP:-

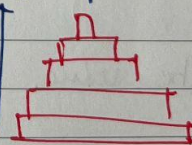
- 2 subproblems share the problem
- ① Overlapping subproblems: subproblems are not independent
 - ② Optimal substructure property: optimal soln can be constructed from optimal solns of subproblems.



① Memoization

② Tabulation.

bottom up.



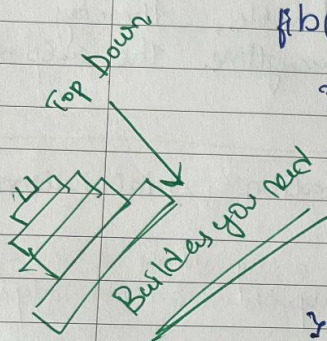
Topdown approach avoids recursion.

Typically done by filling up a lookup table, and computing the soln based on results of the table.

Similar to recursive - Except for it looks for the answer in a lookup before computing the answer.

- + Lookup Table is used to store and ^{DATE} search values.
- + Stores values in an 'array' returned from a computation. and or linear associative.
- + Lookup-table stores (key-values) pairs.

Fibonacci using Memoization: ↓



```
fib(int n, int lookupTable[]) {
```

```
    if (lookupTable[n] == -1) { // if not found
        if (n <= 1) lookupTable[n] = n; O(N) * x5
```

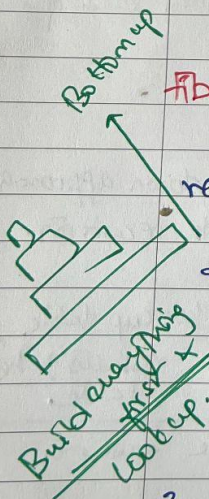
```
    else lookupTable[n] = fib(n-1, lookupTable)
        + fib(n-2, lookupTable);
```

```
    return lookupTable[n]
```

Fibonacci using Tabulation:-

Build the whole table first &

return the value.



```
fib(int n, int[] lookupTable) {
    lookupTable[0] = 0; lookupTable[1] = 1;
```

```
    for (i = 2; i <= n; i++)
```

```
        lookupTable[i] = fib(n-1, lookupTable) +
            fib(n-2, lookupTable);
```


Another Optimal:

Fibonacci requires only last 2 places.

lastSecond = 0

last = 1

current = 1

Initiation

```
for (i=2; i<=n; i++) {
```

```
    current = last + lastSecond;
```

```
    lastSecond = last;
```

```
    last = current;
```

```
}
```

```
return current;
```

reducing Space.

$O(N) = 1$

$O(1) = S.$

0/1 KNAPSACK PROBLEM.

Q. Given two Integer Arrays of weights and profits, implement a knapsack() func, where \rightarrow max wt. cannot exceed "capacity weight" but it should have max. profit.

Each item can be selected only once. either take it/skip it.

A: Example:-