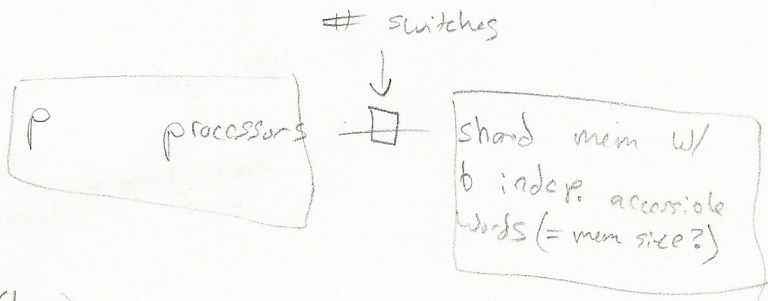


2.8)



$\Theta(bp)$ switches. cost prohibitive for a reasonable memory size. p. 32

UMA

shared-mem is physical org. that can enable logical org. of shared addr space (as can distributed mem). p. 29-30.

1MB w/ 1 cpu = $8 \cdot 1024 = 8192$ sw
 1GB w/ 1 cpu = $8 \cdot 1024^2 = 8.3 \times 10^6$ sw

2.11) $f(p) = \left(\begin{matrix} \vdots \\ \vdots \end{matrix} \right) \times p$

blocking = 2 comms ^{paths.} will access the same link

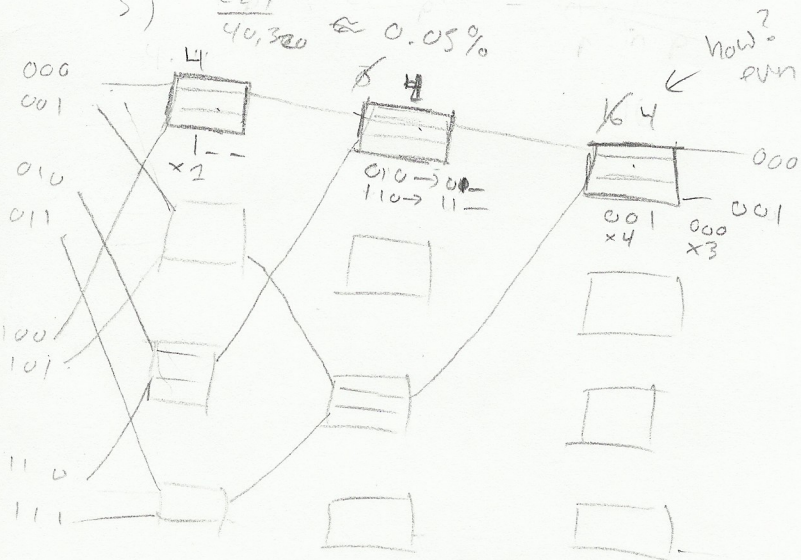
1) $p!$ $p=8 \rightarrow 40,320$

assuming perfect shuffle and $\log p$ switches.

2) $p \sum_{i=1}^{\log_2 p} (2^{(\log_2 p) - i} - 1) \cdot i$
 $p=8 \rightarrow 8 \cdot (4+8+16) = 224$
 $p! - (2^{\log_2 p}) 2^{p-1} = 40,320 - 16,384 = 23,936$
 $p=8 \rightarrow 4 \cdot 3 = 12$

010 \rightarrow 111
 110 \rightarrow 100

3) $\frac{224}{40,320} \approx 0.05\%$



now? not even 16 nodes... overlapping paths.

$\begin{matrix} \times \rightarrow 010 \rightarrow 7 \\ \times \rightarrow 000 \rightarrow 0 \\ 0 \{ 100 \rightarrow 2 \} = 4 \\ 1 \{ 010 \rightarrow 01 = 4 \\ \{ 110 \rightarrow 01 = 4 \} \\ 2 \{ 001 \rightarrow 001 = 4 \\ \{ 101 \rightarrow 001 = 4 \\ \{ 011 \rightarrow 001 = 4 \\ \{ 111 \rightarrow 001 = 4 \\ \times \rightarrow 000 \rightarrow 0 \end{matrix}$

$(p-1) + \left((2^{(\log_2 p) - 1}) \log_2 p \right)$ blocking