

$\Rightarrow$  definition (work some as function)

Ex:- `def name(m, n)`

$\equiv$

$\Rightarrow \text{range}(l, m+1)$

take value from l to m

Mainly used in loop

$\Rightarrow$  append :- to add content. Ex:- `em.append(i)`

$\Rightarrow$  if `in` in stringname :- to find word in string

$\Rightarrow$  `LF = []` to declare a list

`min(n, m)`

$\Rightarrow$  `from filename import *` (to include already written code)

- to find type of string type (variable name)
- to denote last element in array use  $(-1)$  element
- `len(variable)` to get string length
- output of `s[3:1]` is empty because Python not work in backward not give error only empty string
- Output ~~not~~ if max. no. is out of range then it output the upto string & don't give error `s='Hello' S[1:15]` output is Hello
- cannot replace a single character in string by `S[3] = 'P'`
- String are immutable mean cannot change the <sup>character</sup> value of string

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- ⇒ List :- List can include no, character strings or value B can be updated.
- List is Mutable

- Nested List :- List → value contain list.

nest = [1, [2, 3], 4, ["hello"]]

nest[2][0][2] is "1"

nest[0][1:2] is [1, 2]

x = 5      [immutable]

y = x

x ≠ y

here y value 5 because assignment operator creates a copy of a value

- in list if copy or list B change in A list will effect change in list list

- if copy or list then A name Point to one list

- in slice create a new list with some value of list.

$l[:]$  is equivalent to  $l[0: len(l)]$  means full string.

so  $l = [1, 2, 3, 4, 5]$

$k = l[:]$

$l[2] = 8$

then change in l does not effect k list.

- $x = y$  check x by has same value or not  
but " $x = y$ " check x by point to same object or not

list 1 = [1, 2]

list 2 = [1, 2]

list 3 = list 2

list 1 == list 2 True

list 2 == list 3 True

list 1 is list 2 False

list 2 is list 3 True

- we can combine 2 list

list 3 = list 1 + list 2

- list 1 = [1, 2]

list 2 = list 1

list 1 = list 1 + [3]

then list 1 & list 2 no longer Point to same object

⇒ control flow :-

def f1():

def f2():

def f3():

def f4():

def f5():

def f6():

def f7():

def f8():

def f9():

def f10():

def f11():

def f12():

def f13():

def f14():

def f15():

def f16():

def f17():

def f18():

def f19():

def f20():

def f21():

def f22():

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def f170():

def f171():

def f172():

def f173():

def f174():

def f175():

def f176():

def f177():

def f178():

def f179():

def f180():

def f181():

def f182():

def f183():

def f184():

def f185():

def f186():

here execution start from statement 1 from top to bottom but definition of function does not execute it only execute on its use

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- Control flow :- order in which statements are executed

- i) Condition execution

- ii) Repeated execution - loops

- iii) Function definitions



- i) Condition execution :-

- Nonempty value 0, empty sequence "", [ ] is treated as false otherwise will be True

- here for use else & if use at bottom then "elif" is used

- ii) Repeated execution :-

for i in [1, 2, ..., n]:      for fix no. of sequence

while condition :      for don't know no. of sequence

- iii) Function definition :-

- If mutable value (list) is pass as argument then original value is changed.

- Define a function before it invoked

11-09-19

- range()

- range(i, j) mean sequence from i, i+1, i+2, ..., j-1.

ex:- y = range(0, 9) mean 0 to 9 element filled with value 0 to 9.  $y[0]=0$  &  $y[9]=9$

- range(j) mean sequence from zero to  $j-1$

- range(i, j, k) here sequence from i to  $j$  step increment is k ex:- i, i+k, i+2k, ..., i+nk

- Default value of k is 1

$$i+nk < j \leq i+(n+1)k$$

- range(i, j, -1) :- here i > j produce i-1, i-2, ..., j+1

- here step increment i does not cross j

- if k = +ve & i >= j, empty sequence because of infinite loop

- if k = -ve stop "before" j

exp:- range(12, 1, -3) produce 12, 9, 6, 3

- range(0, len(l)) produce fill all element of array l  
value of l  $\Rightarrow$  0 to len(l)-1      range = 0 to len(l)-1

- range(0, 10) is not equal to [0, 1, 2, 3, 4, 5, 6, 7, 8, 9] for python 3  
for python 2 it is true

- list(range(0, 5)) convert range into list

- list l = [1, 3, 5, 6] if we want new list with [2]=5 then

\* list2 = list1[0:2] + [5] + list1[3:]      here len(2) list point to same list

- `list.append(value)` :- to add new element at end (one value at time)
  - `list.extend(list2)` :- to ~~add~~ combine 2 lists & store in list1  
list2 is list1 like  $list1 = list1 + list2$
  - `list1.remove(x)` :- ~~remove~~ remove x ~~from~~ value from list.   
give error if no x in list  
if more than one x is there ~~as~~ first  
x is only remove from beginning
  - `list.append(x)` :- here list is object, append() is function,  
list is object or list is value.
  - Slice is part of list  $list[3:0]$
- $list1 = [1, 3, 5, 7]$   
 $list1[2:] = [9, 10, 11]$   
 $list1 = [1, 3, 9, 10, 11]$   
 $list1[0:2] = [7]$   
 $list1 = [7, 9, 10, 11]$
- if  $x$  in  $l$  :- here if  $x$  element present in list  $l$  then  $l$ .  
✓  $l.remove(x)$  part is executed
- remove  $x$
- while  $x$  in  $l$ :
- remove  $x$
- `l.reverse()` :- reverse ~~elements~~ sequence of list  $list[5:10] = list$
- $l.sort()$  :- sort the list
- `l.index(x)` :- find leftmost position of  $x$  in list  $l$  & give error if no  $x$  in list

- `l.rindex(x)` :- find the rightmost position of value  $x$  in list
- $y=x+y$ , gives error if  $x$  is not assign mentioned
- for iterate over variable as list  $l$  or  $l[1:5]$
- loop :- a set of statements repeated until certain condition is met
- for :- if we know the no. of iteration or used or list
- while :- if we don't know no. of iteration
- $(found, i) = (False, 0)$  here found = "False" &  $i=0$
- if condition1  $\mid$  condition2 condition :-  $\Rightarrow$  condition combine by and

if cond1:  
 $\equiv$   
else:  $\equiv$  or  $\equiv$

- $\Rightarrow$  array :- insertion & deletion is difficult in array as we shift every element below it
- List :- list elements are pointed to its next element not in next memory location so insertion & deletion is easy
- plumbing mean :- add new element in list  $\square \rightarrow \square$
- exchange value is easy in array but difficult in list
- deleting  $\square \rightarrow \square$  list  $\square \rightarrow \square$  array

• time taken by sequential search from start to end time taken by array & list is same for unsorted

- (for sorted) :- go to middle value if find stop search if value is less than then search left side otherwise right side it is called binary search
- only work on arrays

⇒ efficiency :-

\* for Linear Scan is  $O(n)$  for arrays & lists

\* for Binary Search is  $O(\log n)$  for sorted arrays

\* Python can do about  $10^7$  steps in a second

\*

⇒ sorting & sorting :-

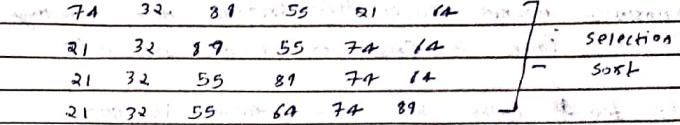
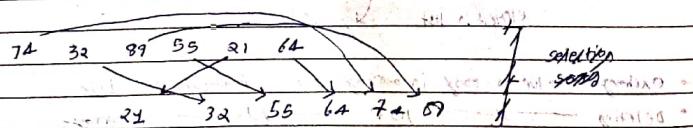
Searching of a value :- unsorted array  $\rightarrow O(n)$

sorted array  $\rightarrow O(\log n)$

by sorting :-

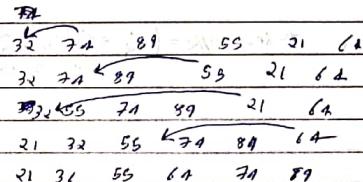
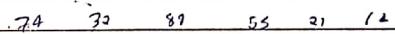
i) easy to find middle value

ii) also check duplicate



Selection sort  $\rightarrow O(n^2) \rightarrow n + (n-1) + (n-2) \dots 1 \rightarrow n^2 + n = n^2$

$\Rightarrow$  insertion sort  $\rightarrow$



Insertion sort  $\rightarrow 1 + 2 + \dots + (n-1) = n(n-1)/2 = O(n^2)$

$\Rightarrow$  recursive

def sum(l, r):

if l == r:

return l

else:

return l + sum(l+1, r)

or

return l + sum(l+1, r)

insertions :- maximum recursion depth exceeded. Mean how many process is pause in execution like out of memory

\* by default Python has maximum recursion is 998

to set recursion limit :-

```
import sys  
sys.setrecursionlimit(100000)
```

→ Merge sort :-

Sort A[0:n/2] # here  $n/2 = \lfloor \frac{n}{2} \rfloor$  take only integer part.

Sort A[n/2:n]

43, 32, 22, 78

[43|32]

[32|22]

[22|78]

[32|32|22|78]

i) if any list is empty .

ii) if non empty & len >= 1

iii) if no, write A in Post

L = A[0:n/2]

R = A[n/2:n]

if not function (A,L,R)

if (L|R <= 1)

+ return (A[L:R])

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def mergesort (A[], L, R)

if (R - L > 1)

Mid = (L+R)/2

L = function (A, L, Mid)

R = function (A, Mid, R)

return (Merge(L, R))

def merge (A, B):

(C, m, n) = (C[], len(A), len(B))

(i, j, k) = (0, 0, 0)

while i+j < m+n:

if i == m:

C.append (B[j])

j=j+1

if j == n:

C.append (A[i])

i=i+1

if B[j] <= A[i]:

C.append (B[j])

j=j+1

if A[i] >= B[j]:

C.append (A[i])

i=i+1

if m > i:

{arr[x] = B[i]}

j=j+1

k=k+1

if A[i] >= B[j]:

C.append (B[j])

j=j+1

i=i+1

if n > j:

{arr[y] = B[i]}

j=j+1

k=k+1

return (C)

efficiency of Merge sort is  $T(n) = O(n \log n)$

④ ⇒ Quick sort

i) choose a Pivot element

ii) divide list w.r.t. to Pivot in 2 halves

worst case is already sorted array that is  $O(n^2)$

average case is  $O(n \log n)$

• Quick sort is not stable because swap operation during partitioning distorts original order

• Merge sort is stable (prove)

i) do not allow element from right to overtake element from left

⇒ Tuples :-

ex:- (age, Name) = (23, "Hari") or (date, (16, 7, 2013))

Tuples are immutable means do not change its value

→ Dictionaries

test1["Dhawal"] = 82, test1["Kohli"] = 56

Dictionaries = {} empty  
TUPLE = () empty tuple

key can be any immutable variable - int, float, bool, str, tuple  
key can not be list or dictionaries

(name) == (name)

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### Dictionaries or Dictionaries

Score ["Test1"] ["Dhawal"] = 82 ⇒ List

Score = { "Test1": { "Dhawal": 82, 33 } ⇒ Dictionaries

first declare outer Dictionaries then inner dict. then assign value score = { 3 }

Score ["Test1"] = { 3 }

Score ["Test1"] ["Dhawal"] = 76

• d.keys() return sequence of key

for k in d.keys():

• to get sorted key -

for k in sorted (d.keys()):

• d.values() is sequence of values in d total = 0

d[0] = 7 → Mean d = [0:73]

• Dictionary allow a flexible association of values to keys

→ function → coordinates (coordinates of a point)

⇒ Default argument

Def → func (a=10, b=85) here if we call func with no parameter then 10 & 85 is used

def f(a, b, c):

g = f  
g(a, b, c) x here f function name is g

→ list comprehension :-

→ filter (P, l) where P is property & l is list  
here filter return sublist of element that satisfied the P property of list l.

→ list(map(f, l)) → here give f as function & l is list  
here each element goes to function as add to result list but map(f, l) does not give list

ex:- [score(i) for i in range(100) if iseven(i)]  
↓ ↓ ↓  
map() generator filter  
Power even value get 0 - 100 get even value

for Pythagorean theorem

[(x, y, z) for x in range(100) for y in range(100) for z in range(100)  
if  $x^2 + y^2 = z^2$ ]

for prime numbers (i=1 to 100) print i if

→ exception handling

Syntax → try:

except IndexError:

action on error

except (NameError, KeyError):

action for all listed errors

except:

action for all errors

else:

odd s

→ Scores[6] = append(s)

except KeyError:

Scores[6] = s

→ Type of error :-

i) SyntaxError

ii) NameError

iii) ZeroDivisionError

iv) IndexError

v) ValueError → when convert english alphabet into integer

→ for, avoid "\n" on print or end use "end=\\"

Print("I am superman", end="")

Print(Round (2/7, 5)) // Print result of Date upto 5 decimal places.

→ on use print("a", "b") there is space b/w a & b to avoid space use sep=" " as end  
ex:- print("a", "b", sep="")

⇒ print("a") is used in Python 2  
but print("a") is used in Python 3

⇒ storage file is read is more slowly than memory

⇒ file handling

hvariable = open("path", "mode") → to connect to file  
string variable = hvariable.read() → to read file (whole)

string variable = hvariable.readline() → to read first line

string variable = hvariable.readlines() → to get list which have each element as line of file

hvariable.seek(n) → Move pointer to n position  
hvariable.read(ya) → to read fixed block of fixed length character

For nextline use "in" &  
externally hvariable.write(s) → to write into file  
in writing hvariable.writeline(s) → to list as write each element by one by one

hvariable.close() → flush output buffer  
hvariable.flush() → write data to file upto now

• whitespace characters which not display like tab, spaces, in

to remove whitespace in string use string.rstrip()  
also remove from left side string.lstrip() & for both sides use string.strip()

s.find(pattern) → return starting position of pattern occur in s  
if not found return -1

s.find(pattern, start, end) → same as above but find in range (start, end)

s.index(pattern) or s.index(pattern, left, right) → same as above but raise ValueError if pattern not find

s.replace(f1, f2) → same as above but replace first f1 by f2 in s

s.replace(f1, f2, n) → same as above but replace first n f1

→ jointString = ""  
(string jointString=join(list), → join list by given character set)

→ convert all to lowercase, both in s & c

s.upper() ⇒ convert all to uppercase

s.title() ⇒ convert first letter in each word to uppercase

s.capitalize() ⇒

s.isalpha() ⇒ return bool. on if s has alphabets or not

s.isnumeric() ⇒ return bool. on if s has numeric ->

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### → format()

a = "first {} second {}" .format(42, 2)

→ pass → to do Nothing

→ function related to list :-

i) sum(list)

ii) list1.extend(list2)

iii) list1.count(x) count occurrence of x

iv) list.index(a) get index

v) min(list) / max(list)

vi) sorted(list)

vii) list.sort(reverse=True)

viii) list.pop(i) delete element by index if not given remove last one

ix) del list[i] delete element by index

x) list.remove(o) → value

→ backtracking

→ for N queen Problem on chess for NxN board

Place N queen as none of them is attack on each other

• 2, 3 is not possible

if so not found go back to previous step & change it

→ for global scope :: - declared → value not needed

global x

→ Set

colours = set()

colours = {red, 'black', 'red', 'green'}

print(colours)

{'black', 'red', 'green'}

• 'black' in colours

this give true

• convert list into set

no. = set([1, 3, 3, 4, 1])

set has {1, 3, 3, 4}

• letters = set('banana')

letters has {n, a, n, s, b, a}

• odd = set([1, 3, 5])

odd prime = set([2, 3, 5])

→ Union | Prime → {1, 2, 3, 5}

⇒ Union

3 ⇒ for intersection

odd - prime → difference

odd prime → {1, 3}

→ exclusive (remove common items)

→ Queue are useful for breadth-first exploration.

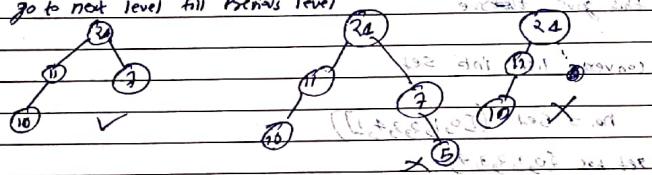
→ Priority Queue :-

delete\_max() → delete of high priority

insert() → insert, maintain heap property

→ balanced tree is called heap

- Parent value is greater than child
- child filled from left to right position is 'sym'.
- before go to next level fill previous level



heap insert take  $O(\log N)$  time

for delete swap Move up to bottom i.e. insert

Behaviors of stack  $\Rightarrow$  Push() & Pop()

queue  $\Rightarrow$  add() & remove()

Heap  $\Rightarrow$  insert() & delete\_max()

- if Heap implement list this not mean `list.append(s)` is legal in Heap

⇒ class :-

int a

every function in class  $\rightarrow$  def classfunctionname(self, Parameter):  
self.a = 10

constructor  $\rightarrow$  \_\_init\_\_( )

•  $\rightarrow$  \_\_str\_\_( )  $\rightarrow$  return str object value

`str(s) = s.__str__( )` both are same

- $P1 + P2 = P1.__add__(P2)$  it will return both value
- $P1 - P2 = P1.__sub__(P2)$

• def append(self, v):

if self.isempty():

self.value = v

elif self.next == None:

newnode = Node(v)

self.next = newnode

else:

self.next.append(v)

return()

⇒ Dictionary :-

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- to get all keys ~~value~~ as list then dict.keys()
  - ———) — values ———) — dict.values()
  - ———) — keys & values as pair then dict.items()
  - for x in my\_dict:
    - Print(x) // Print all keys
  - for y in my\_dict.values():
    - Print(y) // Print all values
  - for z in my\_dict.items():
    - Print(x, ":", y) // Print key-value pair
  - My\_dict ['Dave'] = '004' // if Dave already exist then update  
otherwise create key-value pair
  - del my\_dict ['Dave'] // remove key-pair value
  - my\_dict.pop('Dave') // remove 'Dave' value
  - my\_dict.popitem() // remove last inserted item

$\Rightarrow$  Pointer in C / C++ :-

```

• int Var = 20;
int *IP;
IP = &Var;
cout << &Var;           /* Print address of var variable */
cout << IP;            /* Print address contain by IP = Point */
cout << *IP;           /* Value store in address SPote */
*IP = NULL;
cout << IP;            /* Point "0" by default.

```

\***Ptr** :- Point that contain address of variable.  
\*\***ptr** :- Point that contain address of another Pointer.

- |                            | no. | point. |
|----------------------------|-----|--------|
| $*(\& \text{nums}[0])$     | 16  |        |
| $*(\& \text{nums}+1)$      | 18  |        |
| $*(\& \text{nums}+2)$      | 20  |        |
| $(*(\& \text{num}+1))$     | 25  |        |
| $(*(\& (\text{num}+1)+1))$ | 26  |        |
| $(*(\& (\text{num}+1)+2))$ | 27  |        |

#### • Pointer in linked list :-

```

node_t * head = NULL; // Create node point of LL

head = (node_t*) malloc (sizeof (node_t)); // init address
head->val = 1; // Put value in first list

head->next = (node_t*) malloc (sizeof (node_t)); // another list
head->next->val = 2; // Put value in second list
head->next->next = NULL; // next point to second set to NULL

```

- To know the height of the binary Tree :-

Total Height (Current Node)

if ( $\text{CurrentNode} \rightarrow \text{Next} == \text{NULL}$ )

{ return -1; }

~~Left~~ Node = & Height (currentNode)

~~Print~~ Node = Height (current Node  $\rightarrow$  RNode)

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Q getuse ( $\max(LNode, RNode) + \pm$ )  
}

⇒ Quicksorting algo :-

QVICK (List)

```
if (len(List) > 1)
{ P = List [len(List) - 1]
  for i in List:
    if (i > P):
      RList = RList + [i]
    else:
      LList = LList + [i]
```

Rfinlist = QVICK (RList)

Lfinlist = QVICK (LList)

finlist = Rfinlist + [P] + Lfinlist

return (finlist)

}

⇒ Mergesorting algo :-

Date \_\_\_\_\_

### APPTITUDE :-

Date \_\_\_\_\_

x is  $\frac{1}{2}$

30% =

- i) Divisibility      2) HCF & LCM      3) Profit & Loss
- 4) Simple & compound interest      5) Time, Speed & Distance
- 6) Inverse      7) Log      8) Permutation & combination
- 9) Probability

#### 1) Divisibility :-

2 → last digit is 0, 2, 4, 6, 8

3 → sum of digit is divisible by 3

4 → last 2 digit is divisible by 4

5 → last digit is 0 or 5

6 → if given no. divisible by 2 & 3

7 → subtracting the twice the last digit with remaining digit give a multiple of 7

Ex:- 658 →  $65 - 2 \times 8 = 49$  so 658 is divisible by 7

8 → last 3 digit is divisible by 8

9 → if sum of digit is divisible by 9

10 → if last digit is 0

11 → difference bet the sum of alternate digit is multiple of 11

Ex:- 2343  $\Rightarrow (2+4) - 3+3 = 0$ , so 2343 is divisible of 11

12 → if given no. is divisible by 3 & 4

- if we divisible a given no. by even no. then answer is also a even no.

13 → 4 time the unit digit is odd to remaining no. (remove unit digit) is divisible by 13

15238  $\Rightarrow 2 \times 8 + 1523 + 4 \times 8 = 1555 \Rightarrow 155 + 4 \times 5 = 175 \Rightarrow 17 + 4 \times 5 = 67$   
67 is not divisible by 13

Date.....

$17 \Rightarrow$  5 times the last digit subtract from remaining no. is divisible by 17  
 $2278 \Rightarrow 227 - 8 \times 5 = 187$ , so 2278 is divisible by 17

$19 \Rightarrow$  same case as 13 but 2 times adding so resultant is divisible by 19

$\Rightarrow 28 / HCF / LCM$

$23 \Rightarrow$  same case as 13 but 7 times.

$29 \Rightarrow$  same case as 28 but 3 times

$\Rightarrow HCF \& LCM$

$$\text{Product of 2 NO.} = LCM \times HCF$$

$$LCM \& HCF (20, 30) = 60$$

$$HCF \text{ of } (20, 30) = 10$$

Date.....

SQL languages :-

i) 4 type of languages :-

i) Data Definition Languages :- define the database  
Example:- Create, Drop, Alter, Truncate

ii) Data Manipulation languages :- Manipulate the data present in DB  
Example :- Insert, Update, Delete

iii) Data Control languages :- deal with user Permission & control of DB system. Example:- Grant, Revoke

iv) Transaction Control languages :- deal with the transaction of DB  
Example:- Commit, Rollback

v) DDL/DML languages :- used to fetch data from DB  
Example :- Select

DDL ——————  
a) Comment -- (select all (single line)) /\* \*/ (multiple line comment)

3) Create Table Name AS Select C<sub>1</sub>, C<sub>2</sub> ... C<sub>n</sub> From Emp

4) Truncate table Name; // delete all info. of table but not table

5) Alter table Name ADD C<sub>1</sub> Int; // use to add, deletes, modify tokens  
A —> ———— ALTER C<sub>1</sub> Chg;

6) type of keys :- there are 7 type of key

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i) **Candidate Key** :- all those columns whose tuple are uniquely identify. Mean. all the eligible Primary Keys are Candidate Keys. A table has more than 1 candidate key but Primary key is one.

ii) **Primary Key** :- all tuple of Primary key uniquely identify only one column is known as Primary key.

iii) **Super Key** :- set of all attributes that uniquely identify one tuple. include candidate key, Primary Key & Unique key.

iv) **Unique Key** :- Candidate key column - Primary key column + all attributes that uniquely identify except Primary key + allow one NULL value in column.

v) **Alternate Key** :- Candidate key - Primary key, all candidate key (except Primary key) & not allow NULL value.

vi) **foreign key** :- column that are common in tables

vii) **Composite key** :- combination of 2 or more columns to identify each column uniquely.

~~3) NEED NOT~~

7) **Constraints** :- these are used to specify the rules for data in a table

i) **NOT NULL** :- other column does not contain null value.  
Ex:- Create Table ht (P\_id int NOT NULL);

ii) **UNIQUE** :- ensure all value of column is unique.  
Ex:- Create Table ht (E\_id int NOT NULL UNIQUE);

iii) **CHECK** :- all value in column satisfy a specific condition.  
Ex:- Create Table ht (E\_id tinyint(2) CHECK (Country='India'));

iv) **Default** :- if no value is given then Default value is stored.  
Ex:- Create Table ht (E\_id tinyint(2) DEFAULT 'India');

v) **INDEX** :- This constraints is used to create index in table, by which we can easily create & retrieve data from DB.  
Ex:- CREATE INDEX <sup>id</sup> ON TableName (column, c2);

DDL \_\_\_\_\_

DML \_\_\_\_\_

8) **USE** :- Select the DB on which work is done.  
Ex:- USE DBName;

9) **INSERT INTO** :- insert new record in the table.  
Ex:- Insert INTO Tablename (c1, c2, c3) Value (v1, v2, v3);

10) **UPDATE** :- Modify the already present record.  
Ex:- UPDATE Tablename SET C1 = V2, C2 = V3 WHERE condition;

11) **DELETE** :- delete the existing record.  
Ex:- DELETE FROM Tablename WHERE E\_ID = 2;

12) **DEVELOP SELECT :-** Select data from DB & store in a table  
Known as result-set

Ex:- `SELECT E-ID, E-Name FROM E-Info;`

There are 5 keyword which are used with `SELECT`

i) `SELECT DISTINCT C1, C2 FROM Tname;`

`SELECT DISTINCT` :- return only different value

ii) `ORDER ORDER BY :-` used to sort the required result.

by default it is ascending ascending order. used ASC & DESC

Ex:- `SELECT * FROM E-Info ORDER BY E-ID;`  
ASC/DESC;

iii) `GROUP BY :-` Used to group result set with aggregate function.

Ex:- `SELECT COUNT(E-ID), city FROM E-Info GROUP BY city;`

iv) "Having" :- Having used where "WHERE" is not used.

Ex:- `SELECT COUNT(E-ID), city FROM E-Info GROUP BY city;  
HAVING COUNT(E-ID) > 2 ORDER BY COUNT(E-ID) DESC;`

v) `SELECT INTO :-` used to copy data one table to another.

Ex:- `SELECT * INTO E-backup FROM Emp;`

New table [in external DB]

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13) Operators in SQL:-

i) Arithmetic Operators :- % (Modulus for remainder), /, \*, +, -

ii) Bitwise Operators :- i) " $\wedge$ " (XOR operation) ( $A\bar{B} + \bar{A}B$ ),  $A \wedge B$

ii) " $\vee$ " (OR operation)  $A \vee B$

iii) " $\wedge\wedge$ " (AND operation),  $A \wedge\wedge B$

iii) Comparison Operators :- i)  $\neq$  (not equal to),  $A \neq B$   
/ Relational operators ii)  $<$  (less than or equal to),  $A <= B$

iii)  $>$  (greater than or equal to),  $A >= B$

iv)  $<$  (less than),  $A < B$

v)  $>$  (greater than)  $A > B$

vi)  $=$  (equal to),  $A = B$

iv) Compound Operators :- i) " $\mid\ast$ " (Bitwise OR equals),  $A \mid\ast = B$   
/ Assignment operator ii)  $\wedge\wedge =$  (Bitwise Exclusive equals),  $A \wedge\wedge = B$

iii)  $\wedge\wedge$  (Bitwise AND equals),  $A \wedge\wedge = B$

iv)  $\% =$  (Modulo equals),  $A \% = B$

v)  $/ =$  (divide equals),  $A / = B$

vi)  $\ast =$  (multiply equals),  $A \ast = B$

vii)  $- =$  (subtract equals),  $A - = B$

viii)  $+ =$  (Add equals),  $A + = B$

v) Logical operators :- ~~AND, OR, NOT, BETWEEN, LIKE, IN, EXISTS, ALL, ANY.~~

a) AND Operator :- filters records that fulfill ~~more than~~ required condition.

Ex:- `SELECT * FROM 'E-Table' WHERE city = "Mumbai" AND category = "ST";`

b) OR Operator :- filter record which satisfied any condition

ex:- SELECT \* FROM E-Table WHERE City = "Goa" OR

City = "Mumbai";

c) NOT Operator :- display records which does not satisfy the condition

ex:- SELECT \* FROM E-Table WHERE NOT City = "Goa";

d) ALL combine query :- SELECT \* FROM E-Table WHERE  
NOT City = "Goa" AND (City = "Mumbai" OR  
City = "Bangalore");

e) BETWEEN Operator :- used when you want to select value within a range including both values

ex:- SELECT \* FROM E-Table WHERE Salary BETWEEN  
40000 AND 50000;

f) IN Operator :- used for multiple OR condition

ex:- SELECT \* FROM E-Table WHERE IN City IN  
('Mumbai', 'Goa', 'Bangalore');

g) EXISTS Operator :- test if a record exist or not.

ex:- SELECT City FROM E-Table WHERE EXISTS (SELECT City FROM E-Table WHERE E-ID = 05 AND

City = 'Kolkata');

h) ANY / ALL Operator :- used with WHERE OR HAVING clause  
return true if all of the subquery values meet the condition.

ex:- SELECT E-Name FROM E-Table WHERE E-ID = ANY  
(SELECT E-ID FROM E-Table WHERE City = 'Goa');

i) Aggregate function :-

i) MIN() :- return smallest value of selected column.

ex:- SELECT MIN(E-ID) FROM E-Table;

ii) MAX() :- return largest value of selected column

iii) COUNT() :- return no. of rows which match a specified criteria

ex:- SELECT COUNT(E-ID) FROM E-Table;

iv) SUM() :- sum of a numeric column that you can

ex:- SELECT SUM(SALARY) FROM E-Table;

v) AVG() :- return the average value of a numeric column

ex:- SELECT AVG(Salary) FROM E\_Table;

15) NULL Function :- return an alternative value if an expression is NULL. ISNULL()

ex:- SELECT E\_ID + (Month\_Year\_of\_Salary + ISNULL(Salary, 0))  
FROM E\_Table;

16) Aliases :- Aliases give a column / table a temporary name & only exists for a duration of query

ex:- SELECT E\_ID AS ID, E\_Name AS Name  
FROM E\_Table;

17) Case :- Go through all condition & return a value when first condition is met. So if no condition is TRUE then return else part. If there is no else part then return "NULL"

ex:- select E\_Name, city FROM E\_Table ORDER BY (CASE WHEN city IS NULL THEN 'Country is India by Default'  
ELSE city  
END);

18) JOIN :- Used to combine two or more tables on bases on related column  
There are of 4 types:

i) INNER JOIN :- Join all those records which (A & B) have matching values in both table

Ex:- select Tech.TechID, E\_Table.E\_Name from Tech INNER JOIN E\_Table ON Tech.E\_ID = E\_Table.E\_ID;

ii) FULL JOIN (A+B) :- Join all records which either have a match in left or right table

Ex:- select E\_Table.E\_Name, Tech.TechID From E\_Table  
FULL OUTER JOIN Tech ON Tech.E\_ID = E\_Table.E\_ID;

iii) LEFT JOIN (A+A+B) :- return all record from left table along with record satisfied condition from right table

iv) RIGHT JOIN (A+B+B) :- return all records from right table along with record satisfied condition from left table

19) Set operation :- UNION, INTERSECT, EXCEPT

i) UNION :- Used to combine the result set of a or More statement.

Ex:- select ColumnName from Table1  
UNION  
select ColumnName from Table2

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ii) INTERSECT :- used to show common in two result set.

Ex:- select C<sub>1</sub>, C<sub>2</sub> from Tname where condition  
INTERSECT  
select C<sub>1</sub>, C<sub>2</sub> from Tname1 where condition

iii) Except :- return tuple that present in ~~first~~  
first select statement except those who are present in second select statement

Ex:- SELECT Coloum FROM Tname1  
EXCEPT  
SELECT Coloum FROM Tname2

20) Dates :- follow datatype store date/time value in DB  
i) DATE :- YYYY-MM-DD

ii) DATETIME :- YYYY-MM-DD HH:MI:SS

iii) TIMESTAMP

iv) TIMESTAMP :- Unique number

Ex:- select \* from Tname where J-Date='2019-04-05';

21) Auto Increment :- generate or new unique no. automatically when new record is inserted  
IDENTITY keyword is used.

Ex:- CREATE TABLE Employee\_info\_table (E-ID INT  
IDENTITY(1,1) PRIMARY KEY, E-Name VARCHAR  
(255) NOT NULL);

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22) GRANT :- Provide access or Privileges on the DB & its object to users

Ex:- GRANT PrivilegesName ON ObjectName To { Username | Public | RoleName } [With GRANT Option];

i) PrivilegesName :- Name of the right/access/privilege grant to user

ii) ObjectName :- Name of DB object like Table/View/STORED:

iii) Username :- Name of user who give access

iv) Public :- To grant access to all users

v) RoleName :- Name of a set of privileges grouped together

vi) With GRANT option :- option of Grant

Ex:- GRANT SELECT ON E-Table To User1;

23) REVOKE :- withdraw the user's access Privileges

Ex:- REVOKE PrivilegesNames ON ObjectName  
FROM { Username | PUBLIC | RoleName };

24) Views :- it is a single table which is derived from other tables. So it has field from one

i) CREATE VIEW :- Create a view from table

Ex:- SELECT E\_Name, E\_Phone FROM E\_Table WHERE

City = "Gurgaon";

ii) "CREATE OR REPLACE VIEW" :- used to update a view

Ex:- CREATE VIEW OR REPLACE [View1] AS

SELECT E\_Name, E\_Phone FROM E\_Table WHERE  
City = "Mumbai";

iii) "DROP" View :- used to delete a view

Ex:- DROP VIEW [View1]

25) Stored Procedures :- code which you save & reuse it again.

Ex:- CREATE PROCEDURE [Procedure Name] AS  
GO;

26) Triggers :- Triggers are set of SQL Statement which are stored in DB catalog. the statement execute when an event occur on table so, a trigger can be invoked either BEFORE or AFTER the Data changed by Insert, Update or Delete Statement.

before insert → after insert → before update → after update →  
after delete ← before delete ←

Ex:- CREATE TRIGGER [TriggerName] [Before | After]  
{ INSERT | UPDATE | DELETE } ON [TableName]  
[FOR EACH ROW] [Trigger Body]

DCL

TCL

27) COMMIT :- used to save the transaction into the DB  
Ex:- COMMIT;

28) ROLLBACK :- used to restore the DB to the last committed state.

Ex:- ROLLBACK;

Ex:- ROLL BACK TO SavePoint Name;

29) SAVEPOINT :- used to temporarily save a transaction. or if we wish to rollback to any point the save that point.

Ex:- SAVEPOINT SavePointName;

Ex:- INSERT INTO E\_Table VALUES (05, 'Anushka');  
COMMIT;

UPDATE E\_Table SET Name = 'Akash' WHERE id = '05';

SavePoint S1;

INSERT INTO E\_Table Value (08, 'veena');

SELECT \* FROM E\_Table;

ROLLBACK TO S1;

TCL

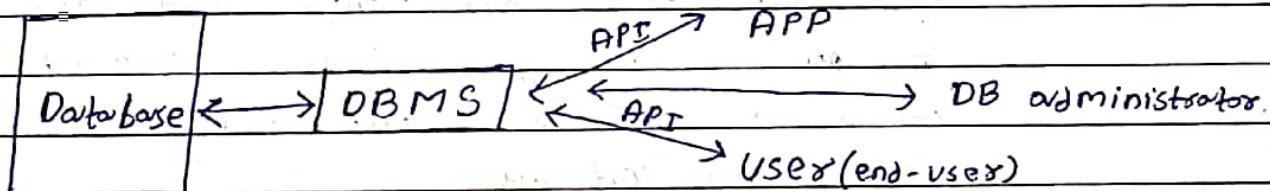
# Data Base Management System.

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⇒ Basic Concept, three-tier architecture, ~~ER Model~~, functional dependencies, normalization, relation Model, stored Procedure & triggers, locking, Clustering

DBMS :- i) DBMS is a system software for creating & managing DB.

ii) DBMS provide a systematic way to handle DB by the request provided through APP API.



Feature of DBMS :-

i) Data Availability :- Data is ready to available whenever required by the user

ii) Data integrity & security :- Protect the existence of DB and maintain the quality of database.

iii) Data independency :- 2 type - Physical data independence  
logical data independence

a) Physical data independence :- Which remain unaffected from changes in storage structures or access method

b) logical data independence :- Which remain unaffected from change in schema

- ⇒ ACID Properties :- (Atomicity, Consistency, Isolation, Durability)  $\Rightarrow$  Characteristics of DBMS :-
- i) Atomicity :- There is only one state for a transaction either commits or abort. If we update a database the all update will commits or none of update commits (All or nothing rule)
  - ii) Consistency :- It refers to the correctness of DB. If we transfer money from one account to other then after deduct money from first account system goes down as amount does not credit in second account then this is called inconsistent DB.
  - iii) Isolation :- This Property ensures that multiple transactions can occur concurrently without leading to inconsistency of DB state.
  - iv) Durability :- This Property ensures that if a transaction is committed then it must be updated in DB even when system or the storage media fail.
- Recovery to the most recent successful commit
- Database :- i) Database is a collection of related data
- ii) data which represent some aspect of real world
  - iii) A DB is designed to be built & populated with data for a specific type
- i) Provides security & removes redundancy
  - ii) Support of multiple views of data
  - iii) Follow ACID Property
  - iv) Supports multi-user environment
- Ex of popular DBMS software :- MySQL, SQLite, dBASE
- Type of DBMS System :-
- i) Hierarchical DBMS :- a) organised in a tree-like structure
  - b) Data is represented using Parent-child relationship
  - c) Parent have many child but child have only one Parent
  - ii) Network Model :- a) allow all child to have multiple parents. (Many-to-many relation)
  - b) Used to address more complex relationship
  - iii) Relation Model :- a) Most widely used DBMS Model
  - b) based on normalizing data in rows & columns.
  - c) having fixed structure & manipulated using SQL
  - IV) Object-oriented Model :- a) stored data in form of object

⇒ 3-Tier architecture :- (Presentation layer, Application layer, Database) ⇒ Component of ER Diagram :-

HTML, CSS, JS   Java, net, Python, C++   MySQL, Oracle  
SQL server

\* A 3-Tier application architecture is client-server architecture consists of 3 tiers. Data store at Data layer, logic handles in application layer, GUI in Presentation layer.

- benefit :- i) each tier can be developed concurrently without affecting other tier.
- ii) easier for enterprise to continuously evolve on application as new need.

⇒ ER Model :- Entity-relationship Model is a design or blue print of DB that can later be implemented on a DB.

ii) 2 component - entity set & relationship set

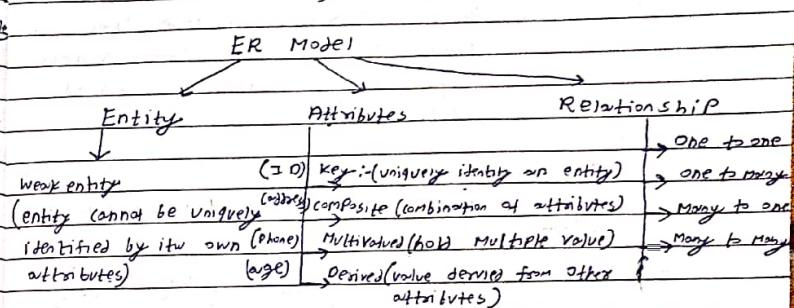
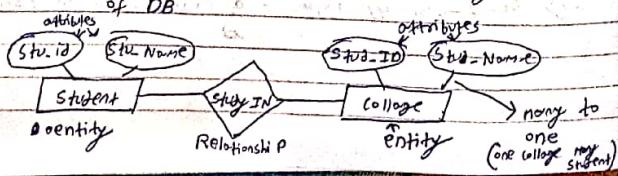
iii) ER Model describe the structure of DB with the help of diagram known as

Entity-relationship diagrams (ER diagram)

⇒ ER Diagram :- i) an entity set is attribute of table in DB

ii) by showing relation among tables & their attributes, ER diagram shows logical structure of DB

Example :-



⇒ functional dependency :- it is a relationship between attributes typically between the Primary key & other non-key attributes within a table.

determinant dependent  
 $X \rightarrow Y$

SIN → Name, Address, DOB

[with SIN we determine Name, Address, DOB]

⇒ Normalization :- i) it is technique in which we use systematically approach of decomposing table to eliminate redundancy (duplicate)

ii) elimination of useless data & ensure data dependencies make sense i.e. data is logically stored.

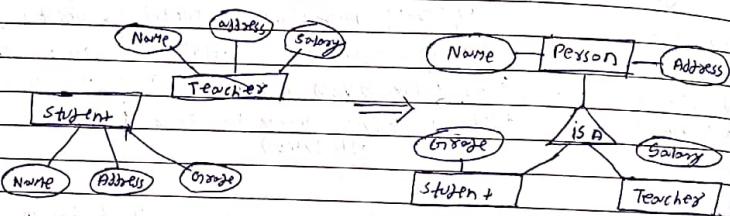
Type of Normal form ⇒ 1NF, 2NF, 3NF & Boyce-Codd NF

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⇒ Generalization :- It is a process in which the common attributes of more than one entity form a new entity & new entity form is called generalized.

before



iv) Boyce-Codd Normal Form :- • must follow 3<sup>rd</sup> NF  
• all table in DB should be only one Primary Key

v) 5<sup>th</sup> NF :-

⇒ Relational Model, Normal form, indexing (B & B+tree)  
⇒ ~~Relational~~, transaction, concurrency, file organisation  
Relational algebra in Relation Model.

⇒ Normal forms :-

- i) 1<sup>st</sup> Normal form :- • no tuple contain 2 value
- ii) 2<sup>nd</sup> Normal form :- • must follow 1<sup>st</sup> Normal form
  - there is No Partial Dependence.
  - if 1 column depend on 2 column  
then it must be separated in 2 different table
- iii) 3<sup>rd</sup> Normal form :- • must follow 2<sup>nd</sup> Normal form
  - there is no transitive dependence
  - if 1 column is derived from 2 columns individually then it must be in 2 different table

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## Struct C ~~Basic~~ Language :-

⇒ Data type in C language :-

range of  
got by (0-79,0-24)

### Data type

Primary Data Type	Derived Data Type	User Defined Data type
→ Void	→ Array	→ Structure
16 bit → int (-32,768 to 32,767)	→ Reference	→ Union
→ long (-2147 Million to 2147 Million)	→ Pointer	→ Enumeration
→ char (0-9 to 18-57 A-Z=65-90)	(A-Z=97-122) (allows -128 to 127)	(-32768 to 32767)
→ float (3.4*10 <sup>-38</sup> to 3.4*10 <sup>38</sup> )		
→ Double (1.7*10 <sup>-38</sup> to 1.7*10 <sup>38</sup> )		
→ short (range same as int)		

• signed :- Means data type act as normal datatype

• unsigned :- Means only → range of no include. (for int unsigned)  
(0 to 65,535)

• Union :- allowing storing various data type in same memory location → P

• Program can define a union with different member or datatype but only → one single member can contain value at a given time.

Ex:- Union Variablename // same as structure definition

```
{ char x;  
  float y;  
 } obj;
```

// here sizeof (obj) => 4 byte as float + byte  
// obj object size is size of largest size variable

• enum or enumeration :- user define data type

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- Enumeration or enum :- • used to assign name to integral constant, the names make a program easy to read & maintain
- keyword 'enum' is used to declare

Example :- #include <stdio.h>

```
enum week { Mon, Tue, wed, Thu, Fri, Sat, Sun };
int main()
{
    enum week day; // week = variable object name
    day = wed;
    printf("%d", day); // output is 2
    return 0;
}
```

- Short :- compact version of int Variable
- range is -32,768 to 32,767
- size is 2 byte for all unsigned, signed

- Pointer :- • it is variable that store address of another variable
- pointer is used to allocate memory dynamically i.e. at runtime

⇒ Variable :- Variable is a name given to a storage location that our program can manipulate

• size of all datatype :-

format specific datatype	size
%d	2 byte (when 16 bit Processor), 4 byte (when 32 bit or 64 bit)
%c	1 byte
%f	4 byte (for 16 bit Processor), 8 byte (when 32 bit, 64 bit)
%f float	4 byte
%lf double	8 byte
Void Pointer	2 byte (16 bit Processor), 4 byte (32 bit), 8 byte (64 bit)
Short	2 byte

• Reference :- When a variable is declared as reference it becomes an alternative name for an existing variable.

• A variable can be declared as reference by putting ' &' in the declaration.

```
int x=10; // 2 variable point to same memory location
int& ref=x; // now if we make any change in any one variable then change is also appear in other
```

⇒ format specifier :- • it is used during input & output • it is a way to tell compiler about the type of does a variable take or point during scan and Point

⇒ Exit(0) ⇒ terminate normally successfully  
exit (-) ⇒ —————— unsuccesfully

⇒ operator in C :-

- i) arithmetic operator :- used in mathematics operation (+,-,\*,/,%,++,-)
- ii) Relational operator :- defined relation between 2 operand ( $=, !=, >, <, \geq, \leq$ ) OR AND NOT
- iii) Logical operator :- used for multiple condition ( $!, \&, \&&$ ) OR AND NOT
- iv) Bitwise operator :- performs manipulation of data at bit level ( $\&, !, ^, \ll, \gg$ )  
Bitwise OR AND Exclusive OR Leftshift Rightshift

v) assign operators :- Used to assigned variable to another variable using other variable or number. ( $=, +=, -=, *=, /=, \%=$ )

vi) conditional operator :-

\* also known as **Ternary Operator**

- \* expression ? expression<sub>1</sub> : expression<sub>2</sub>
- \* same work as if - else but way of writing is different

vii) special operator :- • size of ()

- \*  $\&$   $\Rightarrow$  return address of a variable
- \*  $*$   $\Rightarrow$  point to a variable

loop:- if

$\Rightarrow$  NULL statement :- it is used in loop when there is no statement is loop body.

Ex:- ~~for~~ while ( ; ; for (int i=0; i>2; i++)  
;

loop :- when we have to execute a block of statement for no. of times then we use loop

- while loop (for unknown no. of iteration)
- for loop (for known no. of )
- do while loop (for execute body of loop for at least one time)

$\Rightarrow$  type conversion in C :- conversion of one data type into others.

- \* conversion is done where it is possible (int to char, int to float)

type conversion is of 2 type

- i) implicit type conversion :- perform by compiler automatically
- ii) Explicit type conversion :- perform by user manually

\* When an operation is done between 2 operation then lower preference operators converted to higher preference operator lowest

char  $\rightarrow$  int  $\rightarrow$  long  $\rightarrow$  float  $\rightarrow$  double  $\rightarrow$  long double

$\Rightarrow$  type of variable in C :-

type	storage	initial value	scope	life
auto/automatic (auto)	stack	garbage	within block	end of block
extern/external (extern)	data segment	zero	global	end of program
static (static)	-	zero	within block	-
register(register)	CPU register	garbage	-	end of block
global	data segment	zero	global	end of program
local	stack	garbage	within block	end of block
• Static :- create only one copy if we do call a function multiple time then then it use same variable				

Ex void function()

```
{ static int y=2;  
  y=y+1;  
  printf("%d",y);  
}
```

// if we call this function 3 times then output is 2,3,4  
as y is only one copy

\* Automatic :- all variable declared in a block by default automatic

\* External/global :- declare outside of all the block & can access at any part of program

• Local Variable :- local variable are variable that declare inside a block

- register variable :- it tell the compiler to store the variable in CPU register instead of memory
- frequently used variable are kept in register because of fast accessibility
- we can not get address of these variable

⇒ for string variable we use <sup>ptrs</sup> for i/p & gets for o/p

⇒ Macros :-

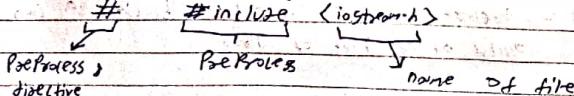
- it is a fragment of code that is given a name. you can
- you can define a Macro using `#define`

Ex:- `#define C 29987 ;`  
here replace `C` with `29987` whenever we use `C`.

⇒ PreProcessor :-

- it is responsible for transforms your program before it is compiled.

- these transform may can be inclusive of header file, macro expansion etc
- all PreProcessor directives begin with `#` symbol

  
PreProcessor directive

• we can also define function Macros

Ex:- `#define CircArea(r) (2*3.14*(r)*(r))`

• ## operator (Token-Pasting operators)

Ex:- `#include <stdio.h>`  
`#define concate(a,b) a##b`  
`int main()`  
`{ int xy = 30;`  
`printf ("%d", concate(x,y)); // replace with printf("%d,%d");`  
`return 0`

3

⇒ dynamic memory allocation (DMA) :-

- use to save memory
- or to utilised our memory

- when we don't know the amount of memory we need then we use DMA

• static memory allocation done on stack  
• dynamic memory allocation done on stack as well as heap

function used in DMA :-

- i) malloc
- ii) calloc
- iii) realloc
- iv) free

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- i) `malloc` function :-
- used for memory allocation
  - reserve memory space of specific size & return pointer which points to memory location
- Syntax  $\Rightarrow$  `Ptr = (datatype*) malloc (byte_size);`

- ii) `calloc` function :-
- used for contiguous allocation
  - used to allocate multiple block of memory
  - Syntax same as above only calloc in place of malloc

- iii) `realloc` function :-
- we can add more memory to size of already allocated memory
  - increase block size & reusing original content as it is
  - some syntax replace malloc with realloc

- iv) `free` function :-
- in DMA you have to deallocate the memory explicitly if not, you may encounter out of memory error
  - Syntax `free (Pointername);`

$\Rightarrow$  type qualifiers in C :-

Type qualifiers are the keyword used to modify the properties of variables. There are 2 type qualifiers :-

- constant
- volatile

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- Constant :- When we define ~~as~~ a volatile constant then whenever we assign any value to it will not change
- use "const" keyword
- Volatile :- volatile is used to create a variable whose values can't be changed in the program explicitly but can change by explicitly by hardware or external device
- for example variable used to store system clock

$\Rightarrow$  Abstract Data type :- it is a special kind of datatype whose behavior is defined by a set of values & set of operations

Some examples of ADT are stack, queue, list (operations)

- $\Rightarrow$  exception :- exceptions are those which can be handled at run time whereas errors cannot.
- Example ~~as~~ exception when we divide only thing by zero
  - error when there is syntax error

$\Rightarrow$  file handling :- for store information & fetch information we use file handling

### • Mode of file-handling :-

- i) r - read      v) x - Create new file, operation fail if file already exist
- ii) w - write     vi) t - open file in text mode
- iii) a - append    vii) b - open file in binary mode
- iv) 'r+' - ~~search file, fail open successfully read & write~~

$\Rightarrow$  threading :- when multiple tasks are performed parallel then it is called threading

$\Rightarrow$  Multithreading in Python :-

- multithread Process share the same data space with the Main thread. can share info & communicate with each other easily.
- there is an instruction Point which keep track of currently running thread

• for already running thread there are 2 state- interrupted or hold

### • Process to start new thread :-

- i) include thread module.
  - ii) cast like `My Thread.start_new_thread(function, args)`
  - iii) define a function for thread implement `Thread class`
  - iv) can create thread using below method
- ~~thread.start\_new\_thread(functionname, [list of arguments])~~

Date \_\_\_\_\_

```

ex:- import thread           // old mode for thread
      import time             // new mode is threading

def Print_time(threadName, delay):
    count = 0
    while count < 5:
        time.sleep(delay)
        count += 1
        print "%s %s" % (threadName, str(time.time()))

try:
    thread.start_new_thread(Print_time, ("Thread-1", 2))
    thread.start_new_thread(Print_time, ("Thread-2", 4))
except:
    print "Error: unable to start thread"
while 1:
    pass

```

### • Method under new thread module in Python :-

- i) `thread.activeCount()` :- no. of active threads
- ii) `thread.currentThread()` :- no. of thread object in callers thread
- iii) `thread.enumerate()` :- return list of active threads

### • Thread class implement following method :-

- i) `run()` :- entry point of thread
- ii) `start()` :- Start a method by calling `run()`
- iii) `isAlive()` :- whether a given thread is still execute or not
- iv) `getName()` :- return name of thread
- v) `setName()` :- set name of thread

\*

- Process to create a thread :-

- i) import threading module

- ii) define a class inheritance of threading.Thread class

- iii) overwrite the \_\_init\_\_ function by initialising the variable

- iv) define Method like run(), start() etc.

- v) define function to perform operation

- vi) declare object of Thread using proper argument

- vii) call start() Method of that class

- viii) thread is exit according to condition

⇒ Synchronizing thread :- • it is the concurrent execution of 2 or more thread that share critical resource

- to avoid the conflict of critical resource by 2 thread then it is necessary that thread must be synchronised

- In Python a ~~thread~~ class Lock is used for locking mechanism

blocking

$\downarrow \rightarrow$  acquire()

$\uparrow \rightarrow$  release()

- the acquire(blocking) method of used for locking a resource by set value of blocking to 1 then it wait for lock to be released by Put its value to 0 by call release() method

# Operating System

Date \_\_\_\_\_

→ Process, thread, inter-Process communication, concurrency, synchronization, Deadlock, CPU scheduling, memory management, virtual memory & file system.

⇒ CPU scheduling :- • it is a process which allow one process to use the CPU while the execution of another process is on hold (in waiting state) due to unavailability of resource - like I/O device

• Term used in scheduling:-

- i) Arrival time :- time at which process arrives in the ready queue
- ii) Completion Time :- time at which process complete its execution
- iii) Burst Time :- Time required by process for CPU execution
- iv) Turnaround Time :- Time difference bet. completion time & arrival time.
- v) Wait time :- Time difference bet. Turnaround Time & burst time

• Type of CPU Scheduling :-

- i) first come first serve :- • simplest scheduling algo in which when a process request the CPU first is allocated the CPU first

- implementation is Perform on queue
  - ii) Shortest job first (SJF) :-
    - Process which have shortest burst time are Schedule first
    - non-preemptive Process means if a Job have same burst time we can not interrupt before then follow first come first serve completion
  - iii) Longest Job first (LJF) :-
    - same as short Job first
    - Non-preemptive in nature
  - iv) Round Robin Scheduling :-
    - Each Process is assigned a fixed time known as Quantum Time in cyclic way
    - Ready queue treated as circular queue
    - If certain CPU burst time is longer than Quantum the complete remain Process in next queue
  - v) Priority scheduling :-
    - Process are scheduling according to their Priorities i.e. highest Priority
    - If 2 Priorities is same then Schedule according to arrival time.
  - ⇒ Deadlock :-
    - Deadlock is a situation where a set of Processes are blocked because each Process is holding a resource & waiting for another resource acquired by some other Process
- method of handling deadlock :-
- i) Deadlock Prevention or avoidance :-
    - we can monitor each process when it is seen that a deadlock will occur we do not complete that process until the required resource is free
  - ii) Deadlock detection & recovery :-
    - Let deadlock occurs after the solution makes to remove deadlock
    - iii) ignore the Problem :-
      - if deadlock is very rare then let it happens & reboot the system
      - this approach followed by windows & Linux
- Process :- Process is instance of a computer program that is being executed by one or many threads

$650 \Rightarrow 100$  hours (3 hours)

Aptitude 50 of 1 marks & 50 of 2 marks  $\Rightarrow 100 \Rightarrow 15M$   
Engg. 250 of 1 marks & 300 of 2 marks  $\Rightarrow 550 \Rightarrow 85M$

Engg. Maths  $\Rightarrow$  75 marks (100) Subject  $\Rightarrow$  70 marks (450)

50 of 1 mark & 50 of 12 mark 250 of 2 marks & 20 of 1 marks

~~i) Engg. Maths  $\Rightarrow$  matrices, system of linear eqns, determinants, linear algebra, decomposition, eigenvalue & eigenvectors, relations, sets~~

~~ii) Computer organization & architecture :-~~

Machine instruction & addressing modes, PLV, Jumper-Patch & control

Unit :- Instruction pipelining, memory hierarchy, cache, main memory & secondary storage, I/O, interface (interrupt & DMA mode)

~~iii) Programming & data structures :-~~

Programming in C, recursion, arrays, stacks, queues, linked lists, trees, binary search trees, binary heap & graphs.

~~iv) Algorithms :-~~

Searching, ~~Sorting~~, Hashing, ~~Asymptotic worst case time~~ & space complexity.

~~v) Algorithm design technology :-~~ Greedy, Dynamic Programming & divide & conquer. Graph search, minimum spanning trees

~~vi) Shortest paths.~~

✓ vi)

Theory of computation :- ~~regular expression & finite~~  $\Rightarrow$

automata, context free grammar & pushdown automata.

Requires context free language, pumping lemma, Turing

$\rightarrow$

Machines & undecidability

vii) Computer design :-

Lexical analysis  $\rightarrow$  Parsing, Syntax-directed translation

Run time environment, intermediate code generation,

X viii) Operation system :- Processes, threads, inter-process communication

$\rightarrow$  concurrency  $\rightarrow$  synchronization, deadlock, CPU scheduling

Memory management, virtual memory, file systems

✓ viii) DBMS - ER model, Relational model: relational algebra

tuple calculus, SQL, integrity constraints, normal forms,

file organisation, indexing (e.g. B&BTrees), ~~test~~

Transaction  $\rightarrow$  consistency

ix) Digital logic

x x) Computer Networks

✓ xi) Optimise :- Numerical computation, numerical estimation, numerical

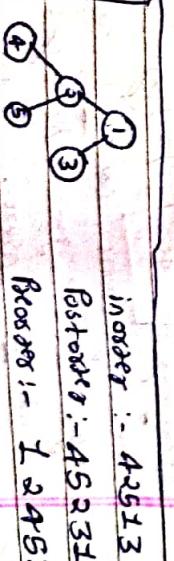
processing  $\rightarrow$  data interpretation

depth

i) 8 - Previous year paper

ii) 7 - full test

iii) 3 - advance test :-



inorder : 4 2 5 1 3

Postorder : 4 5 2 3 1

Preorder : 1 2 4 5 3

+ 8 -

8\*

82

## Programming 28 Data structure:-

Pointers :-

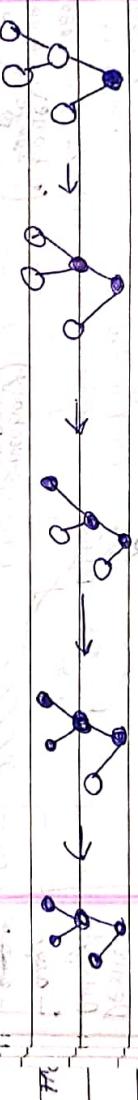
$A = [1, 2, 3, 4, 5, 6]$

\* $P = A$  (Point to  $A$ )

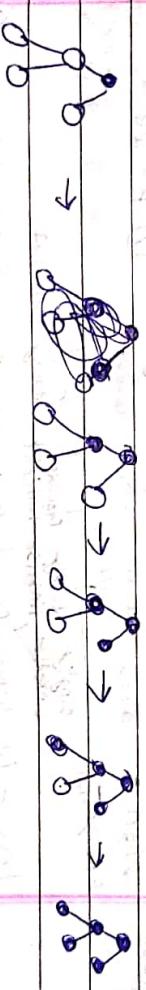
\* $P = P + 4$  (Point to  $B$ )

\* $P[2]$  (Point to  $C$ )

Depth first search is implemented by stack



Breadth first search is implemented by queue.



(operator left to right) → (operator right to left)

Infix, Postfix, Prefix

$a+b*c+d$  →  $ABC * + D +$

(a+b) \* c + d →  $A(B + C) * D$  (operator left to right)

no 2 operators at same priority can not stay together in stack

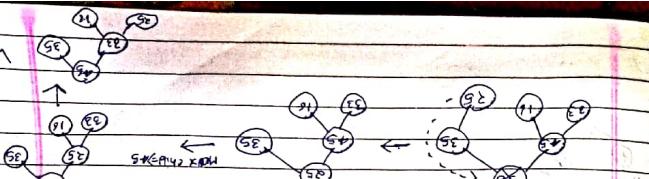
same priority

use same priority →  $a+b*c+d$  (operator left to right)

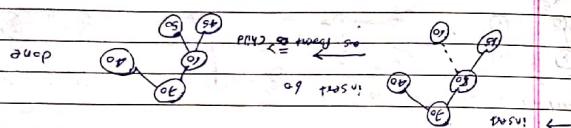
in order : left, root, right (bst)

post -> : left, right, root

pre -> : root, left, right

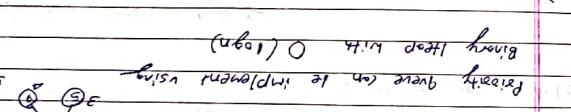


insert & search =  $O(n)$  (the worst)

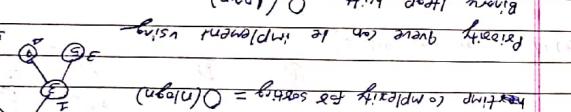


left part having lesser value than ~~the~~ parent node  
having higher value than ~~the~~ parent node

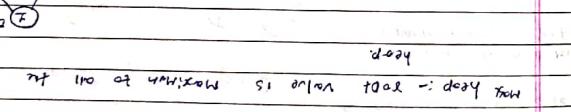
insert (only delete root node)



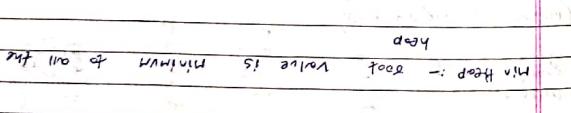
tree - 10-layer (not have other manners i.e. hierarchical view)



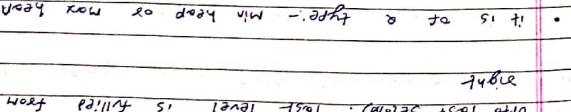
logistics: parallel distributed database system



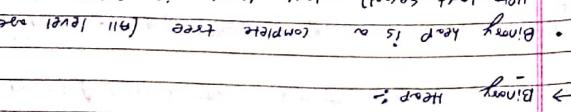
single point failure doesn't affect others



binary search tree i.e. hierarchical view



parallel processing (not all nodes are same)



multiple parallel processes



multiple parallel processes



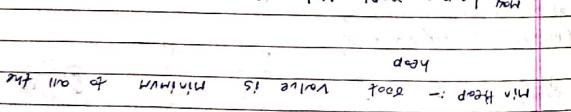
multiple parallel processes



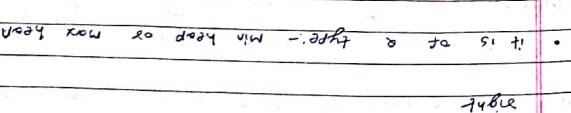
multiple parallel processes



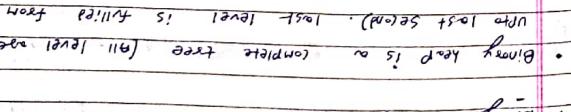
multiple parallel processes



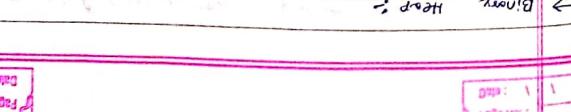
multiple parallel processes



multiple parallel processes



multiple parallel processes



multiple parallel processes



multiple parallel processes



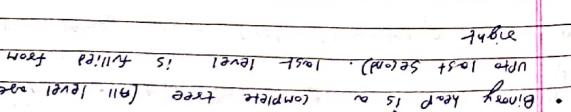
multiple parallel processes



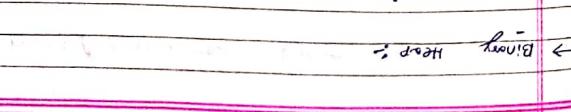
multiple parallel processes



multiple parallel processes



multiple parallel processes



multiple parallel processes



multiple parallel processes



multiple parallel processes



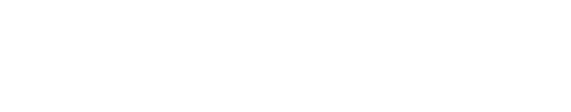
multiple parallel processes



multiple parallel processes



multiple parallel processes



multiple parallel processes



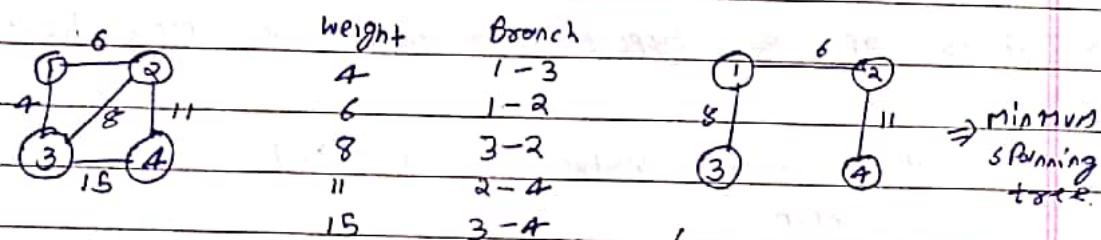
multiple parallel processes



$\Rightarrow$  graphs  $\rightarrow$  consist of (finite set of vertices & edges)

- Kruskal's  
~~Prim's~~ minimum Spanning Tree algo. :- (connected & undirected graphs)

Minimum Spanning Tree :- connected & undirected graphs with no cycle

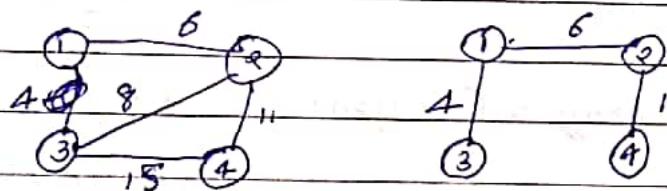


(least weight is selected)

(skip that vertex who create cycle)

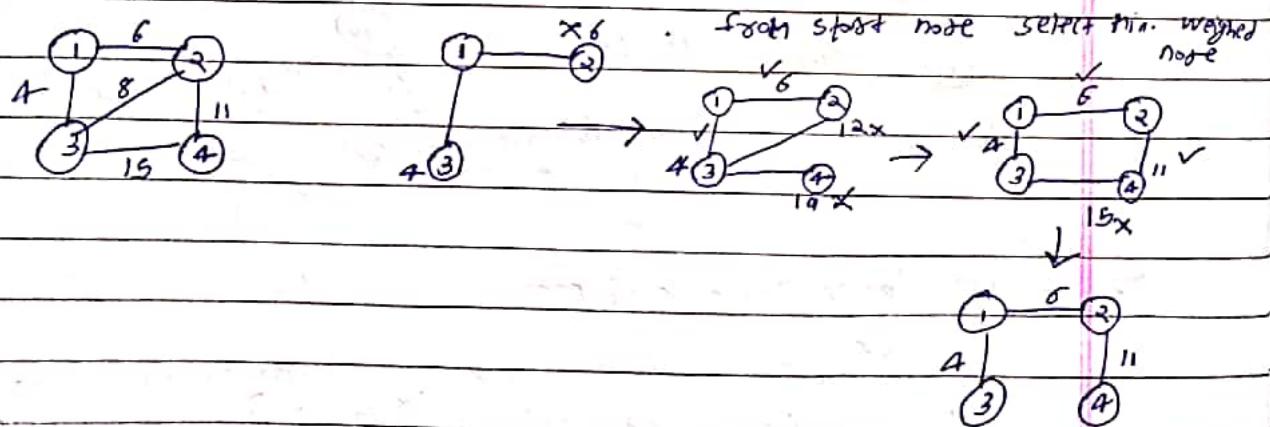
- Prim's minmum spanning tree algo. :-

Start with initial node minimum connect node is selected



(skip that vertex who create cycle)

$\Rightarrow$  Dijkstra's shortest path :-



→ algorithms :-

→ Search :- 2 type Linear search, binary search

- Linear search :-
  - Search Scan one item at a time without jumping to any items.
  - Worst case is  $O(n)$
  - Time increase as element increase =  $O(n)$

- Binary search :-
  - Search Middle element as forward repetitive.
  - Worst case complexity is  $O(\log n)$
  - Check element is less or greater
  - Check only half of list,
  - Time complexity =  $O(\log n)$

• Sorting :-

Type - Bubble sort, Selection sort, Merge sort, Insertion sort, Quick sort, Heap sort.

→ Asymptotic Notations :- 3 type. - theta notation ( $\Theta$ ), Big O notation ( $O$ ), Omega notation ( $\Omega$ )

• Theta notation Big O notation ( $O$ ) :-

- Express upper bound of algo.
- Measure worst case time complexity (Maximum time taken)

• Omega notation ( $\Omega$ ) :-

- Express lower bound of algo.
- Measure best case time complexity (Minimum time taken)

• Theta notation ( $\Theta$ ) :-

- Express lower as well as upper bound
- Measure average case time complexity (Average time taken)

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→ Algo Design techniques :- types:

- Divide & Conquer Approach
- Correct technique
- Dynamic Programming

• Divide & Conquer Approach :-

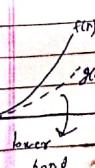
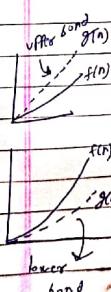
- Divide original Problem into a set of subproblems
- Solve each subproblem individually
- Combine the solution of subproblems recursively into a solution of whole original problem

• Greedy Technique :-

- Used to solve optimization Problem
- Optimization Problem is one in which we are given a set of maximized or minimized. i.e. some condition
- Greedy does not guarantees that solution is optimal but value find is very close to Optimal.

- Make a choice that best of that moment.

• Dynamic Programming :-





# Data Structure

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- ⇒ binary tree :-
- binary tree is part of non-linear data type, i.e. tree
  - in which every node or vertex has either 0 or 1 child or 2 children
  - two ways to represent it
    - i) using array
    - ii) using Linked lists
  - used to represent hierarchical relationship between elements

## Type of binary tree :-

- i) full binary tree (proper binary tree) :- ~~tree in which~~ binary tree in which except root node has exactly 2 children
- ii) complete binary tree :- tree in which all levels are filled & in last level filled left as possible (like heap)

⇒ BST :- Binary Search tree

# Syllabus e-litmus

3 section are there:- (20 each)

- i) Quantitative Ability (8-10) + (3-6) (6-8)
- ii) Logical Reasoning (10-12) (3-6) (6-10)
- iii) Verbal Ability (12-15) (6-10) (10-12)

## i) Quantitative Ability :-

- ✓ i) Geometry - (H)
- ✓ ii) Speed, Time and Distance - (H).
- iii) Mensuration (L)
- iv) Time And Work - (M)
- ✓ v) Number system - (H)
- vi) Probability - (M)
- vii) Permutation and Combinations - (M)
- viii) Few Miscellaneous Question on topics like Progress, Algebra
- ix) Averages →
- x) Co-ordinate Geometry → (Part of Geometry)
- xi) Logarithms →  $\log_b^a = c \quad a = b^c$
- xii) Quadratic eqn → geo.  $x^2 - (\text{sum of roots})x + (\text{Product of roots})$
- xiii) AP, GP, HP (M) →

## ii) Logical Reasoning :-

- ✓ i) Cryptarithmetic - (H)
- ii) Logical Reasoning (M)
- ✓ iii) Data Interpretation (Bar & Pie charts) (H)
- ✓ iv) Data sufficient (Algebra & arithmetic, Geometry, Logr & Misc) (H)
- v) Data Tabulation based questions
- vi) Arrangement Based Problems.

### iii) Verbal Ability

- ✓ i) Reading comprehension (H)
- ✓ ii) Para jumbles (M)
- iii) Sentence completion based on vocabulary (H)
- iv) Sentence completion based on subject-verb agreement (H)
- v) Miscellaneous question related grammar.
- vi) Grammatical Errors based questions
- vii) Paragraph Based Questions.
- viii) Fill in the blanks

### ⇒ Cryptarithmic conversion :-

i) Search for 0 or 9

$$A + B = A \text{ here } A = 0 \text{ or } 9 \text{ (only 2)}$$

ii) Additional 1

$$\text{SEND} + \text{MORE} = \text{MONEY} \text{ here } M=1 \text{ as answer is one more digit}$$

iii) Search for 1 in multiplication or division

$$\text{MAN} \times B = \text{MAN} \text{ here } B \text{ must be 1}$$

iv) Search for 6 or 1 in multiplication

$$A \times B = A \text{ here } B=1 \text{ or } B=6 \text{ (when A is even)}$$

v) search 0 or 5 in multiplication

$$A \times B = A \text{ here } A = 0 \text{ or } 5 (7 \times 5 = 35 \text{ (longer?)})$$

### ⇒ IIT JEE Previous year Paper :-

• Forget last 2 digit of  $7^{145}$

$$\text{Sol} \Rightarrow \text{last 2 digits of } 7^{14} = 07$$

$$7^{12} = 49$$

$$7^{13} = 43$$

$$7^{14} = 07$$

$$7^{15} = 07$$

so,  $7^{145} \Rightarrow 07$  as last digit

• factors common to  $30^1$  &  $20^3$

$$\text{Sol} \Rightarrow \text{HCF of both is } 10^1$$

$$\text{all the factors common to both} = 10^1 = 2^1 \times 5^1$$

$$\text{no. of common factors} = (1+1)(1+1) = 12$$

• a work done in  $\frac{2}{3}$  day with certain worker. If 7 of them are absent  
8 remain is done work in  $\frac{3}{5}$  day so no of employ work on

2nd day is

so, total work done if no of employ is  $x = \frac{2}{3}x$

$$(2-7)x = \frac{2}{3}x$$

$$x = 35$$

$$2^{\text{nd}} \text{ day} = x - 7 = 28$$

• when integer  $n$  is divisible by 8, remainder is 3 then what is reminder on  
is divisible by 8

$$n = 8k+3$$

$$8n = 8(8k+3)$$

$$= 8(6k)+8(2)+2$$

$8n$  is divisible by 8 so remainder is 2

- no of common terms in 2 series  $17, 21, 25, \dots, 412$  &  $16, 21, 26, \dots, 461$  is  
Sol)  $n^{\text{th}}$  term of 1st series =  $17 + (n-1)4 = 4n + 13$   
Min term of 2nd series  $\geq 16 + (n-1)5 = 5n + 11$   
 $4n + 13 = 5n + 11$   
 $4m = 4n + 2$   $[M = 2, 7, 12, \dots]$

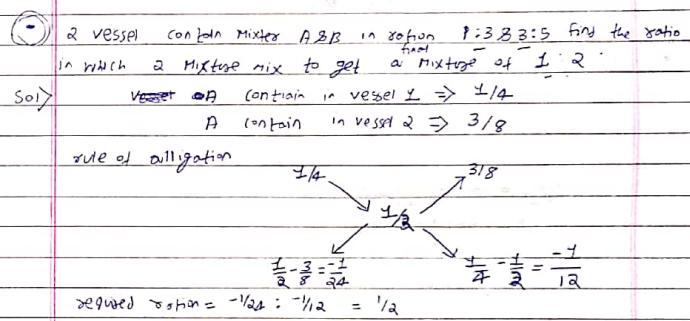
- How many integers are formed using digit 0, 1, 2, 3, 4 but no. is greater than 999 & not greater than 4000.  
Sol)  $\boxed{\quad} \quad \boxed{\quad} \quad \boxed{\quad} \quad \boxed{\quad}$  4000  
 $3 \times 5 \times 5 \times 5 = 375 + 1 = 376$

- What was the distinct term in expansion of  $(a+b+c)^{20}$ ?  
Sol)  $\binom{20+3-1}{3-1} = \frac{22}{2!} = 22! = 11 \times 21 = 231$

- Integers 1, 2, 3, ..., 40 written on board & 39 operation is done. In each repetition one no. is crossed & new no.  $a+b-1$  is written.  
Sol) sum of no.  $1+2+3+\dots+40 = 820$   
no. of operation  $\Rightarrow 39$   
last no. on board is  $= 820-39 = 781$

- Code generated by a digit from 0 to 9 but = 5 times digit except 0  
Let  $V[n+j] = 2V[n] + j$  ( $n=0, 1, 2, \dots$ ),  $V[0] = 0$  then  $V[10] = ?$   
Sol)  $V[0] = 0$ ,  $V[1] = 2 \times 0 + 1 = 1$ ,  $V[2] = 2 \times 1 + 1 = 3$   
 $V[3] = 2 \times 3 + 1 = 7$   $V[4] = 2 \times 7 + 1 = 15$   $V[5] = 2 \times 15 + 1 = 31$   
 $V[6] = 2 \times 31 + 1 = 63$   $V[7] = 2 \times 63 + 1 = 127$   $V[8] = 2 \times 127 + 1 = 255$   
 $V[9] = 2 \times 255 + 1 = 511$   $V[10] = 2 \times 511 + 1 = 1023$   
In  $V[0], V[3], \dots = 2^n - 1$   
 $V[10] = 2^{10} - 1 = 1024 - 1 = 1023$

- Product of 1 to 100 has no of zeroes at end  
Sol)  $\binom{100}{5} + \binom{100}{25} = 24$  as count no of 5's in 100!



- Vessel full of water & kerosene in which 18% is kerosene. Now 8 L of mixture is drawn off & certain vessel full with petrol now 15% of mixture is kerosene so volume of vessel is :-  
Sol) Ratio =  $18:82 \Rightarrow 9:41$   
mixture of kerosene =  $\frac{9x}{50}$   
final kerosene after drawn off  $\Rightarrow \frac{9x}{50} - \frac{9x}{50} = \frac{9x}{50}$   
 $= \frac{(9x-72)}{50}$

$$\frac{15x}{100} = \frac{9x-72}{50} \quad x = 48$$

- a) Mixture of Water & Soda in which Soda cost  $12/L$  is sold  
at a Profit of  $25\%$  with  $13.75/L$  of mixture. Then  
What is Ratio of Water to Soda

$$\text{Sol} \quad \text{Let } \frac{x}{L} \text{ is water} \quad \text{Soda Cost} = 12 \\ x/L - 1 = 12x$$

$$\text{Profit} = 25\% \text{ of } 12x = 3x$$

$$\text{Sell Price} = 12x + 3x = 15x$$

Let  $x$  L be of water in Mixture

$$\text{Selling Price} \Rightarrow 13.75(x+y) = 15x$$

$$x/y = 11/1$$

- The value of  $y$  in  $\log_{10} 13.75y = 3$

$$\log_{10} 13.75y = 3$$

$$\log_{10} 13.75 + \log_{10} y = 3$$

$$\log_{10} 13.75 + 1 = 3$$

$$\log_{10} 37^2 = 2$$

$$2 \log_{10} 37 = 2$$

$$10^2 = 1$$

$$10^2y = 1$$

$$\log_{10} 37 = \log_{10} y$$

$$y = 37$$

- b) In a race of 200m A beat B by 15m to C by 21m. By how much distance will B beat C in a race of 3700m

$\text{Sol} \quad \text{When B at } 185 \text{ m then C at } 179 \text{ m} = 6 \text{ m gap}$

$$Bx20 = 3700 \quad 6 \times 20 = \frac{3700}{6} = 616.67 \text{ m}$$

Distance is of 3700m

$$\text{then gap is } 3700 - 3580 = 120 \text{ m}$$

- c) Two train having lengths  $L_1$  &  $L_2$  having velocities  $v_1$  &  $v_2$  cross each other in  $t$  unit of time when moving in opposite direction to take  $T$  unit time when moving in same direction if ratio of  $T$  to  $t$  is  $11$  then value of  $v_1$  &  $v_2$  is

$$\text{Sol} \quad T = L_1 + L_2 / v_1 - v_2$$

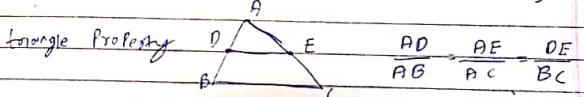
$$t = L_1 + L_2 / v_1 + v_2$$

$$\frac{T}{t} = \frac{v_1 + v_2}{v_1 - v_2} = 11 \quad \frac{v_1}{v_2} = 1.2$$

$$\text{Q) Compound interest} = P \left(1 + \frac{\alpha}{100}\right)^t$$

$$\text{Simple interest} = \frac{P \times Q \times t}{100}$$

$$\text{D) Cubical eqn} \Rightarrow \alpha x^3 + \beta x^2 + (\gamma + \delta) x + \epsilon = 0, \alpha + \beta + \gamma + \delta = -6/a, \gamma BDY = -d/a, \gamma BCYBDP = c/a$$



$$\frac{AD}{AB} = \frac{AE}{AC} = \frac{DE}{BC}$$



Q)  $|x-y| = 1.2$  &  $x^2 - y^2 = 1.6$  Find  $|x| + |y|$

$$\begin{aligned} \Rightarrow |x-y| &= 1.2 & (|x|+|y|)^2 &= x^2+y^2+2xy \\ (x-y)^2 &= (1.2)^2 & &= 1.6 + 0.16 \\ 1.44 &= x^2+y^2-2xy & &= 1.76 \\ 1.44 &= 1.6 - 2xy & |x|+|y| &= \sqrt{1.76} \\ 2xy &= 1.6 - 1.44 & &= 0.26 \\ &= 0.16 \end{aligned}$$

$\Rightarrow AP, GP, HP$

$$1+2+3+\dots+n = \frac{n(n+1)}{2}$$

$$1^2+2^2+3^2+\dots+n^2 = \frac{n(n+1)(2n+1)}{6}$$

$$1^3+2^3+3^3+\dots = \left(\frac{n(n+1)}{2}\right)^2$$

Q) Product of 2 co-prime no. is 153 then LCM is

Ans) HCF of 2 prime no. is 1

$$L.CM \times HCF = 153$$

$$L.CM \times 1 = 153$$

$$L.CM = 153$$

Q) Mehi can complete a work in 12 hrs. If he joined by Jethani who is 50% more efficient, in how many days both can finish the work together.

Mehi in 1 hr. =  $\frac{1}{12}$

$$Jethani in 1 hr = \frac{1}{12} + \frac{50 \times 1}{100} \times \frac{1}{12} = \frac{3}{16}$$

$$\text{Combine} = \frac{1}{12} + \frac{3}{16} = \frac{1}{10} \text{ so 10 days need to}$$

$\Rightarrow$  Trigonometry :-

P	B	P	H	H	B
H	H	B	P	B	P
Sin $\alpha$	Cos $\alpha$	Tan $\alpha$	Cosec $\alpha$	Sec $\alpha$	Cot $\alpha$
Angle	0	$30^\circ$	$45^\circ$	$60^\circ$	$90^\circ$
sin $\alpha$	0	$\frac{1}{2}$	$\frac{1}{\sqrt{2}}$	$\frac{\sqrt{3}}{2}$	1
cos $\alpha$	1	$\frac{\sqrt{3}}{2}$	$\frac{1}{\sqrt{2}}$	$\frac{1}{2}$	0
tan $\alpha$	0	$\frac{1}{\sqrt{3}}$	1	$\sqrt{3}$	$\infty$

$$\sin^2 \alpha + \cos^2 \alpha = 1$$

$$\tan^2 \alpha + \sec^2 \alpha = 1$$

$$\cot^2 \alpha + \operatorname{cosec}^2 \alpha = 1$$

$$\sin(\alpha + \beta) = \sin(\alpha)\cos(\beta) + \cos(\alpha)\sin(\beta)$$

$$\sin(\alpha - \beta) = \sin(\alpha)\cos(\beta) - \cos(\alpha)\sin(\beta)$$

$$\cos(\alpha + \beta) = \cos(\alpha)\cos(\beta) - \sin(\alpha)\sin(\beta)$$

$$\cos(\alpha - \beta) = \cos(\alpha)\cos(\beta) + \sin(\alpha)\sin(\beta)$$

$$\tan(\alpha + \beta) = \frac{\tan(\alpha) + \tan(\beta)}{1 - \tan(\alpha)\tan(\beta)}$$

$$\tan(\alpha - \beta) = \frac{\tan(\alpha) - \tan(\beta)}{1 + \tan(\alpha)\tan(\beta)}$$

If  $A, B, C$  are angles &  $a, b, c$  are sides of triangle then :-

$$\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C} \quad ; \quad C^2 = a^2 + b^2 - 2ab \cos C$$

$$a^2 = b^2 + c^2 - 2bc \cos A$$

$$b^2 = a^2 + c^2 - 2ac \cos B$$

Q) 2 men & 3 women can do a work in 10 days while 3 men & 2 women can do it in 8 days. In how many days can 2 men & 2 women complete the work?

Sol)  $x \Rightarrow$  efficiency of men  $y \Rightarrow$  efficiency of women

$$(2x + 3y) 10 = (3x + 2y) 8$$

$$2x = 7y$$

$$(2x + y) d = (2x + 3y) 10$$

$$(7y + y) d = (7y + 3y) 10$$

$$8yd = 100y$$

$$d = 100/8$$

⇒ logarithms :-

$$\log_b a = y \Rightarrow a = b^y$$

$$\log_b(x \cdot y) = \log_b x + \log_b y$$

$$\log_b(0) = \infty$$

$$\log_b(x/y) = \log_b x - \log_b y$$

$$\log_b(1) = 0$$

$$\log_b(xy) = y \log_b(x)$$

$$\log_b(b) = 1$$

$$\frac{\log_b(x)}{\log_b(y)} = \log_y(x)$$

$$\log_e(2) = 0.69$$

$$\log_e(6)$$

$$\log_e(3) = 1.09$$

.

$$\log_e(4) = 1.38$$

$$\log_e(5) = 1.6$$

$$\log_e(7) = 1.94$$