QUICK SORT

```
algorithm quicksort(A, lo, hi) is
    if lo < hi then</pre>
        p := partition(A, lo, hi)
                                                                   O(n^2)
                                                 Worst case
        quicksort(A, lo, p - 1)
        quicksort(A, p + 1, hi)
                                                 performance
                                                 Best case
                                                                   O(n \log n) (simple partition)
algorithm partition(A, lo, hi) is
                                                 performance
                                                                   or O(n) (three-way partition
    pivot := A[hi]
                                                                   and equal keys)
    i := lo  // place for swapping
    for j := lo to hi - 1 do
                                                 Average case
                                                                   O(n \log n)
        if A[j] ≤ pivot then
                                                 performance
            swap A[i] with A[j]
            i := i + 1
    swap A[i] with A[hi]
    return i
```

INSERTION SORT

```
for i = 1 to length(A) - 1
                                                                   O(n^2) comparisons.
                                            Worst case
   x = A[i]
                                            performance
                                                                   swaps
    j = i - 1
    while j \ge 0 and A[j] > x
                                             Best case performance O(n) comparisons, O(1)
        A[j+1] = A[j]
                                                                   swaps
        j = j - 1
                                                                   O(n^2) comparisons.
                                            Average case
    end while
                                            performance
                                                                   swaps
    A[j+1] = x^{[3]}
 end for
```

HEAP SORT

```
void maxHeapify(long a[], int i, int size)
    int l=leftChild(a,i),r=rightChild(a,i),largest;
    if(l<=size-1 && a[l]>a[i])
        largest =1;
    else largest =i;
    if (r<=size-1 && a[r]>a[largest])
        largest=r;
    if(largest!=i)
        swap(&a[i],&a[largest]);
        maxHeapify(a,largest,size);
}
void buildMaxHeap(long a[],int size)
    int i;
    for (i = (size-1)/2; i >= 0; i--)
        maxHeapify(a,i,size);
}
void heapSort(long a[],int *size)
    buildMaxHeap(a,*size);
    int i;
    for(i=*size-1;i>=1;i--)
        {
             swap(&a[0],&a[i]);
             (*size) --;
            maxHeapify(a,0,*size);
        }
}
```

```
Worst case O(n\log n) performance
Best case performance \Omega(n), O(n\log n)

Average case O(n\log n) performance

Worst case space O(1) auxiliary complexity
```

MERGE SORT

```
i = 0;
\dot{j} = 0;
k = start;
while (i < n1 \&\& j < n2)
      if (Leftarray[i] <= Rightarray[j])</pre>
          arr[k] = Leftarray[i];
          i++;
     }
     else
          arr[k] = Rightarray[j];
          j++;
     k++;
}
while (j < n2)
{
         arr[k++] = Rightarray[j++];
}
while(i<n1)
{
            arr[k++] = Leftarray[i++];
```

```
Worst case performanceO(n \log n)Best case performanceO(n \log n) typical,<br/>O(n) natural variantAverage case performanceO(n \log n)Worst case spaceO(n) total, O(n)<br/>auxiliary
```

SELECTION SORT

```
Worst case O(n^2)
performance

Best case performance O(n^2)

Average case O(n^2)
performance
```

BUBBLE SORT

```
void bubbleSort(float arr[],int size)
{
    if(size==1)
        return;
    int i,j;
    for(i=0;i<size-1;i++)
        if(arr[i]>arr[i+1])
            swap(&arr[i],&arr[i+1]);
        bubbleSort(arr,size-1);
}
```

```
Worst case performance O(n^2)
Best case performance O(n)
Average case performance O(n^2)
Worst case space O(1)
complexity auxiliary
```