

Postlab Exp1

1. What is the easiest trick to win Tic Tac Toe?
- A simple trick to win Tic Tac Toe is put your 'X' in any corner, then whatever may be opponent's move put another 'X' at the last position of same row or column or to the diagonal.
- After the opponents following more place at 'X' at remaining of these empty boxes.
- This is a simple trick to win.

X		X
	O	
O		X

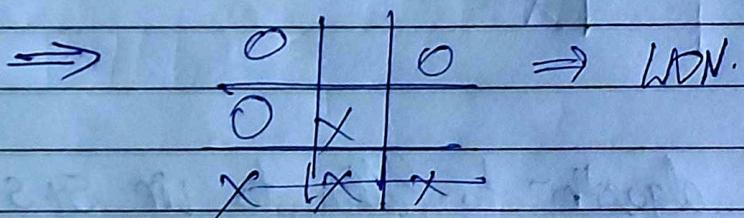
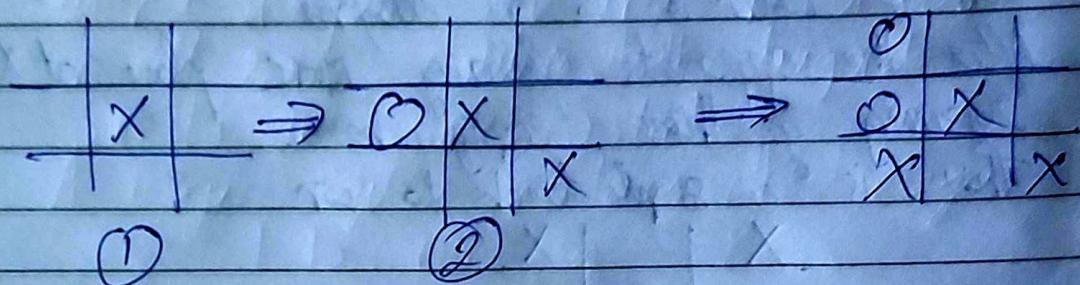
2. What is the algorithm to follow to win 5x5 Tic Tac Toe?

The Advanced Tic Tac Toe strategies of a 5x5 Tactical Approach is the algorithm one can follow to win 5x5 Tic Tac Toe.

Here player should,

- 1) Aim to occupy centre square
- 2) Place next 'X' in attempt to make winning lines
- 3) Block Opponent
- 4) Fork moves
- 5) Play defense
- 6) Adjust your strategies be flexible
- 7) Control the corner situation
- 8) Maintain Pressure play in attack mode.

Q3. Is there any way to never loose at Tic Tac Toe?
 To n St strategies which ensure that you never
 loose a game doesn't ensure you to win it every game.
 1) If you go first
 Place X on centre



Q4. What can Tic Tac Toe help you with?
 In learning and understanding concepts of AI
 Tic Tac Toe is perfect example to illustrate
 basic concepts of game theory.
 If we can understand min max algorithm
 and the decision making involved.
 We can understand & practice search
 algorithms like depth-first search, BFS,
 heuristic ; also brute force method.

Part b) Exp 2

Q1

What is the relationship between tic-tac-toe and magic square?

A the game of tic-tac-toe can also be implemented using the concept of magic square of an 3×3 magic square.

The sum of numbers of all rows, columns, diagonal is always the same = 15.

2	7	6
9	5	1
4	3	8

Hence player can always compute the score of the board to decide the next move to win.

C: Computer H: Human:

eg.	H		C
			C

$$\text{Diff} = 15 - (5+6)$$

$$= 15 - 11 = 4$$

The difference is not should not be greater than 9, hence computer can place C in 3rd cell if it is greater than nine with score 4 and win.

Q2.

		C
	C	C
n		H

$$\text{Diff} = 15 - 6 - 8$$

$$= 1$$

But it is not empty so computer

Ques 1.

Q2. What is magic square of order n?

A Magic Square of order n is an arrangement of n^2 numbers, usually distinct integers such that n numbers in all rows, columns, both diagonals sum to the same constant.

A magic square contains the integers from 1 to n^2 .

Postlab Exp 3

Q1. What is the time complexity of the Water Jug Problem?

1) Using BFS method.

Time Complexity : $O(n+m)$

Space Complexity : $O(n+m)$

where n and m are the quantity of jug 1 and jug 2.

2) Using DFS method:

As each state has six possible next states filling, emptying, pouring each jug so branching factor is 6.

The time complexity is exponential
 $O(3^d)$ where d is the depth of the search tree.

Postlab Exp 5.

Q1

Explain the time complexity of the A* Algorithm.

The time complexity of A* depends on the quality of the heuristic function.
In a worst case, algorithm can be $O(b^d)$ where b is the branching factor i.e. average number of edges from each node and d is the number of nodes on the resulting path.

Q2

What are the limitations of A* Algorithm?

A* Algorithm may consume significant memory and processing resources.

A* heavily relies on the quality of the heuristic function.

If heuristic function is poorly designed and does not account to estimate the distance to

to the goal the algorithm's performance and optimality may be compromised.

A* Algorithm might miss the best path if the guess estimate is overly optimistic.

A* Algorithm can be computationally expensive for very complex problems.

Q3. Discuss A*, BFS, DFS & Dijkstra's algorithm in detail with examples

A* Search Algorithm

A* Search combines the advantages of informed and uninformed search by using a heuristic function to guide the search towards the goal state efficiently.

It evaluates nodes based on a combination of the cost to reach them from the start and an estimate of the cost to reach the goal from them.

Breadth First Search

BFS explores all nodes at the current level before moving to the next level.

It guarantees finding the shortest solution path, but it may require a large amount of memory, especially for deeper

Search trees.

Depth first search

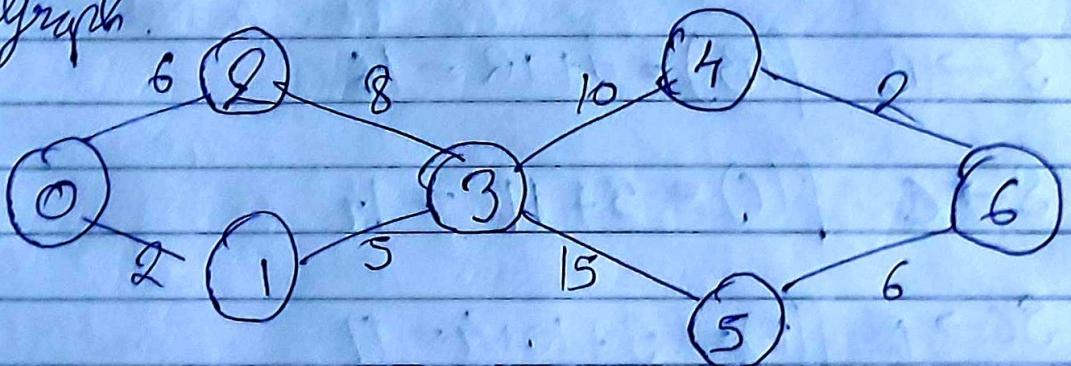
DFS explores as far as possible along each branch before backtracking.
It can be memory efficient, but may not always find the shortest solution, as it can get stuck in deep branches.

Dijkstra's Algorithm

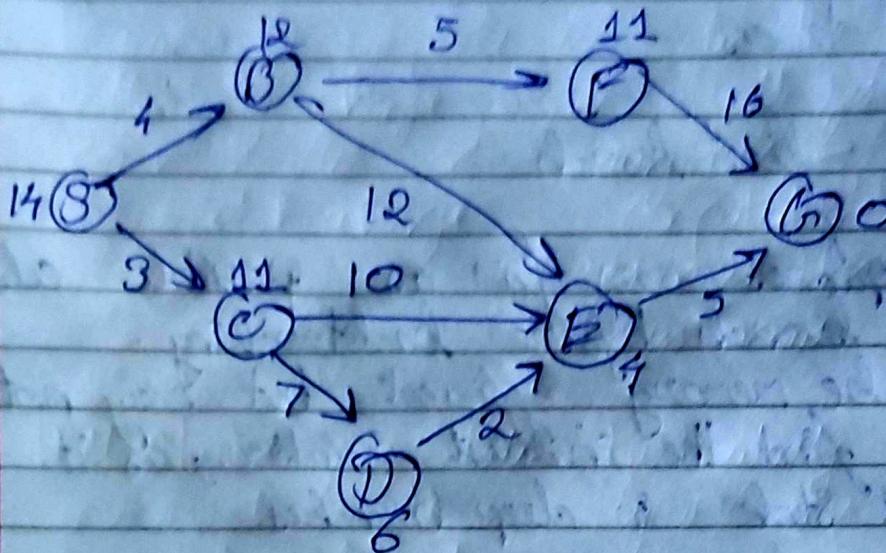
is an algorithm for finding the shortest paths between nodes in a weighted graph.
Dijkstra's original algorithm for finding the shortest paths between nodes in a weighted graph, producing a shortest-path tree.

Ex.

Graph.



Ex. A*



$$f(n) = g(n) + h(n).$$

$g(n)$ = actual cost from start node to n

$h(n)$ = estimated cost.

$$S \rightarrow S \quad f(S) = 0 + 14 =$$

$$S \rightarrow B \quad f(B) = 1 + 12 = 16$$

$$S \rightarrow C \quad f(C) = 3 + 11 = 14$$

$$S \rightarrow E \quad f(E) = 3 + 10 + 4 = 17$$

$$S \rightarrow D \quad f(D) = 3 + 7 + 6 = 16.$$

$$S \rightarrow D \rightarrow E \quad f(E) = 3 + 7 + 2 + 4 = 16$$

$$S \rightarrow D \rightarrow E \rightarrow G \quad f(G) = 3 + 7 + 2 + 5 = 17$$



BFS

S C B D E F G

DPS

S B F O G E E D.

BFS

S C B D E F G

DPS

S B F A G E C D

Postlab Exp 6.

- i). What is the difference between A* and AO*

A*

AO*

- | | |
|--|--|
| 1. Not designed for handling changes in the environment. | 1. Specifically designed to adapt to changes without initiating a new search. |
| 2. Primarily uses AND operation, considering one path at time. | 2. Uses OR and AND operations exploring multiple plans simultaneously. |
| 3. Explores fewer nodes. | 3. May explores more nodes due to adaptability. |
| 4. Less suited for high uncertainty or frequent environmental changes. | 4. Excels in situations with uncertainty, quickly adjusting response to new information. |

Requires a complete restart after an environmental charge
 restart of charge for a full restart
 leads to loss of computation resources when charges occur.

Well suited for static environment with consistent node costs
 Particularly in dynamic environment where cost is in.

Q2

Why AO* algorithm only works when heuristic values are underestimated?
 A* algorithm has a key property: is admissibility: an admissible heuristic is one that never overestimates the cost to reach a goal.

Overestimating means that the heuristic estimate is higher than the actual final path cost.

If we fail for inadmissible heuristic, the algorithm can wind up doing tons of spurious work exploring paths that it should be ignoring, and possibly finding suboptimal paths because of exploring them.

Postlab Exp 7.

Q1. What are the advantages and disadvantages of state space search?

Ans Advantages :

1. State space search systematically explores the possible states and transitions of a problem, ensuring a comprehensive examination of potential solutions.
2. It allows for the representation of complex problems in a structured manner, with states representing configurations and transitions depicting possible actions or moves.
3. State Space Search can be applied to a wide array of problems in robotics & game playing to natural language processing & scheduling.
4. It can adapt to various problem types including deterministic, stochastic, & adversarial scenarios, making it applicable to diverse set of challenges.
5. Informed decision making can be done using heuristic functions by acquiring domain specific knowledge.

Disadvantages :

Exponential growth : State space can grow exponentially with the size of the problem;
Storage and managing a large state space can require significant memory resources.

Search process can be time consuming
specifically state space is large
or if A* algorithm does not employ
efficient heuristics.

Cannot perform optimally in stochastic
or partially observable environments

Problems with a large number of
possible actions for each state
can lead to a high branching factor.

Q2. What are the advantages and
disadvantages of Hill Climbing approach?

Advantages:

1. The algorithm is straightforward to understand and implement.
2. It's memory efficient, maintaining only the current states data.
3. It often converges quickly to a solution, which is beneficial in scenarios where time is critical.

Disadvantages:

1. The algorithm can become stuck at locally optimal solutions that aren't the best overall.
2. Its tendency to focus on the immediate vicinity limits its exploration, potentially overlooking globally optimal solutions.

The quality and effectiveness of the solution found heavily depend on starting point.

Q3

Describe variations of Hill climbing approach?

Ans. a) Simple Hill climbing:

1. Evaluate: If initial state is a good goal state then stop & return success.
Else initial state = current state.
2. Do loop until solution is found:
or there are no new operators present which can be applied

• IF current state = goal state
 // stop
 return success

Else IF Score > current state
 current state = score
 Continue

Else :

 Continue.

b). Steepest Ascent Hill climbing:

IF initial state = goal state
 return success

Else

 in current state = initial state

Loop until solution state of find
no new operators present of which are
applicable

1. Stochastic state \rightarrow current state \Rightarrow New state

IF current state \neq goal state
return success

Else if $Score > Current$
best \cdot $Current\ state = Score$
Continue

Else
continue

exit.
current state = best state

exit

23. Stochastic hill climbing

Evaluate IF initial state == goal state
return success

Else
current state = initial state

Loop

unapplied | = state
better state = successor(state) // generate neighbors

IF better state == goal state
return success

Ex:

current state = better state
Continue

Exit.

Qn. Solve the Block world problem by using the STRIPS method.

Ay Let's consider 3 blocks : A, B, C
Introducing STRIPS operators.

Pickup(x), : picks black x from table

Putdown(x) puts black x on the table

Unstack(x,y) : picks up black x which is stacked on y

Stack(x,y) : stacks block x & y

Representation of
Pickup(x):

act: Pickup(x)

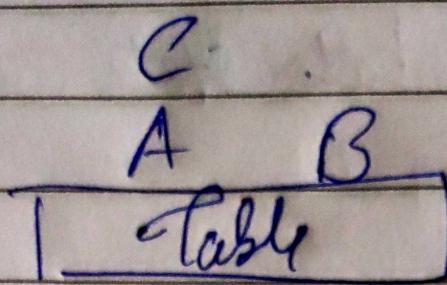
pre: OnTable(x), Clear(x), NotEmpty

add: Holding(x)

del: OnTable(x), Clear(x), NotEmpty.

initial State:

On(C,C), OnTable(A), OnTable(B), Clear(C), Clear(B)



On (A, B) A On (B, C)

