## University of Massachusetts Lowell — Comp 3010: Organization of Programming Languages Assignment 2

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Make sure that the remaining pages of this assignment do not contain any identifying information.

## 1 Small-Step Semantics

(25 points)

Let us define the calculator language below that performs arithmetic and boolean operations. It has three different syntactic classes: (1) arithmetic expressions a (2) boolean expressions b, and (3) final values v. The following questions require you define small-step semantics. That is, first you need to define the configuration. Then, you need to write inference rules that show one configuration small-steps to another configuration.

$$\begin{array}{llll} n & \in & \mathbb{Z} \\ a & ::= & n \mid a_1 + a_2 \mid a_1 \times a_2 \\ b & ::= & \textbf{true} \mid \textbf{false} \mid a = a \mid a \neq a \\ & & \mid a \leq a \mid a > a \mid \neg b \mid b\&\&b \\ v & ::= & n \mid \textbf{true} \mid \textbf{false} \end{array}$$

(a) Write small-step semantics for the syntactic class of arithmetic expressions generated by a.

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(a) Small-step semantics for arithmetic expressions generated by athe rules for arithmetic expressions:

$$a ::= n \mid a_1 + a_2 \mid a_1 \times a_2$$

We define small-step reduction rules for arithmetic expressions:

1.Addition: If the left-hand side of the addition is not a value, reduce it:

$$\frac{a_1 \to a_1'}{a_1 + a_2 \to a_1' + a_2}$$

If the right-hand side is not a value:

$$\frac{a_2 \to a_2'}{n + a_2 \to n + a_2'}$$

When both sides are values:

$$n_1 + n_2 \rightarrow n_1 + n_2$$

2. Multiplication: If the left-hand side of the multiplication is not a value, reduce it:

$$\frac{a_1 \to a_1'}{a_1 \times a_2 \to a_1' \times a_2}$$

If the right-hand side is not a value:

$$\frac{a_2 \to a_2'}{n \times a_2 \to n \times a_2}$$

When both sides are values:

$$n_1 \times n_2 \rightarrow n_1 \times n_2$$

3. For all integer literals:

$$n \to n(where \ x \in Z)$$

(b) Write small-step semantics for the syntactic class of boolean expressions generated by *b*.

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Boolean expressions are generated by the following grammar:

$$b ::= true \mid false \mid a = a \mid a \neq a \mid a \leq a \mid a \rangle a \mid \neg b \mid b \wedge b$$

The small-step semantics for boolean expressions are defined as:

1. Equality check:

$$\begin{aligned} \frac{a_1 \rightarrow a_1'}{a_1 = a_2 \rightarrow a_1' = a_2} \\ \frac{a_2 \rightarrow a_2'}{n = a_2 \rightarrow n = a_2'} \\ n_1 = n_2 \rightarrow \textit{true} \quad \textit{if } n_1 = n_2 \\ n_1 = n_2 \rightarrow \textit{false} \quad \textit{if } n_1 \neq n_2 \end{aligned}$$

2. Inequality check:

$$\begin{aligned} \frac{a_1 \rightarrow a_1'}{a_1 \neq a_2 \rightarrow a_1' \neq a_2} \\ n_1 \neq n_2 \rightarrow \textit{true} \quad \textit{if } n_1 \neq n_2 \\ n_1 \neq n_2 \rightarrow \textit{false} \quad \textit{if } n_1 = n_2 \end{aligned}$$

3. Relational operators ( $\langle , \rangle$ ):

$$egin{aligned} rac{a_1 
ightarrow a_1'}{a_1 \leq a_2 
ightarrow a_1' \leq a_2} \ n_1 \leq n_2 
ightarrow extit{true} & extit{if } n_1 \leq n_2 \ n_1 \leq n_2 
ightarrow extit{false} & extit{if } n_1 
angle n_2 \end{aligned}$$

4. Negation:

$$\dfrac{b 
ightarrow b'}{\lnot b 
ightarrow \lnot b'}$$
 $\lnot true 
ightarrow false$ 
 $\lnot false 
ightarrow true$ 

5. Boolean conjunction (AND):

$$egin{aligned} & b_1 
ightarrow b_1' \ \hline & b_1 \wedge b_2 
ightarrow b_1' \wedge b_2 \ & ext{true} \wedge b_2 
ightarrow b_2 \ & ext{false} \wedge b_2 
ightarrow ext{false} \end{aligned}$$

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